

CODE IN C (INPUT) :

```
print("varun nadar \n 1702")  
a=[4,21,29,64,25,24,3]  
j=0  
print(a)  
search=int(input("enter no:"))  
for i in range (len(a)):  
    if (search==a[i]):  
        print("num found at: ",i+1)  
        j=1  
        break  
if(j==0):  
    print("not found")
```

plcode

OUTPUT :

1.8. unsorted linear search :-
The data is entered in random manner.
user needs to specify the elements to be searched in the entered list.
check the condition that whether the entered number matches with the location of that element.
if all elements are checked one by one and element not found then prompt message number is not found.

python 3.7.4 (tags/v3.7.4:009359112e, Jul 8 2019, 20:34:20) [MSC v.1916 64 bit (AMD64)] on win32
Type "help()", "copyright()", "credits()", or "license()" for more information.
>>> RESTART: C:\Users\Nadar\OneDrive\Desktop\SEM2\DATA STRUCTURE\prac1.py ==>
varun nadar
1702
[4, 21, 29, 64, 25, 24, 3]
enter no: 21
num found at: 2

IS INPUT

```
print("varun nadar (n 1702)")
a=[3,4,21,24,25,29,64]
i=0
print(a)
search=int(input("enter no:"))
if (search-a[i]) or (search-a[6])):
    print("no. doesn't exist")
else:
    for i in range(len(a)):
        if search==a[i]:
            print("num found at:", i+1)
            break
    if i==0:
        print("num not found")
```

p2code

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Blattical - 2

Aim: To search a number as a item using linear search sorted method

Key: Linear search sorted method: searching and sorting the data in ascending order. To average the data in ascending order.

Key: To search elements and to display the location of element

In if linear search sorted method the elements are arranged in ascending or descending order. That is all what we meant by searching through a sorted list.

sorted linear search:

The user is supposed to enter the sorted manner

User has to given an element through the sorted

an element is found display
from updation of value
location

data
number
all element not
does not EXIST.

order list of element
the condition that
checked the number is
can check entered point the last
the starting not can way without any
if we can say number in

IS CODING INPUT

p3code

```
print("varun nadar \n 1702")
a=[3,4,21,24,25,29,64]
print(a)
search=input("enter no to be searched:")
l=0
h=len(a)-1
m=(l+h)/2
if ((search<=a[l])and(search>=a[h])):
    print("does not exist")
else:
    while(l!=h):
        if(search==a[m]):
            print("number found at:",m)
            break
        else:
            if (search<a[m]):
                h=m-1
                m=(l+h)/2
            else:
                l=m+1
                m=(l+h)/2
    if search==a[l]:
        print("number found at:",m)
    else:
        print("not found")
```


reduces to zero
Binary search works :-

Let the following array by sorted
assume that we need to
location of 51

low					high	
0	1	2	3	4	5	6
9	14	27	34	51	78	83

First, determine half of the array
formula

$$\text{mid} = \text{low} + (\text{high} - \text{low}) / 2$$

Here, it is

$$0 + (6 - 0) / 2 = 3$$

The mid = 3

Now, we compare the stored value
location

3, with value being searched
we change our low to mid + 1 and
the new mid value again

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C++ Input :

plcode

```
print("various number\n 1700")
a[4] = 20; a[5] = 24; a[6] = 28;
for p in range(len(a)-1):
    for q in range(p+1, len(a)):
        if(a[p] > a[q]):
            a[p], a[q] = a[q], a[p]
    print(a)
```

Practical u
TO sort given random data by
bubble sort method.

aim
using theory:
bubble sort method
means arranging a set of
in some order. These are different
used to sort the data in
all descending order.

method on element is compared
1st number. If it is found
greater than 1st element then
are interchanged.

the 1st element is compared with
element. If it is found to be greater
than they are interchanged in this way
all the elements are compared with the
next element and interchanged if required.

This is the first iteration and on
completing this iteration the last element
gets placed at the last position.

Similarly in the second iteration the
elements are made till the last but
one and this time the second element

elements gets placed at the
last position in list

As a result after del
the list becomes a sorted list

08/11/2020
coding Input :

p5code

```
print("varun nadar \n 1702")
class stack:
    global tos
    def __init__(self):
        self.l=[0,0,0,0,0,0,0,0]
        self.tos=-1
    def push(self,data):
        n=len(self.l)
        if (self.tos==n-1):
            print("stack is full")
        else:
            self.tos=self.tos+1
            self.l[self.tos]=data
    def pop(self):
        if(self.tos<0):
            print("empty stack")
        else:
            k=self.l[self.tos]
            print("data=",k)
            self.tos=self.tos-1
```

```
s=stack()
s.push(10)
s.push(20)
s.push(30)
s.push(40)
s.push(50)
s.push(60)
s.push(70)
s.push(80)
s.pop()
s.pop()
s.pop()
s.pop()
s.pop()
```

in either case is not the
but the interface is the user
allowed to pop out push
the away with few cells
creations.

CODE INPUT:

```
print("varun nadar\n 1702")
class Queue:
    global f
    def __init__(self):
        self.f = 0
        self.r = 0
        self.data = []
    def add(self, data):
        n = len(self.data)
        if self.f == n:
            self.r = self.f + 1
        else:
            print("queue is full")
    def remove(self):
        n = len(self.data)
        if self.f > 0:
            print(self.data[self.f])
            self.f = self.f + 1
        else:
            print("queue empty")
```

pycode

```
q = Queue()
q.add(30)
q.add(40)
q.add(50)
q.add(60)
q.add(70)
q.add(80)
q.remove()
q.remove()
q.remove()
q.remove()
```

Radical 6:
Aim: adding and deleting data in a queue

Meaning: A queue is a collection of entities that are maintained in a sequence and can use one end of the sequence by which the elements are added is called the back or rear and the end at which elements are removed is called the head or front of the queue.

The operation of adding an element to the rear of the queue is known as Enqueue.

The operation of removing an element from the front is known as Dequeue.

The operations of a queue queue is a first-in first-out (FIFO) structure in FIFO data structure first element added to the queue is the first one to be removed.

it queue as an example of a linear data structure, as it stores sequential collection.

COPYING OUTPUT

python 3.7.4 (tags/v3.7.4:09359112e, Jul 8 2019, 20:34:20) [MSC v.1916 64 bit (AMD64)] on win32
type "help", "copyright", "credits" or "license()" for more
>>> RESTART: C:\Users\Desktop\SEM2\DATA STRUCTURE\prac6.py ==
varun nadar
1702
queue is full
30
40
50
60
70
queue empty
>>>

p7code

```

print("varun nadar \n 1702")
class Queue:
    global r
    global f
    def __init__(self):
        self.r=0
        self.f=0
        self.l=[0,0,0,0,0,0]
    def add(self,data):
        n=len(self.l)
        if self.r<n-1:
            self.l[self.r]=data
            print("data added:",data)
            self.r=self.r+1
        else:
            s=self.r
            self.r=0
            if self.r<self.f:
                self.l[self.r]=data
                self.r=self.r+1
            else:
                self.r=s
                print("queue is full")
    def remove(self):
        n=len(self.l)
        if self.f<n-1:
            print("data removed:",self.l[self.f])
            self.f=self.f+1
        else:
            s=self.f
            self.f=0
            if self.f<self.r:
                print(self.l[self.f])
                self.f=self.f+1
            else:
                print("queue empty")
                self.f=s

q=Queue()
q.add(30)
q.add(40)
q.add(50)
q.add(60)
q.add(70)
q.add(80)
q.remove()

```

As queue data is only the data
between head and tail, hence
data left outside is not a
removed queue anymore, hence
the head and tail pointers will
be initialised to 0 every time
we reach the end of the queue.

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Printout 8 2019, 20:34:20) [MSC v.1916 64 bit
-09359112e, Jul 8 2019, 20:34:20) for more information.
"credits" or "license()" for more information.
C:\Desktop\NEHACL\SEM2\DATA STRUCTURE\pract.py

p8code

```
print("varun nadar \n 1702")
class node:
    global data
    global next
    def __init__(self, item):
        self.data = item
        self.next = None
class linkedlist:
    global s
    def __init__(self):
        self.s = None
    def addl(self, item):
        newnode = node(item)
        if self.s == None:
            self.s = newnode
        else:
            head = self.s
            while head.next != None:
                head = head.next
            head.next = newnode
    def addB(self, item):
        newnode = node(item)
        if self.s == None:
            self.s = newnode
        else:
            newnode.next = self.s
            self.s = newnode
    def display(self):
        head = self.s
        while head.next != None:
            print(head.data)
            head = head.next
        print(head.data)
start = linkedlist()
start.addl(50)
start.addl(60)
start.addl(70)
start.addl(80)
start.addB(40)
start.addB(30)
start.addB(20)
start.display()
```

1) Insertion
2) Deletion
3) Traversal

1) Insertion - Adds an element across the linked list.

2) Deletion - Deletes an element from the linked list.

3) Traversal - Traverses the linked list.

Representation of a linked list:

```

graph LR
    A[A] -- next --> B[B]
    B[B] -- next --> C[C]
    C[C] -- next --> null
  
```

data

ms 3.7.4 (tags/v3.7.4: e09359112e, Jul 8 2019, 20:34:20) [MSC v.1916 64 bit
 x64] on win32
 "help", "copyright", "credits" or "license()" for more information.
 !START: C:/Users/MOHINI/Desktop/NEHACLG/SEM2/DATA STRUCTURE/prac8.py ==
 nadar

Example 9
Evaluate an expression using
Postfix evaluation:

Postfix notation is used to represent an expression. The expression is written from left to right and the operators are placed after the operands. In postfix notation, no parentheses are required to convert infix to postfix.

For evaluation of postfix expression

a stack is used to store operands (or values) of the given expression and do the following:
If the scanned element is a number, push it onto the stack.

If the scanned element is an operator, pop the operands from the stack, perform the operation, and push the result back to the stack.

When the expression is ended, the value left in the stack is the result of the expression.

Example: Evaluate the postfix expression: 9 3 5 * + 9 - 1 1

'9' is a number, so push it to stack.

If it's a number, so push it to the stack. If it's an operator, pop two elements from the stack, apply the operation, and push the result back to the stack. If it's a closing parenthesis, pop until the opening parenthesis is found, then apply the operation and push the result. If it's an opening parenthesis, push it to the stack. If the stack is empty at the end, the expression is valid.

step	symbol	operation	stack	bliss
1	9	push	9	
2	3	push	9, 3	
3	5	push	9, 3, 5	
4	5	pop (3/5)	9, 3	34
5	5	push (15)	9, 15	

Python 3.7.4 (tags/v3.7.4:ef09359112e, Jul 8 2019, 20:34:20) [MSC v.1915 64 bit
 (AMD64)] on win32
 Type "help", "copyright", "credits" or "license()" for more information.
 >>> RESTART: C:/Users/DATA/STRUCTURE/postfix.py ==
 evaluated value is: 62
 varun nadar
 >>>

Practical to
implementing the use of

aim: quick sort
Theory: quick sort is divide and conquer algorithm. It picks an element as pivot and partitions the given array around the picked pivot.
There are many different versions of quick sort that pick pivot in different ways:

Always pick first element as pivot.
Always pick last element as pivot.
Pick a random element as pivot.
Pick median as pivot.

The key process in quicksort in partition() target of partition is, given an array and an element x of array as pivot put x at its current position in sorted array and put all smaller elements before it and put all greater elements after it.

```
print("varun nadar\n 1702")
def quickSort(alist):
    return quickSortHelper(alist, 0, len(alist)-1)
def quickSortHelper(alist, first, last):
    if first < last:
        splitpoint = partition(alist, first, last)
        quickSortHelper(alist, first, splitpoint-1)
        quickSortHelper(alist, splitpoint+1, last)
    return alist
def partition(alist, first, last):
    pivotvalue = alist[first]
    leftmark = first+1
    rightmark = last
    done = False
    while not done:
        while leftmark < rightmark and alist[leftmark] <= pivotvalue:
            leftmark = leftmark + 1
        while alist[rightmark] >= pivotvalue and rightmark >= leftmark:
            rightmark = rightmark - 1
        if rightmark < leftmark:
            done = True
    else:
        temp = alist[leftmark]
        alist[leftmark] = alist[rightmark]
        alist[rightmark] = temp
        temp = alist[first]
        alist[first] = alist[rightmark]
        alist[rightmark] = temp
    return rightmark
alist = [42, 54, 45, 67, 89, 66, 55, 80, 100]
quickSort(alist)
print(alist)
```

output

```
Python 3.8.1 (tags/v3.8.1:1b293b6, Dec 18 2019, 23:11:46) [MSC v.1916 64 bit
Type "help", "copyright", "credits" or "license()" for more information.
>>>
= RESTART: C:/Users/Raj/AppData/Local/Programs/Python/Python38/qu
varun nadar
1702
[42, 45, 54, 55, 66, 89, 67, 80, 100]
>>>
```

Analysis of quick sort :

Time taken by quick sort in general can be written as following

$$T(n) = T(k) + T(n-k-1) + O(n)$$

The first two term are for two recursive calls, the last term is for the partition process. k is the number of elements which are smaller than pivot.

The time taken by quick sort depends upon the input array and partition strategy.

Following are three cases :

Worst case: $T(n) = T(0) + T(n-1) + O(n)$
which is equivalent to

$$T(n) = O(n) + T(n-1)$$

The above recurrence is $O(n^2)$

Best case: $T(n) = 2T(n/2) + O(n)$

The recurrence is $O(n \log n)$

Average case: $T(n) = T(n/10) + T(9n/10) + O(n)$

Practical - II
Aim. Demonstrating the use of mergesort

Theory merge sort

It divides input array in two halves, calls itself for the two halves and then merges the two sorted halves.

The merge() function is used for merging two halves. The merge(arr, l, m, r) is used to merge the two sorted sub-arrays arr[l..m] and arr[m+1..r] into a single sorted sub-array.

Time complexity: Sorting array on different machines. Merge sort is a recursive algorithm and its time complexity can be expressed as follows:

Recurrence relation: $T(n) = 2T(n/2) + O(n)$
Application of merge sort: It is a sorting technique.

```
k+=1
```

```
while i<n1:
```

```
    arr[k]=L[i]
```

```
    i+=1
```

```
    k+=1
```

```
while j<n2:
```

```
    arr[k]=R[j]
```

```
    j+=1
```

```
    k+=1
```

```
def mergesort(arr,l,r):
```

```
    if l<r:
```

```
        m=int((l+(r-1))/2)
```

```
        mergesort(arr,l,m)
```

```
        mergesort(arr,m+1,r)
```

```
        sort(arr,l,m,r)
```

```
arr=[12,11,13,5,6,7,52,47,21]
```

```
print("Before Mergesort\n",arr)
```

```
n=len(arr)
```

```
mergesort(arr,0,n-1)
```

```
print("After Mergesort\n",arr)
```

Practical - 12
Implementing the use of
Binary tree

Binary tree:

A binary tree is a tree which each node has at most two children, which are defined as the left child and the right child.

A recursive definition using set theory notation is a binary tree is a tuple (L, S, R) , where L and R are binary trees and S is a single empty set.

Types of binary trees:

Full binary tree: All levels are filled except possibly the last one and all nodes are filled in as far left as possible.

Complete binary tree: All levels are filled except possibly the last one and all nodes are filled in as far left as possible.

```
print('Varun Nadar/n1702')
class Node:
    global l
    global data
    def __init__(self, l):
        self.l = None
        self.data = l
        self.r = None
```

```
class Tree:
    global root
    def __init__(self):
        self.root = None
    def add(self, val):
        if self.root == None:
            self.root = Node(val)
```

```
        else:
            newnode = Node(val)
            h = self.root
            while True:
                if newnode.data < h.data:
                    if h.l == None:
                        h.l = newnode
                        print(newnode.data, "added on left of", h.data)
                        break
```

```
                    else:
                        h = h.l
                else:
                    h.r = newnode
                    print(newnode.data, "added on right of", h.data)
                    break
```

```
    def preorder(self, start):
        if start != None:
            print(start.data)
            self.preorder(start.l)
            self.preorder(start.r)
```

```
    def inorder(self, start):
        if start != None:
            self.inorder(start.l)
            print(start.data)
            self.inorder(start.r)
```

```
    def postorder(self, start):
        if start != None:
            self.inorder(start.l)
            self.inorder(start.r)
            print(start.data)
```

```
T = Tree()
T.add(100)
T.add(80)
T.add(70)
T.add(85)
```

```
T.add(10)
T.add(78)
T.add(60)
T.add(88)
T.add(15)
T.add(12)
print("preorder")
T.preorder(T.root)
print("inorder")
T.inorder(T.root)
print("postorder")
T.postorder(T.root)
```



```

python 3.8.1 (tags/v3.8.1:1b293b6, Dec 18 2019, 23:11:46) [MSC v.1916 64 bit
(AMD64)] on win32
type "help", "copyright", "credits" or "license()" for more information.
>>>
===== RESTART: C:/Users/Raj/Desktop/vrn.py =====
varun Nadar/n1702
80 added on left of 100
70 added on left of 80
55 added on right of 80
10 added on left of 70
78 added on right of 70
60 added on right of 10
88 added on right of 85
15 added on left of 60
12 added on left of 15
preorder
100
80
70
10
60
15
12
78
85
88
inorder
10
12
15
60
70
78
80
85
88
100
postorder
10
12
15
60
70
78
80
85
88
100
>>>

```