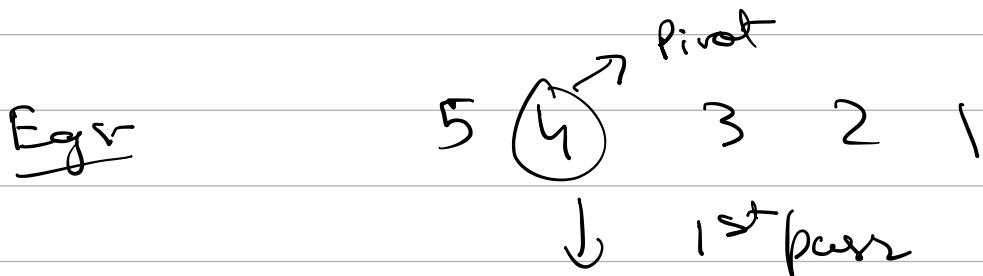
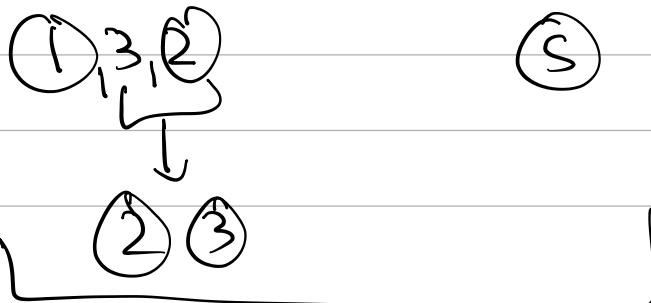
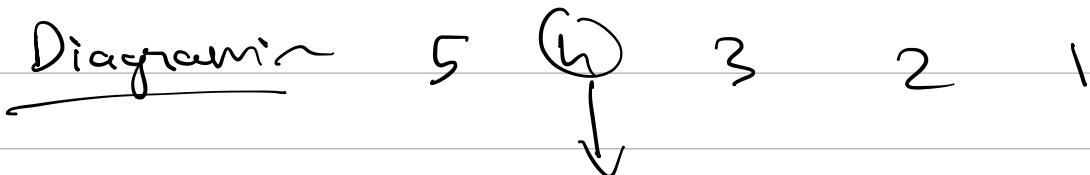


Quick Sort

*). PIVOT:- Choose any element \rightarrow after first pass all the elements $< P$ will be on LHS of P & elements $> P$ will be RHS of P.



Logic: In quick sort, after each pass, the pivot element will be put in the correct position.



Answer:-

1 2 3 4 5

—————
—————

How to put pivot in correct position ?

Logic:-

- ① Pivot should be $\geq \text{start}$. $\uparrow \text{start}++$
- ② Pivot should be $\leq \text{end}$. $\downarrow \text{end}--$
- ③ if any of these conditions are violated. swap start and end and also do $\text{start}++$ and $\text{end}--$.

8
5 \textcircled{h} 3 2 e

\downarrow $P ! > S$ (swap).
 $S++$ $e--$

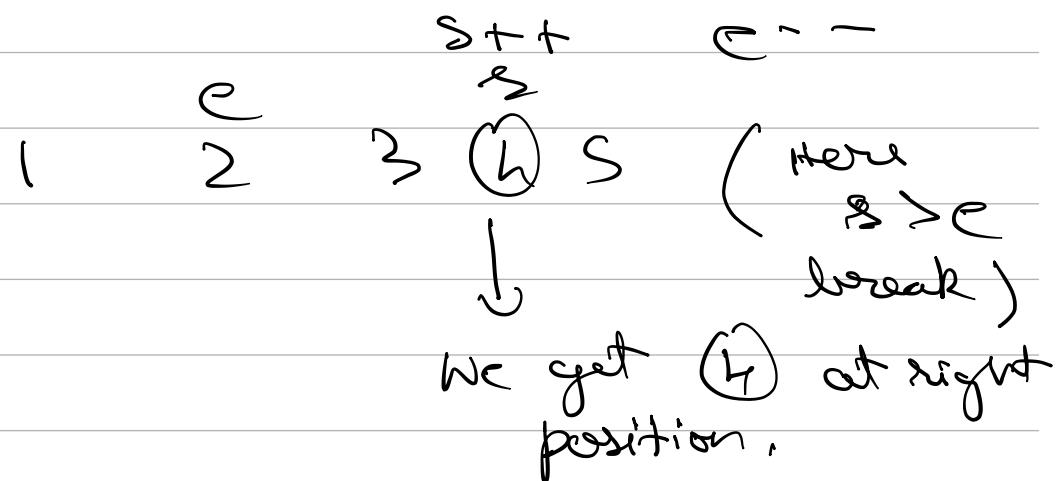
1 \textcircled{h} 3 2 5

\downarrow $P ! > S$ $e ! > P$
(swap)

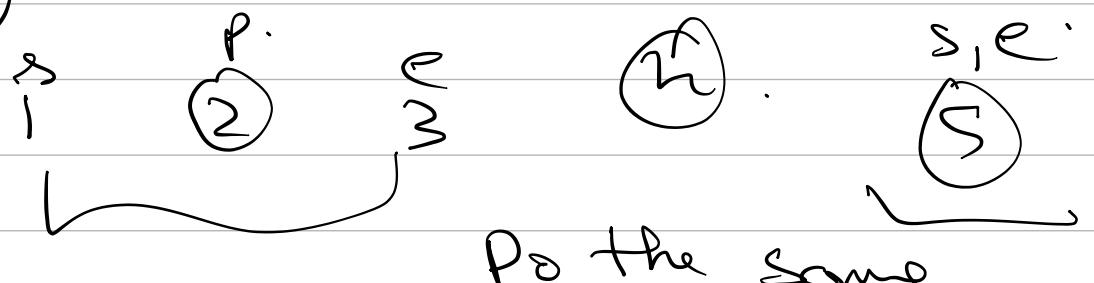
$S++$ $e--$
1 2 $\overset{(se)}{\textcircled{3}}$ \textcircled{h} 5

[

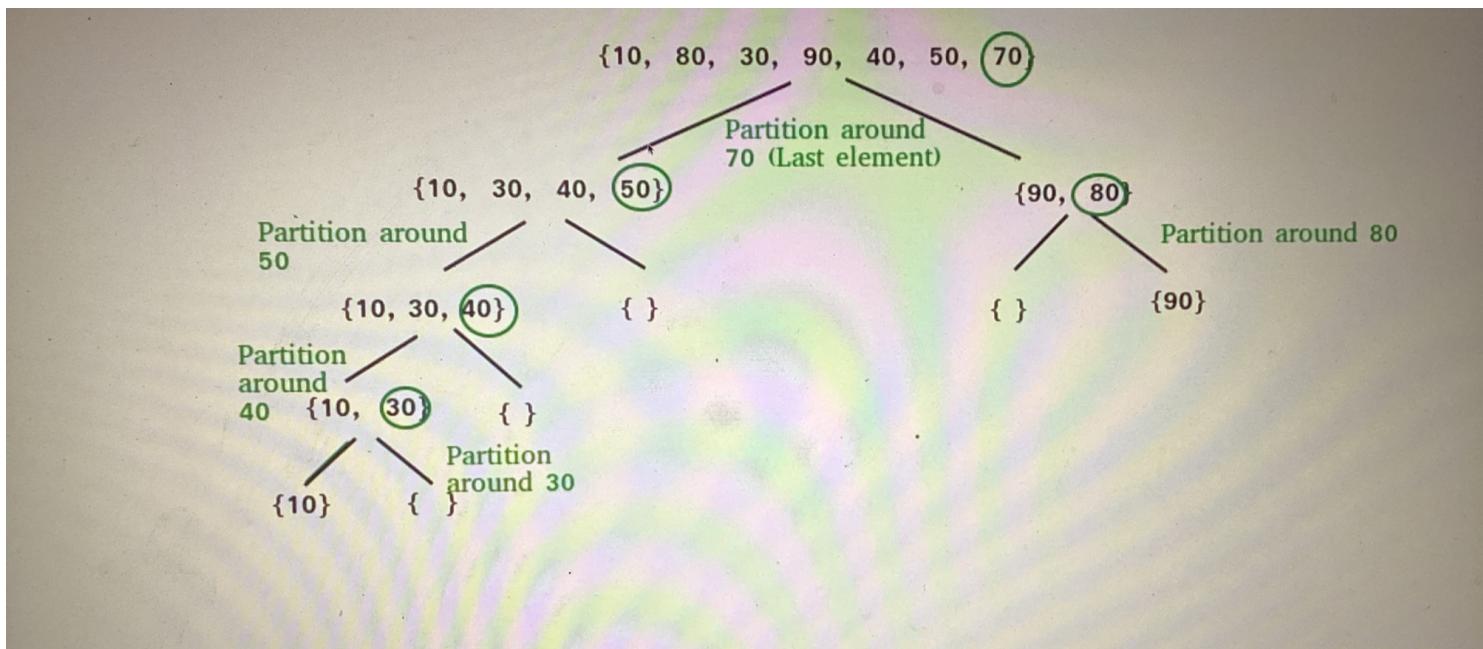
$p > s$ but $e! > p$
(swap)



Now;



ANOTHER EXAMPLE:-

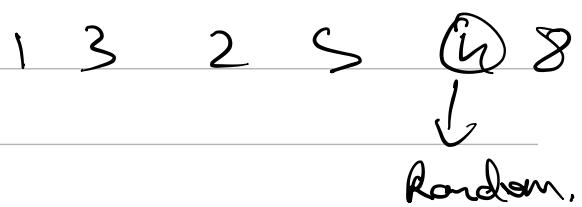


NOTE:-

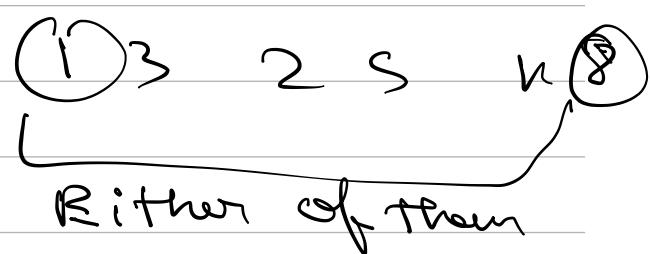
Better choose last element
or pivot

Different positions the pivot can be ?

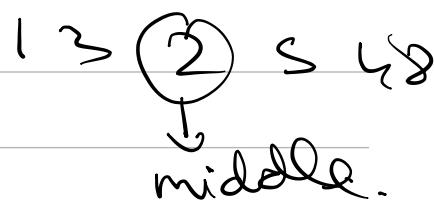
→ Pivot our position :-



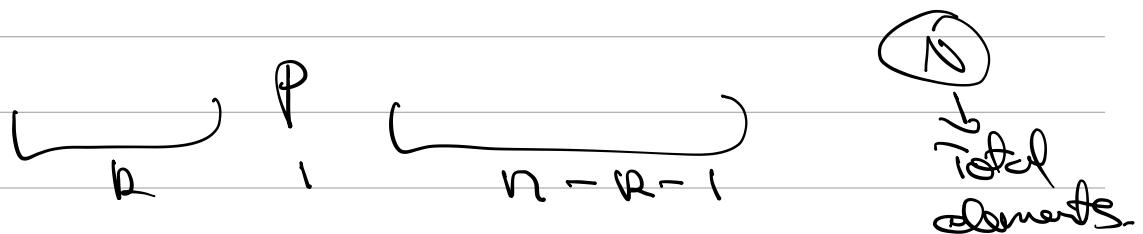
→ Corner element :-



→ Pick the middle element :-



Complexity comparison of different positions



$$T(N) = T(K) + T(N - K - 1) + O(N)$$

Time to sort K elements

Time to sort $N - K - 1$ elements

Time to put pivot at right pos.

↓

Recurrence relation of Quick sort-

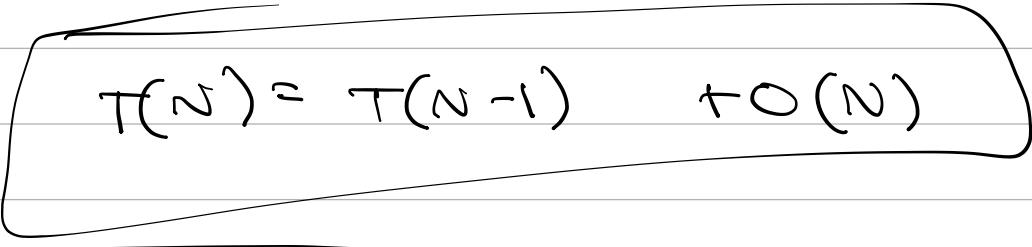
Worst case:-

→ When we pick highest or lowest element as pivot.

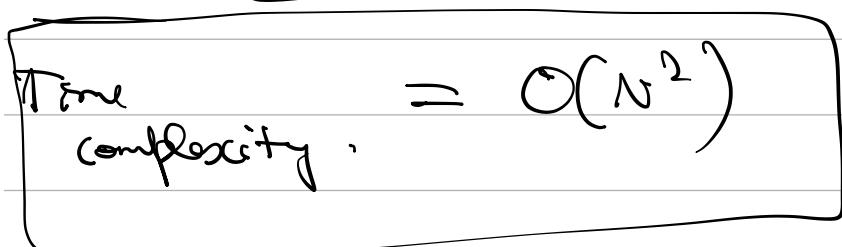
→ One side will be empty.



$$T(n) = T(0) + T(n-1) + O(n).$$



A horizontal bracket is positioned under the term $T(n-1)$ in the equation $T(n) = T(n-1) + O(n)$. An arrow points downwards from the center of this bracket to a rectangular box containing the equation $T(n) = T(n-1) + O(n)$.



A horizontal bracket is positioned under the term $O(n^2)$ in the equation $\text{Time complexity} = O(n^2)$. An arrow points downwards from the center of this bracket to a rectangular box containing the text "Time complexity" followed by an equals sign and $O(n^2)$.

→ check Time and space complexity for derivation.

Best case:-

$$R = \frac{N}{2} \quad \left[\text{we choose middle element} \right]$$

$$T(n) = T\left(\frac{n}{2}\right) + T\left(\frac{n}{2}\right) + O(n)$$

$$T(n) = 2 T\left(\frac{n}{2}\right) + O(n). \quad \Downarrow g(x)$$

Akhila
Beezji-

$$f(x) = x + x \int_1^x \frac{u}{u^{p+1}} du$$

To find P,

$$2 \times \frac{1}{2^P} = 1$$

$$\int P = 1$$

$$f(x) = x + x \int_1^x \frac{u}{u^2} du$$

$$= x + x \int_1^x \frac{1}{u} du$$

$$= x + x \int_1^x \log u du$$

$$= x + x (\log x - \log 1)$$

$$f(x) = x + x \log x$$

→ ignore x,

$$f(x) = x \log x$$

$$\text{Time complexity} = f(N) = O(N \log N)$$

NOTES:-

- *). Not stable (In case of equal elements their relative positions are not in same order as of before sorting)
- *). In-place algorithm (This is better than merge sort as merge sort take $O(N)$ auxiliary space but quicksort is in-place algorithm).
- *). Merge sort is better used in linked list , as unlike array , linked list does not have continuous memory allocation.

Hybrid sorting algorithms

We use combination of sorting algorithms to get better performance .

For example :- Merge sort + Insertion sort

This work well for

partially sorted data.

Code:-

```
public class QuickSort {  
    public static void main(String[] args) {  
        int[] arr = {5, 4, 3, 2, 1};  
        sort(arr, low: 0, hi: arr.length - 1);  
        System.out.println();  
    }  
}
```

```
static void sort(int[] nums, int low, int hi) {  
    if (low >= hi) {  
        return;  
    }  
  
    int s = low;  
    int e = hi;  
    int m = s + (e - s) / 2;  
    int pivot = nums[m];  
  
    while (s <= e) {  
        while (nums[s] < pivot) {  
            s++;  
        }  
        while (nums[e] > pivot) {  
            e--;  
        }  
    }  
}
```

```
if (s <= e) { → again checking inside  
    int temp = nums[s];  
    nums[s] = nums[e];  
    nums[e] = temp;  
    s++;  
    e--;  
}  
} → while loop closed.  
  
// now my pivot is at correct index, please sort two halves now  
sort(nums, low, e); → passes low to e  
sort(nums, s, hi); → s to hi
```

while loop
as s++ and
e-- is
also done
above.

3

low e s high.
1 2 ③ h s

After putting pivot (i.e 3) in
correct position, e-- and s++ is done, so.
passing low to e and s to high will

pass two subarrays.

Internal sorting Algorithms

- Each language uses its own sorting algorithm.
- We can call this sorting function.
- In Java, Arrays.sort(arr) function sorts 'arr' based on Dual pivot quick sort.
- Dual pivot is more efficient than single pivot we done above.
- Dual pivot internal implementation has hybrid algorithm (i.e combination of other algorithms like mergesort, etc).
- Each language uses its own algorithm which we can use by calling few functions.