

Today in Cryptography (5830)

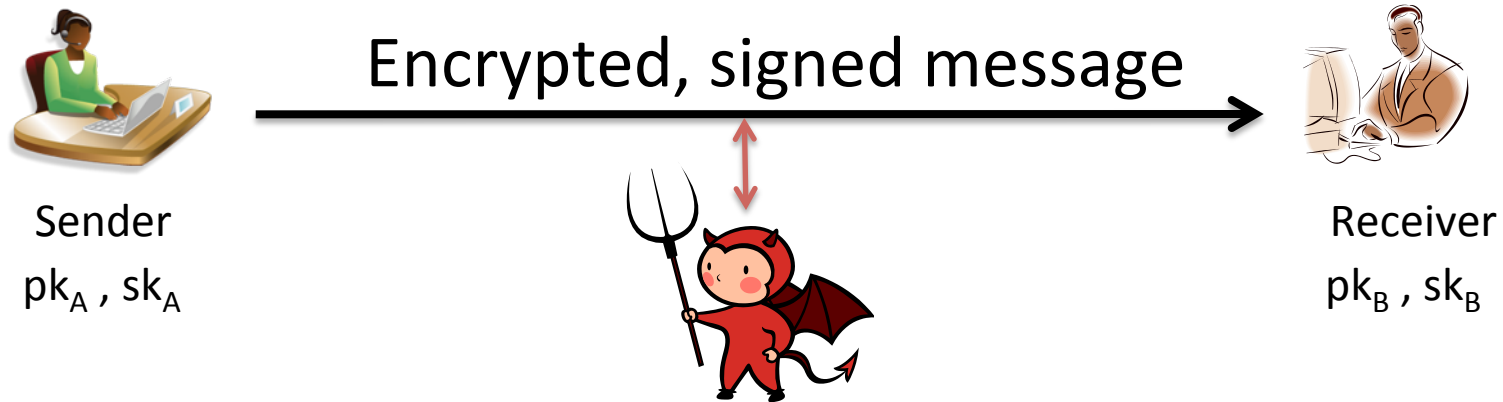
Hybrid encryption recap
OpenPGP standard
TextSecure

Katz-Lindell Chapter 10.3 (Hybrid Encryption)
RFC 4880 (OpenPGP standard)

Application-layer crypto

- So far focused on TLS as running example
 - Transport Layer Security
 - Provides network socket style stream interface
- What about if an application wants to encrypt discrete messages (as opposed to stream)?
 - Email
 - Text messages
 - Etc.

Email encryption

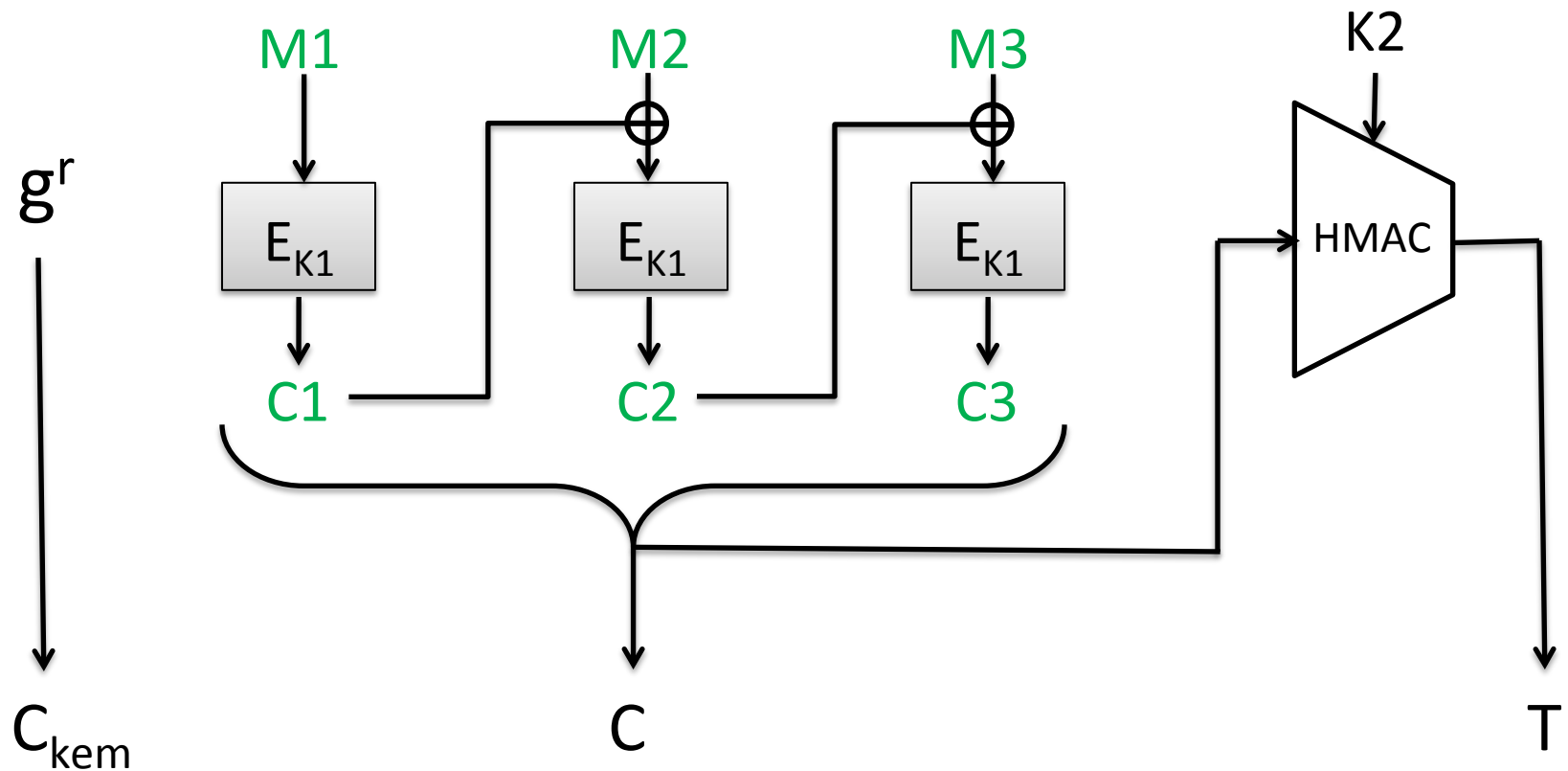


- Message may be large (body of email, PDF of attachments)
- Desire authenticity and confidentiality
- Public-keys delivered out-of-band
 - Websites, key parties, key directory servers

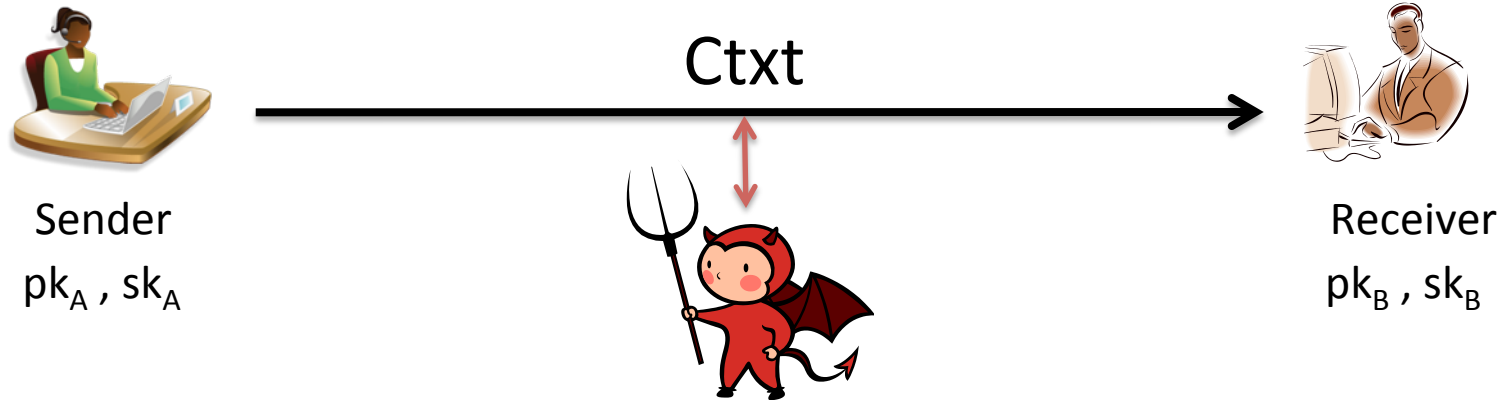
Example hybrid encryption

Enc(X,M):

$$K1 || K2 = \text{SHA256}(g^{xr})$$



Email encryption



- To digitally sign, let $M = \text{Msg} \parallel \text{Sign}(sk_A, \text{Msg})$
- $\text{Ctxt} = \text{Encrypt}(pk_B, M)$

PGP history

- Phil Zimmerman released “Pretty Good Privacy” in 1991 on a USENET post marked as “US only”
- 1993: Criminal investigation by US government for munitions export without a license.
 - Printed PGP source code into a book. First amendment gambit

OpenPGP overview

- Standard for PGP is RFC 4880
- Key encapsulation mechanism:
 - RSA PKCS#1 v1.5 encryption
 - ElGamal over finite field or elliptic curve
- Digital signatures:
 - RSA PKCS#1 v1.5 signatures
 - DSA
- Symmetric encryption:
 - Password-based key derivations using iterated hashing
 - CFB mode using block cipher

OpenPGP overview

- Security problems:
 - Padding oracle attacks against CFB & PKCS#1 v1.5
 - <http://www.ssi.gouv.fr/uploads/2015/05/format-Oracles-on-OpenPGP.pdf>
 - Attacks against home-brewed integrity checks (modification detection check, MDC)
 - Subject lines always in the clear
- Usability problems:
 - Users must manage their own keys
 - Copying private keys to each device
 - “Why Johnny Can’t Encrypt”
<http://www.gaudior.net/alma/johnny.pdf>
 - 90 minutes to send encrypted email, 66% failure rate



How Glenn Greenwald Began Communicating With NSA Whistleblower Edward Snowden

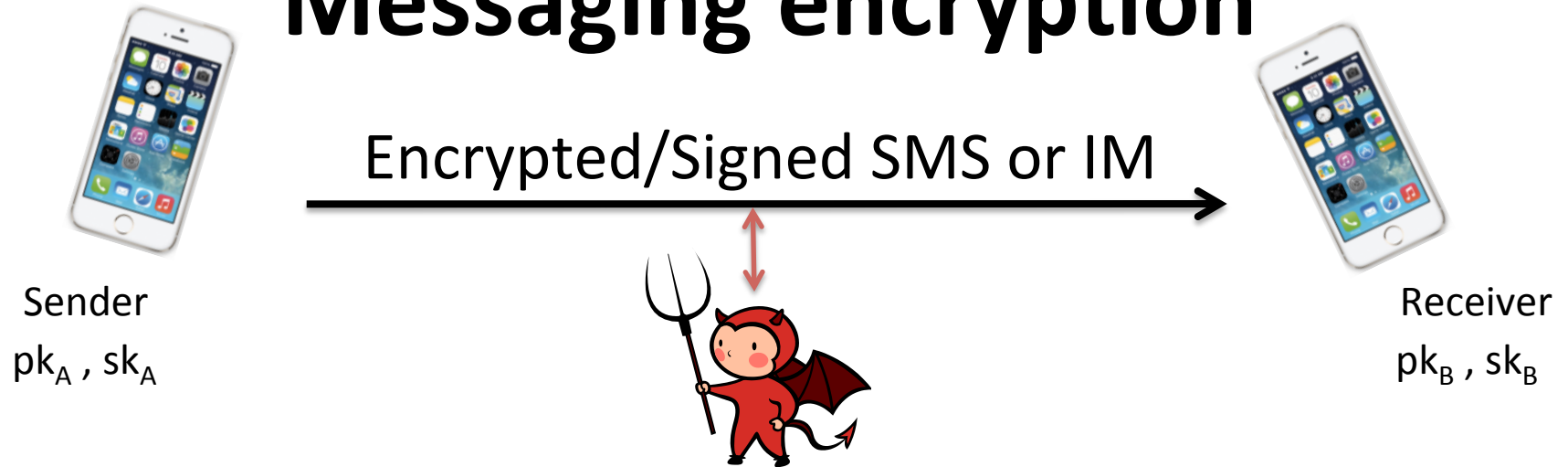
🕒 06/10/2013 04:23 pm ET | Updated Jun 11, 2013

The Switch

Yahoo's plan to get Mail users to encrypt their e-mail: Make it simple

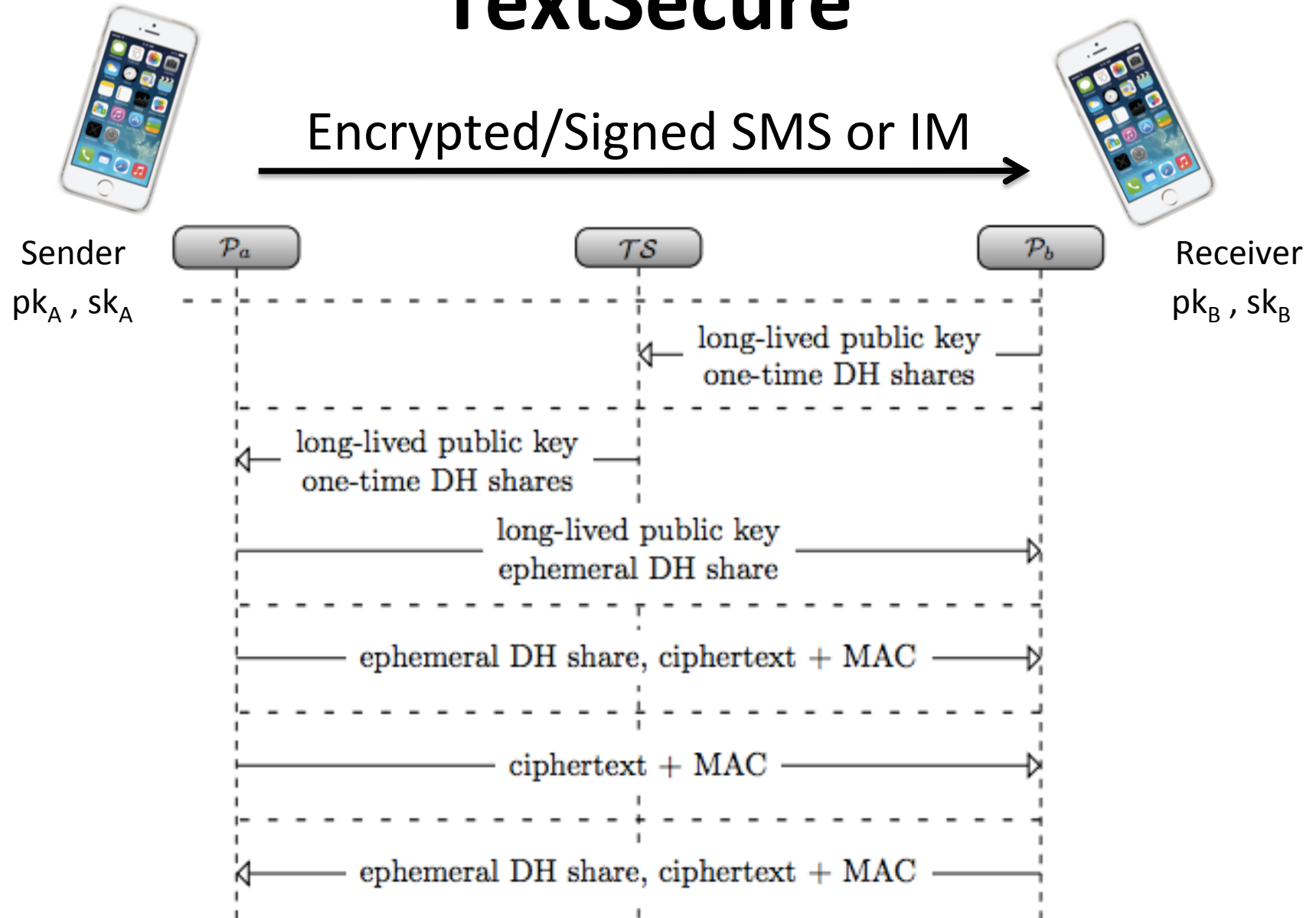
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Messaging encryption



- End-to-end encrypted messaging is a big topic
- TextSecure is protocol adopted by WhatsApp (~1 billion users)

TextSecure

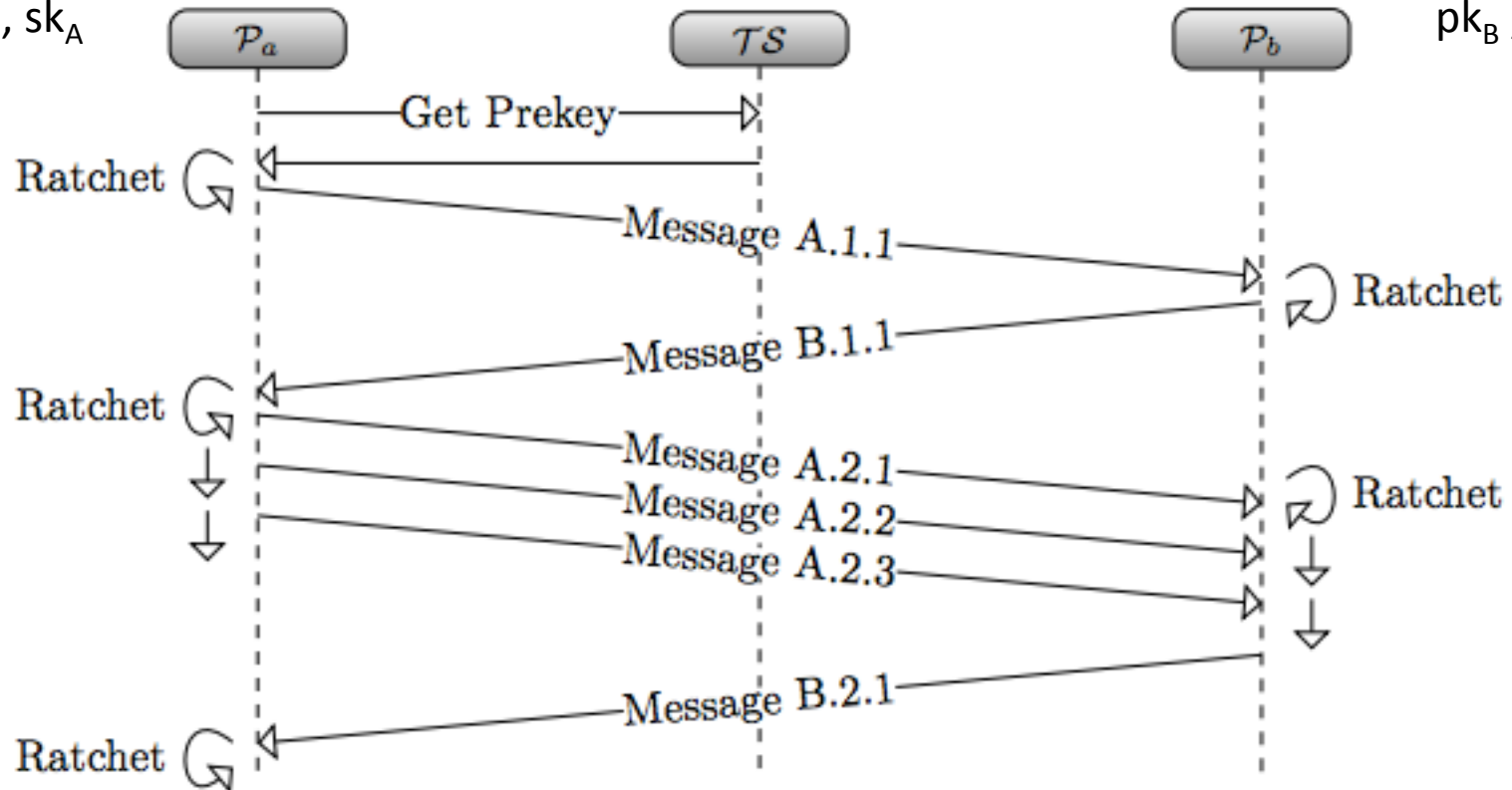


TextSecure

Encrypted/Signed SMS or IM

Sender
 pk_A, sk_A

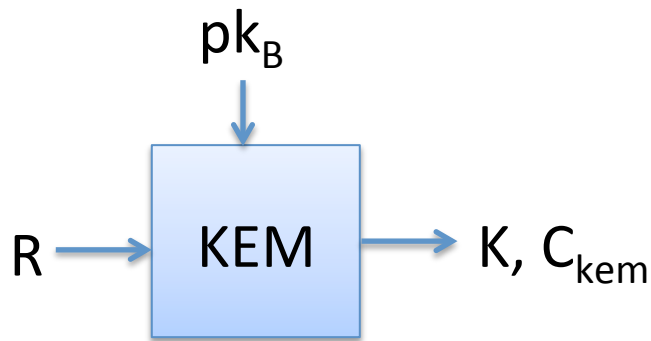
Receiver
 pk_B, sk_B



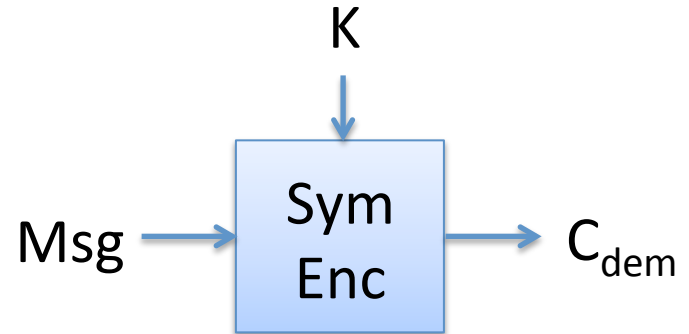
Summary

- Hybrid encryption uses combination of asymmetric and symmetric cryptography
 - Key encapsulation mechanisms (KEM) based on secure PKE, (elliptic curve) Diffie-Hellman
 - Use an authenticated encryption scheme for data encapsulation mechanism (DEM)
- PGP is historical example (and still somewhat widely used)
- End-to-end messaging for IM/chat hotter topic, now widely deployed

Hybrid encryption (KEM/DEM)



KEM = key encapsulation mechanism
Public-key primitive



DEM = data encapsulation mechanism
One-time secure authenticated encryption

HybEnc(pk, M)

Choose randomness R

$K, C_{\text{kem}} \leftarrow \text{KEM}(pk, R)$

$C_{\text{dem}} \leftarrow \text{Enc}(K, M)$

Return $C_{\text{kem}}, C_{\text{dem}}$

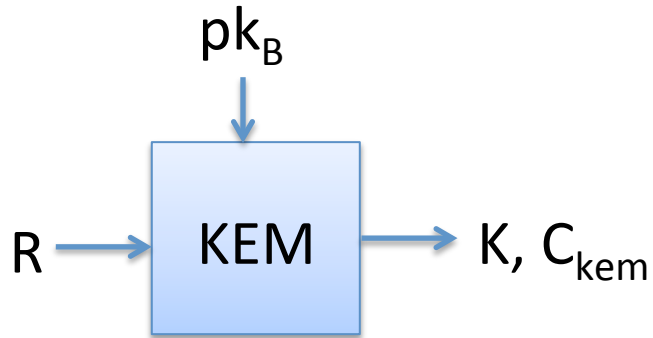
HybDec(sk, C_{kem} , C_{dem})

$K \leftarrow \text{KEM}^{-1}(sk, C_{\text{kem}})$

$M \leftarrow \text{Dec}(K, C_{\text{dem}})$

Return M

KEM from PKE



$KEM(pk, R)$

$C_{kem} \leftarrow \text{PKE-Enc}(pk, R)$

Return $H(R), C_{kem}$

KEM = key encapsulation mechanism
Public-key primitive

ElGamal encryption

Kg outputs $pk = (g, X = g^x)$ and $sk = (g, x)$

g is generator for group of order prime p

Enc((g, X) , M , R)

$r = R \bmod p$

$C1 = g^r$

$C2 = X^r * M$

Return $C1, C2$

Dec((g, x) , $C1, C2$):

Return $C2 * C1^{-x}$

This is only at most chosen-plaintext attack secure. CCA attacks?

ElGamal KEM

Kg outputs $pk = (g, X = g^x)$ and $sk = (g, x)$
 g is generator for group of order prime p

EG-KEM((g,X), R)

$r = R \bmod p$

$C_{\text{kem}} = g^r$

$K = X^r$

Return $H(K)$, C_{kem}

Dec((g,x), C_{kem}):

Return $H(C_{\text{kem}}^x)$

Secure if computational Diffie-Hellman assumption holds in group