## University of Texas at Dallas CS 6322 : Information Retrieval Fall 2012 Homework # 2

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Issued: October 8<sup>th</sup> 2012 Due November 8<sup>th</sup> 2012 before midnight

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# Problem (200 points)

Indexing

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In this assignment you build the index for a simple statistical retrieval system. In the next assignment, you will build the retrieval system itself. For this assignment, index the Cranfield documents used in the last assignment:

A copy of the publicly available Cranfield collection is located on the UTD Apache machine at:

/people/cs/s/sanda/cs6322/Cranfield/

Do not store stop-words in your index. You may use your own stop-word list or the list from the directory on the UTD Apache machine at:

/people/cs/s/sanda/cs6322/Cranfield/resourcesIR.

The terms in your index should be stemmed with the Porter stemmer. The code for the Porter stemmer is available in the same directory on the UTD Apache machine at: /people/cs/s/sanda/cs6322/Cranfield /resourcesIR.

For every term that is indexed, store:

- Document frequency (df): The number of documents that the term occurs in.
- Term frequency (tf): The number of times that the term occurs in each document, and
- The list of documents containing the term.

For each document, store the frequency of the most frequent stem in that document (max\_tf), and the total number of word occurrences in the document, including stopwords (doclen).

Store the inverted lists in your own storage manager. Also, in a second version of your index, compress the inverted lists before storing them, using delta encoding for the document-id and gamma code for the frequency information. A penalty of -100 points will be applied if you do not have also a version of the compressed index, obtaining only max 100 points in this homework.

Delta codes are similar to the gamma codes: they represent a gap by a pair: (length, offset). First the number is represented in binary code. The length of the binary representation is encoded in gamma code, prior to removing the leading 1-bit. After generating the code of the length only, the leading 1-bit is removed and represented in gamma code.

#### <u>Example 1</u>: To write 5 in gamma and delta codes we perform the following operations:

- 1. write 5 in binary as 101
- 2. For the gamma code remove the leading 1-bit to obtain the offset: 01
- 3. The length of the offset is 2:
  - In unary the length is 110
- 4. The code of 5 in gamma is 11001 (or 110,01 to represent [length, offset])
- 5. For the delta code, the length of the offset is 3, because the leading 1-bit is removed afterwards. When writing the length=3 in gamma code it becomes:10,1
- 6. for delta code, the leading 1-bit of the offset is removed now, generating an offset of 10
- 7. The code of 5 in delta is 10101

## Example 2: To write 9 in gamma and delta codes we perform the following operations:

- 1. write 9 in binary as 1001
- 2. for gamma code, remove the leading 1-bit to obtain the offset: 001
- 3. The length of the offset is 3:
  - In unary the length is 1110
- 4. The code of 9 in gamma is 1110001 (or 1110,001 to represent [length, offset])
- 5. for the delta code, the length of the binary representation is 4
- 6. The length is represented in gamma code: 11000
- 7. The leading 1-bit of the binary representation is removed, generating the offset 001
- 8. The code of 9 in delta is 11000001

### <u>Example 3:</u> To write 1 in gamma and delta , we perform the following operations:

- 1. write 1 in binary: 1
- 2. For the gamma code, we remove 1, generating a length=0, which is still 0 in unary code
- 3. in Gamma code, 1 becomes 0
- 4. for the Delta code, the length of the binary representation for 1 is 1. Gamma code for 1 is 0
- 5. The code for 1 in delta is 0

More values for gamma and delta codes are given in the following table:

N	Binary	Gamma code	Delta code
1	1	Len(-)=0; unary(0)=0; Gamma(1)=0	Len(1)=1; Gamma(1)=0; Delta(1)=0
2	10	Len(0)=1; unary(1)=10; $Gamma(2)=100$	Len(10)=2; Gamma(2)=100;
			Delta(2)=1000
3	11	Len(1)=1; unary(1)=10; $Gamma(3)=101$	Len(11)=2; Gamma(2)=100;
			Delta(3)=1001
4	100	Len(00)=2; unary(2)=110;	Len(100)=3; Gamma(3)=101;
		Gamma(4)=11000	Delta(4)=10100
5	101	Len(01)=2; unary(2)=110;	Len(101)=3; Gamma(3)=101;
		Gamma(5)=11001	Delta(5)=10101
6	110	Len(10)=2; unary(2)=110;	Len(110)=3; Gamma(3)=101;
		Gamma(6)=11010	Delta(6)=10110
7	111	Len(11)=2; unary(2)=110;	Len(111)=3; Gamma(3)=101;
		Gamma(7)=11011	Delta(7)=10111
8	1000	Len(000)=3; unary(3)=1110;	Len(1000)=4; Gamma(4)=11000;
		Gamma(8)=111000	Delta(8)=11000000
9	1001	Len(001)=3; unary(3)=1110;	Len(1001)=4; Gamma(4)=11000;
		Gamma(9)=111001	Delta(9)=11000001
10	1010	Len(010)=3; unary(2)=1110;	Len(1010)=4; Gamma(4)=11000;
		Gamma(10)=1110010	Delta(10)=11000010

Production-level IR systems build these compressed indices in about 5 minutes. If your program takes more than an hour, you are doing something wrong.

Turn in your program, a written description of your program, and the following statistics:

- the elapsed time ("wall-clock time") required to build your index,
- the size of the index uncompressed (in bytes),
- the size of the index compressed (in bytes),
- the number of inverted lists in the index, and
- the df, tf, and inverted list length (in bytes) for the terms:

"Reynolds", "NASA", "Prandtl", "flow", "pressure", "boundary", "shock" (or stems that correspond to them).