Greedy Algorithms: Fractional Knapsack

Michael Levin

Higher School of Economics

Algorithmic Toolbox
Data Structures and Algorithms

Outline

Long Hike

Long Hike

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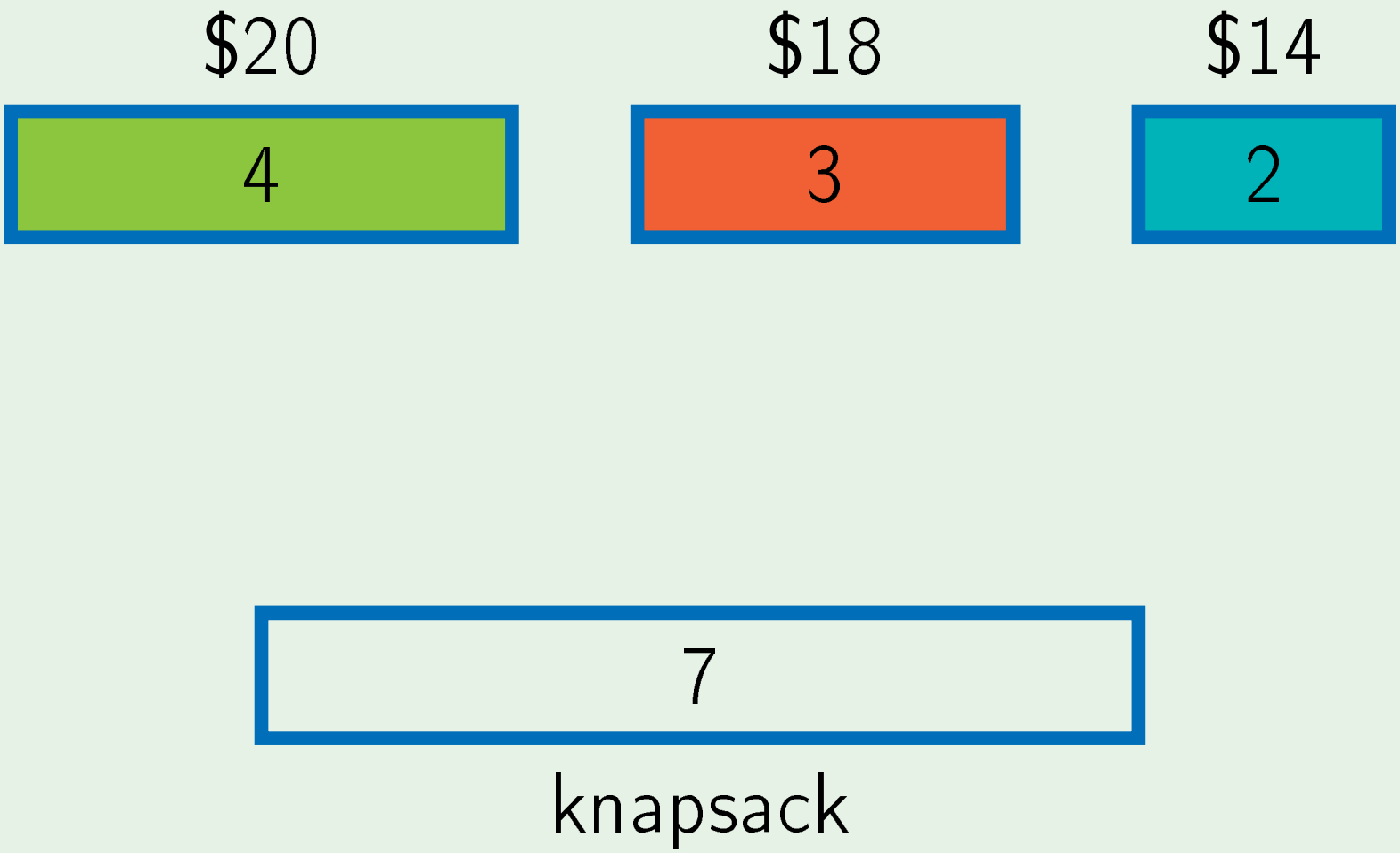
Fractional knapsack

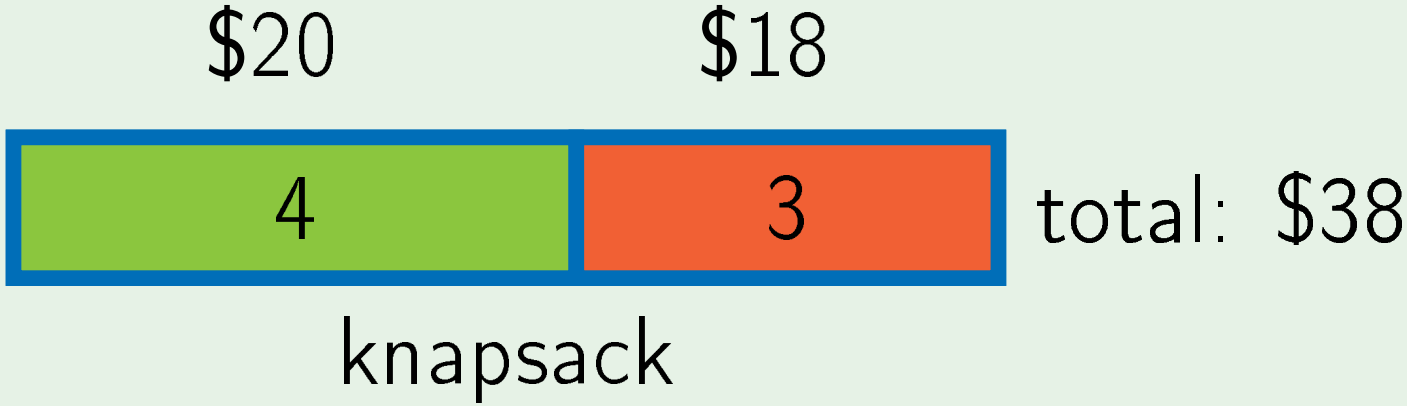
Input: Weights $w_1, \ldots, w_n$ and values $v_1, \ldots, v_n$ of n items; capacity W.

Output: The maximum total value of fractions of items that fit into a bag of capacity W.

Example



Example



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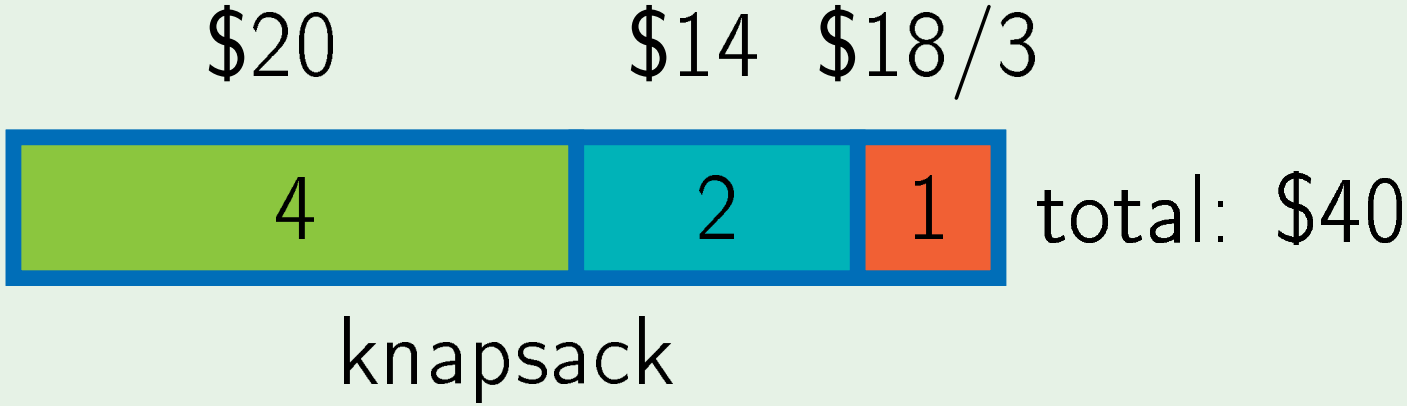
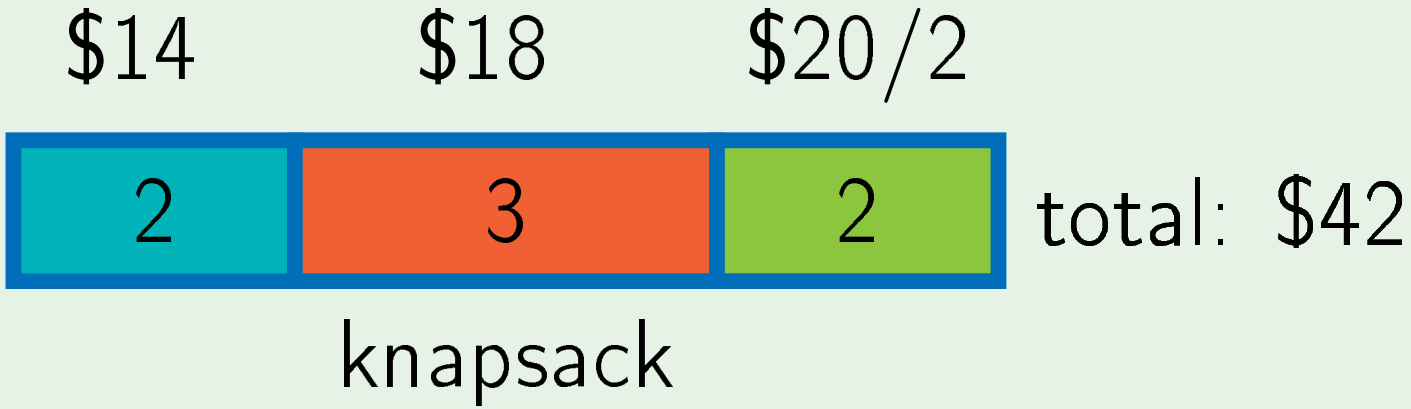
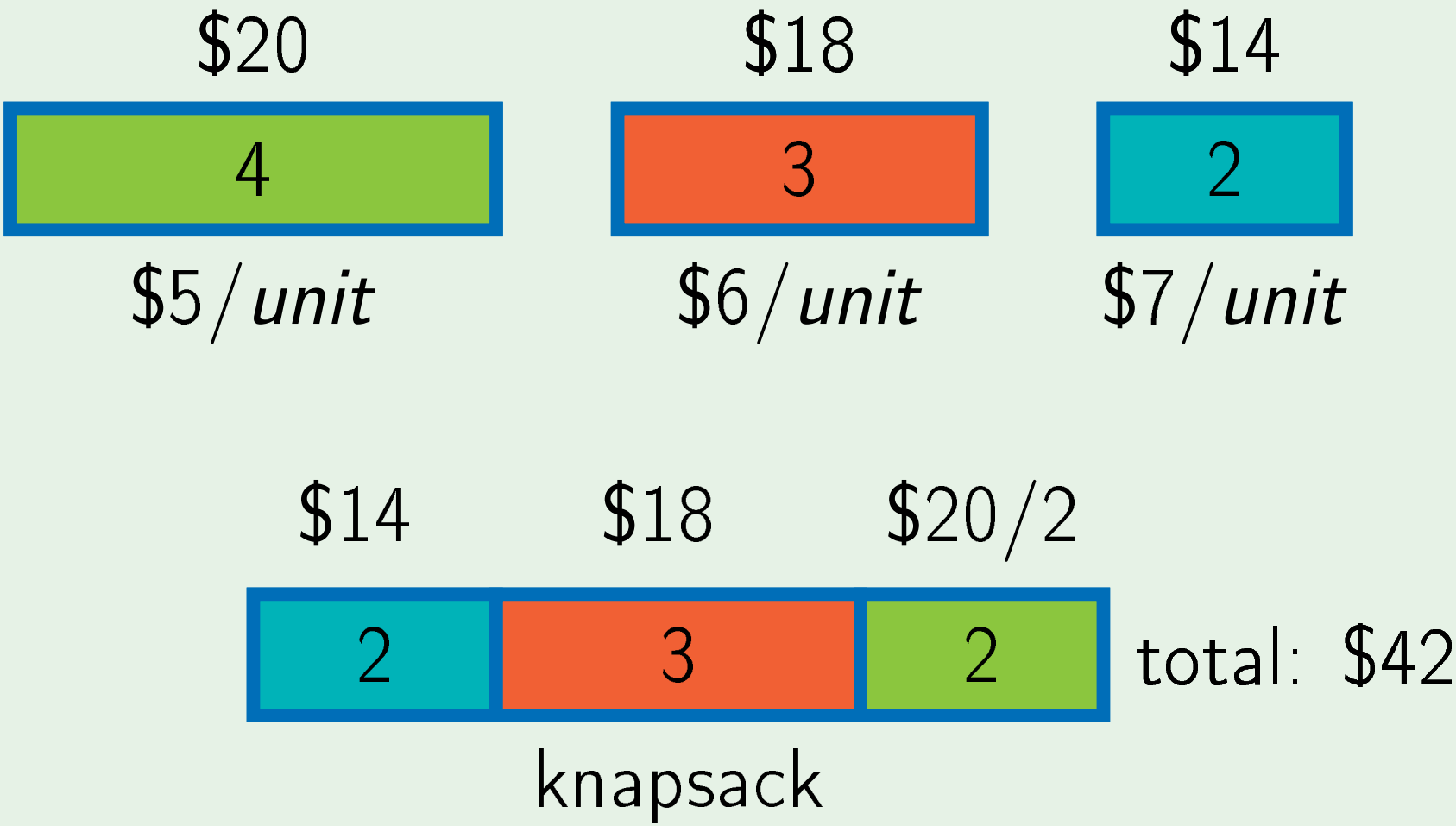


Example



Example

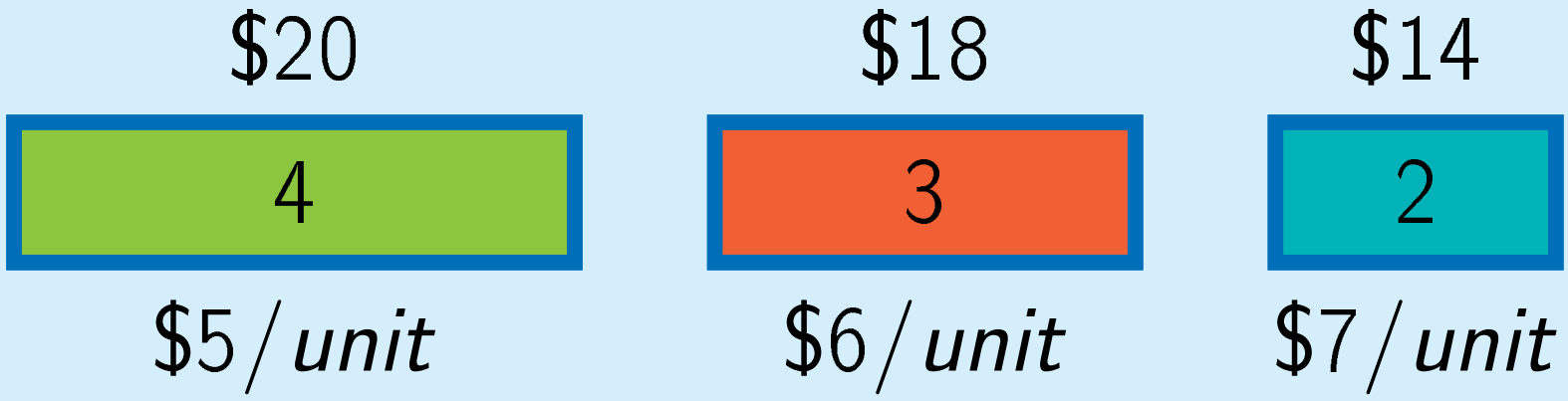


Safe move

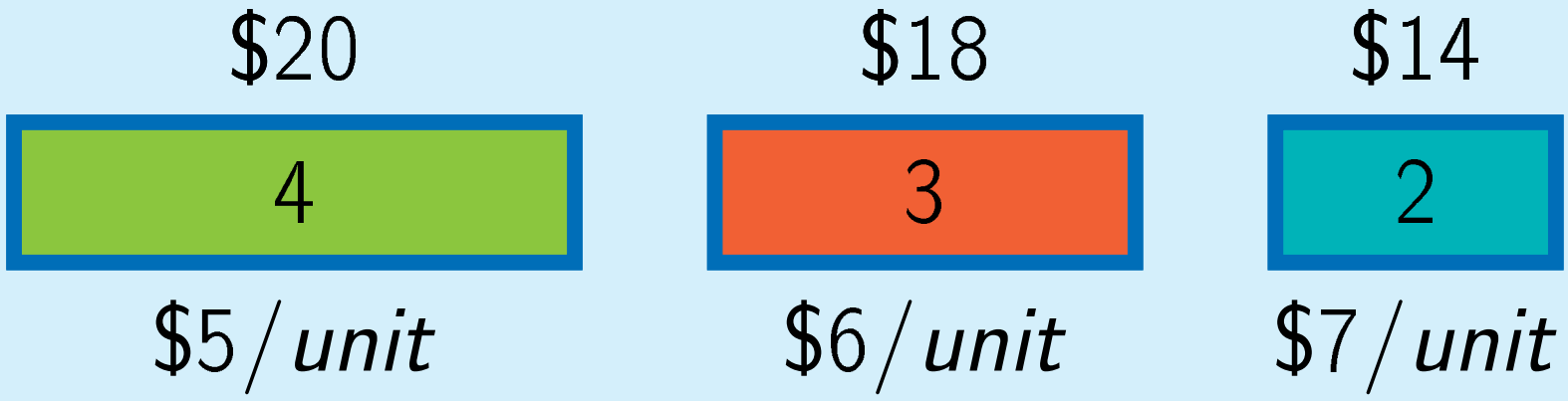
Lemma

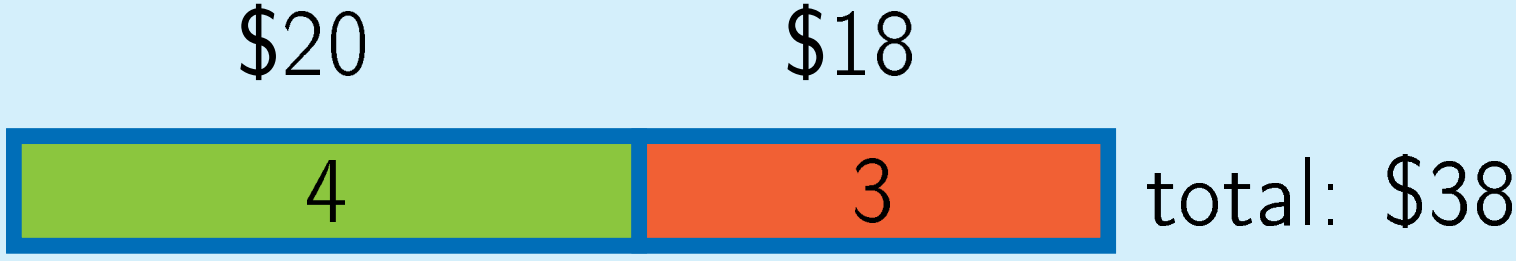
There exists an optimal solution that uses as much as possible of an item with the maximal value per unit of weight.

Proof

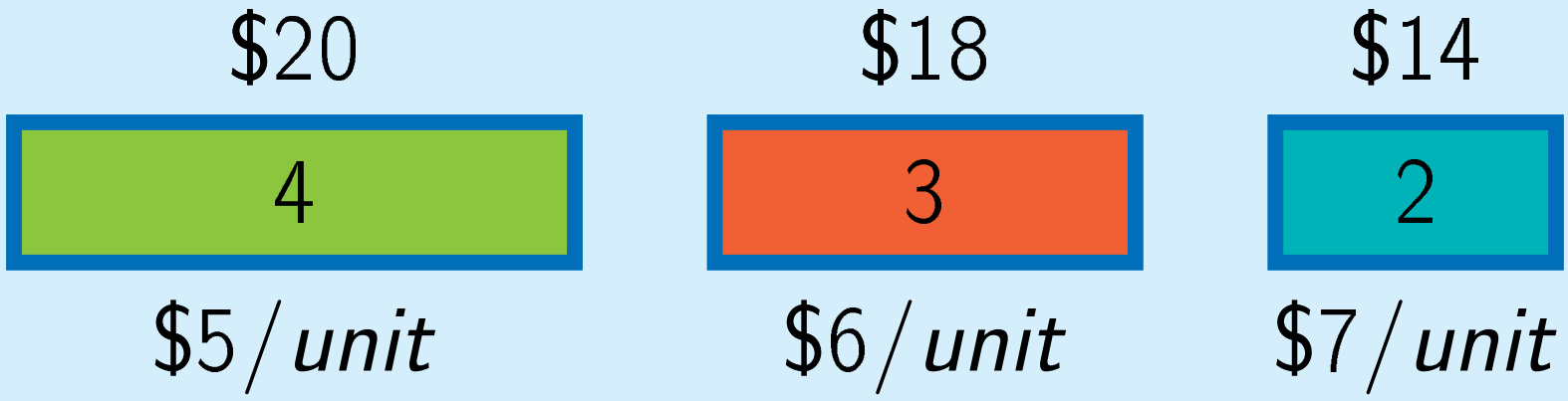


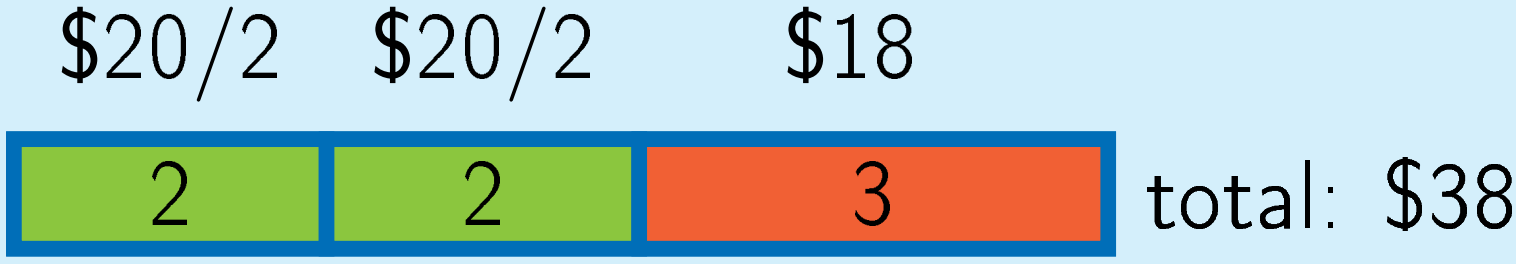
Proof





Proof





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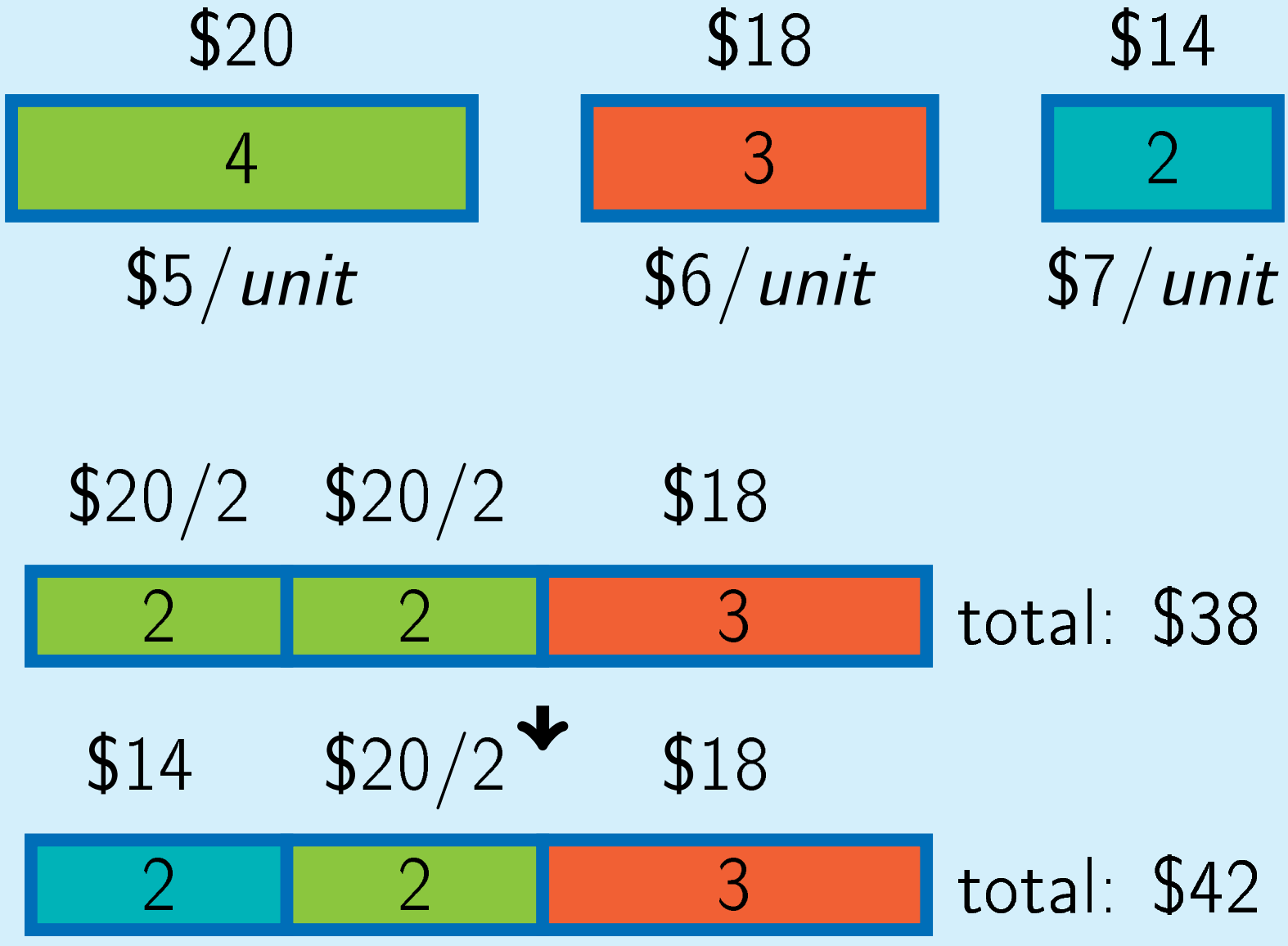
$20
$18
$14
4
3
2
$5/unit
$6/unit
$7/unit
$20/2
$20/2
$18
2
2
3
total: $38
$14
$20/2
$18
2
2
3
total: $42

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Knapsack$(W, w_1, v_1, \ldots, w_n, v_n)$

$A \leftarrow [0, 0, \ldots, 0], V \leftarrow 0$

repeat n times:

  if $W = 0$:

    return (V, A)

  select i with $w_i > 0$ and max $\frac{v_i}{w_i}$

  $a \leftarrow \min(w_i, W)$

  $V \leftarrow V + a\frac{v_i}{w_i}$

  $w_i \leftarrow w_i - a, A[i] \leftarrow A[i] + a, W \leftarrow W - a$

return (V, A)

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The running time of `Knapsack` is $O(n^2)$.

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- Select best item on each step is $O(n)$
- Main loop is executed n times
- Overall, $O(n^2)$ □

Optimization

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- First, sort items by decreasing $\frac{v}{w}$

Assume $\frac{v_1}{w_1} \geq \frac{v_2}{w_2} \geq \cdots \geq \frac{v_n}{w_n}$

Knapsack$(W, w_1, v_1, \ldots, w_n, v_n)$

$A \leftarrow [0, 0, \ldots, 0], V \leftarrow 0$
for i from 1 to n:
  if $W = 0$:
    return (V, A)
  $a \leftarrow \min(w_i, W)$
  $V \leftarrow V + a\frac{v_i}{w_i}$
  $w_i \leftarrow w_i - a, A[i] \leftarrow A[i] + a, W \leftarrow W - a$
return (V, A)

Asymptotics

- Now each iteration is $O(1)$
- `Knapsack` after sorting is $O(n)$
- Sort + `Knapsack` is $O(n \log n)$