

# Views

- In some cases, it is not desirable for all users to see the entire logical model (that is, all the actual relations stored in the database.)
- Consider a person who needs to know an instructors name and department, but not the salary. This person should see a relation described, in SQL, by

```
select ID, name, dept_name  
from instructor
```

- A **view** provides a mechanism to hide certain data from the view of certain users.
- Any relation that is not of the conceptual model but is made visible to a user as a “virtual relation” is called a **view**.

# View Definition

- A view is defined using the **create view** statement which has the form

**create view** *v* **as** <query expression>

where <query expression> is any legal SQL expression. The view name is represented by *v*.

- Once a view is defined, the view name can be used to refer to the virtual relation that the view generates.
- View definition is not the same as creating a new relation by evaluating the query expression
  - Rather, a view definition causes the saving of an expression; the expression is substituted into queries using the view.

# Example Views

- A view of instructors without their salary

```
create view faculty as  
  select ID, name, dept_name  
  from instructor
```

- Find all instructors in the Biology department

```
select name  
from faculty  
where dept_name = 'Biology'
```

- Create a view of department salary totals

```
create view departments_total_salary(dept_name, total_salary) as  
  select dept_name, sum (salary)  
  from instructor  
  group by dept_name;
```

# Views Defined Using Other Views

- **create view** *physics\_fall\_2009* **as**  
    **select** *course.course\_id, sec\_id, building, room\_number*  
    **from** *course, section*  
    **where** *course.course\_id = section.course\_id*  
          **and** *course.dept\_name = 'Physics'*  
          **and** *section.semester = 'Fall'*  
          **and** *section.year = '2009'*;
- **create view** *physics\_fall\_2009\_watson* **as**  
    **select** *course\_id, room\_number*  
    **from** *physics\_fall\_2009*  
    **where** *building = 'Watson'*;

# View Expansion

- Expand use of a view in a query/another view

```
create view physics_fall_2009_watson as  
(select course_id, room_number  
from (select course.course_id, building, room_number  
      from course, section  
      where course.course_id = section.course_id  
           and course.dept_name = 'Physics'  
           and section.semester = 'Fall'  
           and section.year = '2009')  
where building = 'Watson';
```

# Views Defined Using Other Views

- One view may be used in the expression defining another view
- A view relation  $v_1$  is said to *depend directly* on a view relation  $v_2$  if  $v_2$  is used in the expression defining  $v_1$
- A view relation  $v_1$  is said to *depend on* view relation  $v_2$  if either  $v_1$  depends directly to  $v_2$  or there is a path of dependencies from  $v_1$  to  $v_2$
- A view relation  $v$  is said to be *recursive* if it depends on itself.

# View Expansion

- A way to define the meaning of views defined in terms of other views.
- Let view  $v_1$  be defined by an expression  $e_1$  that may itself contain uses of view relations.
- View expansion of an expression repeats the following replacement step:
  - repeat**
    - Find any view relation  $v_i$  in  $e_1$
    - Replace the view relation  $v_i$  by the expression defining  $v_i$
  - until** no more view relations are present in  $e_1$
- As long as the view definitions are not recursive, this loop will terminate

# Update of a View

- Add a new tuple to *faculty* view which we defined earlier

**insert into *faculty* values** ('30765', 'Green', 'Music');

This insertion must be represented by the insertion of the tuple

('30765', 'Green', 'Music', null)

into the *instructor* relation



# Some Updates cannot be Translated Uniquely

- **create view** *instructor\_info* as  
    **select** *ID, name, building*  
    **from** *instructor, department*  
    **where** *instructor.dept\_name= department.dept\_name;*
- **insert into** *instructor\_info* **values** ('69987', 'White', 'Taylor');
  - ▶ which department, if multiple departments in Taylor?
  - ▶ what if no department is in Taylor?
- Most SQL implementations allow updates only on simple views
  - The **from** clause has only one database relation.
  - The **select** clause contains only attribute names of the relation, and does not have any expressions, aggregates, or **distinct** specification.
  - Any attribute not listed in the **select** clause can be set to null
  - The query does not have a **group** by or **having** clause.

# And Some Not at All

- **create view** *history\_instructors* **as**  
    **select** \*  
    **from** *instructor*  
    **where** *dept\_name*= 'History';
- What happens if we insert ('25566', 'Brown', 'Biology', 100000)  
into *history\_instructors*?

# Materialized Views

- **Materializing a view**: create a physical table containing all the tuples in the result of the query defining the view
- If relations used in the query are updated, the materialized view result becomes out of date
  - Need to **maintain** the view, by updating the view whenever the underlying relations are updated.

# Transactions

- Unit of work
- Atomic transaction
  - either fully executed or rolled back as if it never occurred
- Isolation from concurrent transactions
- Transactions begin implicitly
  - Ended by **commit work** or **rollback work**
- But default on most databases: each SQL statement commits automatically
  - Can turn off auto commit for a session (e.g. using API)
  - In SQL:1999, can use: **begin atomic .... end**
    - ▶ Not supported on most databases

# Integrity Constraints

- Integrity constraints guard against accidental damage to the database, by ensuring that authorized changes to the database do not result in a loss of data consistency.
  - A checking account must have a balance greater than \$10,000.00
  - A salary of a bank employee must be at least \$4.00 an hour
  - A customer must have a (non-null) phone number

# Integrity Constraints on a Single Relation

- **not null**
- **primary key**
- **unique**
- **check** (P), where P is a predicate

# Not Null and Unique Constraints

## ■ not null

- Declare *name* and *budget* to be **not null**

*name* **varchar(20) not null**

*budget* **numeric(12,2) not null**

## ■ unique ( $A_1, A_2, \dots, A_m$ )

- The unique specification states that the attributes  $A_1, A_2, \dots, A_m$  form a candidate key.
- Candidate keys are permitted to be null (in contrast to primary keys).

# The check clause

## ■ **check** (P)

where P is a predicate

Example: ensure that semester is one of fall, winter, spring or summer:

```
create table section (  
    course_id varchar (8),  
    sec_id varchar (8),  
    semester varchar (6),  
    year numeric (4,0),  
    building varchar (15),  
    room_number varchar (7),  
    time slot id varchar (4),  
    primary key (course_id, sec_id, semester, year),  
    check (semester in ('Fall', 'Winter', 'Spring', 'Summer'))  
);
```



# Referential Integrity

- Ensures that a value that appears in one relation for a given set of attributes also appears for a certain set of attributes in another relation.
  - Example: If “Biology” is a department name appearing in one of the tuples in the *instructor* relation, then there exists a tuple in the *department* relation for “Biology”.
- Let A be a set of attributes. Let R and S be two relations that contain attributes A and where A is the primary key of S. A is said to be a **foreign key** of R if for any values of A appearing in R these values also appear in S.

# Cascading Actions in Referential Integrity

- **create table** *course* (  
    *course\_id* **char**(5) **primary key**,  
    *title* **varchar**(20),  
    *dept\_name* **varchar**(20) **references** *department*  
)
- **create table** *course* (  
    ...  
    *dept\_name* **varchar**(20),  
    **foreign key** (*dept\_name*) **references** *department*  
        **on delete cascade**  
        **on update cascade**,  
    ...  
)
- alternative actions to cascade: **set null**, **set default**

# Integrity Constraint Violation During Transactions

■ E.g.

```
create table person (  
  ID char(10),  
  name char(40),  
  mother char(10),  
  father char(10),  
  primary key ID,  
  foreign key father references person,  
  foreign key mother references person)
```

- How to insert a tuple without causing constraint violation ?
  - insert father and mother of a person before inserting person
  - OR, set father and mother to null initially, update after inserting all persons (not possible if father and mother attributes declared to be **not null**)
  - OR defer constraint checking (next slide)

# Complex Check Clauses

- **check** (*time\_slot\_id* **in** (**select** *time\_slot\_id* **from** *time\_slot*))
  - why not use a foreign key here?
- Every section has at least one instructor teaching the section.
  - how to write this?
- Unfortunately: subquery in check clause not supported by pretty much any database
  - Alternative: triggers (later)
- **create assertion** <assertion-name> **check** <predicate>;
  - Also not supported by anyone

# Built-in Data Types in SQL

- **date**: Dates, containing a (4 digit) year, month and date
  - Example: **date** '2005-7-27'
- **time**: Time of day, in hours, minutes and seconds.
  - Example: **time** '09:00:30'      **time** '09:00:30.75'
- **timestamp**: date plus time of day
  - Example: **timestamp** '2005-7-27 09:00:30.75'
- **interval**: period of time
  - Example: interval '1' day
  - Subtracting a date/time/timestamp value from another gives an interval value
  - Interval values can be added to date/time/timestamp values

# Index Creation

- **create table** *student*  
(*ID* **varchar** (5),  
*name* **varchar** (20) **not null**,  
*dept\_name* **varchar** (20),  
*tot\_cred* **numeric** (3,0) **default** 0,  
**primary key** (*ID*))
- **create index** *studentID\_index* **on** *student*(*ID*)
- Indices are data structures used to speed up access to records with specified values for index attributes
  - e.g. **select** \*  
    **from** *student*  
    **where** *ID* = '12345'

can be executed by using the index to find the required record, without looking at all records of *student*

*More on indices in Chapter 11*

# User-Defined Types

- **create type** construct in SQL creates user-defined type

**create type** *Dollars* as numeric (12,2) final

- **create table** *department*  
(*dept\_name* **varchar** (20),  
*building* **varchar** (15),  
*budget Dollars*);

# Domains

- **create domain** construct in SQL-92 creates user-defined domain types

**create domain *person\_name* char(20) not null**

- Types and domains are similar. Domains can have constraints, such as **not null**, specified on them.
- **create domain *degree\_level* varchar(10)**  
**constraint *degree\_level\_test***  
**check (value in ('Bachelors', 'Masters', 'Doctorate'));**



# Large-Object Types

- Large objects (photos, videos, CAD files, etc.) are stored as a *large object*:
  - **blob**: binary large object -- object is a large collection of uninterpreted binary data (whose interpretation is left to an application outside of the database system)
  - **clob**: character large object -- object is a large collection of character data
  - When a query returns a large object, a pointer is returned rather than the large object itself.

# Authorization

Forms of authorization on parts of the database:

- **Read** - allows reading, but not modification of data.
- **Insert** - allows insertion of new data, but not modification of existing data.
- **Update** - allows modification, but not deletion of data.
- **Delete** - allows deletion of data.

Forms of authorization to modify the database schema

- **Index** - allows creation and deletion of indices.
- **Resources** - allows creation of new relations.
- **Alteration** - allows addition or deletion of attributes in a relation.
- **Drop** - allows deletion of relations.

# Authorization Specification in SQL

- The **grant** statement is used to confer authorization  
    **grant** <privilege list>  
    **on** <relation name or view name> **to** <user list>
- <user list> is:
  - a user-id
  - **public**, which allows all valid users the privilege granted
  - A role (more on this later)
- Granting a privilege on a view does not imply granting any privileges on the underlying relations.
- The grantor of the privilege must already hold the privilege on the specified item (or be the database administrator).

# Privileges in SQL

- **select**: allows read access to relation, or the ability to query using the view
  - Example: grant users  $U_1$ ,  $U_2$ , and  $U_3$  **select** authorization on the *instructor* relation:  
**grant select on instructor to  $U_1$ ,  $U_2$ ,  $U_3$**
- **insert**: the ability to insert tuples
- **update**: the ability to update using the SQL update statement
- **delete**: the ability to delete tuples.
- **all privileges**: used as a short form for all the allowable privileges