CMPT 409/981: Optimization for Machine Learning

Lecture 3

Sharan Vaswani September 12, 2024

Recap

- For an *L*-smooth function, $f(y) \le f(x) + \langle \nabla f(x), y x \rangle + \frac{L}{2} ||y x||^2$ for all $x, y \in \mathcal{D}$.
- For *L*-smooth functions lower-bounded by f^* , gradient descent with $\eta = \frac{1}{L}$ returns \hat{w} such that $\|\nabla f(\hat{w})\|^2 \le \epsilon$ and requires $T \ge \frac{2L[f(w_0) f^*]}{\epsilon}$ iterations (oracle calls).
- \bullet Importantly, the GD rate does not depend on the dimension of w.
- Lower-Bound: When minimizing a smooth function (without additional assumptions), any first-order algorithm requires $\Omega\left(\frac{1}{\epsilon}\right)$ oracle calls to return a point \hat{w} such that $\|\nabla f(\hat{w})\|^2 \leq \epsilon$.
- Hence, GD is optimal for minimizing smooth functions.

Gradient Descent

- The above results require setting the step-size to $\frac{1}{L}$. In fact, GD with any $\eta \in (0, \frac{2}{L})$ will result in convergence to the stationary point (prove in Assignment 1).
- However, estimating *L* can be difficult as the functions get more complicated.
- Even for simple functions, the theoretically computed *L* is global (the "local" *L* might be much smaller) and often loose in practice. Typically we tend to overestimate *L* resulting in a smaller step-size.
- Instead of setting η according to L, we can "search" for a good step-size η_k in each iteration k. We will study 2 ways to do so:
 - Exact Line-search
 - Backtracking Armijo Line-search

Exact Line-search

Exact line-search: At iteration k, solve the following sub-problem:

$$\eta_k = \operatorname*{arg\,min}_{\eta} f(w_k - \eta \nabla f(w_k)).$$



After computing η_k , do the usual GD update: $w_{k+1} = w_k - \eta_k \nabla f(w_k)$.

- Can adapt to the "local" L, resulting in larger step-sizes and better performance.
- Can solve the sub-problem approximately by doing gradient descent w.r.t η (known as hyper-gradient descent [BCR⁺17]). This is computationally expensive.
- ullet Can compute η_k analytically. This can only be done in special cases such as for quadratics.

Exact Line-search for Linear Regression

Recall linear regression: for $X \in \mathbb{R}^{n \times d}$ and $y \in \mathbb{R}^n$, we aim to solve: $\min_{w \in \mathbb{R}^d} f(w) := \frac{1}{2} \|Xw - y\|^2 = \frac{1}{2} \left[w^\mathsf{T}(X^\mathsf{T}X)w - 2\langle X^\mathsf{T}y, w \rangle + \|y\|^2 \right].$

For the exact line-search, we need to $\min_{\eta} h(\eta) := f(w_k - \eta \nabla f(w_k))$.

Since f is a quadratic, we can directly use the second-order Taylor series:

$$f(w_k - \eta \nabla f(w_k)) = f(w_k) + \langle \nabla f(w_k), -\eta \nabla f(w_k) \rangle + \frac{1}{2} [-\eta \nabla f(w_k)]^\mathsf{T} \nabla^2 f(w_k) [-\eta \nabla f(w_k)]$$

$$\implies \nabla h(\eta_k) = -\|\nabla f(w_k)\|^2 + \eta_k [\nabla f(w_k)]^\mathsf{T} \nabla^2 f(w_k) [\nabla f(w_k)] = 0$$

$$\implies \eta_k = \frac{\|\nabla f(w_k)\|^2}{\|\nabla f(w_k)\|^2_{\nabla^2 f(w_k)}}$$

For linear regression,
$$\nabla^2 f(w_k) = X^\mathsf{T} X$$
 and $\nabla f(w_k) = X^\mathsf{T} (X w_k - y)$. $\implies \eta_k = \frac{\|X^\mathsf{T} (X w_k - y)\|^2}{\|X^\mathsf{T} (X w_k - y)\|_{X^\mathsf{T} X}^2}$. (Implement in Assignment 1)

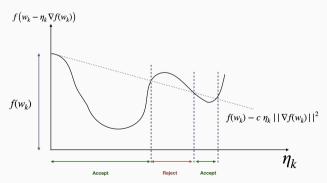
Armijo Condition

Usually, the cost of doing an exact line-search is not worth the computational effort.

Armijo condition for a prospective step-size $\tilde{\eta_k}$:

$$f(w_k - \tilde{\eta}_k \nabla f(w_k)) \le f(w_k) - c \, \tilde{\eta}_k \, \|\nabla f(w_k)\|^2$$

where $c \in (0,1)$ is a hyper-parameter.



Gradient Descent with Backtracking Armijo Line-search

Algorithm GD with Armijo Line-search

- 1: function GD with Armijo line-search(f, w_0 , η_{max} , $c \in (0,1)$, $\beta \in (0,1)$)
- 2: **for** k = 0, ..., T 1 **do**
- 3: $\tilde{\eta}_k \leftarrow \eta_{\text{max}}$
- 4: while $f(w_k \tilde{\eta}_k \nabla f(w_k)) > f(w_k) c \cdot \tilde{\eta}_k \|\nabla f(w_k)\|^2$ do
- 5: $\tilde{\eta}_k \leftarrow \tilde{\eta}_k \beta$
- 6: end while
- 7: $\eta_k \leftarrow \tilde{\eta}_k$
- 8: $w_{k+1} = w_k \eta_k \nabla f(w_k)$
- 9: end for
- 10: **return** w_T

Backtracking Armijo Line-search

Simplification for analysis: Assume that the backtracking line-search procedure returns the largest η that satisfies the Armijo condition. Will be referred to as exact backtracking line-search.

Claim: For *L*-smooth functions, the exact backtracking line-search procedure terminates and returns $\eta_k \geq \min\left\{\frac{2(1-c)}{L}, \eta_{\max}\right\}$.

Proof: For a prospective step-size $\tilde{\eta}_k$, we will use the following two inequalities:

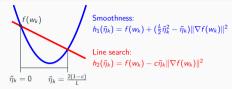
$$f(w_{k} - \tilde{\eta}_{k} \nabla f(w_{k})) \leq \underbrace{f(w_{k}) - \|\nabla f(w_{k})\|^{2} \left(\tilde{\eta}_{k} - \frac{L\tilde{\eta}_{k}^{2}}{2}\right)}_{h_{1}(\tilde{\eta}_{k})} \text{ (Quadratic bound using smoothness)}$$

$$f(w_{k} - \tilde{\eta}_{k} \nabla f(w_{k})) \leq \underbrace{f(w_{k}) - \|\nabla f(w_{k})\|^{2} (c\tilde{\eta}_{k})}_{h_{2}(\tilde{\eta}_{k})} \text{ (Armijo condition)}$$

7

Backtracking Armijo Line-search

Recall that if the Armijo condition is satisfied, the back-tracking line-search procedure terminates.



Case (i) $\eta_{\text{max}} \leq \frac{2(1-c)}{L}$: From smoothness, $f(w_k - \eta_{\text{max}} \nabla f(w_k)) \leq h_1(\eta_{\text{max}})$. For $\eta_{\text{max}} \leq \frac{2(1-c)}{L}$, we know that $h_1(\eta_{\text{max}}) \leq h_2(\eta_{\text{max}})$. Hence, $f(w_k - \eta_{\text{max}} \nabla f(w_k)) \leq h_2(\eta_{\text{max}})$, meaning that the Armijo condition is satisfied for η_{max} . \Longrightarrow if $\eta_{\text{max}} \leq \frac{2(1-c)}{L}$, then the line-search terminates immediately and $\eta_k = \eta_{\text{max}}$.

Case (ii): If $\eta_{\text{max}} > \frac{2(1-c)}{L}$: While backtracking, if $\tilde{\eta}_k = \frac{2(1-c)}{L}$, then $f(w_k - \tilde{\eta}_k \nabla f(w_k)) \leq h_1(\tilde{\eta}_k) = h_2(\tilde{\eta}_k)$, the line-search terminates immediately and $\eta_k = \frac{2(1-c)}{L}$. If the Armijo condition is satisfied for a step-size η_k s.t. $h_2(\eta_k) < h_1(\eta_k)$, then $f(w_k - \eta_k \nabla f(w_k)) \leq h_2(\eta_k) < h_1(\eta_k) \implies c\eta_k \geq \eta_k - \frac{L\eta_k^2}{2} \implies \eta_k \geq \frac{2(1-c)}{L}$.

Putting everything together, the step-size η_k returned by the Armijo line-search satisfies $\eta_k \geq \min\left\{\frac{2(1-c)}{L}, \eta_{\max}\right\}$.

Gradient Descent with Backtracking Armijo Line-search

Claim: For *L*-smooth functions lower-bounded by f^* , gradient descent with exact backtracking Armijo line-search (with c=1/2) returns point \hat{w} such that $\|\nabla f(\hat{w})\|^2 \leq \epsilon$ and requires $T \geq \frac{\max\{2L, 2/\eta_{\max}\}\left[f(w_0) - \min_w f(w)\right]}{\epsilon}$ iterations.

Proof: Since η_k satisfies the Armijo condition and $w_{k+1} = w_k - \eta_k \nabla f(w_k)$,

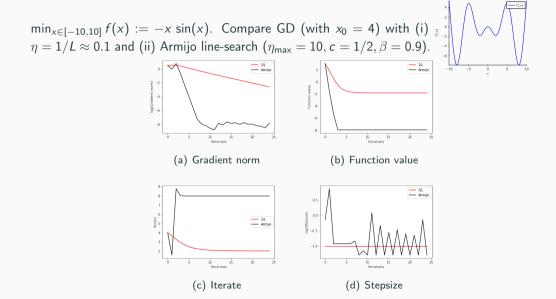
$$\begin{split} f(w_{k+1}) & \leq f(w_k) - c \, \eta_k \, \left\| \nabla f(w_k) \right\|^2 \\ & \leq f(w_k) - \left(\min \left\{ \frac{1}{2L}, \frac{\eta_{\mathsf{max}}}{2} \right\} \right) \, \left\| \nabla f(w_k) \right\|^2 \\ & \qquad \qquad \left(\mathsf{Result from previous slide with } \, c = 1/2) \end{split}$$

Continuing the proof as before,

$$\implies \|\nabla f(\hat{w})\|^2 \le \frac{\max\{2L, 2/\eta_{\max}\} [f(w_0) - f^*]}{T}$$

The claim can be proved by the same reasoning as in Lecture 2.

Gradient Descent with Backtracking Armijo Line-search - Example





Convex Optimization

For smooth functions, GD requires $\Theta(1/\epsilon)$ iterations to converge to an ϵ -approximate stationary point. Alternatively, if we care about global optimization (reach the vicinity of the true minimizer), any algorithm requires $\Omega(1/\epsilon^d)$ iterations.

Convex functions: Class of functions where local optimization can result in convergence to the global minimizer of the function.

In general, convex optimization involves minimizing a convex function over a convex set \mathcal{C} .

Examples of convex optimization in ML

Ridge regression: $\min_{w \in \mathbb{R}^d} \frac{1}{2} \|Xw - y\|^2 + \frac{\lambda}{2} \|w\|^2$.

Logistic regression: $\min_{w \in \mathbb{R}^d} \sum_{i=1}^n \log (1 + \exp(-y_i \langle X_i, w \rangle))$

Support vector machines: $\min_{w \in \mathbb{R}^d} \sum_{i=1}^n \max \{0, 1 - y_i \langle X_i, w \rangle\} + \frac{\lambda}{2} \|w\|^2$

Planning in MDPs in RL: $\max_{\mu \in \mathcal{F}_{\rho}} \langle \mu, r \rangle$ where \mathcal{F}_{ρ} is the flow-polytope.

Convex Sets

A set $\mathcal C$ is convex if every point along the line joining two points in $\mathcal C$ also lies in the set.

For points x, y, the *convex combination* of x, y is $z_{\theta} := \theta x + (1 - \theta)y$ for $\theta \in [0, 1]$.

A set C is convex iff $\forall x, y \in C$, the convex combination $z_{\theta} \in C$ for all $\theta \in [0, 1]$.

Examples of convex sets:

- Positive orthant $\mathbb{R}^d_+: \{x|x \geq 0\}.$
- Hyper-plane: $\{x | Ax = b\}$.
- Half-space: $\{x | Ax \le b\}$.
- Norm-ball: $\{x | \|x\|_p \le r\}$ for $p \ge 1$.
- Norm-cone: $\{(x,r)| \|x\|_p \le r\}$ for $p \ge 1$.

Convex Sets

Q: Prove that the hyper-plane (set of linear equations): $\mathcal{H} := \{x | Ax = b\}$ is a convex set.

If $x, y \in \mathcal{H}$, then, Ax = b and Ay = b. Consider a point $z_{\theta} := \theta x + (1 - \theta)y$ for $\theta \in [0, 1]$.

$$Az_{\theta} = A[\theta x + (1 - \theta)y] = \theta Ax + (1 - \theta)Ay = b.$$

Hence, $z_{\theta} \in \mathcal{H}$ for all $\theta \in [0,1]$ and \mathcal{H} is a convex set.

Q: Prove that the ball of radius r centered at point x_c : $\mathcal{B}(x_c, r) := \{x | \|x - x_c\|_p \le r\}$ for $p \ge 1$ is convex.

If
$$x, y \in \mathcal{B}(x_c, r)$$
, then, $\|x - x_c\|_p \le r$ and $\|y - x_c\|_p \le r$. Consider a point $z_\theta := \theta x + (1 - \theta)y$ for $\theta \in [0, 1]$. $\|z_\theta - x_c\|_p = \|\theta(x - x_c) + (1 - \theta)(y - x_c)\|_p$ $\leq \|\theta(x - x_c)\|_p + \|(1 - \theta)(y - x_c)\|_p$ (Triangle inequality for norms) $\leq \theta \|(x - x_c)\|_p + (1 - \theta) \|(y - x_c)\|_p$ (Homogeneity of norms)

$$\implies \|z - x_c\|_p \le r$$

Hence, $z_{\theta} \in \mathcal{B}(x_c, r)$ for all $\theta \in [0, 1]$ and $\mathcal{B}(x_c, r)$ is a convex set.

Convex Sets

- Q: Prove that the set of symmetric PSD matrices: $S^n_+ = \{X \in \mathbb{R}^{n \times n} | X \succeq 0, X = X^{\mathsf{T}}\}$ is convex.
- ullet Intersection of convex sets is convex \Longrightarrow can prove the convexity of a set by showing that it is an intersection of convex sets.

Example: We know that a half-space: $\langle a_i, x \rangle \leq b_i$ is a convex set. The set of inequalities $Ax \leq b$ is an intersection of half-spaces and is hence convex.



Zero-order definition: A function f is convex iff its domain \mathcal{D} is a convex set, and for all $x,y\in\mathcal{D}$ and $\theta\in[0,1]$,

$$f(\theta x + (1 - \theta)y) \le \theta f(x) + (1 - \theta) f(y)$$

i.e. the function is below the chord between two points.

- Alternatively, f is convex iff the set formed by the area above the function is a convex set.
 Examples of convex functions:
 - All *p*-norms $||x||_p$ with $p \ge 1$.
 - $f(x) = 1/\sqrt{x}$, $f(x) = -\log(x)$, $f(x) = \exp(-x)$
 - Negative entropy: $f(x) = x \log(x)$
 - Logistic loss: $f(x) = \log(1 + \exp(-x))$
 - Linear functions $f(x) = \langle a, x \rangle$

First-order condition: If f is differentiable, it is convex iff its domain \mathcal{D} is a convex set and for all $x, y \in \mathcal{D}$,

$$f(y) \ge f(x) + \langle \nabla f(x), y - x \rangle$$

i.e. the function is above the tangent to the function at any point x.

For a convex f, consider w^* such that $\nabla f(w^*) = 0$, then using convexity, for all $y \in \mathcal{D}$, $f(y) \geq f(w^*)$. If w^* is a stationary point i.e. $\|\nabla f(w^*)\|^2 = 0$, then it is a global minimum. Hence, local optimization to make the gradient zero results in convergence to a global minimum!

Q: For a convex f, if $\nabla f(w^*) = 0$, then is w^* a unique minimizer of f?

Second-order condition: If f is twice differentiable, it is convex iff its domain \mathcal{D} is a convex set and for all $x \in \mathcal{D}$,

$$\nabla^2 f(x) \succeq 0$$

i.e. the Hessian is positive semi-definite ("curved upwards") for all x.

Q: Prove that $f(x) = \max_i x_i$ is a convex function.

$$f\left(\theta x + (1-\theta)y\right) = \max_{i} [\theta x_i + (1-\theta)y_i] \leq \theta \max_{i} x_i + (1-\theta) \max_{i} y_i = \theta f(x) + (1-\theta)f(y)$$

Hence, by using the zero-order definition of convexity, f(x) is convex.

Q: Prove that $f(x) = \frac{1}{2}x^2$ is a convex function.

$$f(y) - f(x) - \langle \nabla f(x), y - x \rangle = \frac{y^2}{2} - \frac{x^2}{2} - x(y - x) = \frac{1}{2} \left[y^2 + x^2 - 2xy \right] = \frac{(x - y)^2}{2} \ge 0$$

Hence, by using the first-order definition of convexity, f(x) is convex.

Q: Prove that $f(x) = \log(1 + \exp(-x))$ is a convex function.

$$f'(x) = \frac{-\exp(-x)}{1 + \exp(-x)} = \frac{-1}{1 + \exp(x)}$$
$$f''(x) = \frac{\exp(x)}{(1 + \exp(x))^2} > 0$$

Hence, by using the second-order definition of convexity, f(x) is convex.

Q: Prove that the ridge regression loss function: $f(w) = \frac{1}{2} \|Xw - y\|^2 + \frac{\lambda}{2} \|w\|^2$ is convex Recall that $\nabla^2 f(w) = X^\mathsf{T} X + \lambda I_d$. For vector v, let us consider $v^\mathsf{T} \nabla^2 f(w) v$,

$$v^{\mathsf{T}}\nabla^{2}f(w)v = v^{\mathsf{T}}[X^{\mathsf{T}}X + \lambda I_{d}]v = v^{\mathsf{T}}[X^{\mathsf{T}}X]v + \lambda v^{\mathsf{T}}v = [Xv]^{\mathsf{T}}[Xv] + \lambda \|v\|^{2} = \|Xv\|^{2} + \lambda \|v\|^{2}$$

$$\implies v^{\mathsf{T}}\nabla^{2}f(w)v \geq 0 \implies \nabla^{2}f(w) \succeq 0.$$

Hence, by using the second-order definition of convexity, f(w) is convex.

Operations that preserve convexity: if f(x) and g(x) are convex functions, then h(x) is convex if,

- $h(x) = \alpha f(x)$ for $\alpha \ge 0$ (Non-negative scaling) E.g. For $w \in R^d$, $f(w) = \|w\|^2$ is convex, and hence $h(w) = \frac{\lambda}{2} \|w\|^2$ for $\lambda \ge 0$ is convex.
- $h(x) = \max\{f(x), g(x)\}$ (Point-wise maximum) E.g. f(w) = 0 and g(w) = 1 - w are convex functions, and hence $h(w) = \max\{0, 1 - w\}$ is convex.
- h(x) = f(Ax + b) (Composition with affine map) E.g.: $f(w) = \max\{0, 1 - w\}$ is convex, and hence $h(w) = \max\{0, 1 - y_i \langle w, x_i \rangle\}$ for $x_i \in \mathbb{R}^d$ and $y_i \in \mathbb{R}$ is convex
- h(x) = f(x) + g(x) (Sum) E.g.: $f(w) = \max\{0, 1 - y_i \langle w, x_i \rangle\}$ is convex, and hence $h(w) = \sum_{i=1}^n \max\{0, 1 - y_i \langle w, x_i \rangle\} + \frac{\lambda}{2} \|w\|^2$ is convex.

Hence, the SVM loss function: $f(w) := \sum_{i=1}^n \max\{0, 1 - y_i \langle X_i, w \rangle\} + \frac{\lambda}{2} \|w\|^2$ is convex.

Q: Prove that ℓ_1 -regularized logistic regression:

$$f(w) := \sum_{i=1}^{n} \log (1 + \exp(-y_i \langle X_i, w \rangle)) + \lambda \|w\|_1$$
 is convex.

We have proved that the logistic loss $f(x) = \log(1 + \exp(-x))$ is convex. Since composition with an affine map is convex, and the sum of convex functions is convex, the first term is convex. Since all norms are convex, and a non-negative scaling of a convex function is convex, the second term is convex. Hence, f(w) is convex.

Another way to prove convexity for logistic regression is to compute the Hessian and show that it is positive semi-definite (In Assignment 1)

Jensen's Inequality

- Recall the zero-order definition of convexity: $\forall x, y \in \mathcal{D}$ and $\theta \in [0, 1]$, $f(\theta x + (1 \theta)y) < \theta f(x) + (1 \theta)f(y)$.
- This can be generalized to *n* points $\{x_1, x_2, \dots, x_n\}$, i.e. for $p_i \ge 0$ and $\sum_i p_i = 1$,

$$f(p_1 x_1 + p_2 x_2 + ... + p_n x_n) \le p_1 f(x_1) + p_2 f(x_2) + ... + p_n f(x_n) \implies f\left(\sum_{i=1}^n p_i x_i\right) \le \sum_{i=1}^n p_i f(x_i)$$

• If X is a discrete r.v. that can take value x_i with probability p_i , and f is convex, then,

$$f\left(\mathbb{E}[X]\right) \leq \mathbb{E}\left[f(X)\right].$$
 (Jensen's inequality)

- Jensen's inequality can be used to prove inequalities like the AM-GM inequality: $\sqrt{ab} \leq \frac{a+b}{2}$.
- Proof: Choose $f(x) = -\log(x)$ as the convex function, and consider two points a and b with $\theta = 1/2$. By Jensen's inequality,

$$-\log\left(\frac{a+b}{2}\right) \leq \frac{-\log(a)-\log(b)}{2} \implies \log\left(\frac{a+b}{2}\right) \geq \log(\sqrt{ab}) \implies \frac{a+b}{2} \geq \sqrt{ab}.$$

Holder's Inequality

Q: Prove Holder's inequality, for $p,q \geq 1$ s.t. $\frac{1}{p} + \frac{1}{q} = 1$ and $x,y \in R^n$, $|\langle x,y \rangle| \leq \|x\|_p \|y\|_q$.

By repeating the AM-GM proof, but for a general $\theta \in [0,1]$, for $a,b \geq 0$, we can prove that,

$$a^{\theta}b^{1-\theta} \leq \theta a + (1-\theta)b$$

Use
$$a = \frac{|x_i|^p}{\sum_{j=1}^n |x_j|^p}$$
, $b = \frac{|y_i|^q}{\sum_{j=1}^n |y_j|^q}$, $\theta = 1/p$, and using the fact that $1 - \theta = 1 - 1/p = 1/q$

$$\left(\frac{|x_i|^p}{\sum_{j=1}^n |x_j|^p}\right)^{1/p} \left(\frac{|y_i|^q}{\sum_{j=1}^n |y_j|^q}\right)^{1/q} \le \frac{1}{p} \frac{|x_i|^p}{\sum_{j=1}^n |x_j|^p} + \frac{1}{q} \frac{|y_i|^p}{\sum_{j=1}^n |y_j|^p}$$

Summing both sides from i=1 to n and using the fact that $\frac{1}{p}+\frac{1}{q}=1$

$$\sum_{i=1}^{n} \frac{|x_{i}|}{\left(\sum_{j=1}^{n} |x_{j}|^{p}\right)^{1/p}} \frac{|y_{i}|}{\left(\sum_{j=1}^{n} |y_{j}|^{q}\right)^{1/q}} \leq 1 \implies \sum_{i} |x_{i}y_{i}| \leq \left(\sum_{i=1}^{n} |x_{i}|^{p}\right)^{1/p} \left(\sum_{i=1}^{n} |y_{i}|^{q}\right)^{1/q}$$

$$\implies |\langle x, y \rangle| \leq ||x||_{p} ||y||_{q} \qquad \text{(Triangle inequality)}$$



References i



Atilim Gunes Baydin, Robert Cornish, David Martinez Rubio, Mark Schmidt, and Frank Wood, *Online learning rate adaptation with hypergradient descent*, arXiv preprint arXiv:1703.04782 (2017).