CMPT 409/981: Optimization for Machine Learning

Lecture 19

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Recap

• Scalar AdaGrad:

$$w_{k+1} = \Pi_C[w_k - \eta_k \nabla f_k(w_k)]$$
 ; $\eta_k = \frac{\eta}{\sqrt{\sum_{s=1}^k \|\nabla f_s(w_s)\|^2}}$

• We proved that if the convex set $\mathcal C$ has diameter D i.e. for all $x,y\in\mathcal C$, $\|x-y\|\leq D$, for an arbitrary sequence of losses such that each f_k is convex, differentiable and G-Lipschitz, scalar AdaGrad with $\eta_k = \frac{\eta}{\sqrt{\sum_{s=1}^k \|\nabla f_s(w_s)\|^2}}$ and $w_1 \in \mathcal C$ has the following regret for all $u \in \mathcal C$,

$$R_T(u) \le \left(\frac{D^2}{2\eta} + \eta\right) G\sqrt{T}$$

- Unlike OGD, scalar AdaGrad does not require the knowledge of G.
- Scalar AdaGrad uses one step-size for each coordinate. In practice, using one step-size per coordinate results in better empirical performance.

- Let us consider the more practical variants of AdaGrad.
- The corresponding update is similar to preconditioned GD with the preconditioner A_k^{-1} :

$$v_{k+1} = w_k - \eta A_k^{-1} \nabla f_k(w_k) \quad ; \quad w_{k+1} = \Pi_{\mathcal{C}}^k[v_{k+1}] := \arg\min_{w \in \mathcal{C}} \frac{1}{2} \|w - v_{k+1}\|_{A_k}^2 \ .$$

$$A_k = \begin{cases} \sqrt{\sum_{s=1}^k \|\nabla f_s(w_s)\|^2} I_d \quad \text{(Scalar AdaGrad)} \\ \operatorname{diag}(G_k^{\frac{1}{2}}) \quad \text{(Diagonal AdaGrad)} \\ G_k^{\frac{1}{2}} \quad \text{(Full-Matrix AdaGrad)} \end{cases}$$

where $G_k \in \mathbb{R}^{d \times d} := \sum_{s=1}^k [\nabla f_s(w_s) \nabla f_s(w_s)^{\mathsf{T}}].$

• For the commonly-used diagonal variant, AdaGrad results in a per-coordinate update, i.e. $\forall i \in [d]$, if $g_{k,i} := [\nabla f_k(w_k)]_i$, then,

$$v_{k+1}[i] = w_k[i] - \eta \frac{g_{k,i}}{\sqrt{\sum_{s=1}^k g_{s,i}^2}} \quad ; \quad w_{k+1} = \arg\min_{w \in \mathcal{C}} \left[\sum_{i=1}^d \sqrt{\sum_{s=1}^k g_{s,i}^2 \left(w[i] - v_{k+1}[i] \right)^2} \right]$$

• We will assume that A_k is invertible (a small ϵI_d can be added to ensure invertibility).

Claim: If the convex set \mathcal{C} has diameter D, for an arbitrary sequence of losses such that each f_k is convex and differentiable, AdaGrad with the general update $w_{k+1} = \Pi_{\mathcal{C}}^k[w_k - \eta A_k^{-1} \nabla f_k(w_k)]$ and $w_1 \in \mathcal{C}$ has the following regret for $u \in \mathcal{C}$,

$$R_{\mathcal{T}}(u) \leq \left(\frac{D^2}{2\eta} + \eta\right) \operatorname{Tr}[A_{\mathcal{T}}]$$

Proof: Starting from the update, $v_{k+1} = w_k - \eta A_k^{-1} \nabla f_k(w_k)$,

$$v_{k+1} - u = w_k - \eta A_k^{-1} \nabla f_k(w_k) - u \implies A_k[v_{k+1} - u] = A_k[w_k - u] - \eta \nabla f_k(w_k)$$

Multiplying the above equations,

$$[v_{k+1} - u]^{\mathsf{T}} A_k [v_{k+1} - u] = [w_k - u - \eta A_k^{-1} \nabla f_k(w_k)]^{\mathsf{T}} [A_k [w_k - u] - \eta \nabla f_k(w_k)]$$

$$\|v_{k+1} - u\|_{A_k}^2 = \|w_k - u\|_{A_k}^2 - 2\eta \langle \nabla f_k(w_k), w_k - u \rangle + \eta^2 [A_k^{-1} \nabla f_k(w_k)]^{\mathsf{T}} [\nabla f_k(w_k)]$$

$$\implies \|v_{k+1} - u\|_{A_k}^2 = \|w_k - u\|_{A_k}^2 - 2\eta \langle \nabla f_k(w_k), w_k - u \rangle + \eta^2 \|\nabla f_k(w_k)\|_{A_k^{-1}}^2$$

Recall that $\|v_{k+1} - u\|_{A_k}^2 = \|w_k - u\|_{A_k}^2 - 2\eta \langle \nabla f_k(w_k), w_k - u \rangle + \eta^2 \|\nabla f_k(w_k)\|_{A_k^{-1}}^2$. Using the update $w_{k+1} = \Pi_{\mathcal{C}}^k[v_{k+1}]$, $u \in \mathcal{C}$ with the non-expansiveness of projections,

$$\|w_{k+1} - u\|_{A_{k}}^{2} = \|\Pi_{\mathcal{C}}[v_{k+1}] - \Pi_{\mathcal{C}}[u]\|_{A_{k}}^{2} \le \|v_{k+1} - u\|_{A_{k}}^{2}$$

$$\implies \|w_{k+1} - u\|_{A_{k}}^{2} \le \|w_{k} - u\|_{A_{k}}^{2} - 2\eta \langle \nabla f_{k}(w_{k}), w_{k} - u \rangle + \eta^{2} \|\nabla f_{k}(w_{k})\|_{A_{k}^{-1}}^{2}$$

$$\le \|w_{k} - u\|_{A_{k}}^{2} - 2\eta [f_{k}(w_{k}) - f_{k}(u)] + \eta^{2} \|\nabla f_{k}(w_{k})\|_{A_{k}^{-1}}^{2} \quad \text{(Convexity)}$$

$$\implies f_{k}(w_{k}) - f_{k}(u) \le \frac{\|w_{k} - u\|_{A_{k}}^{2} - \|w_{k+1} - u\|_{A_{k}}^{2}}{2\eta} + \frac{\eta}{2} \|\nabla f_{k}(w_{k})\|_{A_{k}^{-1}}^{2}$$

Summing from k = 1 to T,

$$\implies R_{T}(u) \leq \frac{1}{2\eta} \underbrace{\sum_{k=1}^{T} \left[\|w_{k} - u\|_{A_{k}}^{2} - \|w_{k+1} - u\|_{A_{k}}^{2} \right]}_{\text{Term (i)}} + \frac{\eta}{2} \sum_{k=1}^{T} \|\nabla f_{k}(w_{k})\|_{A_{k}^{-1}}^{2}$$

Let us now bound Term (i).

Term (i)
$$= \sum_{k=1}^{T} \left[\| w_k - u \|_{A_k}^2 - \| w_{k+1} - u \|_{A_k}^2 \right]$$

$$= \sum_{k=2}^{T} \left[(w_k - u)^T [A_k - A_{k-1}] (w_k - u)] + \| w_1 - u \|_{A_1}^2 - \| w_{T+1} - u \|_{A_T}^2$$

$$\leq \sum_{k=2}^{T} \| w_k - u \|^2 \ \lambda_{\max} [A_k - A_{k-1}] + \| w_1 - u \|_{A_1}^2 \leq \sum_{k=2}^{T} D^2 \lambda_{\max} [A_k - A_{k-1}] + \| w_1 - u \|_{A_1}^2$$

$$(\text{Since } A_{k-1} \leq A_k, \ \lambda_{\max} [A_k - A_{k-1}] \geq 0 \text{ and } \| w_k - u \|^2 \leq D)$$

$$\implies \sum_{k=1}^{T} \left[\| w_k - u \|_{A_k}^2 - \| w_{k+1} - u \|_{A_k}^2 \right] \leq D^2 \sum_{k=2}^{T} \text{Tr}[A_k - A_{k-1}] + \| w_1 - u \|_{A_1}^2$$

$$(\text{For any PSD matrix } B, \lambda_{\max}[B] \leq \text{Tr}[B])$$

Continuing the proof from the previous slide,

$$\begin{aligned} \text{Term (i)} &= \sum_{k=1}^{T} \left[\| w_k - u \|_{A_k}^2 - \| w_{k+1} - u \|_{A_k}^2 \right] \leq D^2 \sum_{k=2}^{T} \text{Tr}[A_k - A_{k-1}] + \| w_1 - u \|_{A_1}^2 \\ &= D^2 \text{ Tr} \left[\sum_{k=2}^{T} [A_k - A_{k-1}] \right] + \| w_1 - u \|_{A_1}^2 \qquad \qquad \text{(Linearity of Trace)} \\ &= D^2 \text{ Tr}[A_T - A_1] + \| w_1 - u \|_{A_1}^2 \leq D^2 \text{ Tr}[A_T - A_1] + \lambda_{\max}[A_1] \| w_1 - u \|^2 \\ \Longrightarrow \text{ Term (i)} \leq D^2 \text{ Tr}[A_T] - D^2 \text{ Tr}[A_1] + D^2 \text{ Tr}[A_1] = D^2 \text{ Tr}[A_T] \end{aligned}$$

Putting everything together,

$$R_{\mathcal{T}}(u) \leq \frac{D^2 \operatorname{Tr}[A_{\mathcal{T}}]}{2\eta} + \frac{\eta}{2} \underbrace{\sum_{k=1}^{T} \left\| \nabla f_k(w_k) \right\|_{A_k^{-1}}^2}_{\operatorname{Term}(ii)}$$

Let us now bound Term (ii).

Claim: Term (ii) = $\sum_{k=1}^{T} \|\nabla f_k(w_k)\|_{A_k^{-1}}^2 \le 2 \operatorname{Tr}[A_T]$

Proof: Let us prove by induction. For convenience, define $g_k := \nabla f_k(w_k)$.

Base case: For k = 1, LHS = $\text{Tr}[g_1^{\mathsf{T}}A_1^{-1}g_1] = \text{Tr}[A_1^{-1}g_1g_1^{\mathsf{T}}] = \text{Tr}[A_1^{-1}A_1A_1] \leq 2 \text{Tr}[A_1] = \text{RHS}.$

Here, we used the cyclic property of trace i.e. Tr[ABC] = Tr[BCA].

Inductive Hypothesis: If the statement is true for T-1, we need to prove it for T.

$$\sum_{k=1}^{T-1} \|g_k\|_{A_k^{-1}}^2 + \|g_T\|_{A_T^{-1}}^2 \le 2\operatorname{Tr}[A_{T-1}] + \|g_T\|_{A_T^{-1}}^2 = 2\operatorname{Tr}[\left(A_T^2 - g_T g_T^{\mathsf{T}}\right)^{1/2}] + \operatorname{Tr}[A_T^{-1} g_T g_T^{\mathsf{T}}]$$

For any $X \succeq Y \succeq 0$, we have [DHS11, Lemma 8], $2 \operatorname{Tr}[(X - Y)^{1/2}] + \operatorname{Tr}[X^{-1/2}Y] \leq 2 \operatorname{Tr}[X^{1/2}]$. Using this for $X = A_T^2$, $Y = g_T g_T^T$, $\sum_{k=1}^T \|g_k\|_{A_k^{-1}}^2 \leq 2 \operatorname{Tr}[A_T]$, which completes the proof.

Putting everything together,

$$R_{\mathcal{T}}(u) \leq \left(\frac{D^2}{2\eta} + \eta\right) \operatorname{Tr}[A_{\mathcal{T}}].$$

Diagonal AdaGrad vs OGD

- We have proved that for both the diagonal and full-matrix variants of AdaGrad, $R_T(u) \leq \left(\frac{D^2}{2n} + \eta\right) \operatorname{Tr}[A_T].$
- By doing a tighter analysis for the diagonal variant, we can prove that the corresponding regret bound is: $R_T(u) \leq \left(\frac{D_\infty^2}{2\eta} + \eta\right) \operatorname{Tr}[A_T]$ where $D_\infty = \max_{x,y \in \mathcal{C}} \|x y\|_\infty$. Setting $\eta = \frac{D_\infty}{\sqrt{2}}$, $R_T(u) \leq \sqrt{2}D_\infty \sum_{i=1}^d \sqrt{\sum_{k=1}^T g_{k,i}^2}$.
- Compare the above bound to the regret for OGD (with $\eta = D/\sqrt{2}G$), $R_T(u) \leq \sqrt{2} D \sqrt{\sum_{i=1}^d \sum_{k=1}^T g_{k,i}^2}$ where $D = \max_{x,y \in \mathcal{C}} \|x y\|_2$.
- If $\mathcal C$ is the unit hypercube, then, $D=\sqrt{d}$ and $D_\infty=1$. If the gradients are sparse (e.g. corresponding to one-hot features for logistic regression), diagonal AdaGrad will result in a better regret bound than OGD.
- For other convex sets, such as the Euclidean ball, and when the gradients are dense, the regret of OGD can be better than that of diagonal AdaGrad.

Recall that $R_T(u) \leq \left(\frac{D^2}{2\eta} + \eta\right) \text{Tr}[A_T]$. In the worst-case, $\text{Tr}[A_T] \leq \sqrt{d} \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2}$.

$$\operatorname{Tr}[A_T] = \operatorname{Tr}[G_T^{\frac{1}{2}}] = \sum_{j=1}^d \sqrt{\lambda_j[G_T]} = d \, \frac{\sum_{j=1}^d \sqrt{\lambda_j[G_T]}}{d} \leq d \, \sqrt{\frac{\sum_{j=1}^d \lambda_j[G_T]}{d}}$$

(Jensen's inequality for \sqrt{x})

$$= \sqrt{d} \sqrt{\sum_{j=1}^{d} \lambda_{j}[G_{T}]} = \sqrt{d} \sqrt{\text{Tr}[G_{T}]} = \sqrt{d} \sqrt{\text{Tr}\left[\sum_{k=1}^{T} \nabla f_{k}(w_{k}) \nabla f_{k}(w_{k})^{\mathsf{T}}\right]}$$

$$\text{Tr}[A_{T}] \leq \sqrt{d} \sqrt{\left[\sum_{k=1}^{T} \text{Tr} \nabla f_{k}(w_{k}) \nabla f_{k}(w_{k})^{\mathsf{T}}\right]} = \sqrt{d} \sqrt{\sum_{k=1}^{T} \|\nabla f_{k}(w_{k})\|^{2}} \quad \text{(Linearity of Trace)}$$

Putting everything together, in the worst-case, the regret can be bounded as:

$$R_T(u) \leq \left(\frac{D^2}{2\eta} + \eta\right) \sqrt{d} \sqrt{\sum_{k=1}^T \left\|\nabla f_k(w_k)\right\|^2}$$

AdaGrad - Convex, Lipschitz functions

Claim: If the convex set C has diameter D, for an arbitrary sequence of losses such that each f_k is convex, differentiable and G-Lipschitz, AdaGrad with the general update

 $w_{k+1} = \Pi_{\mathcal{C}}^k[w_k - \eta A_k^{-1} \nabla f_k(w_k)]$ with $\eta = \frac{D}{\sqrt{2}}$ and $w_1 \in \mathcal{C}$ has the following regret for $u \in \mathcal{C}$,

$$R_T(u) \leq \sqrt{2}DG\sqrt{d}\sqrt{T}$$

Proof: Using the general result for AdaGrad and that each f_k is G-Lipschitz,

$$\begin{aligned} R_T(u) &\leq \left(\frac{D^2}{2\eta} + \eta\right) \sqrt{d} \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2} \leq \left(\frac{D^2}{2\eta} + \eta\right) \sqrt{d} \, G \sqrt{T} \\ R_T(u) &\leq \sqrt{2} DG \sqrt{d} \sqrt{T} \end{aligned} \tag{Setting } \eta = \frac{D}{\sqrt{2}})$$

- Unlike scalar AdaGrad, when using the diagonal or full-matrix variant, the worst-case regret
 has a dimension dependence.
- Similar to scalar AdaGrad, we can derive regret bounds for the strongly-convex Lipschitz and smooth convex losses.



Adaptive Gradient Methods

Update for a generic method: For $k \ge 1$ with $m_0 := 0$, $\beta \ge 0$,

$$w_{k+1} = \Pi_{\mathcal{C}}^{k}[w_{k} - \eta_{k} A_{k}^{-1} m_{k}]; \qquad m_{k} = \beta m_{k-1} + (1 - \beta) \nabla f_{k}(w_{k})$$
 where, $\Pi_{\mathcal{C}}^{k}[v] := \underset{w \in \mathcal{C}}{\arg\min} \frac{1}{2} \|w - v\|_{A_{k}}^{2}$.

Instantiating the generic method:

- **SGD**: $A_k = I_d$, $\beta = 0$. Resulting update: $w_{k+1} = w_k \eta_k \nabla f_k(w_k)$.
- Stochastic Heavy-Ball Momentum: $A_k = I_d$. For $\alpha_k = \eta_k (1 \beta)$ and $\gamma_k = \frac{\beta \eta_k}{\eta_{k-1}}$, Resulting update: $w_{k+1} = w_k \alpha_k \nabla f_k(w_k) + \gamma_k(w_k w_{k-1})$ (Prove in Assignment 4!)
- AdaGrad: $A_k = G_k^{\frac{1}{2}}$ where $G_0 = 0$ and $G_k = G_{k-1} + \nabla f_k(w_k) \nabla f_k(w_k)^{\mathsf{T}}$, $\beta = 0$, $\eta_k = \eta$. Resulting update: $w_{k+1} = w_k \eta A_k^{-1} \nabla f_k(w_k)$.
- Adam: $A_k = G_k^{\frac{1}{2}}$ where $G_0 = 0$ and $G_k = \beta_2 G_{k-1} + (1 \beta_2) \nabla f_k(w_k) \nabla f_k(w_k)^\mathsf{T}$, $\beta = \beta_1$ for $\beta_1, \beta_2 \in (0, 1)$. Resulting update: $w_{k+1} = w_k \eta_k A_k^{-1} m_k$ where $m_k = \beta_1 m_{k-1} + (1 \beta_1) \nabla f_k(w_k)$.

Adam

- Recall the update: $w_{k+1} = \prod_{\mathcal{C}}^k [w_k \eta_k A_k^{-1} m_k]$; $m_k = \beta m_{k-1} + (1 \beta) \nabla f_k(w_k)$.
- For Adam, $G_k = (1 \beta_2) \sum_{i=1}^k \beta_2^{k-i} [\nabla f_i(w_i) \nabla f_i(w_i)^{\mathsf{T}}]$ and $m_k = (1 \beta_1) \sum_{i=1}^k \beta_1^{k-i} [\nabla f_i(w_i)].$

Hence, the influence of the past gradients is decayed exponentially which ensures that G_k and m_k are both primarily influenced by the most recent gradient $\nabla f_k(w_k)$. This results in better empirical performance.

• Consider scalar Adam for which $G_k = (1 - \beta_2) \sum_{i=1}^k \beta_2^{k-i} \|\nabla f_i(w_i)\|^2$. Unlike scalar AdaGrad (for which $G_k = \sum_{i=1}^k \|\nabla f_i(w_i)\|^2$), G_k is not guaranteed to increase monotonically (i.e. $G_{k+1} > G_k$). Hence the "effective step-size" $\tilde{\eta}_k$ equal to $\frac{\eta}{\sqrt{G_k}}$ is not guaranteed to decrease.

Hence, to ensure convergence, Adam requires $\eta_k = \tilde{\eta_k} \alpha_k$ for some decreasing sequence α_k . The original paper [KB14] claimed convergence for $\eta_k = O(1/\sqrt{k})$, $\beta_2 \in [0,1)$ and $\beta_1 \in [0,1)$.

• However, the non-monotonic behaviour of G_k can result in non-convergence of Adam even with an explicitly decreasing sequence of η_k , constant $\beta_2 \in (0,1)$ and $\beta_1 = 0$ (no momentum).

- We will construct an example on which Adam can result in linear regret in the online setting (and is hence not guaranteed to converge to the minimizer in the stochastic setting) [RKK19].
- For C>2, run Adam with $\beta_1=0$ (no momentum), $\beta_2=\frac{1}{1+C^2}$ and $\eta_k=\frac{\eta}{\sqrt{k}}$ such that $\eta<\sqrt{1-\beta_2}$ on the following problem:
- ullet Consider $\mathcal{C} = [-1,1]$ and the following sequence of linear functions.

$$f_k(w) = \begin{cases} C & w \text{ for } k \text{ mod } 3 = 1 \\ -w & \text{otherwise} \end{cases}$$

Update: $w_1 = 1$ and for $k \ge 1$,

$$v_{k+1} := w_k - \frac{\eta_k}{\sqrt{\beta_2 G_{k-1} + (1 - \beta_2) \left\| \nabla f_k(w_k) \right\|^2}} \nabla f_k(w_k) \text{ and } w_{k+1} = \Pi_{[-1,1]}[v_{k+1}]$$

• We will compare Adam to the "best" fixed decision (w^*) that minimizes the regret. To compute w^* , consider the sequence of 3 functions from iteration 3k to 3k + 2 for $k \ge 0$. In this case,

$$w^* := \arg\min_{[-1,1]} \left[f_{3k}(w) + f_{3k+1}(w) + f_{3k+2}(w) \right] = \arg\min_{[-1,1]} \left[(C-2)w \right] = -1 \quad \text{(Since } C > 2)$$

Claim: For Adam's iterates, for $k \ge 0$, for all $i \le [3k+1]$, $w_i > 0$ and $w_{3k+1} = 1$.

Proof: Let us prove the statement by induction. Base case: For k = 0, $w_{3k+1} = w_1 = 1$.

Inductive hypothesis: Assume that for $i \leq [3k+1]$, $w_i > 0$ and $w_{3k+1} = 1$. We need to prove that (a) $w_{3k+2} > 0$, (b) $w_{3k+3} > 0$ and (c) $w_{3k+4} = 1$.

In order to show this, note that $\nabla f_i(w) = C$ for i mod 3 = 1 and $\nabla f_i(w) = -1$ otherwise.

Consider the update at iteration (3k+1). By the induction hypothesis, we know that $w_{3k+1}=1$.

$$\begin{aligned} v_{3k+2} &= w_{3k+1} - \left[\frac{\eta_{3k+1}}{\sqrt{\beta_2 \, G_{3k} + (1 - \beta_2) \, \|\nabla f_{3k+1}(w_{3k+1})\|^2}} \, \nabla f_{3k+1}(w_{3k+1}) \right] \\ &= 1 - \left[\frac{C\eta}{\sqrt{(3k+1) \, (\beta_2 \, G_{3k} + (1 - \beta_2) \, C^2)}} \right] \qquad \text{(Using the value of } \eta_{3k+1}) \\ &\geq 1 - \left[\frac{C\eta}{\sqrt{(3k+1) \, (1 - \beta_2) \, C^2}} \right] = 1 - \left[\frac{\eta}{\sqrt{(3k+1) \, (1 - \beta_2)}} \right] \quad \text{(Since } G_{3k} \geq 0) \\ \implies v_{3k+2} \geq 1 - \frac{1}{\sqrt{3k+1}} > 0 \qquad \qquad \text{(Since } \eta < \sqrt{1 - \beta_2} \text{ and } k \geq 1) \end{aligned}$$

Since
$$\left[\frac{C\eta}{\sqrt{(3k+1)(\beta_2 \ G_{3k}+(1-\beta_2)C^2)}}\right] > 0$$
, $v_{3k+2} < 1$. Since $v_{3k+2} \in (0,1)$, $w_{3k+2} = v_{3k+2} < 1$ which proves (a).

• For the update at iteration (3k+2), since $\nabla f_{3k+2}(w) = -1$ for all w,

$$v_{3k+3} = w_{3k+2} + \left[\frac{\eta}{\sqrt{(3k+2)(\beta_2 G_{3k+1} + (1-\beta_2))}} \right]$$

Since $w_{3k+2} \in (0,1)$ and $\frac{\eta}{\sqrt{(3k+2)(\beta_2 G_{3k+1} + (1-\beta_2))}} > 0$, $v_{3k+3} > 0$ and hence $w_{3k+3} > 0$ which proves (b).

• In order to prove (c), consider iteration 3k + 3. Since $\nabla f_{3k+3}(w) = -1$ for all w,

$$v_{3k+4} = w_{3k+3} + \left[\frac{\eta}{\sqrt{(3k+3)(\beta_2 G_{3k+2} + (1-\beta_2))}} \right]$$

From the above update, we can conclude that $v_{3k+4} > w_{3k+3}$.

To prove (c), we will show that $v_{3k+4} \ge 1$ and hence $w_{3k+4} = \Pi_{[-1,1]}v_{3k+4} = 1$. For this, we consider two cases – when $v_{3k+3} \ge 1$ or when $v_{3k+3} < 1$.

Case 1: When $v_{3k+3} \ge 1 \implies w_{3k+3} = 1 \implies v_{3k+4} \ge 1 \implies w_{3k+4} = 1$.

Case 2: When $v_{3k+3} \le 1 \implies w_{3k+3} = v_{3k+3} \le 1$. Combining iterations (3k+4) and (3k+3),

$$v_{3k+4} = v_{3k+3} + \left[\frac{\eta}{\sqrt{(3k+3)(\beta_2 G_{3k+2} + (1-\beta_2))}} \right]$$

$$= w_{3k+2} + \left[\frac{\eta}{\sqrt{(3k+2)(\beta_2 G_{3k+1} + (1-\beta_2))}} \right] + \left[\frac{\eta}{\sqrt{(3k+3)(\beta_2 G_{3k+2} + (1-\beta_2))}} \right]$$

$$= 1 - \left[\frac{C\eta}{\sqrt{(3k+1)(\beta_2 G_{3k} + (1-\beta_2)C^2)}} \right] \qquad \text{(Since } v_{3k+2} = w_{3k+2} \text{ and } w_{3k+1} = 1)$$

$$+ \left[\frac{\eta}{\sqrt{(3k+2)(\beta_2 G_{3k+1} + (1-\beta_2))}} \right] + \left[\frac{\eta}{\sqrt{(3k+3)(\beta_2 G_{3k+2} + (1-\beta_2))}} \right]$$

In order to show that $v_{3k+4} \ge 1$, it is sufficient to show that $T_1 \le T_2$.

Recall from Slide 6, $T_1 \leq \left[\frac{\eta}{\sqrt{(3k+1)(1-\beta_2)}}\right]$. Let us lower-bound T_2 .

$$T_{2} := \left[\frac{\eta}{\sqrt{(3k+2)(\beta_{2} G_{3k+1} + (1-\beta_{2}))}} \right] + \left[\frac{\eta}{\sqrt{(3k+3)(\beta_{2} G_{3k+2} + (1-\beta_{2}))}} \right]$$

$$\geq \left[\frac{\eta}{\sqrt{(3k+2)(\beta_{2} C^{2} + (1-\beta_{2}))}} \right] + \left[\frac{\eta}{\sqrt{(3k+3)(\beta_{2} C^{2} + (1-\beta_{2}))}} \right]$$
(Since $G_{k} \leq C^{2}$ for all k)

$$= \frac{\eta}{\sqrt{(\beta_2 C^2 + (1 - \beta_2))}} \left[\sqrt{\frac{1}{3k + 2}} + \sqrt{\frac{1}{3k + 3}} \right]$$

$$\geq \frac{\eta}{\sqrt{(\beta_2 C^2 + (1 - \beta_2))}} \left[\sqrt{\frac{1}{2(3k + 1)}} + \sqrt{\frac{1}{2(3k + 1)}} \right] = \frac{\sqrt{2}\eta}{\sqrt{(\beta_2 C^2 + (1 - \beta_2))}} \left[\frac{1}{\sqrt{3k + 1}} \right]$$

$$\implies T_2 \ge \left\lceil \frac{\eta}{\sqrt{(3k+1)(1-\beta_2)}} \right\rceil \ge T_1 \qquad \text{(Since } \beta_2 = \frac{1}{1+C^2} \implies \frac{\beta_2 C^2 + (1-\beta_2)}{2} = 1 - \beta_2 \text{)}$$

Since we have proved that $T_2 \ge T_1$, $v_{3k+4} = 1 - T_1 + T_2 \ge 1 \implies w_{3k+4} = 1$. This completes the induction proof.

Hence, for the Adam iterates, for $k \ge 0$, for all $i \le [3k+1]$, $w_i > 0$ and $w_{3k+1} = 1$. Now that we have bounds on the Adam iterates, let us compute its regret $R_{[3k \to 3k+2]}(w^*)$ w.r.t $w^* = -1$ for iterations 3k to 3k + 2.

$$R_{[3k \to 3k+2]}(w^*) = [f_{3k}(w_{3k}) - f_{3k}(-1)] + [f_{3k+1}(w_{3k+1}) - f_{3k+1}(-1)] + [f_{3k+2}(w_{3k+2}) - f_{3k+2}(-1)]$$

$$= [-w_{3k} - 1] + [C w_{3k+1} + C] + [-w_{3k+2} - 1] > 2C - 4 > 0$$
(Since w_{3k} and w_{3k+2} are in $(0, 1)$, $w_{3k+1} = 1$ and $C > 2$)

- Hence for every three functions, Adam has a regret > 2C 4 and hence $R_T(w^*) = O(T)$.
- Both OGD and AdaGrad achieve sublinear regret when run on this example.

• The example takes advantage of the non-monotonicity in the Adam step-sizes – resulting in smaller updates for $k=1 \mod 3$ (when the gradient is positive and will push the iterates towards -1) and larger updates for the other k (when the gradient is negative and will push the iterates towards 1).

The example can be modified [RKK19] to consider:

- Updates of the form $w_{k+1} = w_k \frac{\eta_k}{\sqrt{G_k + \epsilon}}$ for $\epsilon > 0$.
- Constant η_k (rather than $O(1/\sqrt{k})$).
- Stochastic setting (rather than the more general online convex optimization setup).
- Decreasing, non-zero β_1 (the momentum parameter).
- To bypass such examples where Adam fails to converge, AMSGrad [RKK19] modifies the update to ensure monotonically decreasing step-sizes and prove convergence.
- In the example, as C>2 increases, the regret increases, $\beta_2=\frac{1}{1+C^2}\to 0$. [ZCS⁺22] show that using a "large" β_2 and ensuring that $\beta_1\leq \sqrt{\beta_2}$ (often the choice in practice) can bypass the lower-bound resulting in convergence for Adam (without modifying the update).



AMSGrad – fixing the convergence of Adam

• Since the non-decreasing step-size for Adam is problematic, AMSGrad [RKK19] fixes this issue by making a small modification (in red) to Adam. It has the following update – for $\beta_1,\beta_2\in(0,1)$,

$$G_{k} = \beta_{2} G_{k-1} + (1 - \beta_{2}) \operatorname{diag} \left[\nabla f_{k}(w_{k}) \nabla f_{k}(w_{k})^{\mathsf{T}} \right] ; \quad A_{k} = \max \{ G_{k}^{\frac{1}{2}}, A_{k-1} \}$$

$$w_{k+1} = \Pi_{\mathcal{C}}^{k} [w_{k} - \eta_{k} A_{k}^{-1} m_{k}]; \quad ; \quad m_{k} = \beta_{1} m_{k-1} + (1 - \beta_{1}) \nabla f_{k}(w_{k})$$

$$\Pi_{\mathcal{C}}^{k} [v_{k+1}] := \underset{w \in \mathcal{C}}{\operatorname{arg min}} \frac{1}{2} \| w - v_{k+1} \|_{A_{k}}^{2} ,$$

where $C = \max\{A, B\}$ for diagonal matrices A and B implies that for all $i \in [d]$, $C_{i,i} = \max\{A_{i,i}, B_{i,i}\}$.

• The AMSGrad update ensures that $A_k \succeq A_{k-1}$ and hence the step-sizes η_k are non-increasing, which guarantees convergence.

References i

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