CMPT 409/981: Optimization for Machine Learning

Lecture 3

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September 15, 2022

Recap

For *L*-smooth functions lower-bounded by f^* , gradient descent with $\eta = \frac{1}{L}$ returns an ϵ -approximate stationary point and requires $\Theta\left(\frac{1}{\epsilon}\right)$ iterations.

Importantly, the GD rate does not depend on the dimension of w.

In practice, we can set η_k in an adaptive manner using an exact line-search:

$$\eta_k = \arg\min_{\eta} f(w_k - \eta \nabla f(w_k)).$$

An exact line-search can adapt to the "local" *L*, resulting in larger step-sizes and better performance.

However, we can compute η_k analytically only in special cases, whereas solving the sub-problem approximately to set η_k can be expensive.

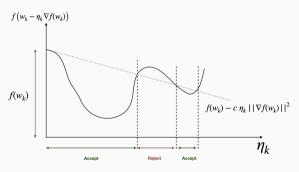
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Usually, the cost of doing an exact line-search is not worth the computational effort.

Armijo condition for a prospective step-size $\tilde{\eta_k}$:

$$f(w_k - \tilde{\eta}_k \nabla f(w_k)) \le f(w_k) - c \, \tilde{\eta}_k \, \|\nabla f(w_k)\|^2$$

where $c \in (0,1)$ is a hyper-parameter.



Algorithm GD with Armijo Line-search

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1: function GD with Armijo line-search(f, w_0, \eta_{\max}, c \in (0,1), \beta \in (0,1))
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2: **for**
$$k = 0, ..., T - 1$$
 do

3:
$$\tilde{\eta}_k \leftarrow \eta_{\text{max}}$$

4: while
$$f(w_k - \tilde{\eta}_k \nabla f(w_k)) > f(w_k) - c \cdot \eta \|\nabla f(w_k)\|^2$$
 do

5:
$$\tilde{\eta}_k \leftarrow \tilde{\eta}_k \beta$$

7:
$$\eta_k \leftarrow \tilde{\eta}_k$$

8:
$$w_{k+1} = w_k - \eta_k \nabla f(w_k)$$

- 9: end for
- 10: **return** w_T

Claim: The (exact) backtracking procedure terminates and returns $\eta_k \ge \min \left\{ \frac{2(1-c)}{L}, \eta_{\text{max}} \right\}$. **Proof**:

$$f(w_{k} - \tilde{\eta}_{k} \nabla f(w_{k})) \leq \underbrace{f(w_{k}) - \|\nabla f(w_{k})\|^{2} \left(\tilde{\eta}_{k} - \frac{L\tilde{\eta}_{k}^{2}}{2}\right)}_{h_{1}(\tilde{\eta}_{k})} \text{ (Quadratic bound using smoothness)}$$

$$f(w_{k} - \tilde{\eta}_{k} \nabla f(w_{k})) \leq \underbrace{f(w_{k}) - \|\nabla f(w_{k})\|^{2} \left(c\tilde{\eta}_{k}\right)}_{\text{(Armijo condition)}} \text{ (Armijo condition)}$$

If the Armijo condition is satisfied, the back-tracking line-search procedure terminates.

Case (i): For
$$\eta_{\max} \leq \frac{2(1-c)}{L}$$
, $f(w_k - \eta_{\max} \nabla f(w_k)) \leq h_1(\eta_{\max}) \leq h_2(\eta_{\max})$ \Longrightarrow if $\eta_{\max} \leq \frac{2(1-c)}{L}$, then the line-search terminates immediately and $\eta_k = \eta_{\max}$.



Case (ii): If
$$\eta_{\text{max}} > \frac{2(1-c)}{L}$$
 and the Armijo condition is satisfied for step-size η_k , then $f(w_k - \eta_k \nabla f(w_k)) \le h_2(\eta_k) \le h_1(\eta_k) \implies c\eta_k \ge \eta_k - \frac{L\eta_k^2}{2} \implies \eta_k \ge \frac{2(1-c)}{L}$.

Putting the two cases together, the step-size η_k returned by the Armijo line-search satisfies $\eta_k \geq \min\left\{\frac{2\,(1-c)}{L},\eta_{\max}\right\}$.

Claim: Gradient Descent with (exact) backtracking Armijo line-search (with c=1/2) returns point \hat{w} such that $\|\nabla f(\hat{w})\|^2 \leq \epsilon$ and requires $T \geq \frac{\max\{2L, \frac{2}{\eta_{\max}}\} [f(w_0) - \min_w f(w)]}{\epsilon}$ oracle calls or iterations.

Proof: Since η_k satisfies the Armijo condition and $w_{k+1} = w_k - \eta_k \nabla f(w_k)$,

$$\begin{split} f(w_{k+1}) & \leq f(w_k) - c \, \eta_k \, \left\| \nabla f(w_k) \right\|^2 \\ & \leq f(w_k) - \left(\min \left\{ \frac{1}{2L}, \frac{\eta_{\mathsf{max}}}{2} \right\} \right) \, \left\| \nabla f(w_k) \right\|^2 \\ & \qquad \qquad \left(\mathsf{Result from previous slide with } \, c = 1/2) \end{split}$$

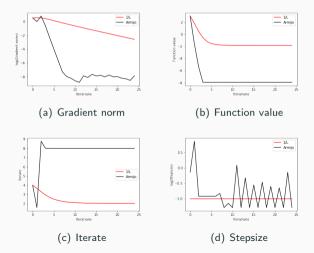
Continuing the proof as before,

$$\implies \|\nabla f(\hat{w})\|^2 \le \frac{\max\{2L, \frac{2}{\eta_{\mathsf{max}}}\}\left[f(w_0) - \min_w f(w)\right]}{T}$$

The claim is proved by the same reasoning as in Lecture 2.

Gradient Descent with Line-search - Examples

 $\min_{x \in [-10,10]} f(x) := -x \sin(x)$. Compare GD (with $x_0 = 4$) with (i) $\eta = 1/L \approx 0.1$ and (ii) Armijo line-search with $\eta_{\max} = 10, c = 1/2, \beta = 0.9$.





Convex Optimization

We have seen that we require $\Theta(1/\epsilon)$ iterations to converge to an ϵ -approximate stationary point for smooth functions. Alternatively, if we care about global optimization (reach the vicinity of the true minimizer) of Lipschitz functions, we require $\Theta(1/\epsilon^d)$ iterations.

Convex functions: Class of functions where local optimization can result in convergence to the global minimizer of the function.

In general, convex optimization involves minimizing a convex function over a convex set \mathcal{C} .

Examples of convex optimization in ML

Ridge regression: $\min_{w \in \mathbb{R}^d} \frac{1}{2} \|Xw - y\|^2 + \frac{\lambda}{2} \|w\|^2$.

Logistic regression: $\min_{w \in \mathbb{R}^d} \sum_{i=1}^n \log (1 + \exp(-y_i \langle X_i, w \rangle))$

Support vector machines: $\min_{w \in \mathbb{R}^d} \sum_{i=1}^n \max\{0, 1 - y_i \langle X_i, w \rangle\} + \frac{\lambda}{2} \|w\|^2$

Planning in MDPs in RL: $\max_{\mu \in \mathcal{F}_{\rho}} \langle \mu, r \rangle$ where \mathcal{F}_{ρ} is the flow-polytope.

Convex Sets

A set $\mathcal C$ is convex if a point along the line joining two points in $\mathcal C$ also lies in the set.

For points x, y, the *convex combination* of x, y is $z := \theta x + (1 - \theta)y$ for $\theta \in [0, 1]$.

A set C is convex iff $\forall x, y \in C$, the convex combination $z \in C$.

Examples of convex sets:

- Positive orthant $\mathbb{R}^d_+: \{x|x \geq 0\}.$
- Hyper-plane: $\{x|Ax = b\}$.
- Half-space: $\{x | Ax \le b\}$.
- Norm-ball: $\{x | \|x\|_p \le r\}$.
- Norm-cone: $\{(x,r)| ||x||_p \le r\}$.

Convex Sets

Q: Prove that the hyper-plane (set of linear equations): $\mathcal{H} := \{x | Ax = b\}$ is a convex set.

If $x, y \in \mathcal{H}$, then, Ax = b and Ay = b. Consider a point $z := \theta x + (1 - \theta)y$ for $\theta \in [0, 1]$.

$$Az = A[\theta x + (1 - \theta)y] = \theta Ax + (1 - \theta)Ay = b.$$

Hence, $z \in \mathcal{H}$ and \mathcal{H} is a convex set.

Q: Prove that the ball of radius r centered at point x_c : $\mathcal{B}(x_c, r) := \{x | \|x - x_c\|_p \le r\}$ is convex.

If $x, y \in \mathcal{B}(x_c, r)$, then, $||x - x_c||_p \le r$ and $||y - x_c||_p \le r$. Consider a point $z := \theta x + (1 - \theta)y$ for $\theta \in [0, 1]$.

$$\begin{split} \left\|z-x_{c}\right\|_{p} &= \left\|\theta(x-x_{c})+(1-\theta)(y-x_{c})\right\|_{p} \\ &\leq \left\|\theta(x-x_{c})\right\|_{p} + \left\|(1-\theta)(y-x_{c})\right\|_{p} & \text{(Triangle inequality for norms)} \\ &\leq \theta\left\|(x-x_{c})\right\|_{p} + (1-\theta)\left\|(y-x_{c})\right\|_{p} & \text{(Homogeneity of norms)} \end{split}$$

$$\implies \|z - x_c\|_p \le r$$

Hence, $z \in \mathcal{B}(x_c, r)$ and $\mathcal{B}(x_c, r)$ is a convex set.

Convex Sets

Q: Prove that the set of symmetric PSD matrices: $S^n_+ = \{X \in \mathbb{R}^{n \times n} | X \succeq 0\}$ is convex.

Ans: If $X \in S_+^n$, for any vector v, $v^T X v \ge 0$. Consider $X, Y \in S_+^n$, and let $Z = \theta X + (1 - \theta)Y$, then, $v^T Z v = \theta \ v^T X v + (1 - \theta) \ v^T Y v \ge 0$, hence $Z \in S_+^n$ and S_+^n is a convex set.

Intersection of convex sets is convex \implies can prove the convexity of a set by showing that it is an intersection of convex sets.

Example: We know that a half-space: $\langle a_i, x \rangle \leq b_i$ is a convex set. The set of inequalities $Ax \leq b$ is an intersection of half-spaces and is hence convex.



Zero-order definition: A function f is convex iff its domain \mathcal{D} is a convex set, and for all $x,y\in\mathcal{D}$ and $\theta\in[0,1]$,

$$f(\theta x + (1 - \theta)y) \le \theta f(x) + (1 - \theta)f(y)$$

i.e. the function is below the chord between two points.

Alternatively, f is convex iff the set formed by the area above the function is a convex set.

Examples of convex functions:

- All norms $||x||_p$
- $f(x) = 1/\sqrt{x}$, $f(x) = -\log(x)$, $f(x) = \exp(-x)$
- Negative entropy: $f(x) = x \log(x)$
- Logistic loss: $f(x) = \log(1 + \exp(-x))$
- Linear functions $f(x) = \langle a, x \rangle$

First-order condition: If f is differentiable, it is convex iff its domain \mathcal{D} is a convex set and for all $x, y \in \mathcal{D}$,

$$f(y) \ge f(x) + \langle \nabla f(x), y - x \rangle$$

i.e. the function is above the tangent to the function at any point x.

For a convex f, consider w^* such that $\nabla f(w^*) = 0$, then using convexity, for all $y \in \mathcal{D}$, $f(y) \geq f(w^*)$. If w^* is a stationary point i.e. $\|\nabla f(w^*)\|^2 = 0$, then it is a global minimum. Hence, local optimization to make the gradient zero results in convergence to a global minimum!

Q: For a convex f, if $\nabla f(w^*) = 0$, then is w^* a unique minimizer of f?

Ans: No, there might many minimizers that all have the same function value

Second-order condition: If f is twice differentiable, it is convex iff its domain \mathcal{D} is a convex set and for all $x \in \mathcal{D}$,

$$\nabla^2 f(x) \succeq 0$$

i.e. the Hessian is positive semi-definite ("curved upwards") for all x.

Q: Prove that $f(x) = \max_i x_i$ is a convex function

$$f\left(\theta x + (1-\theta)y\right) = \max_{i} [\theta x_i + (1-\theta)y_i] \leq \theta \max_{i} x_i + (1-\theta) \max_{i} y_i = \theta f(x) + (1-\theta)f(y)$$

Hence, by using the zero-order definition of convexity, f(x) is convex.

Q: Prove that $f(x) = \frac{1}{2}x^2$ is a convex function

$$f(y) - f(x) - \langle \nabla f(x), y - x \rangle = \frac{y^2}{2} - \frac{x^2}{2} - x(y - x) = \frac{1}{2} \left[y^2 + x^2 - 2xy \right] = \frac{(x - y)^2}{2} \ge 0$$

Hence, by using the first-order condition of convexity, f(x) is convex.

Q: Prove that $f(x) = \log(1 + \exp(-x))$ is a convex function

$$f'(x) = \frac{-\exp(-x)}{1 + \exp(-x)} = \frac{-1}{1 + \exp(x)}$$
$$f''(x) = \frac{\exp(x)}{(1 + \exp(x))^2} > 0$$

Hence, by using the second-order condition of convexity, f(x) is convex.

Q: Prove that the ridge regression loss function: $f(w) = \frac{1}{2} \|Xw - y\|^2 + \frac{\lambda}{2} \|w\|^2$ is convex Recall that $\nabla^2 f(w) = X^\mathsf{T} X + \lambda I_d$. For vector v, let us consider $v^\mathsf{T} \nabla^2 f(w) v$,

$$v^{\mathsf{T}}\nabla^{2}f(w)v = v^{\mathsf{T}}[X^{\mathsf{T}}X + \lambda I_{d}]v = v^{\mathsf{T}}[X^{\mathsf{T}}X]v + \lambda v^{\mathsf{T}}v = [Xv]^{\mathsf{T}}[Xv] + \lambda \|v\|^{2} = \|Xv\|^{2} + \lambda \|v\|^{2}$$

$$\implies v^{\mathsf{T}}\nabla^{2}f(w)v \geq 0 \implies \nabla^{2}f(w) \succeq 0.$$

Hence, by using the second-order condition of convexity, f(w) is convex.

Operations that preserve convexity: if f(x) and g(x) are convex functions, then h(x) is convex if,

- $h(x) = \alpha f(x)$ for $\alpha \ge 0$ (Non-negative scaling) E.g. For $w \in R^d$, $f(w) = \|w\|^2$ is convex, and hence $h(w) = \frac{\lambda}{2} \|w\|^2$ for $\lambda \ge 0$ is convex.
- $h(x) = \max\{f(x), g(x)\}$ (Point-wise maximum) E.g. f(w) = 0 and g(w) = 1 - w are convex functions, and hence $h(w) = \max\{0, 1 - w\}$ is convex.
- h(x) = f(Ax + b) (Composition with affine map) E.g.: $f(w) = \max\{0, 1 - w\}$ is convex, and hence $h(w) = \max\{0, 1 - y_i \langle w, x_i \rangle\}$ for $x_i \in \mathbb{R}^d$ and $y_i \in \mathbb{R}$ is convex
- h(x) = f(x) + g(x) (Sum) E.g.: $f(w) = \max\{0, 1 - y_i \langle w, x_i \rangle\}$ is convex, and hence $h(w) = \sum_{i=1}^n \max\{0, 1 - y_i \langle w, x_i \rangle\} + \frac{\lambda}{2} \|w\|^2$ is convex.

Hence, the SVM loss function: $f(w) := \sum_{i=1}^n \max\{0, 1 - y_i \langle X_i, w \rangle\} + \frac{\lambda}{2} \|w\|^2$ is convex.

Q: Prove that ℓ_1 -regularized logistic regression:

$$f(w) := \sum_{i=1}^{n} \log (1 + \exp(-y_i \langle X_i, w \rangle)) + \lambda \|w\|_1$$
 is convex

We have proved that the logistic loss $f(x) = \log(1 + \exp(-x))$ is convex. Since composition with an affine map is convex, and the sum of convex functions is convex, the first term is convex. Since all norms are convex, and a non-negative scaling of a convex function is convex, the second term is convex. Hence, f(w) is convex.

Another way to prove convexity for logistic regression is to compute the Hessian and show that it is positive semi-definite (In Assignment 1!)

Jensen's Inequality

Recall the zero-order definition of convexity: $\forall x, y \in \mathcal{D}$ and $\theta \in [0, 1]$, $f(\theta x + (1 - \theta)x) < \theta f(x) + (1 - \theta)f(y)$.

This can be generalized to *n* points $\{x_1, x_2, \dots, x_n\}$, i.e. for $p_i \ge 0$ and $\sum_i p_i = 1$,

$$f(p_1 x_1 + p_2 x_2 + \ldots + p_n x_n) \leq p_1 f(x_1) + p_2 f(x_2) + \ldots + p_n f(x_n) \implies f\left(\sum_{i=1}^n p_i x_i\right) \leq \sum_{i=1}^n p_i f(x_i)$$

i.e. if X is a discrete r.v. that can take value x_i with probability p_i , and f is convex, then,

$$f\left(\mathbb{E}[X]\right) \leq \mathbb{E}\left[f(X)\right].$$
 (Jensen's inequality)

Can be used to prove inequalities like the AM-GM inequality: $\sqrt{ab} \le \frac{a+b}{2}$.

Choose $f(x) = -\log(x)$ as the convex function, and consider two points a and b with $\theta = 1/2$. By Jensen's inequality,

$$-\log\left(\frac{a+b}{2}\right) \le \frac{-\log(a) - \log(b)}{2} \implies \log\left(\frac{a+b}{2}\right) \ge \log(\sqrt{ab})$$

Holder's Inequality

Q: Prove Holder's inequality, for p,q>1 s.t. $\frac{1}{p}+\frac{1}{q}=1$ and $x,y\in R^d$, $\langle x,y\rangle\leq \|x\|_p$ $\|y\|_q$

By repeating the AM-GM proof, but for a general $\theta \in [0,1]$, for $a,b \geq 0$,

$$a^{\theta}b^{1-\theta} \leq \theta a + (1-\theta)b$$

Use $a = \frac{|x_i|^p}{\sum_{j=1}^n |x_j|^p}$, $b = \frac{|y_i|^q}{\sum_{j=1}^n |y_j|^q}$, $\theta = 1/p$, and using the fact that $1 - \theta = 1 - 1/p = 1/q$

$$\left(\frac{|x_i|^p}{\sum_{j=1}^n |x_j|^p}\right)^{1/p} \left(\frac{|y_i|^q}{\sum_{j=1}^n |y_j|^q}\right)^{1/q} \le \frac{1}{p} \frac{|x_i|^p}{\sum_{j=1}^n |x_j|^p} + \frac{1}{q} \frac{|y_i|^p}{\sum_{j=1}^n |y_j|^p}$$

Summing both sides from i = 1 to n,

$$\sum_{i=1}^{n} \frac{|x_i|}{\left(\sum_{j=1}^{n} |x_j|^p\right)^{1/p}} \frac{|y_i|}{\left(\sum_{j=1}^{n} |y_j|^q\right)^{1/q}} \leq 1 \implies \sum_{i} x_i y_i \leq \left(\sum_{i=1}^{n} |x_i|^p\right)^{1/p} \left(\sum_{i=1}^{n} |y_i|^q\right)^{1/q}$$

