

CMPT 409/981: Optimization for Machine Learning

Lecture 3

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Recap

For L -smooth functions lower-bounded by f^* , gradient descent with $\eta = \frac{1}{L}$ returns an ϵ -approximate stationary point and requires $\Theta\left(\frac{1}{\epsilon}\right)$ iterations.

Importantly, the GD rate does not depend on the dimension of w .

In practice, we can set η_k in an adaptive manner using an exact line-search:

$$\eta_k = \arg \min_{\eta} f(w_k - \eta \nabla f(w_k)).$$

Exact line-search can adapt to the “local” L , resulting in larger step-sizes and better performance.

Can compute η_k analytically only in special cases, whereas solving the sub-problem approximately to set η_k can be expensive.

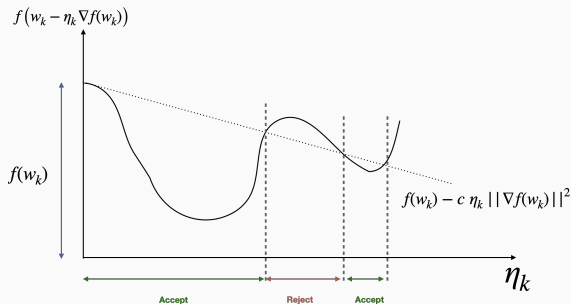
Gradient Descent with Line-search

Usually, the cost of doing an exact line-search is not worth the computational effort.

Armijo condition for a prospective step-size $\tilde{\eta}_k$:

$$f(w_k - \tilde{\eta}_k \nabla f(w_k)) \leq f(w_k) - c \tilde{\eta}_k \|\nabla f(w_k)\|^2$$

where $c \in (0, 1)$ is a hyper-parameter.



Backtracking line-search: At iteration k , starting with an initial “guess” of the step-size η_{\max} , check the Armijo condition for a prospective step-size $\tilde{\eta}_k$.

- If $\tilde{\eta}_k$ satisfies the Armijo condition, set $\eta_k = \tilde{\eta}_k$ and do the usual GD update.
- Else, decrease $\tilde{\eta}_k$ by a multiplicative factor $\beta \in (0, 1)$ and check the Armijo condition for the new prospective step-size equal to $\tilde{\eta}_k\beta$.
- Keep “backtracking” on $\tilde{\eta}_k$ until the Armijo condition is satisfied.
- Do the usual GD step: $w_{k+1} = w_k - \eta_k \nabla f(w_k)$ using the η_k for which the Armijo condition is satisfied.

Gradient Descent with Line-search

Claim: The (exact) backtracking procedure terminates and returns $\eta_k \geq \min \left\{ \frac{2(1-c)}{L}, \eta_{\max} \right\}$.

Proof:

$$f(w_k - \tilde{\eta}_k \nabla f(w_k)) \leq \underbrace{f(w_k) - \|\nabla f(w_k)\|^2 \left(\eta_k - \frac{L\eta_k^2}{2} \right)}_{h_1(\tilde{\eta}_k)} \quad (\text{Quadratic bound using smoothness})$$

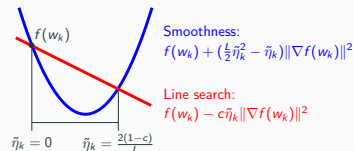
$$f(w_k - \tilde{\eta}_k \nabla f(w_k)) \leq \underbrace{f(w_k) - \|\nabla f(w_k)\|^2 (c\tilde{\eta}_k)}_{h_2(\tilde{\eta}_k)} \quad (\text{Armijo condition})$$

If the Armijo condition is satisfied, the back-tracking line-search procedure terminates.

Case (i): For $\eta_{\max} \leq \frac{2(1-c)}{L}$,

$$f(w_k - \eta_{\max} \nabla f(w_k)) \leq h_1(\eta_{\max}) \leq h_2(\eta_{\max})$$

\implies if $\eta_{\max} \leq \frac{2(1-c)}{L}$, then the line-search terminates immediately and $\eta_k = \eta_{\max}$.



Case (ii): If $\eta_{\max} > \frac{2(1-c)}{L}$ and the Armijo condition is satisfied for step-size η_k , then

$$f(w_k - \eta_k \nabla f(w_k)) \leq h_2(\eta_k) \leq h_1(\eta_k) \implies c\eta_k \geq \eta_k - \frac{L\eta_k^2}{2} \implies \eta_k \geq \frac{2(1-c)}{L}.$$

Putting the two cases together, the step-size η_k returned by the Armijo line-search satisfies

$$\eta_k \geq \min \left\{ \frac{2(1-c)}{L}, \eta_{\max} \right\}.$$

Gradient Descent with Line-search

Claim: Gradient Descent with (exact) backtracking Armijo line-search (with $c = 1/2$) returns point \hat{w} such that $\|\nabla f(\hat{w})\|^2 \leq \epsilon$ and requires $T = \frac{2L[f(w_0) - \min_w f(w)]}{\epsilon}$ oracle calls or iterations.

Proof: Since η_k satisfies the Armijo condition and $w_{k+1} = w_k - \eta_k \nabla f(w_k)$,

$$\begin{aligned} f(w_{k+1}) &\leq f(w_k) - c \eta_k \|\nabla f(w_k)\|^2 \\ &\leq f(w_k) - \left(\min \left\{ \frac{1}{2L}, \eta_{\max} \right\} \right) \|\nabla f(w_k)\|^2 \\ &\quad \text{(Result from previous slide with } c = 1/2) \end{aligned}$$

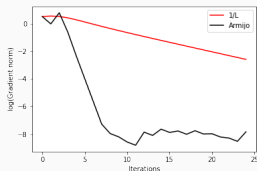
Continuing the proof as before,

$$\Rightarrow \|\nabla f(\hat{w})\|^2 \leq \frac{\max\{2L, 1/\eta_{\max}\} [f(w_0) - \min_w f(w)]}{T}$$

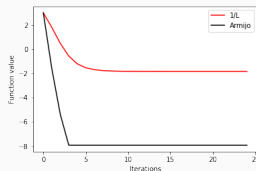
The claim is proved by reasoning as before.

Gradient Descent with Line-search – Examples

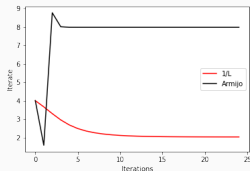
$\min_{x \in [-10, 10]} f(x) := -x \sin(x)$. Compare GD (with $x_0 = 4$) with (i) $\eta = 1/L \approx 0.1$ and (ii) Armijo line-search with $\eta_{\max} = 10, c = 1/2, \beta = 0.9$.



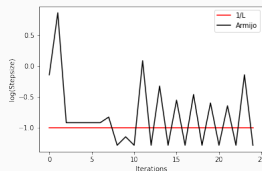
(a) Gradient norm



(b) Function value



(c) Iterate



(d) Stepsize

Questions?

Convex Optimization

We have seen that we require $\Theta(1/\epsilon)$ iterations to converge to an ϵ -approximate stationary point for smooth functions. Alternatively, if we care about global optimization (reach the vicinity of the true minimizer) of Lipschitz functions, we require $\Theta(1/\epsilon^d)$ iterations.

Convex functions: Class of functions where local optimization can result in convergence to the global minimizer of the function.

In general, convex optimization involves minimizing a convex function over a convex set \mathcal{C} : $\min_{w \in \mathcal{C}} f(w)$.

Examples of convex optimization in ML

Ridge regression: $\min_{w \in \mathbb{R}^d} \frac{1}{2} \|Xw - y\|^2 + \frac{\lambda}{2} \|w\|^2$.

Logistic regression: $\min_{w \in \mathbb{R}^d} \sum_{i=1}^n \log(1 + \exp(-y_i \langle X_i, w \rangle))$

Support vector machines: $\min_{w \in \mathbb{R}^d} \sum_{i=1}^n \max\{0, 1 - y_i \langle X_i, w \rangle\} + \frac{\lambda}{2} \|w\|^2$

Planning in MDPs in RL: $\max_{\mu \in \mathcal{F}_\rho} \langle \mu, r \rangle$ where \mathcal{F}_ρ is the flow-polytope.

A set \mathcal{C} is convex if a point along the line joining two points in \mathcal{C} also lies in the set.

For points x, y , the *convex combination* of x, y is $z = \theta x + (1 - \theta)y$ for $\theta \in [0, 1]$.

A set \mathcal{C} is convex iff $\forall x, y \in \mathcal{C}$, the convex combination $z \in \mathcal{C}$.

Examples of convex sets:

- Positive orthant $\mathbb{R}_+^d : \{x | x \geq 0\}$.
- Hyper-plane: $\{x | Ax = b\}$.
- Half-space: $\{x | Ax \leq b\}$.
- Norm-ball: $\{x | \|x\|_p \leq r\}$.
- Norm-cone: $\{(x, r) | \|x\|_p \leq r\}$.

Q: Prove that the hyper-plane (set of linear equations): $\mathcal{H} := \{x | Ax = b\}$ is a convex set.

If $x, y \in \mathcal{H}$, then, $Ax = b$ and $Ay = b$. Consider a point $z := \theta x + (1 - \theta)y$ for $\theta \in [0, 1]$.

$$Az = A[\theta x + (1 - \theta)y] = \theta Ax + (1 - \theta)Ay = b.$$

Hence, $z \in \mathcal{H}$ and \mathcal{H} is a convex set.

Q: Prove that the ball of radius r centered at point x_c : $\mathcal{B}(x_c, r) := \{x | \|x - x_c\|_p \leq r\}$ is convex.

If $x, y \in \mathcal{B}(x_c, r)$, then, $\|x - x_c\|_p \leq r$ and $\|y - x_c\|_p \leq r$. Consider a point $z := \theta x + (1 - \theta)y$ for $\theta \in [0, 1]$.

$$\begin{aligned}\|z - x_c\|_p &= \|\theta(x - x_c) + (1 - \theta)(y - x_c)\|_p \\ &\leq \|\theta(x - x_c)\|_p + \|(1 - \theta)(y - x_c)\|_p && \text{(Triangle inequality for norms)} \\ &\leq \theta \|x - x_c\|_p + (1 - \theta) \|y - x_c\|_p && \text{(Homogeneity of norms)}\end{aligned}$$

$$\implies \|z - x_c\|_p \leq r$$

Hence, $z \in \mathcal{B}(x_c, r)$ and $\mathcal{B}(x_c, r)$ is a convex set.

Q: Prove that the set of symmetric PSD matrices: $S_+^n = \{X \in \mathbb{R}^{n \times n} | X \succeq 0\}$ is convex.

Intersection of convex sets is convex \implies can prove the convexity of a set by showing that it is an intersection of convex sets.

Example: We know that a half-space: $\langle a_i, x \rangle \leq b_i$ is a convex set. The set of inequalities $Ax \leq b$ is an intersection of half-spaces and is hence convex.

Questions?

Convex Functions

Zero-order definition: A function f is convex iff its domain \mathcal{D} is a convex set, and for all $x, y \in \mathcal{D}$, then for all $\theta \in [0, 1]$,

$$f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y)$$

i.e. the function is below the chord between two points.

Alternatively, f is convex iff the set formed by the area above the function is a convex set.

Examples of convex functions:

- All norms $\|x\|_p$
- $f(x) = 1/\sqrt{x}$, $f(x) = -\log(x)$, $f(x) = \exp(-x)$
- Negative entropy: $f(x) = x \log(x)$
- Logistic loss: $f(x) = \log(1 + \exp(-x))$
- Linear functions $f(x) = \langle a, x \rangle$

Convex Functions

First-order condition: If f is differentiable, it is convex iff its domain \mathcal{D} is a convex set and for all $x, y \in \mathcal{D}$,

$$f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle$$

i.e. the function is above the tangent to the function at any point x .

For a convex f , consider w^* such that $\nabla f(w^*) = 0$, then using convexity, for all $y \in \mathcal{D}$, $f(y) \geq f(w^*)$. If w^* is a stationary point i.e. $\|\nabla f(w^*)\|^2 = 0$, then it is a global minimum. Hence, local optimization to make the gradient zero results in convergence to a global minimum!

Q: For a convex f , if $\nabla f(w^*) = 0$, then is w^* a unique minimizer of f ?

Second-order condition: If f is twice differentiable, it is convex iff its domain \mathcal{D} is a convex set and for all $x \in \mathcal{D}$,

$$\nabla^2 f(x) \succeq 0$$

i.e. the Hessian is positive semi-definite (“curved upwards”) for all x .

Q: Prove that $f(x) = \max_i x_i$ is a convex function

$$f(\theta x + (1 - \theta)y) = \max_i [\theta x_i + (1 - \theta)y_i] \leq \theta \max_i x_i + (1 - \theta) \max_i y_i = \theta f(x) + (1 - \theta)f(y)$$

Hence, by using the zero-order definition of convexity, $f(x)$ is convex.

Q: Prove that $f(x) = \frac{1}{2}x^2$ is a convex function

$$f(y) - f(x) - \langle \nabla f(x), y - x \rangle = \frac{y^2}{2} - \frac{x^2}{2} - x(y - x) = \frac{1}{2} [y^2 + x^2 - 2xy] = \frac{(x - y)^2}{2} \geq 0$$

Hence, by using the first-order condition of convexity, $f(x)$ is convex.

Convex Functions

Q: Prove that $f(x) = \log(1 + \exp(-x))$ is a convex function

$$f'(x) = \frac{-\exp(-x)}{1 + \exp(-x)} = \frac{-1}{1 + \exp(x)}$$
$$f''(x) = \frac{\exp(x)}{(1 + \exp(x))^2} > 0$$

Hence, by using the second-order condition of convexity, $f(x)$ is convex.

Q: Prove that the ridge regression loss function: $f(w) = \frac{1}{2} \|Xw - y\|^2 + \frac{\lambda}{2} \|w\|^2$ is convex

Recall that $\nabla^2 f(w) = X^T X + \lambda I_d$. For vector v , let us consider $v^T \nabla^2 f(w) v$,

$$v^T \nabla^2 f(w) v = v^T [X^T X + \lambda I_d] v = v^T [X^T X] v + \lambda v^T v = [Xv]^T [Xv] + \lambda \|v\|^2 = \|Xv\|^2 + \lambda \|v\|^2$$
$$\implies v^T \nabla^2 f(w) v \geq 0 \implies \nabla^2 f(w) \succeq 0.$$

Hence, by using the second-order condition of convexity, $f(w)$ is convex.

Convex Functions

Operations that preserve convexity: if $f(x)$ and $g(x)$ are convex functions, then $h(x)$ is convex if,

- $h(x) = \alpha f(x)$ for $\alpha \geq 0$ (Non-negative scaling)

E.g: For $w \in \mathbb{R}^d$, $f(w) = \|w\|^2$ is convex, and hence $h(w) = \frac{\lambda}{2} \|w\|^2$ for $\lambda \geq 0$ is convex.

- $h(x) = \max\{f(x), g(x)\}$ (Point-wise maximum)

E.g: $f(w) = 0$ and $g(w) = 1 - w$ are convex functions, and hence $h(w) = \max\{0, 1 - w\}$ is convex.

- $h(x) = f(Ax + b)$ (Composition with affine map)

E.g.: $f(w) = \max\{0, 1 - w\}$ is convex, and hence $h(w) = \max\{0, 1 - y_i \langle w, x_i \rangle\}$ for $x_i \in \mathbb{R}^d$ and $y_i \in \mathbb{R}$ is convex

- $h(x) = f(x) + g(x)$ (Sum)

E.g.: $f(w) = \max\{0, 1 - y_i \langle w, x_i \rangle\}$ is convex, and hence $h(w) = \sum_{i=1}^n \max\{0, 1 - y_i \langle w, x_i \rangle\} + \frac{\lambda}{2} \|w\|^2$ is convex.

Hence, the SVM loss function: $f(w) := \sum_{i=1}^n \max\{0, 1 - y_i \langle X_i, w \rangle\} + \frac{\lambda}{2} \|w\|^2$ is convex.

Q: Prove that ℓ_1 -regularized logistic regression:

$$f(w) := \sum_{i=1}^n \log(1 + \exp(-y_i \langle X_i, w \rangle)) + \lambda \|w\|_1 \text{ is convex}$$

We have proved that the logistic loss $f(x) = \log(1 + \exp(-x))$ is convex. Since composition with an affine map is convex, and the sum of convex functions is convex, the first term is convex. Since all norms are convex, and a non-negative scaling of a convex function is convex, the second term is convex. Hence, $f(w)$ is convex.

Another way to prove convexity for logistic regression is to compute the Hessian and show that it is positive semi-definite (In Assignment 1!)

Jensen's Inequality

Recall that the zero-order definition of convexity: $\forall x, y \in \mathcal{D}$ and $\theta \in [0, 1]$,
 $f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y)$.

This can be generalized to n points $\{x_1, x_2, \dots, x_n\}$, i.e. for $p_i \geq 0$ and $\sum_i p_i = 1$,

$$f(p_1 x_1 + p_2 x_2 + \dots + p_n x_n) \leq p_1 f(x_1) + p_2 f(x_2) + \dots + p_n f(x_n) \implies f\left(\sum_{i=1}^n p_i x_i\right) \leq \sum_{i=1}^n p_i f(x_i)$$

i.e. if X is a random variable that can take value x_i with probability p_i , and f is convex, then,

$$f(\mathbb{E}[X]) \leq \mathbb{E}[f(X)]. \quad (\text{Jensen's inequality})$$

Can be used to prove useful inequalities like the AM-GM inequality: $\sqrt{ab} \leq \frac{a+b}{2}$.

Choose $f(x) = -\log(x)$ as the convex function, and consider two points a and b with $\theta = 1/2$.

By Jensen's inequality,

$$-\log\left(\frac{a+b}{2}\right) \leq \frac{-\log(a) - \log(b)}{2} \implies \log\left(\frac{a+b}{2}\right) \geq \log(\sqrt{ab})$$

Holder's Inequality

Q: Prove Holder's inequality, for $p, q > 1$ s.t. $\frac{1}{p} + \frac{1}{q} = 1$ and $x, y \in R^d$, $\langle x, y \rangle \leq \|x\|_p \|y\|_q$

By repeating the AM-GM proof, but for a general $\theta \in [0, 1]$, for $a, b \geq 0$,

$$a^\theta b^{1-\theta} \leq \theta a + (1 - \theta)b$$

Use $a = \frac{|x_i|^p}{\sum_{j=1}^n |x_j|^p}$, $b = \frac{|y_i|^q}{\sum_{j=1}^n |y_j|^q}$, $\theta = 1/p$, and using the fact that $1 - \theta = 1 - 1/p = 1/q$

$$\left(\frac{|x_i|^p}{\sum_{j=1}^n |x_j|^p} \right)^{1/p} \left(\frac{|y_i|^q}{\sum_{j=1}^n |y_j|^q} \right)^{1/q} \leq \frac{1}{p} \frac{|x_i|^p}{\sum_{j=1}^n |x_j|^p} + \frac{1}{q} \frac{|y_i|^q}{\sum_{j=1}^n |y_j|^q}$$

Summing both sides from $i = 1$ to n ,

$$\sum_{i=1}^n \frac{|x_i|}{\left(\sum_{j=1}^n |x_j|^p \right)^{1/p}} \frac{|y_i|}{\left(\sum_{j=1}^n |y_j|^q \right)^{1/q}} \leq 1 \implies \sum_i x_i y_i \leq \left(\sum_{i=1}^n |x_i|^p \right)^{1/p} \left(\sum_{i=1}^n |y_i|^q \right)^{1/q}$$

Questions?

Minimizing Smooth, Convex Functions

Recall that for convex functions, minimizing the gradient norm results in finding the minimizer.

Let us analyze the convergence of GD on smooth, convex functions: $\min_{w \in \mathbb{R}^d} f(w)$.

Claim: For L -smooth, convex functions GD with $\eta = \frac{1}{L}$ requires $T = \frac{2L \|w_0 - w^*\|^2}{\epsilon}$ iterations to obtain point w_T that is ϵ -suboptimal in the sense that $f(w_T) \leq f(w^*) + \epsilon$.

Proof: For L -smooth functions, $\forall x, y \in \mathcal{D}$, $f(y) \leq f(x) + \langle \nabla f(x), y - x \rangle + \frac{L}{2} \|y - x\|^2$. Similar to Lecture 2, using GD: $w_{k+1} = w_k - \frac{1}{L} \nabla f(w_k)$ yields

$$f(w_{k+1}) - f(w^*) \leq f(w_k) - f(w^*) - \frac{1}{2L} \|\nabla f(w_k)\|^2 \quad (1)$$

Using $y = w^*$, $x = w_k$ in the first-order condition for convexity: $f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle$,

$$f(w_k) - f(w^*) \leq \langle \nabla f(w_k), w_k - w^* \rangle \leq \|\nabla f(w_k)\| \|w_k - w^*\| \quad (\text{Cauchy Schwarz})$$

$$\implies \|\nabla f(w_k)\| \geq \frac{f(w_k) - f(w^*)}{\|w_k - w^*\|} \quad (2)$$

Minimizing Smooth, Convex Functions

In addition to descent on the function, when minimizing smooth, convex functions, GD decreases the distance to a minimizer w^* .

Claim: For GD with $\eta = \frac{1}{L}$, $\|w_{k+1} - w^*\|^2 \leq \|w_k - w^*\|^2$.

Proof:

$$\|w_{k+1} - w^*\|^2 = \|w_k - \eta \nabla f(w_k) - w^*\|^2 = \|w_k - w^*\|^2 - 2\eta \langle \nabla f(w_k), w_k - w^* \rangle + \eta^2 \|\nabla f(w_k)\|^2$$

Using $y = w^*$, $x = w_k$ in the first-order condition for convexity: $f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle$,

$$\|w_{k+1} - w^*\|^2 \leq \|w_k - w^*\|^2 - 2\eta[f(w_k) - f(w^*)] + \eta^2 \|\nabla f(w_k)\|^2$$

For convex functions, L -smoothness is equivalent to

$$f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle + \frac{1}{2L} \|\nabla f(x) - \nabla f(y)\|^2. \text{ Using } x = w^*, y = w_k,$$

$$\leq \|w_k - w^*\|^2 - 2\eta[f(w_k) - f(w^*)] + 2L\eta^2[f(w_k) - f(w^*)]$$

$$\implies \|w_{k+1} - w^*\|^2 \leq \|w_k - w^*\|^2 \quad (\text{By setting } \eta = \frac{1}{L})$$

Minimizing Smooth, Convex Functions

Combining Eq. 2 with the result of the previous claim,

$$\|\nabla f(w_k)\| \geq \frac{f(w_k) - f(w^*)}{\|w_k - w^*\|} \geq \frac{f(w_k) - f(w^*)}{\|w_0 - w^*\|}$$

Combining the above inequality with Eq. 1,

$$f(w_{k+1}) - f(w^*) \leq f(w_k) - f(w^*) - \frac{1}{2L} \|\nabla f(w_k)\|^2 \leq f(w_k) - f(w^*) - \frac{1}{2L} \frac{[f(w_k) - f(w^*)]^2}{\|w_0 - w^*\|^2}$$

Dividing by $[f(w_k) - f(w^*)][f(w_k) - f(w^*)]$

$$\begin{aligned} \frac{1}{f(w_k) - f(w^*)} &\leq \frac{1}{f(w_{k+1}) - f(w^*)} - \frac{1}{2L} \frac{f(w_k) - f(w^*)}{\|w_0 - w^*\|^2} \frac{1}{f(w_{k+1}) - f(w^*)} \\ \Rightarrow \frac{1}{2L \|w_0 - w^*\|^2} \underbrace{\frac{f(w_k) - f(w^*)}{f(w_{k+1}) - f(w^*)}}_{\geq 1} &\leq \left[\frac{1}{f(w_{k+1}) - f(w^*)} - \frac{1}{f(w_k) - f(w^*)} \right] \end{aligned} \quad (3)$$

Minimizing Smooth, Convex Functions

Summing Eq. 3 from $k = 0$ to $T - 1$,

$$\begin{aligned} \sum_{k=0}^{T-1} \left[\frac{1}{2L \|w_0 - w^*\|^2} \right] &\leq \sum_{k=0}^{T-1} \left[\frac{1}{f(w_{k+1}) - f(w^*)} - \frac{1}{f(w_k) - f(w^*)} \right] \\ \frac{T}{2L \|w_0 - w^*\|^2} &\leq \frac{1}{f(w_T) - f(w^*)} - \frac{1}{f(w_0) - f(w^*)} \leq \frac{1}{f(w_T) - f(w^*)} \\ \implies f(w_T) - f(w^*) &\leq \frac{2L \|w_0 - w^*\|^2}{T} \end{aligned}$$

The suboptimality, $f(w_T) - f(w^*)$ decreases at an $O\left(\frac{1}{T}\right)$ rate, i.e. the function value at iterate w_T approaches the minimum function value $f(w^*)$.

In order to obtain a function value ϵ close to the optimal function value, GD requires $T = \frac{2L \|w_0 - w^*\|^2}{\epsilon}$ iterations.

Minimizing Smooth, Convex Functions

Recall that GD was optimal (amongst first-order methods with no dependence on the dimension) when minimizing smooth (possibly non-convex) functions.

Is GD also optimal when minimizing smooth, convex functions, or can we do better?

Lower Bound: For any initialization, there exists a smooth, convex functions such that any first-order method requires $\Omega\left(\frac{1}{\sqrt{\epsilon}}\right)$ iterations.

Possible reasons for the discrepancy between the $O(1/\epsilon)$ upper-bound for GD, and the $\Omega(1/\sqrt{\epsilon})$ lower-bound:

- 1 Our upper-bound analysis of GD is loose, and GD actual matches the lower-bound.
- 2 The lower-bound is loose, and there is a function that requires $\Omega(1/\epsilon)$ iterations to optimize.
- 3 Both the upper and lower-bounds are tight, and GD is sub-optimal. There exists another algorithm that has an $O(1/\epsilon)$ upper-bound and is hence optimal.

Option [3] is correct – GD is sub-optimal for minimizing smooth, convex functions. Using Nesterov acceleration is optimal and requires $\Theta(1/\epsilon)$ iterations.

Questions?