# CMPT 409/981: Optimization for Machine Learning

Lecture 15

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### Recap

#### Online Optimization

- 1: Online Optimization ( $w_0$ , Algorithm  $\mathcal{A}$ , Convex set  $\mathcal{C}$ )
- 2: **for** k = 1, ..., T **do**
- 3: Algorithm  $\mathcal{A}$  chooses point (decision)  $w_k \in \mathcal{C}$
- 4: Environment chooses and reveals the (potentially adversarial) loss function  $f_k:\mathcal{C}\to\mathbb{R}$
- 5: Algorithm suffers a cost  $f_k(w_k)$
- 6: end for

**Regret**: For any fixed decision  $u \in C$ ,  $R_T(u) := \sum_{k=1}^T [f_k(w_k) - f_k(u)]$ .

1

Online Gradient Descent (OGD): At iteration k, OGD chooses  $w_k$ . After the loss function  $f_k$  is revealed, OGD uses the function to compute

$$w_{k+1} = \Pi_C[w_k - \eta_k \nabla f_k(w_k)] \text{ where } \Pi_C[x] = \operatorname*{arg\,min}_{y \in \mathcal{C}} \frac{1}{2} \|y - x\|^2.$$

If the convex set  $\mathcal C$  has a diameter D i.e. for all  $x,y\in\mathcal C$ ,  $\|x-y\|^2\leq D$ , for an arbitrary sequence losses such that each  $f_k$  is convex, differentiable and G-Lipschitz, OGD with  $\eta_k=\frac{D}{\sqrt{2}\,G\,\sqrt{k}}$  and  $w_1\in\mathcal C$ , has regret  $R_T(u)\leq\sqrt{2}DG\,\sqrt{T}$ .

Additionally, if each  $f_k$  is  $\mu_k$  strongly-convex, OGD with  $\eta_k = \frac{1}{\sum_{i=1}^k \mu_i}$  has regret  $R_T(u) \leq \frac{G^2}{2\mu} \ (1 + \log(T))$ .

2

### Recap

**Follow the Leader** (FTL): At iteration k, FTL chooses the point  $w_k$ . After the loss function  $f_k$  is revealed, FTL uses it to compute

$$w_{k+1} = \arg\min_{w \in \mathcal{C}} \sum_{i=1}^k f_i(w).$$

Running FTL on a quadratic lower-bound for the loss recovers OGD in the strongly-convex case.

For strongly-convex, G-Lipschitz losses, FTL has regret  $R_T(u) \leq \frac{G^2}{2\mu} (1 + \log(T))$  that matches OGD, but does not require knowledge of  $\mu$  (Proof today).

If the losses are not necessarily strongly-convex, then FTL can result in O(T) regret.

#### Recap

Idea: Add an explicit regularization to fix FTL for a convex sequence of losses.

Follow the Regularized Leader (FTRL): At iteration  $k \ge 0$ , FTRL chooses the point  $w_k$ . After the loss function  $f_k$  is revealed, FTRL uses it to compute

$$w_{k+1} = \underset{w \in C}{\operatorname{arg \, min}} \sum_{i=1}^{k} \left[ f_i(w) + \frac{\sigma_i}{2} \|w - w_i\|^2 \right] + \frac{\sigma_0}{2} \|w\|^2 ,$$

where  $\sigma_i \geq 0$  is the regularization strength. If we set  $\sigma_i = 0$  for all i, FTRL reduces to FTL.

Running FTRL on a linear lower-bound for the loss recovers OGD in the convex case.

FTRL has the following regret for a general sequence of convex losses,

$$R_{T}(u) \leq \sum_{k=1}^{T} \left[ \frac{1}{2\lambda_{k+1}} \left\| \nabla f_{k}(w_{k}) \right\|^{2} \right] + \sum_{k=1}^{T} \frac{\sigma_{k}}{2} \left\| u - w_{k} \right\|^{2} + \frac{\sigma_{0}}{2} \left\| u \right\|^{2} \text{ where } \lambda_{k} = \sum_{i=1}^{k-1} [\mu_{i}] + \sum_{i=0}^{k} [\sigma_{i}].$$

For convex, G-Lipschitz losses, FTRL has regret  $R_T(u) \leq \sqrt{2}\sqrt{D^2 + \|u\|^2} G\sqrt{T}$ .

# Follow the Leader - Strongly-Convex, Lipschitz functions

**Claim**: If the convex set  $\mathcal C$  has diameter D, for an arbitrary sequence losses such that each  $f_k$  is  $\mu_k$  strongly-convex (s.t.  $\mu:=\min_{k=1}^T \mu_k>0$ ), G-Lipschitz and differentiable, then FTL with  $w_1\in\mathcal C$  satisfies the following regret bound for all  $u\in\mathcal C$ ,

$$R_T(u) \leq \frac{G^2}{2\mu} \ (1 + \log(T))$$

**Proof**: Using the general result for FTRL, for  $\lambda_{k+1} = \sum_{i=1}^k \mu_i + \sum_{i=0}^k \sigma_i$ . Since  $f_k$  is  $\mu_k$  strongly-convex, we will set  $\sigma_i = 0$  for all i. Hence,  $\lambda_{k+1} = \sum_{i=1}^k \mu_i \geq \mu_i k$ .

$$R_{T}(u) \leq \sum_{k=1}^{T} \left[ \frac{1}{2\lambda_{k+1}} \left\| \nabla f_{k}(w_{k}) \right\|^{2} \right] + \sum_{i=1}^{T} \frac{\sigma_{i}}{2} \left\| u - w_{i} \right\|^{2} + \frac{\sigma_{0}}{2} \left\| u \right\|^{2} \leq \frac{G^{2}}{2\mu} \sum_{k=1}^{T} \left[ \frac{1}{k} \right]$$
(Since  $f_{k}$  is  $G$ -Lipschitz)

$$\implies R_T(u) \leq \frac{G^2(1+\log(T))}{2\mu}$$

Hence, FTL matches the regret for OGD for strongly-convex, Lipschitz functions, but does not require knowledge of  $\mu$ .



# Adaptive step-sizes

Recall the claim we proved in Lecture 14 (Slide 6): If the convex set  $\mathcal C$  has diameter D, for an arbitrary sequence of losses such that each  $f_k$  is convex and differentiable, OGD with the update  $w_{k+1} = \Pi_{\mathcal C}[w_k - \eta_k \nabla f_k(w_k)]$  such that  $\eta_k \leq \eta_{k-1}$  and  $w_1 \in \mathcal C$  has the following regret for  $u \in \mathcal C$ ,

$$R_T(u) \leq \frac{D^2}{2\eta_T} + \sum_{k=1}^T \frac{\eta_k}{2} \|\nabla f_k(w_k)\|^2 = \frac{D^2}{2\eta} + \frac{\eta}{2} \sum_{k=1}^T \|\nabla f_k(w_k)\|^2 \quad \text{ (If } \eta_k = \eta \text{ for all } k\text{)}$$

In order to find the optimal  $\eta$ , differentiating the RHS w.r.t  $\eta$  and setting it to zero,

$$-\frac{D^2}{2\eta^2} + \frac{1}{2} \sum_{k=1}^{T} \|\nabla f_k(w_k)\|^2 = 0 \implies \eta^* = \frac{D}{\sqrt{\sum_{k=1}^{T} \|\nabla f_k(w_k)\|^2}}$$

Since the second derivative equal to  $\frac{2D^2}{\eta^3} > 0$ ,  $\eta^*$  minimizes the RHS. Setting  $\eta = \eta^*$ ,

$$R_T(u) \leq D \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2}$$

6

# Adaptive step-sizes

Choosing  $\eta = \eta^* = \frac{D}{\sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2}}$  minimizes the upper-bound on the regret. However, this is not practical since setting  $\eta$  requires knowing  $\nabla f_k(w_k)$  for all  $k \in [T]$ .

To approximate  $\eta^*$  to have a practical algorithm, we can set  $\eta_k$  as follows:

$$\eta_k = \frac{D}{\sqrt{\sum_{s=1}^k \|\nabla f_s(w_s)\|^2}}$$

Hence, at iteration k, we only use the gradients upto that iteration.

Algorithmically, we only need to maintain the running sum of the squared gradient norms.

Moreover, this choice of step-size ensures that  $\eta_k \leq \eta_{k-1}$  (since we are accumulating gradient norms in the denominator so the step-size cannot increase) and hence we can use our general result for bounding the regret.

#### Scalar AdaGrad

Hence, we have the following update for any  $\eta > 0$ ,

$$w_{k+1} = \Pi_C[w_k - \eta_k \nabla f_k(w_k)]$$
 ;  $\eta_k = \frac{\eta}{\sqrt{\sum_{s=1}^k \|\nabla f_s(w_s)\|^2}}$ 

This is exactly the AdaGrad update without a per-coordinate scaling and is referred to as scalar AdaGrad or AdaGrad Norm [WWB20].

For a sequence of convex, differentiable losses, using the general result,

$$R_{T}(u) \leq \frac{D^{2}}{2\eta_{T}} + \sum_{k=1}^{T} \frac{\eta_{k}}{2} \|\nabla f_{k}(w_{k})\|^{2} = \frac{D^{2}}{2\eta} \sqrt{\sum_{k=1}^{T} \|\nabla f_{k}(w_{k})\|^{2}} + \frac{\eta}{2} \sum_{k=1}^{T} \frac{\|\nabla f_{k}(w_{k})\|^{2}}{\sqrt{\sum_{s=1}^{k} \|\nabla f_{s}(w_{s})\|^{2}}}$$

In order to bound the regret for AdaGrad, we need to bound the last term.

### Scalar AdaGrad

We prove the following general claim and will use it for  $a_s = \|\nabla f_s(w_s)\|^2$ .

**Claim**: For all T and  $a_s \ge 0$ ,  $\sum_{k=1}^T \frac{a_k}{\sqrt{\sum_{s=1}^k a_s}} \le 2\sqrt{\sum_{k=1}^T a_k}$ .

**Proof**: Let us prove by induction. **Base case**: For T = 1, LHS =  $\sqrt{a_1} < 2\sqrt{a_1} = \text{RHS}$ .

**Inductive Hypothesis**: If the statement is true for T-1, we need to prove it for T.

$$\sum_{k=1}^{T} \frac{a_k}{\sqrt{\sum_{s=1}^{k} a_s}} = \sum_{k=1}^{T-1} \frac{a_k}{\sqrt{\sum_{s=1}^{k} a_s}} + \frac{a_T}{\sqrt{\sum_{s=1}^{T} a_s}} \le 2\sqrt{\sum_{s=1}^{T-1} a_s} + \frac{a_T}{\sqrt{\sum_{s=1}^{T} a_s}} = 2\sqrt{Z - x} + \frac{x}{\sqrt{Z}}$$

$$(x := a_T, Z := \sum_{s=1}^{T} a_s)$$

The derivative of the RHS w.r.t to x is  $-\frac{1}{\sqrt{Z-x}} + \frac{1}{\sqrt{Z}} < 0$  for all  $x \ge 0$  and hence the RHS is maximized at x = 0. Setting x = 0 completes the induction proof.

$$\implies \sum_{k=1}^{T} \frac{a_k}{\sqrt{\sum_{s=1}^{k} a_s}} \le 2\sqrt{Z} = 2\sqrt{\sum_{s=1}^{T} a_s}$$

#### Scalar AdaGrad

Recall that  $R_T(u) \leq \frac{D^2}{2\eta} \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2} + \frac{\eta}{2} \sum_{k=1}^T \frac{\|\nabla f_k(w_k)\|^2}{\sqrt{\sum_{s=1}^k \|\nabla f_s(w_s)\|^2}}$ . Using the claim in the previous slide with  $a_s := \|\nabla f_s(w_s)\|^2 \geq 0$ ,

$$R_T(u) \leq \frac{D^2}{2\eta} \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2} + \eta \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2} = \left(\frac{D^2}{2\eta} + \eta\right) \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2}.$$

The step-size that minimizes the above bound is equal to  $\eta^* = \frac{D}{\sqrt{2}}$ . With this choice,

$$R_T(u) \leq \sqrt{2}D \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2}$$

Comparing to the regret for the optimal (impractical) constant step-size on Slide 3,

$$R_T(u) \leq \sqrt{2} \min_{\eta} \left[ \frac{D^2}{2\eta} + \frac{\eta}{2} \sum_{k=1}^{T} \left\| \nabla f_k(w_k) \right\|^2 \right]$$

Hence, AdaGrad is only sub-optimal by  $\sqrt{2}$  when compared to the best constant step-size!

# Scalar AdaGrad - Convex, Lipschitz functions

**Claim**: If the convex set  $\mathcal{C}$  has diameter D i.e. for all  $x,y\in\mathcal{C}$ ,  $\|x-y\|^2\leq D$ , for an arbitrary sequence losses such that each  $f_k$  is convex, differentiable and G-Lipschitz, scalar AdaGrad with  $\eta_k = \frac{\eta}{\sqrt{\sum_{i=1}^k \|\nabla f_i(w_i)\|^2}}$  and  $w_1\in\mathcal{C}$  has the following regret for all  $u\in\mathcal{C}$ ,

$$R_T(u) \le \left(\frac{D^2}{2\eta} + \eta\right) G\sqrt{T}$$

**Proof**: Using the general result from the previous slide,

$$R_{T}(u) \leq \left(\frac{D^{2}}{2\eta} + \eta\right) \sqrt{\sum_{k=1}^{T} \|\nabla f_{k}(w_{k})\|^{2}} \leq \left(\frac{D^{2}}{2\eta} + \eta\right) \sqrt{G^{2}T} = \left(\frac{D^{2}}{2\eta} + \eta\right) G\sqrt{T}$$
(Since each  $f_{k}$  is  $G$ -Lipschitz)

With  $\eta = \frac{D}{\sqrt{2}}$ ,  $R_T(u) \le \sqrt{2} D G \sqrt{T}$ . Hence, for convex, Lipschitz functions, AdaGrad achieves the same regret as OGD but is adaptive to G.

# Scalar AdaGrad - Strongly-Convex, Lipschitz functions

**Claim**: If the convex set  $\mathcal C$  has diameter D i.e. for all  $x,y\in \mathcal C$ ,  $\|x-y\|^2\leq D$ , for an arbitrary sequence losses such that each  $f_k$  is  $\mu$  strongly-convex, differentiable and G-Lipschitz, scalar AdaGrad with  $\eta_k = \frac{G^2/\mu}{1+\sum_{s=1}^k \|\nabla f_s(w_s)\|^2}$  and  $w_1\in \mathcal C$  has the following regret for all  $u\in \mathcal C$ ,

$$R_T(u) = \frac{G^2}{2\mu} \left[ 1 + \log(1 + G^2 T) \right]$$

Though AdaGrad can achieve logarithmic regret for strongly-convex, Lipschitz functions similar to OGD and FTL, it requires knowledge of  ${\it G}$  and  ${\it \mu}$  and is not adaptive to these quantities.

Proof: Need to prove this in Assignment 4!



Let us consider a more general and practical variant of AdaGrad that uses a per-coordinate step-size. The corresponding update is:

$$\begin{aligned} v_{k+1} &= w_k - \eta \, A_k^{-1} \nabla f_k(w_k) \quad ; \quad w_{k+1} &= \Pi_{\mathcal{C}}^k[v_{k+1}] := \arg\min_{w \in \mathcal{C}} \frac{1}{2} \left\| w - v_{k+1} \right\|_{A_k}^2 \,. \\ A_k &= \begin{cases} \sqrt{\sum_{s=1}^k \left\| \nabla f_s(w_s) \right\|^2} \, I_d \quad \text{(Scalar AdaGrad)} \\ \operatorname{diag}(G_k^{\frac{1}{2}}) \quad \text{(Diagonal AdaGrad)} \\ G_k^{\frac{1}{2}} \quad \text{(Full-Matrix AdaGrad)} \end{cases} \end{aligned}$$

where 
$$G_k \in \mathbb{R}^{d \times d} := \sum_{s=1}^k [\nabla f_s(w_s) \nabla f_s(w_s)^{\mathsf{T}}].$$

Claim: If the convex set  $\mathcal{C}$  has diameter D, for an arbitrary sequence of losses such that each  $f_k$  is convex and differentiable, AdaGrad with the general update  $w_{k+1} = \Pi_{\mathcal{C}}^k[w_k - \eta A_k^{-1}\nabla f_k(w_k)]$  and  $w_1 \in \mathcal{C}$  has the following regret for  $u \in \mathcal{C}$ ,

$$R_T(u) \leq \left(\frac{D^2}{2\eta} + \eta\right) \sqrt{d} \sqrt{\sum_{k=1}^T \left\|\nabla f_k(w_k)\right\|^2}$$

**Proof**: Starting from the update,  $v_{k+1} = w_k - \eta A_k^{-1} \nabla f_k(w_k)$ ,

$$v_{k+1} - u = w_k - \eta A_k^{-1} \nabla f_k(w_k) - u \implies A_k[v_{k+1} - u] = A_k[w_k - u] - \eta \nabla f_k(w_k)$$

Multiplying the above equations,

$$[v_{k+1} - u]^{\mathsf{T}} A_k [v_{k+1} - u] = [w_k - u - \eta A_k^{-1} \nabla f_k(w_k)]^{\mathsf{T}} [A_k [w_k - u] - \eta \nabla f_k(w_k)]$$

$$\|v_{k+1} - u\|_{A_k}^2 = \|w_k - u\|_{A_k}^2 - 2\eta \langle \nabla f_k(w_k), w_k - u \rangle + \eta^2 [A_k^{-1} \nabla f_k(w_k)]^{\mathsf{T}} [\nabla f_k(w_k)]$$

$$\implies \|v_{k+1} - u\|_{A_k}^2 = \|w_k - u\|_{A_k}^2 - 2\eta \langle \nabla f_k(w_k), w_k - u \rangle + \eta^2 \|\nabla f_k(w_k)\|_{A_k^{-1}}^2$$

Recall that  $\|v_{k+1} - u\|_{A_k}^2 = \|w_k - u\|_{A_k}^2 - 2\eta \langle \nabla f_k(w_k), w_k - u \rangle + \eta^2 \|\nabla f_k(w_k)\|_{A_k^{-1}}^2$ . Using the update  $w_{k+1} = \Pi_{\mathcal{C}}^k[v_{k+1}]$ ,  $u \in \mathcal{C}$  with the non-expansiveness of projections,

$$\|w_{k+1} - u\|_{A_{k}}^{2} = \|\Pi_{\mathcal{C}}[v_{k+1}] - \Pi_{\mathcal{C}}[u]\|_{A_{k}}^{2} \le \|v_{k+1} - u\|_{A_{k}}^{2}$$

$$\implies \|w_{k+1} - u\|_{A_{k}}^{2} \le \|w_{k} - u\|_{A_{k}}^{2} - 2\eta \langle \nabla f_{k}(w_{k}), w_{k} - u \rangle + \eta^{2} \|\nabla f_{k}(w_{k})\|_{A_{k}^{-1}}^{2}$$

$$\le \|w_{k} - u\|_{A_{k}}^{2} - 2\eta [f_{k}(w_{k}) - f_{k}(u)] + \eta^{2} \|\nabla f_{k}(w_{k})\|_{A_{k}^{-1}}^{2} \qquad \text{(Convexity)}$$

$$\implies f_{k}(w_{k}) - f_{k}(u) \le \frac{\|w_{k} - u\|_{A_{k}}^{2} - \|w_{k+1} - u\|_{A_{k}}^{2}}{2\eta} + \frac{\eta}{2} \|\nabla f_{k}(w_{k})\|_{A_{k}^{-1}}^{2}$$

Summing from k = 1 to T,

$$\implies R_T(u) \leq \frac{1}{2\eta} \sum_{k=1}^T \left[ \|w_k - u\|_{A_k}^2 - \|w_{k+1} - u\|_{A_k}^2 \right] + \frac{\eta}{2} \sum_{k=1}^T \|\nabla f_k(w_k)\|_{A_k^{-1}}^2$$

Let us now bound the first term in the above expression.

$$\sum_{k=1}^{T} \left[ \| w_{k} - u \|_{A_{k}}^{2} - \| w_{k+1} - u \|_{A_{k}}^{2} \right]$$

$$= \sum_{k=2}^{T} \left[ (w_{k} - u)^{\mathsf{T}} [A_{k} - A_{k-1}] (w_{k} - u)] + \| w_{1} - u \|_{A_{1}}^{2} - \| w_{T+1} - u \|_{A_{T}}^{2}$$

$$\leq \sum_{k=2}^{T} \| w_{k} - u \|^{2} \lambda_{\max} [A_{k} - A_{k-1}] + \| w_{1} - u \|_{A_{1}}^{2} \leq \sum_{k=2}^{T} D^{2} \lambda_{\max} [A_{k} - A_{k-1}] + \| w_{1} - u \|_{A_{1}}^{2}$$

$$\left( \text{Since } A_{k-1} \leq A_{k}, \lambda_{\max} [A_{k} - A_{k-1}] \geq 0 \text{ and } \| w_{k} - u \|^{2} \leq D \right)$$

$$\implies \sum_{k=1}^{T} \left[ \| w_{k} - u \|_{A_{k}}^{2} - \| w_{k+1} - u \|_{A_{k}}^{2} \right] \leq D^{2} \sum_{k=2}^{T} \text{Tr}[A_{k} - A_{k-1}] + \| w_{1} - u \|_{A_{1}}^{2}$$

$$\left( \text{For any PSD matrix } B, \lambda_{\max}[B] \leq \text{Tr}[B] \right)$$

Continuing the proof from the previous slide,

$$\sum_{k=1}^{T} \left[ \| w_k - u \|_{A_k}^2 - \| w_{k+1} - u \|_{A_k}^2 \right] \le D^2 \sum_{k=2}^{T} \operatorname{Tr}[A_k - A_{k-1}] + \| w_1 - u \|_{A_1}^2$$

$$= D^2 \operatorname{Tr} \left[ \sum_{k=2}^{T} [A_k - A_{k-1}] \right] + \| w_1 - u \|_{A_1}^2 \qquad \text{(Linearity of Trace)}$$

$$= D^2 \operatorname{Tr}[A_T - A_1] + \| w_1 - u \|_{A_1}^2 \le D^2 \operatorname{Tr}[A_T - A_1] + \lambda_{\max}[A_1] \| w_1 - u \|^2$$

$$\sum_{k=2}^{T} \left[ \| w_k - u \|_{A_k}^2 - \| w_{k+1} - u \|_{A_k}^2 \right] \le D^2 \operatorname{Tr}[A_T] - D^2 \operatorname{Tr}[A_1] + D^2 \operatorname{Tr}[A_1] = D^2 \operatorname{Tr}[A_T]$$

Putting everything together,

$$R_T(u) \leq \frac{D^2 \operatorname{Tr}[A_T]}{2\eta} + \frac{\eta}{2} \sum_{k=1}^I \|\nabla f_k(w_k)\|_{A_k^{-1}}^2$$

Let us now bound the second term in the above expression.

Claim:  $\sum_{k=1}^{T} \|\nabla f_k(w_k)\|_{A_k^{-1}}^2 \le 2 \operatorname{Tr}[A_T]$ 

**Proof**: Let us prove by induction. For convenience, define  $\nabla_k := \nabla f_k(w_k)$ .

RHS. Here, we used the cyclic property of trace i.e. Tr[ABC] = Tr[BCA].

**Inductive Hypothesis**: If the statement is true for T-1, we need to prove it for T.

$$\sum_{k=1}^{T-1} \|\nabla_k\|_{A_k^{-1}}^2 + \|\nabla_T\|_{A_T^{-1}}^2 \le 2\operatorname{Tr}[A_{T-1}] + \|\nabla_T\|_{A_T^{-1}}^2 = 2\operatorname{Tr}[\left(A_T^2 - \nabla_T\nabla_T^{\mathsf{T}}\right)^{1/2}] + \operatorname{Tr}[A_T^{-1}\nabla_T\nabla_T^{\mathsf{T}}]$$

For any  $X \succeq Y \succeq 0$ , we have [DHS11, Lemma 8],  $2 \operatorname{Tr}[(X - Y)^{1/2}] + \operatorname{Tr}[X^{-1/2}Y] \le 2 \operatorname{Tr}[X^{1/2}]$ . Using this for  $X = A_T^2$ ,  $Y = \nabla_T \nabla_T^\mathsf{T}$ ,  $\sum_{k=1}^T \|\nabla_k\|_{A_k^{-1}}^2 \le 2 \operatorname{Tr}[A_T]$ , which completes the proof.

Putting everything together,

$$R_{\mathcal{T}}(u) \leq \left(\frac{D^2}{2\eta} + \eta\right) \operatorname{Tr}[A_{\mathcal{T}}].$$

Recall that  $R_T(u) \leq \left(\frac{D^2}{2\eta} + \eta\right) \operatorname{Tr}[A_T]$ . Bounding  $\operatorname{Tr}[A_T]$ 

$$\operatorname{Tr}[A_T] = \operatorname{Tr}[G_T^{\frac{1}{2}}] = \sum_{j=1}^d \sqrt{\lambda_j[G_T]} = d \, \frac{\sum_{j=1}^d \sqrt{\lambda_j[G_T]}}{d} \leq d \, \sqrt{\frac{\sum_{j=1}^d \lambda_j[G_T]}{d}}$$

$$= \sqrt{d} \sqrt{\sum_{i=1}^{d} \lambda_{j}[G_{T}]} = \sqrt{d} \sqrt{\mathsf{Tr}[G_{T}]} = \sqrt{d} \sqrt{\mathsf{Tr}\left[\sum_{k=1}^{T} \nabla f_{k}(w_{k}) \nabla f_{k}(w_{k})^{\mathsf{T}}\right]}$$

$$\operatorname{Tr}[A_T] \leq \sqrt{d} \, \sqrt{\sum_{k=1}^T \operatorname{Tr}\left[\nabla f_k(w_k) \nabla f_k(w_k)^\mathsf{T}\right]} = \sqrt{d} \, \sqrt{\sum_{k=1}^T \left\|\nabla f_k(w_k)\right\|^2} \qquad \text{(Linearity of Trace)}$$

Putting everything together,

$$R_T(u) \leq \left(\frac{D^2}{2\eta} + \eta\right) \sqrt{d} \sqrt{\sum_{k=1}^T \left\| \nabla f_k(w_k) \right\|^2}$$

(Jensen's inequality for  $\sqrt{x}$ )

# AdaGrad - Convex, Lipschitz functions

**Claim**: If the convex set C has diameter D, for an arbitrary sequence of losses such that each  $f_k$  is convex, differentiable and G-Lipschitz, AdaGrad with the general update

 $w_{k+1} = \Pi_{\mathcal{C}}^k[w_k - \eta A_k^{-1} \nabla f_k(w_k)]$  with  $\eta = \frac{D}{\sqrt{2}}$  and  $w_1 \in \mathcal{C}$  has the following regret for  $u \in \mathcal{C}$ ,

$$R_{\mathcal{T}}(u) \leq \left(\frac{D^2}{2\eta} + \eta\right) \sqrt{d} \sqrt{\sum_{k=1}^{T} \left\|\nabla f_k(w_k)\right\|^2}$$

**Proof**: Using the general result from the previous slide and that each  $f_k$  is G-Lipschitz,

$$R_{T}(u) \leq \left(\frac{D^{2}}{2\eta} + \eta\right) \sqrt{d} \sqrt{\sum_{k=1}^{T} \left\|\nabla f_{k}(w_{k})\right\|^{2}} \leq \left(\frac{D^{2}}{2\eta} + \eta\right) \sqrt{d} G \sqrt{T}$$

$$R_{T}(u) \leq \sqrt{2}DG d \sqrt{T} \qquad \qquad \text{(Setting } \eta = \frac{D}{\sqrt{2}}\text{)}$$

Unlike scalar AdaGrad, when using the diagonal or full-matrix variant, the regret depends on the dimension d.

### AdaGrad - Convex, Smooth functions

Recall that for convex functions, the regret for AdaGrad is bounded as:

$$R_T(u) \leq \left(\frac{D^2}{2\eta} + \eta\right) \sqrt{d} \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2}.$$

In order to bound the regret for smooth functions, we define  $\zeta^2$  such that  $f_k(u) - f_k^* \leq \zeta^2$ . Hence, if the learner is competing against a fixed decision u that minimizes each  $f_k$ , then  $\zeta^2 = 0$ .  $\zeta^2$  characterizes the analog of interpolation in the online setting.

Using L-smoothness of  $f_k$  to bound the gradient norm term (for each k) in the regret expression,

$$\|\nabla f_{k}(w_{k})\|^{2} \leq 2L[f_{k}(w_{k}) - f_{k}^{*}] = 2L[f_{k}(w_{k}) - f_{k}(u)] + 2L[f_{k}(u) - f_{k}^{*}] \leq 2L[f_{k}(w_{k}) - f_{k}(u)] + 2L\zeta^{2}$$

$$\implies \sum_{k=1}^{T} \|\nabla f_{k}(w_{k})\|^{2} \leq 2L\sum_{k=1}^{T} [f_{k}(w_{k}) - f_{k}(u)] + 2L\sum_{k=1}^{T} \zeta^{2} = 2L[R_{T}(u) + \zeta^{2}T]$$

$$R_{T}(u) \leq \left(\frac{D^{2}}{2\eta} + \eta\right) \sqrt{d}\sqrt{2L[R_{T}(u) + \zeta^{2}T]}$$

## AdaGrad - Convex, Smooth functions

Recall that  $R_T(u) \leq \left(\frac{D^2}{2\eta} + \eta\right) \sqrt{d} \sqrt{2L[R_T(u) + \zeta^2 T]}$ . Squaring this expression,

$$[R_{T}(u)]^{2} \leq 2dL \left(\frac{D^{2}}{2\eta} + \eta\right)^{2} \underbrace{[R_{T}(u) + \zeta^{2}T]}_{:=x}$$

$$\implies x^{2} \leq \alpha(x+\beta) \implies x \leq \frac{\alpha + \sqrt{\alpha^{2} + 4\alpha\beta}}{2} \leq \alpha + \sqrt{\alpha\beta}$$

$$\implies R_{T}(u) \leq 2dL \left(\frac{D^{2}}{2\eta} + \eta\right)^{2} + \sqrt{2dL} \left(\frac{D^{2}}{2\eta} + \eta\right) \zeta\sqrt{T}$$

Note that the above bound holds for all  $\eta > 0$  and AdaGrad does not need to know  $\zeta$  or L. The regret depends on  $\zeta^2$ , the upper-bound on  $\max_{k \in [T]} [f_k(u) - f_k^*]$ . Such bounds that depend on the fixed decision that we are comparing against are called *first-order regret bounds*.

For example, when  $u = w^* := \arg\min_w \sum_{k=1}^T f_k(w)$  and  $\zeta = 0$ , then AdaGrad only incurs a constant regret that is independent of T. This observation has been used to explain the good performance of IL algorithms when using over-parameterized (convex) models [YBC20, LVS22].



#### References i

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