

# CMPT 409/981: Optimization for Machine Learning

## Lecture 3

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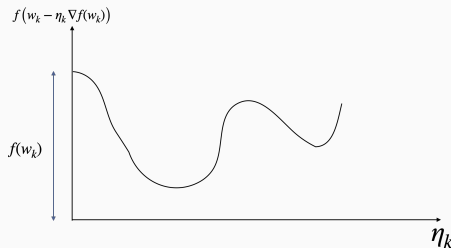
- For an  $L$ -smooth function,  $f(y) \leq f(x) + \langle \nabla f(x), y - x \rangle + \frac{L}{2} \|y - x\|^2$  for all  $x, y \in \mathcal{D}$ .
- For  $L$ -smooth functions lower-bounded by  $f^*$ , gradient descent with  $\eta = \frac{1}{L}$  returns  $\hat{w}$  such that  $\|\nabla f(\hat{w})\|^2 \leq \epsilon$  and requires  $T \geq \frac{2L[f(w_0) - f^*]}{\epsilon}$  iterations (oracle calls).
- Importantly, the GD rate does not depend on the dimension of  $w$ .
- *Lower-Bound*: When minimizing a smooth function (without additional assumptions), any *first-order* algorithm requires  $\Omega\left(\frac{1}{\epsilon}\right)$  oracle calls to return a point  $\hat{w}$  such that  $\|\nabla f(\hat{w})\|^2 \leq \epsilon$ .
- Hence, GD is optimal for minimizing smooth functions.

- The above results require setting the step-size to  $\frac{1}{L}$ . In fact, GD with any  $\eta \in (0, \frac{2}{L})$  will result in convergence to the stationary point (prove in Assignment 1).
- However, estimating  $L$  can be difficult as the functions get more complicated.
- Even for simple functions, the theoretically computed  $L$  is global (the “local”  $L$  might be much smaller) and often loose in practice. Typically we tend to overestimate  $L$  resulting in a smaller step-size.
- Instead of setting  $\eta$  according to  $L$ , we can “search” for a good step-size  $\eta_k$  in each iteration  $k$ . We will study 2 ways to do so:
  - Exact Line-search
  - Backtracking Armijo Line-search

# Exact Line-search

**Exact line-search:** At iteration  $k$ , solve the following sub-problem:

$$\eta_k = \arg \min_{\eta} f(w_k - \eta \nabla f(w_k)).$$



After computing  $\eta_k$ , do the usual GD update:  $w_{k+1} = w_k - \eta_k \nabla f(w_k)$ .

- Can adapt to the “local”  $L$ , resulting in larger step-sizes and better performance.
- Can solve the sub-problem approximately by doing gradient descent w.r.t  $\eta$  (known as *hyper-gradient descent* [BCR<sup>+</sup>17]). This is computationally expensive.
- Can compute  $\eta_k$  analytically. This can only be done in special cases such as for quadratics.

# Exact Line-search for Linear Regression

Recall linear regression: for  $X \in \mathbb{R}^{n \times d}$  and  $y \in \mathbb{R}^n$ , we aim to solve:

$$\min_{w \in \mathbb{R}^d} f(w) := \frac{1}{2} \|Xw - y\|^2 = \frac{1}{2} \left[ w^\top (X^\top X) w - 2 \langle X^\top y, w \rangle + \|y\|^2 \right].$$

For the exact line-search, we need to  $\min_\eta h(\eta) := f(w_k - \eta \nabla f(w_k))$ .

Since  $f$  is a quadratic, we can directly use the second-order Taylor series:

$$\begin{aligned} f(w_k - \eta \nabla f(w_k)) &= f(w_k) + \langle \nabla f(w_k), -\eta \nabla f(w_k) \rangle + \frac{1}{2} [-\eta \nabla f(w_k)]^\top \nabla^2 f(w_k) [-\eta \nabla f(w_k)] \\ \implies \nabla h(\eta_k) &= -\|\nabla f(w_k)\|^2 + \eta [\nabla f(w_k)]^\top \nabla^2 f(w_k) [\nabla f(w_k)] = 0 \\ \implies \eta_k &= \frac{\|\nabla f(w_k)\|^2}{\|\nabla f(w_k)\|_{\nabla^2 f(w_k)}^2} \end{aligned}$$

For linear regression,  $\nabla^2 f(w_k) = X^\top X$  and  $\nabla f(w_k) = X^\top (Xw_k - y)$ .

$$\implies \eta_k = \frac{\|X^\top (Xw_k - y)\|^2}{\|X^\top (Xw_k - y)\|_{X^\top X}^2}. \quad (\text{Implement in Assignment 1})$$

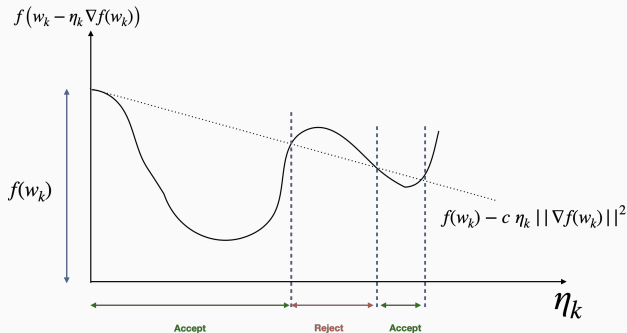
# Armijo Condition

Usually, the cost of doing an exact line-search is not worth the computational effort.

**Armijo condition** for a prospective step-size  $\tilde{\eta}_k$ :

$$f(w_k - \tilde{\eta}_k \nabla f(w_k)) \leq f(w_k) - c \tilde{\eta}_k \|\nabla f(w_k)\|^2$$

where  $c \in (0, 1)$  is a hyper-parameter.



# Gradient Descent with Backtracking Armijo Line-search

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**Algorithm** GD with Armijo Line-search

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1: function GD with Armijo line-search( $f, w_0, \eta_{\max}, c \in (0, 1), \beta \in (0, 1)$ )
2:   for  $k = 0, \dots, T - 1$  do
3:      $\tilde{\eta}_k \leftarrow \eta_{\max}$ 
4:     while  $f(w_k - \tilde{\eta}_k \nabla f(w_k)) > f(w_k) - c \cdot \tilde{\eta}_k \|\nabla f(w_k)\|^2$  do
5:        $\tilde{\eta}_k \leftarrow \tilde{\eta}_k \beta$ 
6:     end while
7:      $\eta_k \leftarrow \tilde{\eta}_k$ 
8:      $w_{k+1} = w_k - \eta_k \nabla f(w_k)$ 
9:   end for
10: return  $w_T$ 
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# Backtracking Armijo Line-search

**Simplification for analysis:** Assume that the backtracking line-search procedure returns the largest  $\eta$  that satisfies the Armijo condition. Will be referred to as *exact backtracking line-search*.

**Claim:** For  $L$ -smooth functions, the exact backtracking line-search procedure terminates and returns  $\eta_k \geq \min \left\{ \frac{2(1-c)}{L}, \eta_{\max} \right\}$ .

**Proof:** For a prospective step-size  $\tilde{\eta}_k$ , we will use the following two inequalities:

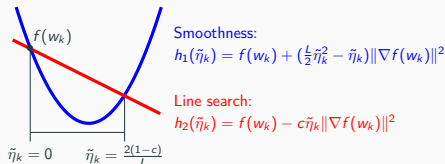
$$\underbrace{f(w_k - \tilde{\eta}_k \nabla f(w_k)) \leq f(w_k) - \|\nabla f(w_k)\|^2 \left( \tilde{\eta}_k - \frac{L\tilde{\eta}_k^2}{2} \right)}_{h_1(\tilde{\eta}_k)} \quad (\text{Quadratic bound using smoothness})$$

$$\underbrace{f(w_k - \tilde{\eta}_k \nabla f(w_k)) \leq f(w_k) - \|\nabla f(w_k)\|^2 (c\tilde{\eta}_k)}_{h_2(\tilde{\eta}_k)} \quad (\text{Armijo condition})$$



# Backtracking Armijo Line-search

If the Armijo condition is satisfied, the backtracking line-search procedure terminates.



**Case (i)**  $\eta_{\max} \leq \frac{2(1-c)}{L}$ : From smoothness,  $f(w_k - \eta_{\max} \nabla f(w_k)) \leq h_1(\eta_{\max})$ . For  $\eta_{\max} \leq \frac{2(1-c)}{L}$ , we know that  $h_1(\eta_{\max}) \leq h_2(\eta_{\max})$ . Hence,  $f(w_k - \eta_{\max} \nabla f(w_k)) \leq h_2(\eta_{\max})$ , meaning that the Armijo condition is satisfied for  $\eta_{\max}$ .  $\implies$  if  $\eta_{\max} \leq \frac{2(1-c)}{L}$ , then the line-search terminates immediately and  $\eta_k = \eta_{\max}$ .

**Case (ii):** If  $\eta_{\max} > \frac{2(1-c)}{L}$ : While backtracking, if  $\tilde{\eta}_k = \frac{2(1-c)}{L}$ , then  $f(w_k - \eta_{\max} \nabla f(w_k)) \leq h_1(\tilde{\eta}_k) = h_2(\tilde{\eta}_k)$ , the line-search terminates immediately and  $\eta_k = \frac{2(1-c)}{L}$ . If the Armijo condition is satisfied for a step-size  $\eta_k$  s.t.  $h_2(\eta_k) < h_1(\eta_k)$ , then  $f(w_k - \eta_k \nabla f(w_k)) \leq h_2(\eta_k) < h_1(\eta_k) \implies c\eta_k \geq \eta_k - \frac{L\eta_k^2}{2} \implies \eta_k \geq \frac{2(1-c)}{L}$ .

Putting everything together, the step-size  $\eta_k$  returned by the Armijo line-search satisfies  $\eta_k \geq \min \left\{ \frac{2(1-c)}{L}, \eta_{\max} \right\}$ .

# Gradient Descent with Backtracking Armijo Line-search

**Claim:** For  $L$ -smooth functions lower-bounded by  $f^*$ , gradient descent with exact backtracking Armijo line-search (with  $c = 1/2$ ) returns point  $\hat{w}$  such that  $\|\nabla f(\hat{w})\|^2 \leq \epsilon$  and requires  $T \geq \frac{\max\{2L, 2/\eta_{\max}\} [f(w_0) - \min_w f(w)]}{\epsilon}$  iterations.

**Proof:** Since  $\eta_k$  satisfies the Armijo condition and  $w_{k+1} = w_k - \eta_k \nabla f(w_k)$ ,

$$\begin{aligned} f(w_{k+1}) &\leq f(w_k) - c \eta_k \|\nabla f(w_k)\|^2 \\ &\leq f(w_k) - \left( \min \left\{ \frac{1}{2L}, \frac{\eta_{\max}}{2} \right\} \right) \|\nabla f(w_k)\|^2 \end{aligned}$$

(Result from previous slide with  $c = 1/2$ )

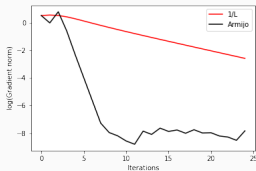
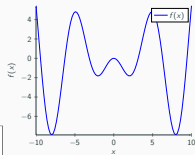
Continuing the proof as before,

$$\implies \|\nabla f(\hat{w})\|^2 \leq \frac{\max\{2L, 2/\eta_{\max}\} [f(w_0) - \min_w f(w)]}{T}$$

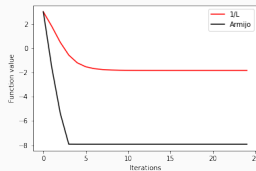
The claim can be proved by the same reasoning as in Lecture 2.

# Gradient Descent with Backtracking Armijo Line-search – Example

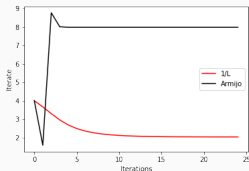
$\min_{x \in [-10, 10]} f(x) := -x \sin(x)$ . Compare GD (with  $x_0 = 4$ ) with (i)  $\eta = 1/L \approx 0.1$  and (ii) Armijo line-search ( $\eta_{\max} = 10, c = 1/2, \beta = 0.9$ ).



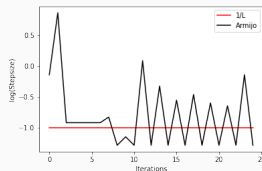
(a) Gradient norm



(b) Function value



(c) Iterate



(d) Step size

Questions?

# Convex Optimization

For smooth functions, GD requires  $\Theta(1/\epsilon)$  iterations to converge to an  $\epsilon$ -approximate stationary point. Alternatively, if we care about global optimization (reach the vicinity of the true minimizer), any algorithm requires  $\Theta(1/\epsilon^d)$  iterations.

**Convex functions:** Class of functions where local optimization can result in convergence to the global minimizer of the function.

In general, convex optimization involves minimizing a convex function over a convex set  $\mathcal{C}$ .

*Examples of convex optimization in ML*

**Ridge regression:**  $\min_{w \in \mathbb{R}^d} \frac{1}{2} \|Xw - y\|^2 + \frac{\lambda}{2} \|w\|^2$ .

**Logistic regression:**  $\min_{w \in \mathbb{R}^d} \sum_{i=1}^n \log(1 + \exp(-y_i \langle X_i, w \rangle))$

**Support vector machines:**  $\min_{w \in \mathbb{R}^d} \sum_{i=1}^n \max\{0, 1 - y_i \langle X_i, w \rangle\} + \frac{\lambda}{2} \|w\|^2$

**Planning in MDPs in RL:**  $\max_{\mu \in \mathcal{F}_\rho} \langle \mu, r \rangle$  where  $\mathcal{F}_\rho$  is the flow-polytope.

A set  $\mathcal{C}$  is convex if every point along the line joining two points in  $\mathcal{C}$  also lies in the set.

For points  $x, y$ , the *convex combination* of  $x, y$  is  $z_\theta := \theta x + (1 - \theta)y$  for  $\theta \in [0, 1]$ .

A set  $\mathcal{C}$  is convex iff  $\forall x, y \in \mathcal{C}$ , the convex combination  $z_\theta \in \mathcal{C}$  for all  $\theta \in [0, 1]$ .

*Examples of convex sets:*

- Positive orthant  $\mathbb{R}_+^d : \{x | x \geq 0\}$ .
- Hyper-plane:  $\{x | Ax = b\}$ .
- Half-space:  $\{x | Ax \leq b\}$ .
- Norm-ball:  $\{x | \|x\|_p \leq r\}$  for  $p \geq 1$ .
- Norm-cone:  $\{(x, r) | \|x\|_p \leq r\}$  for  $p \geq 1$ .

**Q:** Prove that the hyper-plane (set of linear equations):  $\mathcal{H} := \{x \mid Ax = b\}$  is a convex set.

If  $x, y \in \mathcal{H}$ , then,  $Ax = b$  and  $Ay = b$ . Consider a point  $z_\theta := \theta x + (1 - \theta)y$  for  $\theta \in [0, 1]$ .

$$Az_\theta = A[\theta x + (1 - \theta)y] = \theta Ax + (1 - \theta)Ay = b.$$

Hence,  $z_\theta \in \mathcal{H}$  for all  $\theta \in [0, 1]$  and  $\mathcal{H}$  is a convex set.

**Q:** Prove that the ball of radius  $r$  centered at point  $x_c$ :  $\mathcal{B}(x_c, r) := \{x \mid \|x - x_c\|_p \leq r\}$  for  $p \geq 1$  is convex.

If  $x, y \in \mathcal{B}(x_c, r)$ , then,  $\|x - x_c\|_p \leq r$  and  $\|y - x_c\|_p \leq r$ . Consider a point  $z_\theta := \theta x + (1 - \theta)y$  for  $\theta \in [0, 1]$ .

$$\begin{aligned}\|z_\theta - x_c\|_p &= \|\theta(x - x_c) + (1 - \theta)(y - x_c)\|_p \\ &\leq \|\theta(x - x_c)\|_p + \|(1 - \theta)(y - x_c)\|_p && \text{(Triangle inequality for norms)} \\ &\leq \theta \|x - x_c\|_p + (1 - \theta) \|y - x_c\|_p && \text{(Homogeneity of norms)}\end{aligned}$$

$$\implies \|z - x_c\|_p \leq r$$

Hence,  $z_\theta \in \mathcal{B}(x_c, r)$  for all  $\theta \in [0, 1]$  and  $\mathcal{B}(x_c, r)$  is a convex set.

**Q:** Prove that the set of symmetric PSD matrices:  $S_+^n = \{X \in \mathbb{R}^{n \times n} | X \succeq 0, X = X^T\}$  is convex.

- Intersection of convex sets is convex  $\implies$  can prove the convexity of a set by showing that it is an intersection of convex sets.

*Example:* We know that a half-space:  $\langle a_i, x \rangle \leq b_i$  is a convex set. The set of inequalities  $Ax \leq b$  is an intersection of half-spaces and is hence convex.



Questions?

# Convex Functions

**Zero-order definition:** A function  $f$  is convex iff its domain  $\mathcal{D}$  is a convex set, and for all  $x, y \in \mathcal{D}$  and  $\theta \in [0, 1]$ ,

$$f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y)$$

i.e. the function is below the chord between two points.

- Alternatively,  $f$  is convex iff the set formed by the area above the function is a convex set.

*Examples of convex functions:*

- All  $p$ -norms  $\|x\|_p$  with  $p \geq 1$ .
- $f(x) = 1/\sqrt{x}$ ,  $f(x) = -\log(x)$ ,  $f(x) = \exp(-x)$
- Negative entropy:  $f(x) = x \log(x)$
- Logistic loss:  $f(x) = \log(1 + \exp(-x))$
- Linear functions  $f(x) = \langle a, x \rangle$

# Convex Functions

**First-order condition:** If  $f$  is differentiable, it is convex iff its domain  $\mathcal{D}$  is a convex set and for all  $x, y \in \mathcal{D}$ ,

$$f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle$$

i.e. the function is above the tangent to the function at any point  $x$ .

For a convex  $f$ , consider  $w^*$  such that  $\nabla f(w^*) = 0$ , then using convexity, for all  $y \in \mathcal{D}$ ,  $f(y) \geq f(w^*)$ . If  $w^*$  is a stationary point i.e.  $\|\nabla f(w^*)\|^2 = 0$ , then it is a global minimum. Hence, local optimization to make the gradient zero results in convergence to a global minimum!

**Q:** For a convex  $f$ , if  $\nabla f(w^*) = 0$ , then is  $w^*$  a unique minimizer of  $f$ ?

**Second-order condition:** If  $f$  is twice differentiable, it is convex iff its domain  $\mathcal{D}$  is a convex set and for all  $x \in \mathcal{D}$ ,

$$\nabla^2 f(x) \succeq 0$$

i.e. the Hessian is positive semi-definite (“curved upwards”) for all  $x$ .

**Q:** Prove that  $f(x) = \max_i x_i$  is a convex function.

$$f(\theta x + (1 - \theta)y) = \max_i [\theta x_i + (1 - \theta)y_i] \leq \theta \max_i x_i + (1 - \theta) \max_i y_i = \theta f(x) + (1 - \theta)f(y)$$

Hence, by using the zero-order definition of convexity,  $f(x)$  is convex.

**Q:** Prove that  $f(x) = \frac{1}{2}x^2$  is a convex function.

$$f(y) - f(x) - \langle \nabla f(x), y - x \rangle = \frac{y^2}{2} - \frac{x^2}{2} - x(y - x) = \frac{1}{2} [y^2 + x^2 - 2xy] = \frac{(x - y)^2}{2} \geq 0$$

Hence, by using the first-order definition of convexity,  $f(x)$  is convex.

# Convex Functions

**Q:** Prove that  $f(x) = \log(1 + \exp(-x))$  is a convex function.

$$f'(x) = \frac{-\exp(-x)}{1 + \exp(-x)} = \frac{-1}{1 + \exp(x)}$$
$$f''(x) = \frac{\exp(x)}{(1 + \exp(x))^2} > 0$$

Hence, by using the second-order definition of convexity,  $f(x)$  is convex.

**Q:** Prove that the ridge regression loss function:  $f(w) = \frac{1}{2} \|Xw - y\|^2 + \frac{\lambda}{2} \|w\|^2$  is convex

Recall that  $\nabla^2 f(w) = X^T X + \lambda I_d$ . For vector  $v$ , let us consider  $v^T \nabla^2 f(w) v$ ,

$$v^T \nabla^2 f(w) v = v^T [X^T X + \lambda I_d] v = v^T [X^T X] v + \lambda v^T v = [Xv]^T [Xv] + \lambda \|v\|^2 = \|Xv\|^2 + \lambda \|v\|^2$$
$$\implies v^T \nabla^2 f(w) v \geq 0 \implies \nabla^2 f(w) \succeq 0.$$

Hence, by using the second-order definition of convexity,  $f(w)$  is convex.

# Convex Functions

Operations that preserve convexity: if  $f(x)$  and  $g(x)$  are convex functions, then  $h(x)$  is convex if,

- $h(x) = \alpha f(x)$  for  $\alpha \geq 0$  (Non-negative scaling)

E.g: For  $w \in \mathbb{R}^d$ ,  $f(w) = \|w\|^2$  is convex, and hence  $h(w) = \frac{\lambda}{2} \|w\|^2$  for  $\lambda \geq 0$  is convex.

- $h(x) = \max\{f(x), g(x)\}$  (Point-wise maximum)

E.g:  $f(w) = 0$  and  $g(w) = 1 - w$  are convex functions, and hence  $h(w) = \max\{0, 1 - w\}$  is convex.

- $h(x) = f(Ax + b)$  (Composition with affine map)

E.g.:  $f(w) = \max\{0, 1 - w\}$  is convex, and hence  $h(w) = \max\{0, 1 - y_i \langle w, x_i \rangle\}$  for  $x_i \in \mathbb{R}^d$  and  $y_i \in \mathbb{R}$  is convex

- $h(x) = f(x) + g(x)$  (Sum)

E.g.:  $f(w) = \max\{0, 1 - y_i \langle w, x_i \rangle\}$  is convex, and hence  $h(w) = \sum_{i=1}^n \max\{0, 1 - y_i \langle w, x_i \rangle\} + \frac{\lambda}{2} \|w\|^2$  is convex.

Hence, the SVM loss function:  $f(w) := \sum_{i=1}^n \max\{0, 1 - y_i \langle X_i, w \rangle\} + \frac{\lambda}{2} \|w\|^2$  is convex.

**Q:** Prove that  $\ell_1$ -regularized logistic regression:

$f(w) := \sum_{i=1}^n \log(1 + \exp(-y_i \langle X_i, w \rangle)) + \lambda \|w\|_1$  is convex.

We have proved that the logistic loss  $f(x) = \log(1 + \exp(-x))$  is convex. Since composition with an affine map is convex, and the sum of convex functions is convex, the first term is convex. Since all norms are convex, and a non-negative scaling of a convex function is convex, the second term is convex. Hence,  $f(w)$  is convex.

Another way to prove convexity for logistic regression is to compute the Hessian and show that it is positive semi-definite (In Assignment 1)

# Jensen's Inequality

- Recall the zero-order definition of convexity:  $\forall x, y \in \mathcal{D}$  and  $\theta \in [0, 1]$ ,  
 $f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y)$ .
- This can be generalized to  $n$  points  $\{x_1, x_2, \dots, x_n\}$ , i.e. for  $p_i \geq 0$  and  $\sum_i p_i = 1$ ,

$$f(p_1 x_1 + p_2 x_2 + \dots + p_n x_n) \leq p_1 f(x_1) + p_2 f(x_2) + \dots + p_n f(x_n) \implies f\left(\sum_{i=1}^n p_i x_i\right) \leq \sum_{i=1}^n p_i f(x_i)$$

- If  $X$  is a discrete r.v. that can take value  $x_i$  with probability  $p_i$ , and  $f$  is convex, then,

$$f(\mathbb{E}[X]) \leq \mathbb{E}[f(X)]. \quad (\text{Jensen's inequality})$$

- Jensen's inequality can be used to prove inequalities like the AM-GM inequality:  
 $\sqrt{ab} \leq \frac{a+b}{2}$ .
- Proof:* Choose  $f(x) = -\log(x)$  as the convex function, and consider two points  $a$  and  $b$  with  $\theta = 1/2$ . By Jensen's inequality,

$$-\log\left(\frac{a+b}{2}\right) \leq \frac{-\log(a) - \log(b)}{2} \implies \log\left(\frac{a+b}{2}\right) \geq \log(\sqrt{ab}) \implies \frac{a+b}{2} \geq \sqrt{ab}.$$



# Holder's Inequality

**Q:** Prove Holder's inequality, for  $p, q \geq 1$  s.t.  $\frac{1}{p} + \frac{1}{q} = 1$  and  $x, y \in R^d$ ,  $\langle x, y \rangle \leq \|x\|_p \|y\|_q$ .

By repeating the AM-GM proof, but for a general  $\theta \in [0, 1]$ , for  $a, b \geq 0$ , we can prove that,

$$a^\theta b^{1-\theta} \leq \theta a + (1 - \theta)b$$


Use  $a = \frac{|x_i|^p}{\sum_{j=1}^n |x_j|^p}$ ,  $b = \frac{|y_i|^q}{\sum_{j=1}^n |y_j|^q}$ ,  $\theta = 1/p$ , and using the fact that  $1 - \theta = 1 - 1/p = 1/q$

$$\left( \frac{|x_i|^p}{\sum_{j=1}^n |x_j|^p} \right)^{1/p} \left( \frac{|y_i|^q}{\sum_{j=1}^n |y_j|^q} \right)^{1/q} \leq \frac{1}{p} \frac{|x_i|^p}{\sum_{j=1}^n |x_j|^p} + \frac{1}{q} \frac{|y_i|^q}{\sum_{j=1}^n |y_j|^q}$$

Summing both sides from  $i = 1$  to  $n$ ,

$$\begin{aligned} \sum_{i=1}^n \frac{|x_i|}{\left( \sum_{j=1}^n |x_j|^p \right)^{1/p}} \frac{|y_i|}{\left( \sum_{j=1}^n |y_j|^q \right)^{1/q}} &\leq 1 \implies \sum_i x_i y_i \leq \left( \sum_{i=1}^n |x_i|^p \right)^{1/p} \left( \sum_{i=1}^n |y_i|^q \right)^{1/q} \\ &\implies \langle x, y \rangle \leq \|x\|_p \|y\|_q \end{aligned}$$

Questions?

-  Atilim Gunes Baydin, Robert Cornish, David Martinez Rubio, Mark Schmidt, and Frank Wood, *Online learning rate adaptation with hypergradient descent*, arXiv preprint arXiv:1703.04782 (2017).