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```
1 Pre-work
1.1 Linux
         //show the file
mkdir test //make a directory - test
cd test //go to test directory
cd ..
         //back to last directory
mv test test2 //if test2 exist, move test to test2, else
   rename test to test2
rm test //delete test
vim a.cpp //write code
{
       //insert mode(type)
      //normal mode(command)
//save and leave
  esc
  :wq
  normal + gg + G + d //delete all
  normal +the line your mouse stop+d //delete the
     Line
  normal + gg + V + G + " = " //indent(tab)
  normal + u //return to last step
  the leter your mouse stop + s //delete the letter and
      start insert mode
  !!DO NOT press ctrl + s!! //if so ,use ctrl+q
g++ a.cpp -o a.out
                        //compile
./a.out //run
1.2 Python
 cat > a.in //copy input to a.in,use ctrl+d to
      finish
                        //enter input in a.in and run
  . / a.out < a.in
  . / a.out<a.in>a.o //enter output to a.o
  cat a.o //show up the content in a.o
}
                 //should use python
  from random import*
                     //produce a random int in a-b
   randint(a, b)
    random(a, b)
                   //maybe produce a float
ctrl + c
           //force to stop
```

//run diff a.out b.out //look difference between two file

python3 test.py

echo "\$i"

//match

n = stdin.readline()

s.strip() //刪頭尾空格

i=ord('B') - ord('A') x=bin(x) //2(以0b 開頭) x=oct(x) //8(以00開頭) x=hex(x) //16(以0x開頭) x=int(x,2) //2轉10 x=int(x,8) //8轉10 x=int(x,16) //16轉10

for s in stdin: //連續輸入

print("%g" %(R)) //去除小數的0 輸出法

i='A' //字元轉數字 i=ord(i) #65

a,b = b,a #swap

do

}

for ((i = 0;; i++))

python3 gen.py > input . / ac < input > ac.out . / wa < input > wa.out diff ac.out wa.out || break

from sys import stdin //讀整行

a,b = map(int,input().split().) //映射輸入

1.3 Header

```
#include<bits/stdc++.h>
#define PI acos(-1.0)
#define e 2.7182818
#define LL long long
#define lowbit(x) (x&-x)
using namespace std;
ios::sync_with_stdio(0);
cin.tie(0); cout.tie(0);
cout << fixed << setprecision(10) <<;

1.4 Int

int -2,147,483,648 ~ 2,147,483,647

unsigned 0 ~ 4,294,967,295

long long -9,223,372,036,854,775,808 ~
9,223,372,036,854,775,807</pre>
```

unsigned long long 0 ~ 18,446,744,073,709,551,615

2 String

2.1 String to int

```
int num = atoi(str.c_str());
long long num = atoll(str.c_str());
```

2.2 Int to string

| to_string(num);

2.3 Lexicographically minimum string rotation

2.4 Hash

2.5 Trie

```
//create new trie
trie* root = new trie();
// 0(|s|)
void insert(string& s) {
    trie* now = root; //start at root everytime
    for (auto i : s) {
        now->sz++;
        if (now->nxt[i - 'a'] == NULL) {
            now->nxt[i - 'a'] = new trie();
        now = now->nxt[i - 'a']; //go to next letter
    now->cnt++;
    now->sz++;
// O(|s|)
int query_prefix(string& s) { //search how many prefix
    trie* now = root;
                          //start at root everytime
    for (auto i : s) {
        if (now->nxt[i - 'a'] == NULL) {
            return 0;
        now = now -> nxt[i - 'a'];
    return now->sz;
int query_count(string& s) { //search how many times s
     appear
    trie* now = root;
                          //start at root everytime
    for (auto i : s) {
        if (now->nxt[i - 'a'] == NULL) {
            return 0;
        now = now - > nxt[i - 'a'];
    return now->cnt;
```

2.6 Bitwise Trie

```
struct trie {
    trie* nxt[2];
                //how many number end at this node
    int cnt;
                //how many number's prefix include this
    int sz;
         node
    trie() :cnt(0), sz(0) {
    memset(nxt, 0, sizeof(nxt));
};
//create new trie
trie* root = new trie();
void insert(int x){
    trie *now = root;
                        //start at root everytime
    for(int i=30;i>=0;i--){
         now - > sz++
         if(now->nxt[x>>i&1] == NULL){
             now->nxt[x>>i&1] = new trie();
        now = now->nxt[x>>i&1]; / go to next
    }
    now->cnt++;
    now->sz++;
}
```

2.7 Double Hash

```
const ll P1 = 75577;
const ll P2 = 12721;
const ll MOD = 998244353;
pair<ll, ll> Hash[MXN];
void build(const string& s) {
    pair<ll, ll> val = make_pair(0, 0);
    for (int i = 0; i < s.size(); i++) {
        val.first = (val.first * P1 + s[i]) % MOD;
        val.second = (val.second * P2 + s[i]) % MOD;
        Hash[i] = val;
    }
}</pre>
```

2.8 Suffix Array

```
//S = "babd"
//而 SA[i] 存的就是第i小的後綴是第幾個字元開始的字串
//SA[0] = 1 ( "abd"
//SA[1] = 0 ( "babd"
//SA[2] = 2 ( "bd" )
//SA[3] = 3 ( "d" )
//Height 數組
//裡面儲存的值為第 i 小的字串
//跟第 i-1 小的字串共同前綴長度(LCP)
//H[0] = 0 (第0個字串沒有前一個)
//H[1] = 0
//H[2] = 1 ( "b" )
//應用:LCS最長共同子字串
//S = A + '@' + B
//就看i , i+1 字串為不同字串, 取 H[i+1] 的 max
//應用:相異字串數量
//全部子字串數量 - H[0~s.size()]即為答案
const int N = 300010; //字串的2倍大小
struct SA{
#define REP(i,n) for ( int i=0; i<int(n); i++ )</pre>
#define REP1(i,a,b) for ( int i=(a); i<=int(b); i++ )</pre>
  bool _t[N*2];
  int _s[N*2], _sa[N*2], _c[N*2], x[N], _p[N], _q[N*2],
       hei[N], r[N];
  int operator [] (int i){ return _sa[i]; }
  void build(int *s, int n, int m){
    memcpy(_s, s, sizeof(int) * n);
    sais(_s, _sa, _p, _q, _t, _c, n, m);
    mkhei(n);
  void mkhei(int n){
    REP(i,n) r[_sa[i]] = i;
    hei[0] = 0;
    REP(i,n) if(r[i]) {
      int ans = i>0 ? max(hei[r[i-1]] - 1, 0) : 0;
      while(_s[i+ans] == _s[_sa[r[i]-1]+ans]) ans++;
     hei[r[i]] = ans;
   }
 }
  void sais(int *s, int *sa, int *p, int *q, bool *t,
     int *c, int n, int z){
    bool uniq = t[n-1] = true, neq;
    int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
        lst = -1;
#define MSO(x,n) memset((x),0,n*sizeof(*(x)))
#define MAGIC(XD) MS0(sa, n); \
    memcpy(x, c, sizeof(int) * z); \
    XD; \
    memcpy(x + 1, c, sizeof(int) * (z - 1)); \
    REP(i,n) if(sa[i] && !t[sa[i]-1]) sa[x[s[sa[i
        ]-1]]++] = sa[i]-1; \setminus
    memcpy(x, c, sizeof(int) * z); \
    for(int i = n - 1; i >= 0; i--) if(sa[i] && t[sa[i
        ]-1]) sa[--x[s[sa[i]-1]]] = sa[i]-1;
    MS0(c, z);
    REP(i,n) uniq \&= ++c[s[i]] < 2;
    REP(i,z-1) c[i+1] += c[i];
    if (uniq) { REP(i,n) sa[--c[s[i]]] = i; return; }
    for(int i = n - 2; i >= 0; i--) t[i] = (s[i] == s[i]
        +1] ? t[i+1] : s[i] < s[i+1]);
    MAGIC(REP1(i,1,n-1) \  \, \textbf{if}(t[i] \  \, \& \  \, !t[i-1]) \  \, sa[--x[s[i
        ]]]=p[q[i]=nn++]=i);
    REP(i, n) if (sa[i] && t[sa[i]] && !t[sa[i]-1]) {
     neq=1st<0 \mid memcmp(s+sa[i],s+lst,(p[q[sa[i]]+1]-sa
          [i])*sizeof(int));
     ns[q[lst=sa[i]]]=nmxz+=neq;
    sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz
         + 1);
    MAGIC(for(int i = nn - 1; i >= 0; i--) sa[--x[s[p[
        nsa[i]]]]] = p[nsa[i]]);
 }
}sa;
int H[ N ], SA[ N ];
//ip[i]=(int)str[i]
void suffix_array(int* ip, int len) {
// should padding a zero in the back
```

```
// ip is int array, len is array length
// ip[0..n-1] != 0, and ip[len] = 0
ip[len++] = 0;
sa.build(ip, len, 128);
for (int i=0; i<len; i++) {
    H[i] = sa.hei[i + 1];
    SA[i] = sa._sa[i + 1];
}
// resulting height, sa array \in [0,len)
}
int ip[N];
for (int i = 0; i < str.size(); i++) {
    ip[i] = int(str[i]);
}
suffix_array(ip,str.size());</pre>
```

2.9 KMP

```
/*len - failure[k]:
在R結尾的情況下,這個子字串可以由開頭
長度為(Len - failure[k])的部分重複出現來表達
failure[k] :
    failure[k]為次長相同前綴後綴
    如果我們不只想求最多,而且以O - base做為考量
    , 那可能的長度由大到小會是
    failuer[k] \( failure[failuer[k] - 1]
     ` failure[failure[failuer[k] - 1] - 1]..
    直到有值為0為止*/
int failure[MXN];
void KMP(string& t, string& p)
{
    if (p.size() > t.size()) return;
    for (int i = 1, j = failure[0] = -1; i < p.size();</pre>
       ++i)
       while (j >= 0 \&\& p[j + 1] != p[i])
           j = failure[j];
        if (p[j + 1] == p[i]) j++;
        failure[i] = j;
    for (int i = 0, j = -1; i < t.size(); ++i)</pre>
        while (j \ge 0 \&\& p[j + 1] != t[i])
           j = failure[j];
        if (p[j + 1] == t[i]) j++;
        if (j == p.size() - 1)
           cout << i - p.size() + 1 << " ";</pre>
           j = failure[j];
       }
    }
```

2.10 Manacher

2.11 Palindromic-Tree 回文樹

```
// Len[state[i]]是對應的回文長度 Len[fail[i]]
// num[state[i]]是有幾個回文後綴
// cnt[state[i]]是這個回文子字串在整個字串中的出現次數
// fail[i]是他長度次長的回文後綴, aba的fail是a
// state[i] 以i為結尾的最長回文
const int MXN = 1000010;
struct PalT{
```

```
int nxt[MXN][26],fail[MXN],len[MXN];
  int tot,lst,n,state[MXN],cnt[MXN],num[MXN];
  int diff[MXN],sfail[MXN],fac[MXN],dp[MXN];
  char s[MXN]={-1};
  int newNode(int 1,int f){
    len[tot]=1,fail[tot]=f,cnt[tot]=num[tot]=0;
    memset(nxt[tot],0,sizeof(nxt[tot]));
    diff[tot]=(1>0?1-len[f]:0);
    sfail[tot]=(1>0&&diff[tot]==diff[f]?sfail[f]:f);
    return tot++;
  int getfail(int x){
    while(s[n-len[x]-1]!=s[n]) x=fail[x];
    return x;
  int getmin(int v){
    dp[v]=fac[n-len[sfail[v]]-diff[v]];
    if(diff[v]==diff[fail[v]])
        dp[v]=min(dp[v],dp[fail[v]]);
    return dp[v]+1;
  int push(){
    int c=s[n]-'a',np=getfail(lst);
    if(!(lst=nxt[np][c])){
      lst=newNode(len[np]+2,nxt[getfail(fail[np])][c]);
      nxt[np][c]=lst; num[lst]=num[fail[lst]]+1;
    fac[n]=n;
    for(int v=lst;len[v]>0;v=sfail[v])
        fac[n]=min(fac[n],getmin(v));
    return ++cnt[lst],lst;
  void init(const string _s){
    tot=1st=n=0;
    newNode(0,1),newNode(-1,1);
    for(;_s[n];) s[n+1]=_s[n],++n,state[n-1]=push();
    for(int i=tot-1;i>1;i--) cnt[fail[i]]+=cnt[i];
}palt;
```

3 STL

3.1 Vector

```
vec.insert(vec.begin()+i,a); //insert a in vec[i]
//erase vec[i]~vec[j-1]
vec.erase(vec.begin()+i,vec.begin()+j);
{
  vector<int>::iterator low, up;
  low = std::lower_bound(v.begin(), v.end(), 20);
  up = std::upper_bound(v.begin(), v.end(), 20);
}
//erase repeat number
{
  sort(vec.begin(), vec.end());
  vec.erase(unique(vec.begin(), vec.end()), iter, vec.end())}
}
```

3.2 Set

```
//sort in small to big,O(lg n)
              //get first value
*st.begin()
              //get last value
*st.rbegin()
              //return 0 represent not found
st.count(x)
st.find(x)
             //return st.end() represent not found
             //x can be value or iterator
st.erase(x)
st.clear() st.empty() st.size()
st.insert(st2.begin(),st2.end()) //insert st2[begin,
    end) to st
/*dont have operator[]*/
//customize compare
struct cmp {
  bool operator()(int lhs, int rhs)/*const*/{
    return lhs > rhs;
set<int, cmp> st;
```

3.3 Map

3.4 Priority queue

```
priority_queue<int> //number from big to small
  //number from small to big
priority_queue<int, vector<int>, greater<int> >num;
//function
empty size push pop top
```

3.5 Queue

```
empty size front back push pop
```

3.6 Stack

```
empty size back push pop top
```

3.7 Next permutation

```
//ex:123->132->213->231->312->321
int a[] = { 0,1,2,3,4,5,6,7,8,9 };
while (next_permutation(a, a + 10)){}
```

4 Divide and conquer

4.1 Merge sort

```
vector<int> merge_sort(vector<int> vec) {
  if (vec.size() == 1)return vec;
  vector<int> left(vec.begin(), vec.begin() + vec.size
      ()/2),
    right(vec.begin() + vec.size()/2, vec.end());
  left = merge_sort(left);
  right = merge_sort(right);
  vector<int> re;
  int i, j;
  for (i = 0, j=0; i < left.size() && j<right.size();) {</pre>
    if (left[i] < right[j]) re.push_back(left[i++]);</pre>
    else if(left[i] > right[j]) re.push_back(right[j
        ++]);
    else {
      re.push_back(left[i++]);
      re.push_back(right[j++]);
  while (i < left.size())re.push_back(left[i++]);</pre>
  while (j < right.size())re.push_back(right[j++]);</pre>
  return re;
```

4.2 Quick sort

```
vector<int> quick_sort(vector<int> vec) {
  if (vec.size() <= 1)return vec;</pre>
  srand(time(NULL));
  int pivot =rand() % vec.size();
  vector<int> left, right;
  for (int i = 0; i < vec.size(); i++) {</pre>
    if (i == pivot)continue;
    if (vec[i] > vec[pivot])right.push_back(vec[i]);
    else left.push_back(vec[i]);
  left = quick_sort(left);
  right = quick_sort(right);
  left.push_back(vec[pivot]);
  for (int i = 0; i < right.size(); i++)left.push_back(</pre>
      right[i]);
  return left;
}
```

4.3 Closest pair

```
struct Node{
 long int x, y;
bool cmp_y(Node a,Node b){
  if (a.y != b.y)
   return a.y < b.y;</pre>
    return a.x < b.x:
bool cmp_x(Node a, Node b) {
 if (a.x != b.x)
    return a.x < b.x;</pre>
  else
    return a.y < b.y;</pre>
void update_dis(Node a, Node b){
  mindis = min(mindis, sqrt((a.x - b.x) * (a.x - b.x) +
       (a.y - b.y) * (a.y - b.y));
vector<Node> node;
void closest(int l,int r) {
  if (r-1 <= 3) {//點少暴力拆解
    for (int i = 1; i < r; i++) {</pre>
      for (int j = i + 1; j < r;</pre>
                                    j++) {
        update_dis(node[i], node[j]);
    }
    return;
  }
  //點不少就拆2邊(分,治)
  int mid = (1 + r) \gg 1, midx = node[mid].x;
  closest(1,mid); closest(mid,r);
  //合
  vector<Node> all;
  int k = mid;
  while (midx - node[k].x < mindis) {</pre>
    all.push_back(node[k]);
    k - - :
    if (k < 1)break;</pre>
  k = mid + 1;
  while (node[k].x - midx < mindis) {</pre>
    all.push_back(node[k]);
    k++;
    if (k >= r)break;
  }
  sort(all.begin(), all.end(), cmp_y);
  for (int i = 0; i < all.size(); i++) {</pre>
    for (int j = i+1; j < all.size(); j++) {</pre>
      if (all[j].y - all[i].y > mindis)break;
      if ((all[i].x >= node[mid].x && all[j].x <= node[</pre>
          mid].x) \mid \mid (all[j].x >= node[mid].x && all[i]
           ].x <= node[mid].x)) {
        update_dis(all[i], all[j]);
   }
 }
int main() {
 ios::sync_with_stdio(0);
  cin.tie(0);
  cout.tie(0);
  int n;
 cin >> n;
 node.resize(n);
 for (int i = 0; i < n; i++) {</pre>
   cin >> node[i].x >> node[i].y;
  sort(node.begin(), node.end(), cmp_x);
 closest(0,n);
  cout << fixed << setprecision(4) << mindis << "\n";</pre>
```

5 DP

5.1 Coin

```
void findway() {
  way[0] = 1;
  int price[5] = { 1,5,10,25,50 };
```

```
for (int i = 0; i < 5; i++) {
   for (int j = price[i]; j < MXN; j++) {
     way[j] += way[j - price[i]];
   }
}</pre>
```

5.2 Knapsack (01backpack)

```
void knapsack(int n, int w){
                                             //--2dim
     memset(c, 0, sizeof(c));
for (int i = 0; i < n; ++i)</pre>
          for (int j = 0; j <= w; ++j)</pre>
               if (j - weight[i] < 0)</pre>
                    c[i+1][j] = c[i][j];
                    c[i+1][j] = max(c[i][j],c[i][j - weight]
                         [i]] + cost[i]);
     cout << "max value:" << c[n][w];</pre>
void knapsack(int n, int w) {
                                          //--1dim
     memset(c, 0, sizeof(c));
for (int i = 0; i < n; ++i)</pre>
          for (int j = w; j - weight[i] >= 0; --j)
               c[j] = max(c[j], c[j - weight[i]] + cost[i]
                    ]);
     cout << "max value:" << c[w];</pre>
}
```

5.3 LCS

5.4 LIS

```
int LIS(const vector<int>& arr){
   vector<int> lis;
   for(auto i:arr){
      if(lis.empty() || lis.back() < i) lis.
            push_back(i);
      else      *lower_bound(lis.begin(),lis.end(),i) =
            i;
   }
   return lis.size();
}</pre>
```

5.5 Max-sum of sub-array

dp[i]=max(dp[i-1],0)+a[i]

6 Flow

6.1 Flow

```
#define PB push_back
#define SZ(x) (int)x.size()
using namespace std;
const int MXN = 10000;
struct Dinic {
  struct Edge { int v, f, re; };
  int n, s, t, level[MXN];
  vector<Edge> E[MXN];
  void init(int _n, int _s, int _t) {
    n = _n; s = _s; t = _t;
    for (int i = 0; i < n; i++) E[i].clear();</pre>
  void add_edge(int u, int v, int f) {
    E[u].PB({ v,f,SZ(E[v]) });
    E[v].PB({ u,0,SZ(E[u]) - 1 });
  bool BFS() {
    for (int i = 0; i < n; i++) level[i] = -1;</pre>
    queue<int> que;
    que.push(s);
    level[s] = 0;
    while (!que.empty()) {
      int u = que.front(); que.pop();
      for (auto it : E[u]) {
        if (it.f > 0 && level[it.v] == -1) {
```

```
level[it.v] = level[u] + 1;
          que.push(it.v);
     }
    return level[t] != -1;
  int DFS(int u, int nf) {
    if (u == t) return nf;
    int res = 0;
    for (auto &it : E[u]) {
      if (it.f > 0 && level[it.v] == level[u] + 1) {
        int tf = DFS(it.v, min(nf, it.f));
        res += tf; nf -= tf; it.f -= tf;
        E[it.v][it.re].f += tf;
        if (nf == 0) return res;
      }
    if (!res) level[u] = -1;
    return res;
  int flow(int res = 0) {
    while (BFS())
      res += DFS(s, 2147483647);
    return res:
}flow;
6.2 KM
```

```
//perfect match, return max value
           //找最小值時,代入-W
struct KM{
  static const int MXN = 2001; // 1-based
  static const 11 INF = 0x3f3f3f3f;
  int n, mx[MXN], my[MXN], pa[MXN];
  11 g[MXN][MXN], lx[MXN], ly[MXN], sy[MXN];
  bool vx[MXN], vy[MXN];
  void init(int _n) {
    n = _n;
    for(int i=1; i<=n; i++)</pre>
  fill(g[i], g[i]+n+1, 0);//找最小值時, 0變-INF
  void addEdge(int x, int y, ll w) {g[x][y] = w;}
  void augment(int y) {
    for(int x, z; y; y = z)
      x=pa[y], z=mx[x], my[y]=x, mx[x]=y;
  void bfs(int st) {
    for(int i=1; i<=n; ++i) sy[i]=INF, vx[i]=vy[i]=0;</pre>
    queue<int> q; q.push(st);
    for(;;) {
      while(q.size()) {
        int x=q.front(); q.pop(); vx[x]=1;
        for(int y=1; y<=n; ++y) if(!vy[y]){</pre>
          11 t = 1x[x]+1y[y]-g[x][y];
          if(t==0){
            pa[v]=x:
            if(!my[y]){augment(y);return;}
             vy[y]=1, q.push(my[y]);
          }else if(sy[y]>t) pa[y]=x,sy[y]=t;
        }
      11 cut = INF;
      for(int y=1; y<=n; ++y)</pre>
        if(!vy[y]&&cut>sy[y]) cut=sy[y];
      for(int j=1; j<=n; ++j){</pre>
        if(vx[j]) lx[j] -= cut;
        if(vy[j]) ly[j] += cut;
        else sy[j] -= cut;
      for(int y=1; y<=n; ++y) if(!vy[y]&&sy[y]==0){</pre>
        if(!my[y]){augment(y);return;}
        vy[y]=1, q.push(my[y]);
      }
   }
  11 solve(){
    fill(mx, mx+n+1, 0); fill(my, my+n+1, 0);
fill(ly, ly+n+1, 0); fill(lx, lx+n+1, -INF);
    for(int x=1; x<=n; ++x) for(int y=1; y<=n; ++y)</pre>
      lx[x] = max(lx[x], g[x][y]);
    for(int x=1; x<=n; ++x) bfs(x);</pre>
```

```
11 \text{ ans} = 0;
     for(int y=1; y<=n; ++y) ans += g[my[y]][y];</pre>
  }
}graph;
```

6.3 MinCostFlow

```
struct MinCostMaxFlow {
   typedef int Tcost;
   static const int MAXV = 20010;
   static const int INFf = 1000000;
   static const Tcost INFc = 1e9;
   struct Edge {
       int v, cap;
       Tcost w:
       int rev:
       Edge() {}
       Edge(int t2, int t3, Tcost t4, int t5)
           : v(t2), cap(t3), w(t4), rev(t5) {}
   int V, s, t;
   vector<Edge> g[MAXV];
                 //sum, start, end
   void init(int n, int _s, int _t) {
    V = n; s = _s; t = _t;
       for (int i = 0; i <= V; i++) g[i].clear();</pre>
                       //start,end,capacity,cost
   void addEdge(int a, int b, int cap, Tcost w) {
       g[a].push_back(Edge(b, cap, w, (int)g[b].size())
       g[b].push_back(Edge(a, 0, -w, (int)g[a].size() -
             1));
   Tcost d[MAXV];
   int id[MAXV], mom[MAXV];
   bool inqu[MAXV];
   queue<int> q;
   pair<int, Tcost> solve() {
       int mxf = 0; Tcost mnc = 0;
       while (1) {
           fill(d, d + 1 + V, INFc);
            fill(inqu, inqu + 1 + V, 0);
            fill(mom, mom + 1 + V, -1);
           mom[s] = s;
           d[s] = 0;
           q.push(s); inqu[s] = 1;
            while (q.size()) {
                int u = q.front(); q.pop();
                inqu[u] = 0;
                for (int i = 0; i < (int)g[u].size(); i</pre>
                    ++) {
                    Edge& e = g[u][i];
                    int v = e.v;
                    if (e.cap > 0 && d[v] > d[u] + e.w)
                        d[v] = d[u] + e.w;
                        mom[v] = u;
                        id[v] = i;
                        if (!inqu[v]) q.push(v), inqu[v]
                              = 1;
                    }
                }
            if (mom[t] == -1) break;
           int df = INFf;
            for (int u = t; u != s; u = mom[u])
                df = min(df, g[mom[u]][id[u]].cap);
           for (int u = t; u != s; u = mom[u]) {
                Edge& e = g[mom[u]][id[u]];
                e.cap -= df;
                g[e.v][e.rev].cap += df;
           mxf += df;
           mnc += df * d[t];
       return { mxf,mnc };
}flow;
```

7 Graph

7.1 BCC

```
//return a 2dim array,if size=2 represent this is a
    bridge
//else if size>=3 represent this is a group
#define MXN (int)(1e5+5)
#define PB push_back
#define REP(i,x) for(int i=0;i<x;i++)</pre>
struct BccVertex {
  int n,nScc,step,dfn[MXN],low[MXN];
  vector<int> E[MXN],sccv[MXN];
  int top,stk[MXN];
  void init(int _n) {
    n = _n; nScc = step = 0;
for (int i=0; i<n; i++) E[i].clear();</pre>
  void addEdge(int u, int v)
  { E[u].PB(v); E[v].PB(u); }
  void DFS(int u, int f) {
    dfn[u] = low[u] = step++;
    stk[top++] = u;
    for (auto v:E[u]) +
      if (v == f) continue;
if (dfn[v] == -1) {
        DFS(v,u);
        low[u] = min(low[u], low[v]);
        if (low[v] >= dfn[u]) {
           int z;
           sccv[nScc].clear();
           do {
             z = stk[--top];
             sccv[nScc].PB(z);
           } while (z != v);
           sccv[nScc++].PB(u);
        }
      }
      else
        low[u] = min(low[u],dfn[v]);
    }
  }
  vector<vector<int>> solve() {
    vector<vector<int>> res;
    for (int i=0; i<n; i++)</pre>
      dfn[i] = low[i] = -1;
    for (int i=0; i<n; i++)</pre>
      if (dfn[i] == -1) {
        top = 0;
        DFS(i,i);
    REP(i,nScc) res.PB(sccv[i]);
    return res;
}graph;
```

7.2 Max Clique

```
struct MaxClique{
  static const int MAXN=105;
  int N,ans;
  int g[MAXN][MAXN], dp[MAXN], stk[MAXN][MAXN];
  int sol[MAXN], tmp[MAXN]; //sol[0~ans-1]is answer
  void init(int n){
   N=n;
            //0-base
    memset(g,0,sizeof(g));
  void add_edge(int u,int v){
    g[u][v]=g[v][u]=1;
  int dfs(int ns,int dep){
    if(!ns){
      if(dep>ans){
        ans=dep;
        memcpy(sol,tmp,sizeof tmp);
        return 1;
      }else return 0;
    for(int i=0;i<ns;++i){</pre>
      if(dep+ns-i<=ans)return 0;</pre>
      int u=stk[dep][i],cnt=0;
      if(dep+dp[u]<=ans)return 0;</pre>
```

7.3 SCC

```
#define PB push_back
#define MXN 10000
#define FZ(x) memset(x,0,sizeof(0))
struct Scc {
  int n, nScc, vst[MXN], bln[MXN]; // 最後每個點所屬的
      連通分量存在bLn陣列
  vector<int> E[MXN], rE[MXN], vec;
  void init(int _n) { //先初始化點的數量
   n = _n;
for (int i = 0; i < MXN; i++)</pre>
      E[i].clear(), rE[i].clear();
  void addEdge(int u, int v) { // 加有向邊
    E[u].PB(v); rE[v].PB(u);
  void DFS(int u) {
    vst[u] = 1;
    for (auto v : E[u]) if (!vst[v]) DFS(v);
    vec.PB(u);
  void rDFS(int u) {
    vst[u] = 1; bln[u] = nScc;
    for (auto v : rE[u]) if (!vst[v]) rDFS(v);
  void solve() { // 跑 kosaraju
   nScc = 0;
    vec.clear();
    FZ(vst);
    for (int i = 0; i < n; i++)</pre>
     if (!vst[i]) DFS(i);
    reverse(vec.begin(), vec.end());
    FZ(vst);
    for (auto v : vec)
      if (!vst[v]) {
        rDFS(v); nScc++;
  }
}scc;
```

7.4 SPFA 差分約束

```
//判斷負環,差分約束
//差分約束:
//xj - xi <= k
//連接一條邊 連接一條邊(i,j), 權重為 k
//最後再設置一個起點 S,連向所有邊邊權為 O
//從起點 s,跑 SPFA,若出現負環則代表這組不等式無解
bool spfa(){
   deque<int> dq;
   dis[0]=0;
   dq.push_back(0);
   inq[0]=1;
   while(!dq.empty()){
       int u=dq.front();
       dq.pop_front();
       inq[u]=0;
       for(auto i:edge[u]){
          if(dis[i.first]>i.second+dis[u]){
              dis[i.first]=i.second+dis[u];
```

7.5 歐拉路徑-迴路

```
//存在歐拉的條件
//歐拉迴路
//無向圖:所有點的度數為偶數
//有向圖:所有點入度等於出度
//歐拉路徑
//無向圖:度數為奇數的點數量不超過2
//有向圖:全部點的入度出度一樣,
//
        或剛好一個點出度-1=入度,
//
        另一點入度-1=出度,
        其他點入度等於出度
// 且圖連通!!!! <反例> a->b b->a c->d d->c
vector<int> path;
void dfs(int x){
   while(!edge[x].empty()){
      int u = edge[x].back();
       edge[x].pop_back();
      dfs(u);
   path.push_back(x);
}
int main(){
   build();
   dfs(st);
   reverse(path.begin(),path.end());
   for(int i:path)
                   cout<<i<<'
   cout<<endl;
}
```

7.6 MinimumMeanCycle

```
#include<cfloat> //for DBL_MAX
int dp[MAXN][MAXN]; // 1-base,O(NM)
vector<tuple<int, int, int>> edge;
double mmc(int n) {//allow negative weight
  const int INF = 0x3f3f3f3f;
  for (int t = 0; t < n; ++t) {</pre>
    memset(dp[t + 1], 0x3f, sizeof(dp[t + 1]));
    for (const auto& e : edge) {
      int u, v, w;
      tie(u, v, w) = e;
      dp[t + 1][v] = min(dp[t + 1][v], dp[t][u] + w);
  double res = DBL_MAX;
 for (int u = 1; u <= n; ++u) {</pre>
   if (dp[n][u] == INF) continue;
    double val = -DBL_MAX;
    for (int t = 0; t < n; ++t)
      val = max(val, (dp[n][u] - dp[t][u]) * 1.0 / (n - val)
           t));
    res = min(res, val);
 return res; //if there are no cycle return DBL_MAX
```

7.7 Minimum Spanning Tree(kruskal)

```
struct Edge{
  int u,v,w;
  friend bool operator<(const Edge& lhs,const Edge&
      rhs){
      return lhs.w<rhs.w;
  }</pre>
```

```
};
vector<Edge> graph;
void kruskal(){
    int sum=0;
    sort(graph.begin(),graph.end());
    for(auto i:graph){
        if(Find(i.u)!=Find(i.v)){
            Union(find(i.u),find(i.v));
            sum+=i.w;
        }
    }
    cout<<sum<<endl;
}</pre>
```

7.8 Minimum Spanning Tree(prim)

```
void prim(){
    v.clear();v.resize(n);
    priority_queue<pair<ll,int>,vector<pair<ll,int>>,
        greater<pair<11,int>>> pq;
    pq.push({0,0});
    11 sum=0;
    while(!pq.empty()){
        auto u=pq.top();pq.pop();
        if(v[u.second]) continue;
        v[u.second]=1;
        sum+=u.first;
        for(auto i:edge[u.second]){
            if(!v[i.first]){
                 pq.push({i.second,i.first});
        }
    cout<<sum<<endl;</pre>
}
```

7.9 Single Source Shortest Paths(dijkstra)

```
void dijkstra(int startPoint,int endPoint){
    priority_queue<pair<ll,int>, vector<pair<ll,int>>,
        greater<pair<11,int>>> pq;
    v.clear();v.resize(n);
    dis.clear();dis.resize(n,INF);
    dis[startPoint]=0;
    pq.push({dis[startPoint],startPoint});//push
        startpoint into pq
    while(!pq.empty()){//if pq is not empty then
        continue
       auto u=pq.top();pq.pop(); //pop the point that
           is closet to startpoint everytime
       if(v[u.second]) continue; //if the point is
           visited represent there already have shorter
            path and dont have to walk again
       v[u.second]=1;
                            //set the point visited
       for(auto i:edge[u.second]){
          if(dis[i.first]>u.first+i.second){//determine
               whether it can relax
              dis[i.first]=u.first+i.second;
              pq.push({dis[i.first],i.first});//connect
                   the path that can relax
       }
    cout<<dis[endPoint]<<endl;</pre>
```

7.10 Union find

7.11 Maxium General Weighted Matching

```
//一般圖帶權匹配
//满足最大匹配情況下最大化權重
struct WeightGraph {
 static const int INF = INT_MAX;
  static const int N = 514;
  struct edge{
    int u,v,w; edge(){}
    edge(int ui,int vi,int wi)
      :u(ui),v(vi),w(wi){}
  int n,n_x;
  edge g[N*2][N*2];
  int lab[N*2];
  int match[N*2],slack[N*2],st[N*2],pa[N*2];
  int flo_from[N*2][N+1],S[N*2],vis[N*2];
  vector<int> flo[N*2];
  queue<int> q;
  int e_delta(const edge &e){
   return lab[e.u]+lab[e.v]-g[e.u][e.v].w*2;
  void update_slack(int u,int x){
    if(!slack[x]||e_delta(g[u][x])<e_delta(g[slack[x]][</pre>
        x]))slack[x]=u;
  void set_slack(int x){
    slack[x]=0;
    for(int u=1;u<=n;++u)</pre>
      if(g[u][x].w>0&&st[u]!=x&&S[st[u]]==0)
        update_slack(u,x);
  void q_push(int x){
    if(x<=n)q.push(x);</pre>
    else for(size_t i=0;i<flo[x].size();i++)</pre>
      q_push(flo[x][i]);
  void set_st(int x,int b){
    st[x]=b;
    if(x>n)for(size_t i=0;i<flo[x].size();++i)</pre>
      set_st(flo[x][i],b);
  int get_pr(int b,int xr){
    int pr=find(flo[b].begin(),flo[b].end(),xr)-flo[b].
        begin();
    if(pr%2==1){
      reverse(flo[b].begin()+1,flo[b].end());
      return (int)flo[b].size()-pr;
    }else return pr;
  void set_match(int u,int v){
    match[u]=g[u][v].v;
    if(u<=n) return;</pre>
    edge e=g[u][v];
    int xr=flo_from[u][e.u],pr=get_pr(u,xr);
    for(int i=0;i<pr;++i)set_match(flo[u][i],flo[u][i</pre>
        ^1]);
    set_match(xr,v);
    rotate(flo[u].begin(),flo[u].begin()+pr,flo[u].end
  void augment(int u,int v){
    for(;;){
      int xnv=st[match[u]];
      set_match(u,v);
      if(!xnv)return;
      set_match(xnv,st[pa[xnv]]);
      u=st[pa[xnv]],v=xnv;
  int get_lca(int u,int v){
    static int t=0;
    for(++t;u||v;swap(u,v)){
      if(u==0)continue;
      if(vis[u]==t)return u;
      vis[u]=t;
      u=st[match[u]];
      if(u)u=st[pa[u]];
    return 0:
  void add_blossom(int u,int lca,int v){
    int b=n+1;
```

```
while(b<=n_x&&st[b])++b;</pre>
  if(b>n_x)++n_x;
  lab[b]=0,S[b]=0;
  match[b]=match[lca];
  flo[b].clear();
  flo[b].push_back(lca);
  for(int x=u,y;x!=lca;x=st[pa[y]])
    flo[b].push_back(x),flo[b].push_back(y=st[match[x
        ]]),q push(y);
  reverse(flo[b].begin()+1,flo[b].end());
  for(int x=v,y;x!=lca;x=st[pa[y]])
    flo[b].push_back(x),flo[b].push_back(y=st[match[x
        ]]),q_push(y);
  set_st(b,b);
  for(int x=1;x<=n_x;++x)g[b][x].w=g[x][b].w=0;</pre>
  for(int x=1;x<=n;++x)flo_from[b][x]=0;</pre>
  for(size_t i=0;i<flo[b].size();++i){</pre>
    int xs=flo[b][i];
    for(int x=1;x<=n_x;++x)</pre>
      if(g[b][x].w==0||e_delta(g[xs][x])< e_delta(g[b])
           ][x]))
        g[b][x]=g[xs][x],g[x][b]=g[x][xs];
    for(int x=1;x<=n;++x)</pre>
      if(flo_from[xs][x])flo_from[b][x]=xs;
  set_slack(b);
void expand_blossom(int b){
  for(size_t i=0;i<flo[b].size();++i)</pre>
    set_st(flo[b][i],flo[b][i]);
  int xr=flo_from[b][g[b][pa[b]].u],pr=get_pr(b,xr);
  for(int i=0;i<pr;i+=2){</pre>
    int xs=flo[b][i],xns=flo[b][i+1];
    pa[xs]=g[xns][xs].u;
    S[xs]=1,S[xns]=0;
    slack[xs]=0,set_slack(xns);
    q_push(xns);
  S[xr]=1,pa[xr]=pa[b];
  for(size_t i=pr+1;i<flo[b].size();++i){</pre>
    int xs=flo[b][i];
    S[xs]=-1,set_slack(xs);
  st[b]=0;
bool on_found_edge(const edge &e){
  int u=st[e.u],v=st[e.v];
  if(S[v]==-1){
    pa[v]=e.u,S[v]=1;
    int nu=st[match[v]];
    slack[v]=slack[nu]=0;
    S[nu]=0,q_push(nu);
  }else if(S[v]==0){
    int lca=get_lca(u,v);
    if(!lca)return augment(u,v),augment(v,u),true;
    else add_blossom(u,lca,v);
  }
  return false;
bool matching(){
  memset(S+1,-1,sizeof(int)*n_x);
  memset(slack+1,0,sizeof(int)*n_x);
  a=queue<int>();
  for(int x=1;x<=n_x;++x)</pre>
    if(st[x]==x&&!match[x])pa[x]=0,S[x]=0,q_push(x);
  if(q.empty())return false;
  for(;;){
    while(q.size()){
      int u=q.front();q.pop();
      if(S[st[u]]==1)continue;
      for(int v=1; v<=n; ++v)</pre>
        if(g[u][v].w>0&&st[u]!=st[v]){
          if(e_delta(g[u][v])==0){
            if(on_found_edge(g[u][v]))return true;
          }else update_slack(u,st[v]);
    int d=INF;
    for(int b=n+1;b<=n_x;++b)</pre>
      if(st[b]==b&&S[b]==1)d=min(d,lab[b]/2);
    for(int x=1;x<=n_x;++x)</pre>
      if(st[x]==x&&slack[x]){
        if(S[x]==-1)d=min(d,e_delta(g[slack[x]][x]));
```

```
else if(S[x]==0)d=min(d,e_delta(g[slack[x]][x
               ])/2);
      for(int u=1;u<=n;++u){</pre>
        if(S[st[u]]==0){
           if(lab[u]<=d)return 0;</pre>
           lab[u]-=d;
        }else if(S[st[u]]==1)lab[u]+=d;
       for(int b=n+1;b<=n_x;++b)</pre>
        if(st[b]==b){
           if(S[st[b]]==0)lab[b]+=d*2;
           else if(S[st[b]]==1)lab[b]-=d*2;
      q=queue<int>();
       for(int x=1;x<=n_x;++x)</pre>
        if(st[x]==x&&slack[x]&&st[slack[x]]!=x&&e_delta
             (g[slack[x]][x])==0)
           if(on_found_edge(g[slack[x]][x]))return true;
      for(int b=n+1;b<=n x;++b)</pre>
        if(st[b]==b&&S[b]==1&&lab[b]==0)expand_blossom(
             b);
    return false;
  }
  pair<long long,int> solve(){
    memset(match+1,0,sizeof(int)*n);
    n x=n;
    int n_matches=0;
    long long tot_weight=0;
    for(int u=0;u<=n;++u)st[u]=u,flo[u].clear();</pre>
    int w_max=0;
    for(int u=1;u<=n;++u)</pre>
      for(int v=1;v<=n;++v){</pre>
        flo_from[u][v]=(u==v?u:0);
        w_max=max(w_max,g[u][v].w);
    for(int u=1;u<=n;++u)lab[u]=w_max;</pre>
    while(matching())++n_matches;
    for(int u=1;u<=n;++u)</pre>
      if(match[u]&&match[u]<u)</pre>
         tot_weight+=g[u][match[u]].w;
    return make_pair(tot_weight,n_matches);
  void add_edge( int ui , int vi , int wi ){ //無向圖
    g[ui][vi].w = g[vi][ui].w = wi;
  void init( int _n ){ //1-base
    n = _n;
for(int u=1;u<=n;++u)</pre>
       for(int v=1; v<=n; ++v)</pre>
        g[u][v]=edge(u,v,0);
} graph;
```

7.12 Maxium General graph Matching

```
//一般圖匹配 用在無向圖
//<cf> flow 只能用在2分匹配
// should shuffle vertices and edges 要打亂輸入的邊和
mt19937 gen(chrono::steady_clock::now().
    time_since_epoch().count());
for(int i=1;i<=n;i++)</pre>
 shuffle(edge[i].begin(),edge[i].end(),gen);
把點打亂: shuffle(ind.begin(),ind.end(),gen);
                                               //ind是
edge[ind[a]].push_back(ind[b]);
edge[ind[b]].push_back(ind[a]);
const int N=100005, E=(2e5)*2+40;
struct Graph{ // 1-based; match: i <-> lnk[i]
  int to[E],bro[E],head[N],e,lnk[N],vis[N],stp,n;
  void init(int _n){
   stp=0; e=1; n=_n;
   for(int i=1;i<=n;i++) head[i]=lnk[i]=vis[i]=0;</pre>
  void add_edge(int u,int v){
   to[e]=v,bro[e]=head[u],head[u]=e++;
```

```
to[e]=u,bro[e]=head[v],head[v]=e++;
  bool dfs(int x){
    vis[x]=stp;
    for(int i=head[x];i;i=bro[i]){
      int v=to[i];
      if(!lnk[v]){ lnk[x]=v,lnk[v]=x; return true; }
    for(int i=head[x];i;i=bro[i]){
      int v=to[i];
      if(vis[lnk[v]]<stp){</pre>
        int w=lnk[v]; lnk[x]=v,lnk[v]=x,lnk[w]=0;
        if(dfs(w)) return true;
        lnk[w]=v, lnk[v]=w, lnk[x]=0;
      }
    }
    return false;
  int solve(){
    int ans=0;
    for(int i=1;i<=n;i++) if(!lnk[i]) stp++,ans+=dfs(i)</pre>
    return ans;
  }
}graph;
```

8 Data structure

8.1 Point Modify

build(0,0,n-1);

```
#define cl(x) (x<<1)+1
#define cr(x) (x << 1) + 2
int arr[MXN];
int tree[MXN*4];
void build(int index,int left,int right){
    if( left == right ){
        tree[index] = arr[left];
        return:
    int mid=(left + right)/2;
    build(cl(index),left,mid);
    build(cr(index),mid+1,right);
    tree[index] = max(tree[cl(x)] ,tree[cr(x)]);
int query(int index,int left,int right,int query_left,
    int query_right){
    if( query_left <= left && right <= query_right){</pre>
        return tree[index];
    int mid=(left + right)/2;
    int ans=-INF;
    if(query_left <= mid){</pre>
        ans = max(ans, query(cl(index),left,mid,
            query_left,query_right));
    if(query_right > mid){
        ans = max(ans, query(cr(index),mid+1,right,
             query_left,query_right));
    return ans;
void update(int index,int left,int right,int position,
    int value){
    if(left == right){
        tree[index] = value;
        return;
    int mid=(left+right)/2;
    if(position <= mid){</pre>
        update(cl(index),left,mid,position,value);
    }
    esle{
        update(cr(index),mid+1,right,position,value);
    tree[index] = max(tree[cl(index)], tree[cr(index)]);
int main(){
```

```
query(0,0,n-1,2,7);
8.2 Interval Modify
#define cl(x) (x << 1) + 1
#define cr(x) (x << 1) + 2
#define INF 1e9
struct seg_tree {//0-base
  static const int MXN = 1e5 + 5, NO_TAG = 0; //to be
  ll a[MXN], val[MXN * 4], tag[MXN * 4], v;
  int n, ql, qr;
  void push(int i, int l, int r) {
    if (tag[i] != NO_TAG) {
      val[i] += tag[i]; //update by tag
      if (1 != r) {
        tag[cl(i)] += tag[i]; //push
        tag[cr(i)] += tag[i]; //push
      tag[i] = NO_TAG;
    }
  void pull(int i, int l, int r) {
    int mid = (1 + r) \gg 1;
    push(cl(i), l, mid); push(cr(i), mid + 1, r);
    val[i] = max(val[cl(i)], val[cr(i)]); //pull
  void build(int i, int l, int r) {
    if (1 == r) {
      val[i] = a[1]; //set value
    int mid = (1 + r) >> 1;
    build(cl(i), l, mid); build(cr(i), mid + 1, r);
    pull(i, 1, r);
  void update(int i, int l, int r) {
    push(i, 1, r);
    if (ql == 1 && r == qr) {
      tag[i] += v; //update tag
      return:
    int mid = (1 + r) >> 1;
    if (ql <= mid) update(cl(i), l, mid);</pre>
    if (qr > mid) update(cr(i), mid + 1, r);
    pull(i, 1, r);
  void query(int i, int l, int r) {
    push(i, 1, r);
    if (ql <= 1 && r <= qr) {</pre>
      v = max(v, val[i]); //update answer
      return;
    int mid = (1 + r) >> 1;
    if (ql <= mid) query(cl(i), l, mid);</pre>
    if (qr > mid) query(cr(i), mid + 1, r);
  int Query(int _ql, int _qr) { //傳 入 詢 問 區 間 v = -INF, ql = _ql, qr = _qr;
    query(0, 0, n - 1);
    return v;
  void Update(int _v, int _ql, int _qr) { //傳 入 更 新
       值,區間
    v = v, ql = ql, qr = qr;
    update(0, 0, n - 1);
  void init() {
    memset(tag, 0, sizeof(tag));
    memset(val, 0, sizeof(val));
}tree;
8.3 BIT
```

```
void update(int x){
  for (;x < MXN;x+= lowbit(x)) {
    BIT[x] += 1;
  }
}
ll query(int x){</pre>
```

```
11 \text{ ans} = 0;
  for (;x > 0; x \rightarrow lowbit(x)) {
   ans += BIT[x];
  }
  return ans;
map<ll, int> v_idx; //點對應的idx(離散化)
for (int i = 1; i <= n; i++)</pre>
    cin >> arr[i];
    v_idx[arr[i]] = 1;
    v_idx[k * arr[i]] = 1;
  map<ll, int>::iterator iter;
  int idx = 0;
  for (iter = v_idx.begin(); iter != v_idx.end(); iter
      ++) //儲存idx
    iter->second = ++idx;
  for (int i = 1; i <= n; i ++)</pre>
    ans += query(idx) - query(v_idx[k*arr[i]]);//所有已
        加進去的-符合條件的=不符合的
    //或ans += i-1 - query(v_idx[k*arr[i]]);
    update(v_idx[arr[i]]);
```

9 Tree

9.1 LCA

```
vector<vector<int>> tree;
vector<vector<int>> anc;
vector<int> timeIn, timeOut;
int ti = 0;
void build(int x, int fa) {
    anc[x].resize(__lg(n)+10);//tle改+1
    for (int i = 0; i < __lg(n) + 10; i++) {</pre>
        anc[x][i] = fa;
        fa = anc[fa][i];
    }
void dfs(int x, int fa) {
    timeIn[x] = ti++;
    build(x, fa);
    for (int i = 0; i < tree[x].size(); i++) {</pre>
        if (tree[x][i] == fa) continue;
        dfs(tree[x][i], x);
    timeOut[x] = ti++;
bool isAnc(int a, int b) {
    if (timeIn[a] <= timeIn[b] && timeOut[a] >= timeOut
        [b])return 1; //a是祖先
    return 0:
int query(int a, int b) {
   if (isAnc(a, b))return a;
    if (isAnc(b, a))return b;
    for (int i = __lg(n) + 10 - 1; i >= 0; i--) {
        if (!isAnc(anc[a][i], b)) a = anc[a][i];
    return anc[a][0];
int nlca(int x, int y) {
                            //x's yth anc
    if (y == 0)return x;
    int tp = log2(y);
    if (y == (1 << tp)) return anc[x][tp];</pre>
    else return nlca(anc[x][tp], y - (1 << tp));</pre>
```

9.2 Persistent Segment Tree 持久化

```
      struct node{
      11 val;

      node *1, *r;
      );

      vector<node *> version;
      //用一個vector紀錄全部版本的根節點

      //線段樹
```

```
void build(node *now_version, 1, r);
11 query(node *now_version, 1, r, q1, qr);
node *update_version(node *pre_version,int 1,int r,int
    pos, int v); //回傳新建的節點
void add_version(int x,int v){
                                  //修改位置 x 的值為 v
    version.push_back(update_version(version.back(), 0,
         n-1, x, v));
}
node *update_version(node *pre_version,int 1,int r,int
    pos, int v){
    node *x = new node();
                           //當前位置建立新節點
    if(1 == r){
       x \rightarrow val = v;
        return x;
    int mid = (l+r)>>1;
    if(pos <= mid){ //更新左邊
       x->l = update(pre_version->l, l, mid, pos, v);
            //左邊節點連向新節點
        x->r = pre->version->r;
           //右邊連到原本的右邊
    else{ //更新右邊
       x->1 = pre->version->1;
                                     //左邊連到原本的左
        x->r = update(pre_version->r, r, mid, pos, v);
             //右邊節點連向新節點
    x->val = x->l->val + x->r->val;
   return x:
//並查集
void build(node* now, int left, int right) {
 if (left == right) {
   now->fa = left;
   now->sz = 1;
    return;
 int mid = (left + right) / 2;
 now->cl = new node;
 now->cr = new node;
 build(now->cl, left, mid);
 build(now->cr, mid + 1, right);
node* update_fa(node* pre, int left, int right, int pos
    . int val) {
  node* x = new node;
  if (left == right) {
   x \rightarrow fa = val;
   x \rightarrow sz = 1;
   return x;
 int mid = (left + right) >> 1;
 if (pos <= mid) {</pre>
   x->cl = update_fa(pre->cl, left, mid, pos, val);
   x \rightarrow cr = pre \rightarrow cr;
  else {
   x->cl = pre->cl;
   x\rightarrow cr = update_fa(pre\rightarrow cr, mid + 1, right, pos, val
        );
 return x;
void update_sz(node* now, int left, int right, int pos,
     int val) {
  if (left == right) {
   now->sz += val;
    return;
 int mid = (left + right) >> 1;
  if (pos <= mid) update_sz(now->cl, left, mid, pos,
      val);
  else update_sz(now->cr, mid + 1, right, pos, val);
pair<int,int> query(node* now, int left, int right, int
     pos) {
  if (left == right) return { now->fa,now->sz};
 int mid = (left + right) >> 1;
```

```
if (pos <= mid) return query(now->cl, left, mid, pos)
  else return query(now->cr, mid+1, right, pos);
pair<int, int> find_(node* now_version,int m,int x) {
  pair<int,int> tp = query(now_version, 1, m, x);
  if (x == tp.first) return tp;
  else return find_(now_version, m, tp.first);
void union_(node* now_version,int new_,int n,int x,int
  pair<int, int> X = find_(now_version, n, x);
  pair<int, int> Y = find_(now_version, n, y);
  if (X.first != Y.first) {
    if (X.second < Y.second) {</pre>
      version[new_] = update_fa(now_version, 1, n, X.
          first, Y.first);
      update_sz(version[new_], 1, n, Y.first, X.second)
    }
    else {
      version[new_] = update_fa(now_version, 1, n, Y.
         first, X.first);
      update_sz(version[new_], 1, n, X.first, Y.second)
  }else version[new_] = now_version;
9.3 Treap
struct Treap {
  int key, pri, sz, tag;
  Treap* cl, * cr;
                   //左右子樹
  Treap() {}
  Treap(int key_) {
    key = key_;
    sz = 1;
    tag = 0;
    pri = rand();
    cl = cr = nullptr;
 }
int Size(Treap* x) { return x ? x->sz : 0; }
void pull(Treap* x) { x->sz = Size(x->cl) + Size(x->cr)}
     + 1; }
Treap* merge(Treap* a, Treap* b) {
  push(a);
  push(b);
  if (!a || !b)return a ? a : b; //其中一個子樹為空則回
      傳另一個
  if (a->pri > b->pri) {
                                //如果a的pri比較大則a
     比較上面
    push(a->cr);
    push(b);
    a->cr = merge(a->cr, b); //將a的右子樹跟b合併
    pull(a);
    return a;
  else {
    push(b->cl);
    push(a);
    b->cl = merge(a, b->cl); //如果b的pri比較大則b比
        較上面
    pull(b);
                              // 將b的左子樹根a合併
    return b;
  }
void splitByKey(Treap* x, int k, Treap*& a, Treap*& b)
    {
  //將一棵Treap分成雨棵,
  //key小於等於k的分到左邊那棵a,其他分到右邊那棵b
  push(x);
  if (!x) { a = b = nullptr; }
  else if (x->key <= k) {</pre>
    a = x;
    splitByKey(x->cr, k, a->cr, b);
    pull(a);
  else {
```

b = x;

```
splitByKey(x->cl, k, a, b->cl);
   pull(b);
 }
void splitByKth(Treap* x, int k, Treap*& a, Treap*& b)
  //將一棵Treap分成雨棵,
  //左邊那棵a的節點數有k個,右邊那棵b節點數為n-k個
 push(x);
 if (!x) { a = b = nullptr; }
else if (Size(x->cl) + 1 <= k) {</pre>
   a = X; //如果左子樹+自己的SiZe小於等於k則左子樹跟
       自己為k個以內
   splitByKth(x->cr, k - Size(x->cl) - 1, a->cr, b);
   pull(a);
 else {
   b = x; //如果左子樹+自己的size大於R個則右子樹跟自
        己會分到右邊
   splitByKth(x->cl, k, a, b->cl);
   pull(b);
 }
void insert(int val){
                                //新增一個值為val的元
   Treap *x = new Treap(val);
                               //設 — 個 trean 節 點
   Treap *1,*r;
   splitByKey(root, val, 1, r); //找到新節點要放的位
   root = merge(merge(1,x),r);
                               //合併到原本的treap裡
void erase(int val){
                                //移除所有值為val的元
   Treap *1,*mid,*r;
   splitByKey(root, val, l, r); //把小於等於val的丟到
   splitByKey(1, val-1, 1, mid); //小於val的丟到L,等於
       val的就會在mid裡
   root = merge(1,r);
                               // 將除了val以外的值合
       併
//翻轉
void push(Treap* x) {
 if (!x)return;
  if (x->tag) {
   swap(x->cl, x->cr);
   if (x->cl)x->cl->tag ^= 1;
   if (x->cr)x->cr->tag ^= 1;
   x->tag ^= 1;
 }
}
9.4 樹壓平
vector<vector<int>>edge;
vector<pair<int,int>>times;
int ti = 0;
void dfs(int x,int fa){
 times[x].fir = times[x].sec = ti++;
 for(int i:edge[x]){
   if(i == fa)continue;
   dfs(i,x);
   times[x].sec = max(times[i].sec,times[x].sec);
 }
}
```

9.5 樹上全點對距離總和

```
for(int i=0;i<edge[x].size();i++){
    if(edge[x][i] == fa) continue;
    ll tp = sum + dp[x] - (dp[edge[x][i]] + sz[edge
        [x][i]]) + (n-sz[edge[x][i]]); //從x到edge[
        x][i]的距離總和
    dfs2(x,edge[x][i],tp);
    }
}
dfs1(1,1);
dfs2(1,1,0); //答案是ans/2
```

10 Geometry

10.1 Convex hull 凸包

```
struct Pt{
    int x,y;
    Pt(){}
    Pt(int _x,int _y){
        x=_x, y=_y;
    friend bool operator<(const Pt& lhs,const Pt& rhs){</pre>
        return lhs.x==rhs.x?lhs.y<rhs.y:lhs.x<rhs.x;</pre>
    friend Pt operator-(const Pt& lhs,const Pt& rhs){
        return Pt(rhs.x-lhs.x,rhs.y-lhs.y);
    friend int cross(const Pt& o,const Pt& a,const Pt&
        b){
        Pt lhs = o-a, rhs = o-b;
        return lhs.x*rhs.y - lhs.y*rhs.x;
};
vector<Pt> convex_hull(vector<Pt> hull){ //回傳凸包陣列
    sort(hull.begin(),hull.end());
    int top=0;
    vector<Pt> stk;
    for(int i=0;i<hull.size();i++){</pre>
        while(top>=2&&cross(stk[top-2],stk[top-1],hull[
             i])<=0)
            stk.pop_back(),top--;
        stk.push_back(hull[i]);
        top++;
    for(int i=hull.size()-2,t=top+1;i>=0;i--){
        while(top>=t&&cross(stk[top-2],stk[top-1],hull[
            i])<=0)
            stk.pop_back(),top--;
        stk.push_back(hull[i]);
        top++;
    stk.pop_back();
    return stk;
double FarthestPair(vector<Pt> arr) { //回傳最遠點對
    double ret = 0;
    for (int i = 0, j = i + 1; i < arr.size(); i++) {</pre>
        while (distance(i, j) < distance(i, (j + 1) %</pre>
            arr.size())) {
            j = (j + 1) % arr.size();
        ret = max(ret, distance(i, j));
    return ret;
vector<Pt> tui;
double distance(int i, int j) {
    double a = (tui[i].x - tui[j].x) * (tui[i].x - tui[
        j].x),
           b = (tui[i].y - tui[j].y) * (tui[i].y - tui[
                j].y);
    return sqrt(a + b);
}
```

10.2 Stack Square Area

```
#include<bits/stdc++.h>
#include<unordered_map>
using namespace std;
int c, n, m, min_area, head, tail, mid, f, area;
bool check;
unordered_map<int, int>done;
```

```
vector<vector<int>>G. h:
vector<int>element;
vector<pair<int, int>>Stack;
pair<long long, long long>ans;
int main() {
  ios::sync_with_stdio(0);
  cin.tie(0);
  cout.tie(0);
  cin >> c;
  while (c--) {
    h.clear();
    G.clear();
    element.clear();
    cin >> n >> m >> min_area;
    G.resize(n);
    done.clear();
    for (int i = 0; i < n; i++) {</pre>
      G[i].resize(m);
      for (int j = 0; j < m; j++) {
        cin >> G[i][j];
        element.push_back(G[i][j]);
      }
    sort(element.begin(), element.end());
    ans.first = *element.begin() - 1;
    head = mid = 0;
    tail = element.size();
    while (mid!=((head+tail)>>1)||!done.count(element[
        mid])) {
      mid = (head + tail) >> 1;
      if (!done.count(element[mid])) {
        area = 0;
        h.clear();
        h.resize(n);
        for (int i = 0; i < n; i++) {</pre>
          h[i].resize(m);
          for (int j = 0; j < m; j++) {</pre>
            h[i][j] = (G[i][j] >= element[mid] ? 1 : 0)
            if (i != 0 && h[i][j] != 0) h[i][j] += h[i
                 - 1][j];
          }
        for (int i = 0; i < n; i++) {
          Stack.clear();
          for (int j = 0; j < m; j++) {
            f = j;
            while (!Stack.empty() && h[i][j] < Stack.</pre>
              back().second) {
area = max(area, (j - Stack.back().first)
                   *Stack.back().second);
              f = Stack.back().first;
              Stack.pop_back();
            if (h[i][j] != 0 && (Stack.empty() || h[i][
                 j] != Stack.back().second))Stack.
                 push_back(make_pair(f, h[i][j]));
          while (!Stack.empty()) {
            area = max(area, (m - Stack.back().first)*
                 Stack.back().second);
            Stack.pop_back();
        }
      if (done.count(element[mid])) area = done[element
          [mid]];
      if (area < min_area) tail = mid;</pre>
      else {
        done[element[mid]] = area;
        ans.first = element[mid];
        ans.second = area;
        head = mid;
      }
    cout << ans.first << " " << ans.second << "\n";</pre>
  return 0;
```

11 Sqrt

11.1 分塊

```
struct blk{
    vector<int> local;
                         //每塊的全部元素
    int global;
                         //儲存每塊的總和
    int tag;
                         //儲存整塊一起更新的值
                         //初始化
    blk(){
        local.clear();
                         //清空區間元素
        tag = global = 0; //將區間總和先設為0
    }
};
vector<blk> b;
void build(){
    int len=sqrt(n),num=(n+len-1)/len;
    for(int i=0;i<n;i++){ //第i個元素分在第 i/Len 塊
        cin>>x;
        //存入區間中
        b[i/len].local.push_back(x);
        //更新區間總和
        b[i/len].global += x;
    }
void update(int ql,int qr,int v){
    int blk_l=ql/len,blk_r=qr/len,ret=0;
    if(blk_1 == blk_r){
                         //如果都在同一塊直接一個一個
        跑過去就好
        for(int i=ql;i<=qr;i++)</pre>
            b[blk_l].local[i%len]+=v;
        b[blk_1].global+=(qr-ql+1)*v;
    for(int i=ql;i<(blk_l+1)*len;i++){ //最左的那一塊
        b[blk 1].local[i%len]+=v;
        b[blk_1].global+=v;
    for(int i=blk_l+1;i<blk_r;i++){ //中間每塊
        b[i].tag+=v;
        b[i].global+=v*len;
    for(int i=blk_r*len;i<=qr;i++){ //最右的那一塊
        b[blk_r].local[i%len]+=v;
        b[blk_r].global+=v;
    }
 int query(int ql,int qr){
    int blk_l=ql/len,blk_r=qr/len,ret=0;
    if(blk_1 == blk_r){
                        //如果都在同一塊直接一個一個
        跑過去就好
        for(int i=ql;i<=qr;i++)</pre>
            ret+=b[blk_1].local[i%len]+b[blk_1].tag;
        return ret:
    for(int i=ql;i<(blk_l+1)*len;i++)</pre>
                                       //最左的那一塊
        ret+=b[blk_1].local[i%len]+b[blk_1].tag;
    for(int i=blk_l+1;i<blk_r;i++)</pre>
                                    //中間每塊的總和
        ret+=b[i].global;
    for(int i=blk_r*len;i<=qr;i++)</pre>
                                     //最右的那一塊
        ret+=b[blk_r].local[i%len]+b[blk_r].tag;
    return ret;
```

11.2 Mo's 莫隊

```
int n,k = sqrt(n);
                   //每塊大小為k
struct query{
   int l,r,id;
                   //詢問的左界右界 以及 第幾筆詢問
   friend bool operator<(const query& lhs,const query&</pre>
        rhs){
       return lhs.1/k==rhs.1/k ? lhs.r<rhs.r : lhs.l</pre>
           rhs.1;
        //先判斷是不是在同一塊 不同塊的話就比較塊的順
       序,否則比較右界r
 int num = 0;
int cnt[1'000'005], ans[30005];
vector<query> q;
void add(int index){ ... }
                           //新增元素到區間內
void sub(int index){ ... }
                           //從區間內移除元素
void solve(){
```

```
sort(q.begin(),q.end());
for(int i=0,l=-1,r=0;i<n;i++){</pre>
   while(1>q[i].1)
                       add(--1);
   while(r<q[i].r)</pre>
                       add(++r);
                                    //記得要先做新
        增元素的
    while(1<q[i].1)</pre>
                       sub(1++);
                                    //再做移除元素
        的
   while(r>q[i].r)
                      sub(r--);
   ans[q[i].id] = num;
                                    //移到區間後儲
        存答案
```

12 Math

12.1 Miller rabin 找質數

```
// n < 4,759,123,141
                            3 :
                                2, 7, 61
// n < 1,122,004,669,633
                                 2, 13, 23, 1662803
                                  6 : pirmes <= 13
// n < 3,474,749,660,383
                                  7 :
// n < 2^64
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
// Make sure testing integer is in range [2, n-2] if
// you want to use magic.
//<ex> magic[3] ={2,7,61};
#define ull unsigned long long
ull magic[]={}
ull mul(ull a, ull b, ull c) { //快速乘
 ull ans = 0;
 while (b) {
   if (b & 1) ans = (ans + a) % c;
   a = (a + a) \% c;
   b >>= 1;
 }
 return ans;
ull mypow(ull a, ull u, ull n) {
 ull ans = 1;
 while (u) {
   if (u & 1) ans = mul(ans,a,n);
   a = mul(a,a,n);
   u >>= 1;
 return ans;
bool witness(ull a, ull n, ull u, ull t) {
 if (!a) return 0;
  ull x = mypow(a, u, n); //a^u \% n
 for (int i = 0; i < t; i++) {</pre>
   ull nx = mul(x, x, n);
    if (nx == 1 && x != 1 && x != n-1) return 1;
   x = nx;
 }
 return x != 1;
bool miller_rabin(ull n) {
 int s = (magic numbers size);
    // iterate s times of witness on n
    if(n < 2) return 0;
 if (!(n & 1)) return n == 2;
 ull u = n-1; ull t = 0;
 // n-1 = u*2^t
 while (!(u & 1)) u >>= 1, t++;
 while (s--) {
   ull a = magic[s] % n;
   if (witness(a, n, u, t)) return 0;
 return 1;
```

12.2 Chinese remainder theorem

```
LL x[N],m[N];
LL CRT(LL x1, LL m1, LL x2, LL m2) {
    LL g = __gcd(m1, m2);
    if((x2 - x1) % g) return -1;// no sol
    m1 /= g; m2 /= g;
    pair<LL,LL> p = gcd(m1, m2);
    LL lcm = m1 * m2 * g;
    LL res = p.first * (x2 - x1) * m1 + x1;
    return (res % lcm + lcm) % lcm;
}
LL solve(int n){ // n>=2,be careful with no solution
```

12.4 矩陣快速冪

```
#define MOD 1'000'000'007
#define ll long long
vector<vector<ll>>> operator*(const vector<vector<ll>>>&
    lhs,const vector<vector<ll>>% rhs){
    vector<vector<ll>> ret(lhs.size(),vector<ll>(rhs
        [0].size(),0));
    for(int i=0;i<lhs.size();i++){</pre>
        for(int j=0;j<rhs[0].size();j++){</pre>
            for(int k=0;k<rhs.size();k++){</pre>
                ret[i][j] += lhs[i][k] * rhs[k][j] %
                     MOD:
                ret[i][j] %= MOD;
            }
        }
    return ret;
vector<vector<ll>> init_value={{1},{0}}; //第\theta,1項
vector<vector<ll>> base={{1,1},{1,0}};
                                           //費式數列轉
vector<vector<ll>> matrix={{1,0},{0,1}}; //單位矩陣
while(y){
    if(y&1){
        matrix = matrix * base;
    base = base * base;
    y >>= 1;
matrix = matrix * init_value;
cout<< matrix[0][0] << endl;</pre>
12.5 模逆元
```

```
}
inv[MXN] = mypow(f[MXN], (MOD - 2), MOD);
for (long long i = MXN - 1; i >= 0; i--) { //模逆元
    inv[i] = inv[i + 1] * (i + 1) % MOD;
}

ll C(int n, int m) { //用途再算C幾取幾 % MOD
    return f[n] * inv[m] % MOD * inv[n - m] % MOD;
}

12.6 判斷線交叉

double Direction(point Bi point Bi point Bk) {
```

```
double Direction(point Pi, point Pj, point Pk){
    return (Pj.x-Pi.x)*(Pk.y-Pi.y)-(Pk.x-Pi.x)*(Pj.y-Pi.y)
    );
}

bool isIntersect(line p, line q){
    double d1, d2, d3, d4;
    d1 = Direction(p.a, p.b, q.a);
    d2 = Direction(p.a, p.b, q.b);
    d3 = Direction(q.a, q.b, p.a);
    d4 = Direction(q.a, q.b, p.b);
    if(d1*d2<0 && d3*d4<0) { return true;} //規範相交
    //非規範相交
    else return false;
```

12.7 Theorem

```
• Wilson's theorem : (p-1)! \equiv -1 (mod \ p) \quad \text{(p is a prime)}
```

• Fermat's little theorem : $a^p\equiv a(mod~p)~\text{(p is a prime and gcd(a,p)=1)}\\ \frac{1}{a}(mod~p)\equiv a^{p-2}~\text{(with fast pow)}$





