

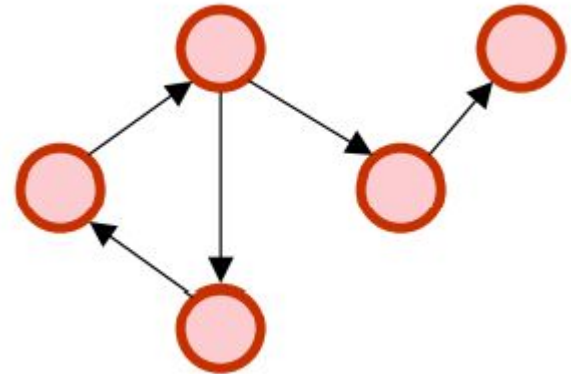
Topological Sort

Topological Sort

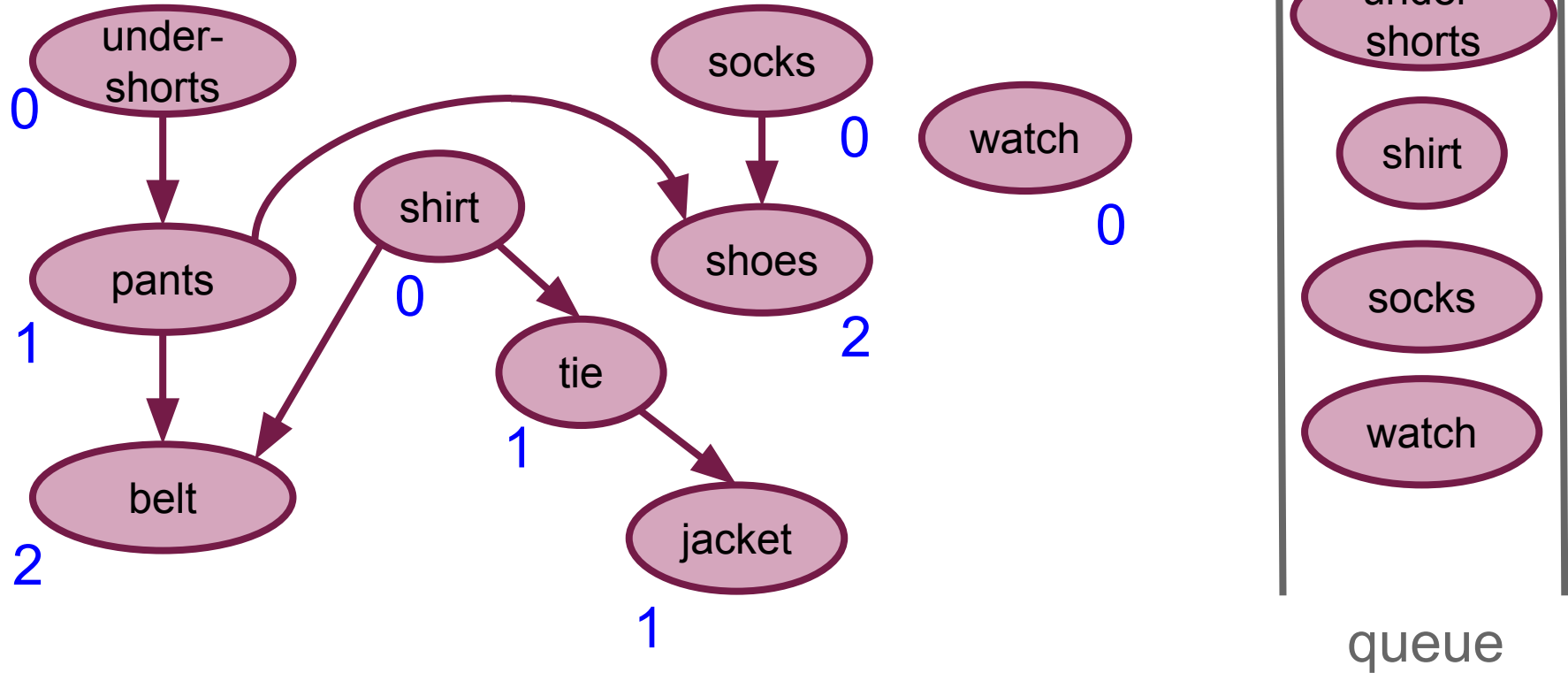
- Given a directed graph G
- Try to find an order of all vertices in G , such that for any edge (u, v) , u appears before v in the ordering.

Topological Sort

- We can find a topological order if G is acyclic (contains no cycle)
- If G contains a cycle, then it is impossible to find such an ordering. (Proof by contradiction)

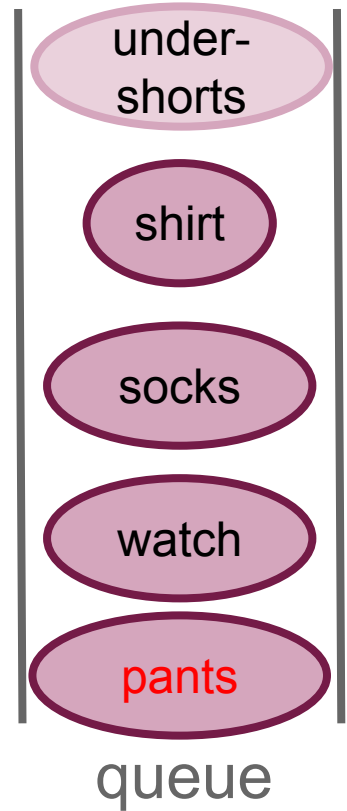
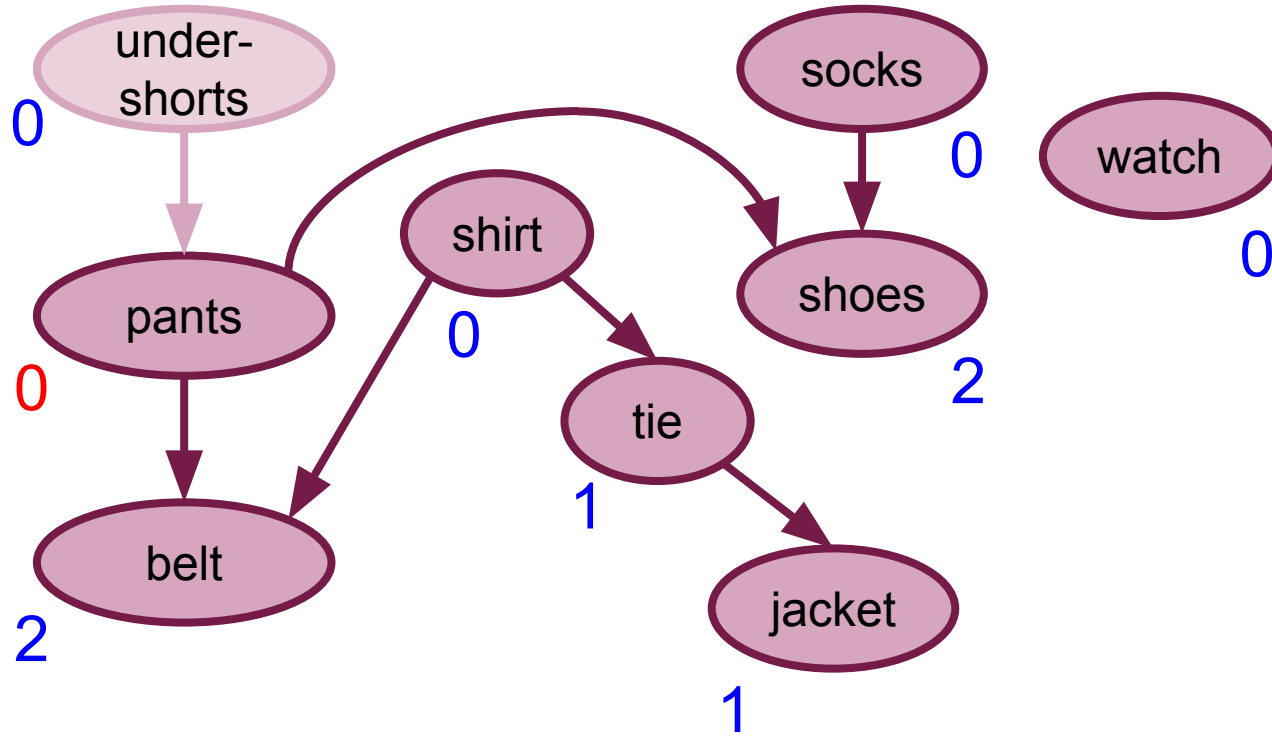


Topological Sort - Example



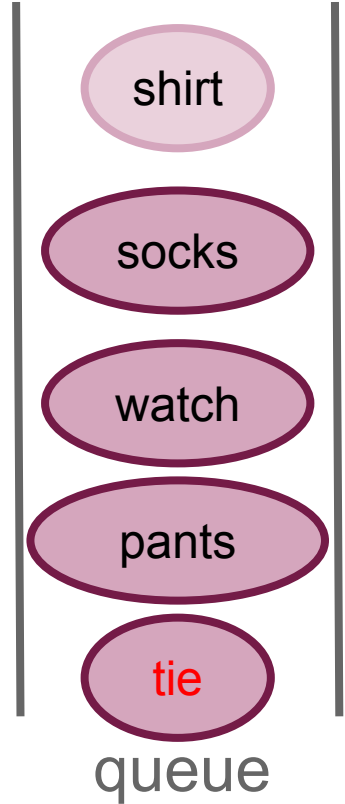
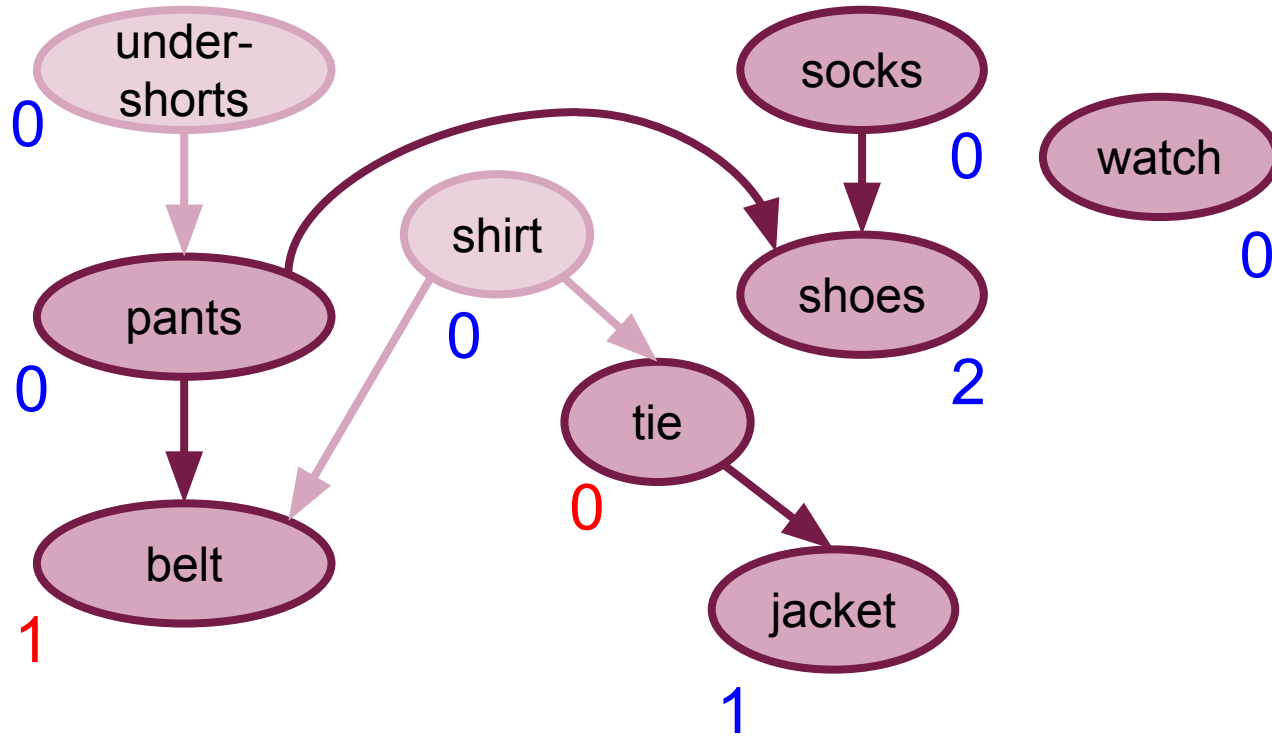
Topological order: (empty)

Topological Sort - Example



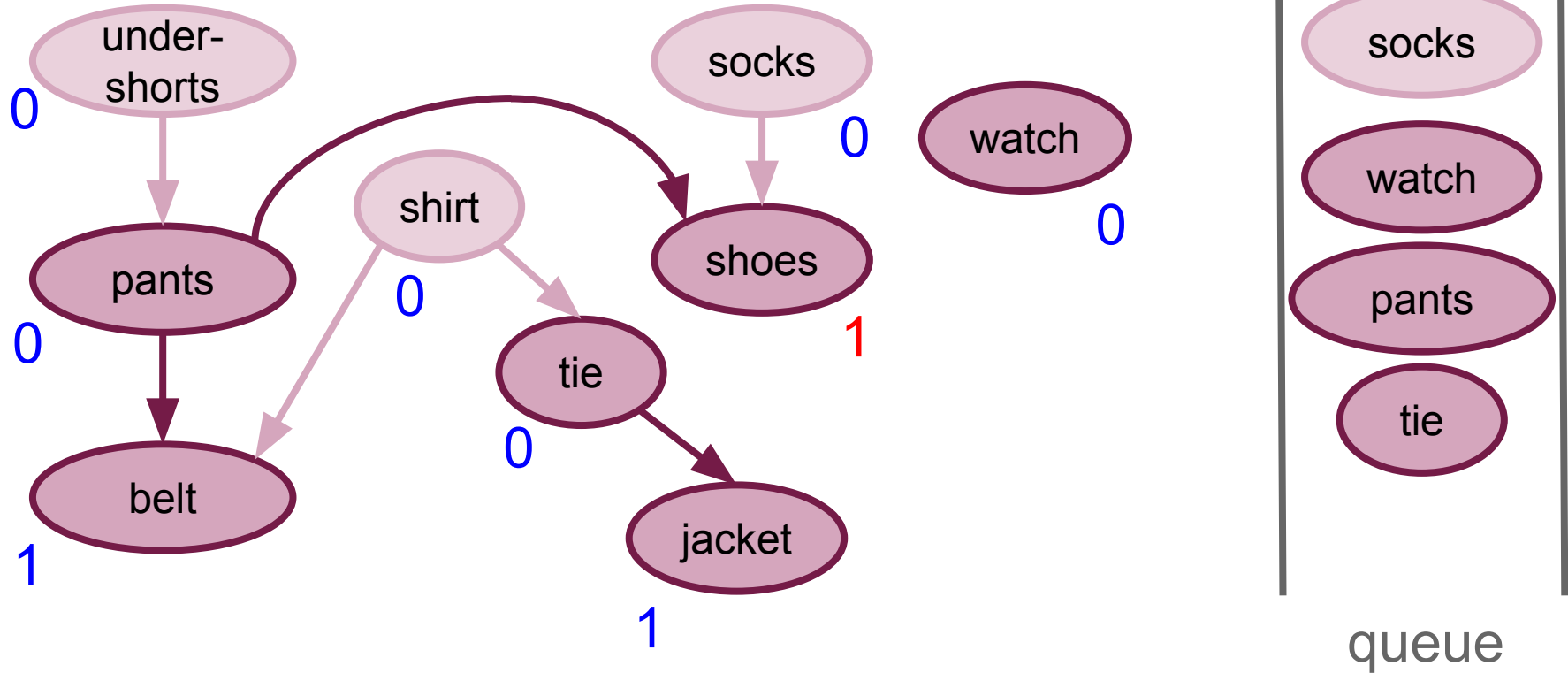
Topological order: **under-shorts**

Topological Sort - Example



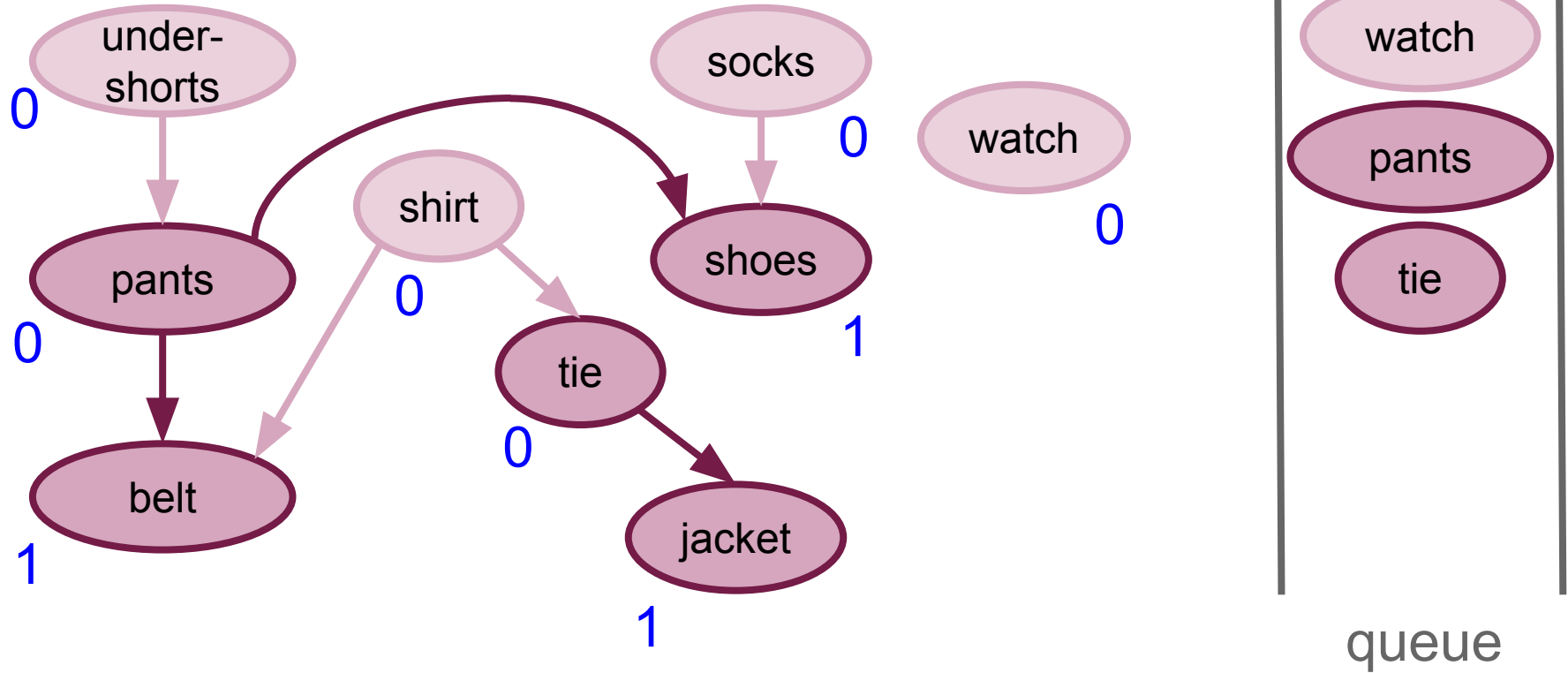
Topological order: under-shorts, **shirt**

Topological Sort - Example



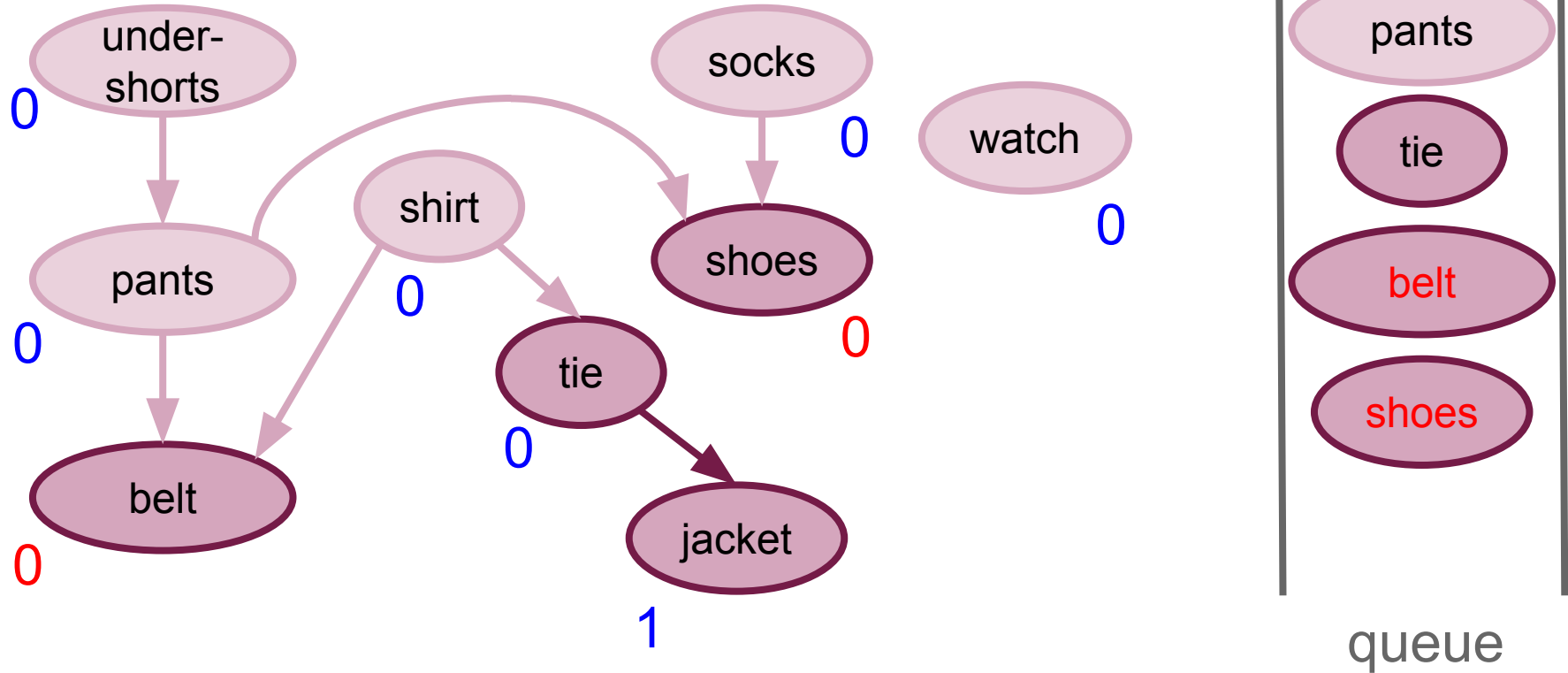
Topological order: under-shorts, shirt, **socks**

Topological Sort - Example



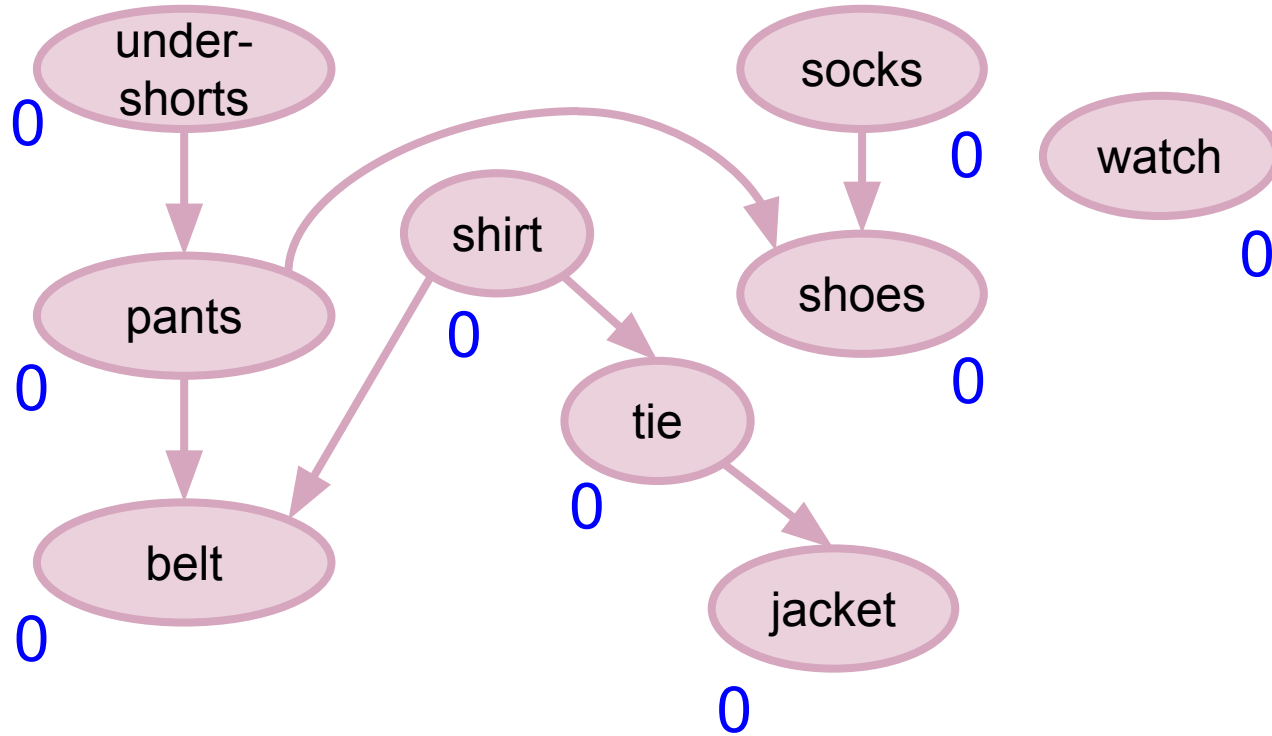
Topological order: under-shorts, shirt, socks, **watch**

Topological Sort - Example



Topological order: under-shorts, shirt, socks, watch, pants

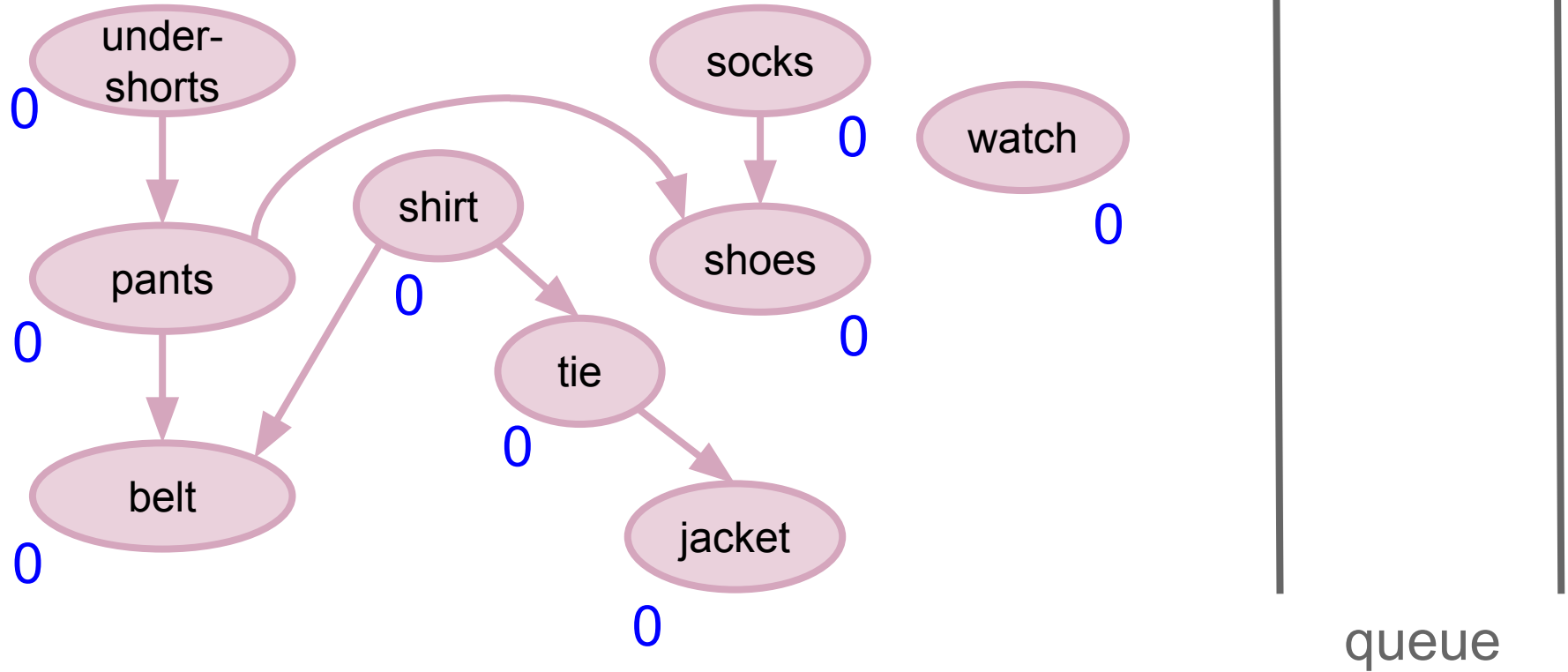
Topological Sort - Example (result)



queue

Topological order: under-shorts, shirt, socks, watch, pants, tie, belt, shoes, jacket

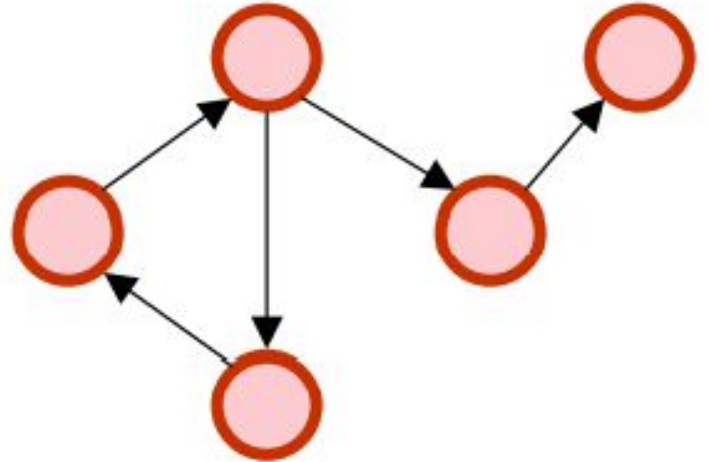
Topological Sort - Example (result)



[shirt, tie, jacket, watch, socks, under-shorts, pants, belt, shoes] can also be a Topological order.

How to check it's DAG or not?

```
topological_sort(G, indegree);  
if(ans.size() == n) cout<<"Yes\n";  
else cout<<"No\n";
```



Topological Sort - Alternative Implementation

- Using out-degree instead of in-degree
- Call DFS
 - For every vertex, save the timestamp of finish time
 - Output vertices in decreasing order of their timestamp

in class sample codes

In-degree:

<https://ide.usaco.guide/O4F1WE2aSueZryJYWYQ>

dfs:

https://ide.usaco.guide/O4F1_8WCsrRDQho7zHn

Lab18 Is Dag

題意：

給你一張無向圖，請回答該圖是否是一張 DAG。

Lab18 Is Dag (cont.)

- 直接用跑拓樸排序檢查是否為 DAG 即可
- <https://gist.github.com/rurutoria7/53fc6da52d0c7db91b7db7732bf4b776>

Lab18 Spreadsheet

題意：

給你一張試算表，每一格的公式為某些其他格的值的和。請求出每一格的值是多少，或回答無法求出。

Lab18 Spreadsheet (cont.)

- 令第 u 格加上的值為 $a[u]$, 若 u 參考到 v 則建邊 (u,v)
 - 令 $f(u)$ 為第 u 格的值
 - 若 $f(u)$ 已經被計算完, 直接 return
 - 若 $f(u)$ 尚未計算, 則 $f(u)$ 為 $a[u]$ 加上所有參考點的值 $f(v)$
 - 若 $f(u)$ 仍在等待答案, 則詢問非法
- <https://gist.github.com/rurutoria7/4156b6a9d22a3d3de257295b3035ab05>