# WRITING-TASK

# Part 2

2. sendFrame() requires the caller to provide the destination MAC address when sending IP packets, but users of IP layer won't provide the address to you. Explain how you addressed this problem when implementing IP protocol.

Use standard ARP protocol. Each host has a globol map for ip and MAC address. Any time the host would like to send a packet to an IP address whose MAC address unknown, the host will send an ARP request to each host in the same subset. And device with the specific IP address will send an ARP reply to the host.

3. Describe your routing algorithm.

Self-Destruct Protocol



SDP(self-destruct protocol) is a routing protocol. Instread of RIP using UDP, the SDP is based on link layer. It is designed for the reason that the routing table stores the *MAC* address of nexthop for the task function setRoutingTable in our lab.

The *format* of SDP frame is shown below.

```
0
           1
                     2
                                3
\begin{smallmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 0 & 1 \\ \end{smallmatrix}
Length
           Flag
+-+-+-+-+-+-+-+-+-+-+-+-+
                      MAC address
Route Address Prefix
itemFlag
                      Distance
```

- **Length**: the number of SDP item. The max number is 255.
- Flag: basic information of the frame. The defination of each bit will be explained latter.

```
[rrrrvnui]
    r (Reserved)
    i Increment : is an increment routing information.
    u Unfinished : have more packet because of length limited
    n isNew : is a new device added into network
    v Verify. : is a verify frame
```

MAC address: MAC address of the sender.

#### A SDP item:

- Route Address Prefix: Route address prefix such as 10.100.1.0.
- plen: The length of prefix such as 24.
- o **Distance**: Distance to the receiver.
- **itemFlag**: Information of the item.

```
[rrrrrrd]
   r (Reserved)
   d Deleted : delete the item
```

### **Protocol**

SDP uses the algorithm of distance vector.

 Once added into network, a new host will initial its own routing table according to its devices' IP address and subnet mask. For example, ns2 has the routing table like this

| IP Prefix     | Device  | Nexthop MAC      | Distance |    |
|---------------|---------|------------------|----------|----|
|               | veth2-1 | (mac of veth2-1) | 0 (dev)  | -1 |
| 10.100.2.0/24 | veth2-3 | (mac of veth2-3) | 0 (dev)  | -1 |

Then the new host will send its routing table to all its neighbours with flag **isNew** called **Request Frame**. (We will discuss matric later)

• Whenever receiving an SDP frame, the host will update its own routing table according to the distance (+1 for each hop) and added it into updateSDPItems. Then the host send updateSDPItems to its neighbours except the src called **Update Frame**. If the SDP frame has the flag isNew, the host will send its own routing table to the new host (src) called **Reply Frame**.

#### Matric:

matric of items in routing table with via will always be -1.

- Every 30~40 seconds, matric of each item in the routing table will +1.
- Every 30~40 seconds, every host will send its whole routing table to its neighbours to verify called **Verify Frame**.
- When receiving a verify frame, matric of items in routing table, whose <a href="nexthop">nexthop</a> is src, will be reset to 0. It means "You are my uplink, so I believe that the routing is still effective!"
- Whenever a host is deleted from the network, the matric will be added without update. Once the matric is greater than or equal to 2, the host will send a update frame to tell neighbours delete this item with itemFlag Deleted. As getting an updating item with Deleted, the metric will be set to 8 directly.
- But, but the host won't delete the item from routing table immediately. It will be reserved for about 1-2 period(s), or a routing frame will be added again from downlink.
- 4. To implement routing properly, you need to detect other hosts/let other hosts know about you. In this lab, you are not required to detect hosts not running your protocol stack automatically/let them know about you, but you must not make them complain about strange incoming packets. Describe how your IP implementation achieved this goal.

Considering SDP is based on link layer, SDP uses  $0 \times 2333$  as EtherType. Other host or router will get a frame with an unregistered EtherType, it will ignore the frame.

# Part 3

2. Describe how you correctly handled TCP state changes.

Using **DFA**. We have two threads for a TCP socket. One of them is mainly used to accept the segments. The other one is used to send segments. So, considering a TCP state change, mostly, *listening thread* handles the recv procedure, and main thread handles the send procedure. Also, listening thread will change criticalst as critical state to promise the protocol **thread safe** if the main thread have things to do.

We promise that listening thread will not be blocked any time. But, listening thread will be blocked if criticalst != st which means main thread does not finish its work for thread safe.

### 3. Describe how you implemented in-order data transferring.

As sender, the potocol send only one segment until it receive an ACK related.

As receiver, the algorithm is similiar with Go-Back-N of sliding window. If we receive a packet with sequence number greater than wanted (rcv\_nxt), the protocol will send a **duplicate ACK** to tell other side retransmit the packet with correct sequence number. If the sequence number less than rcv\_nxt, the protocol will send an **ACK** acknowledged actually.

#### 4. corner cases

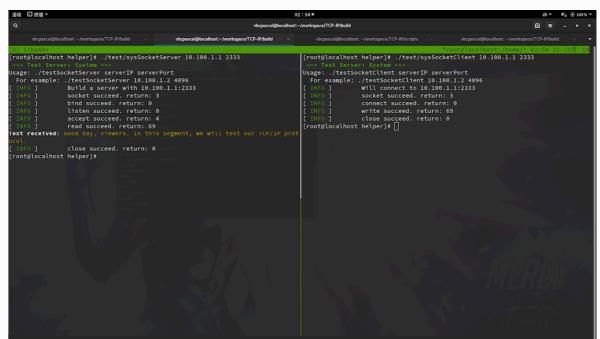
It can work. But it does not have enough robustness actually. For example, real Linux machine will **always** send segments with  $\boxed{\mathtt{PSH}}$  and  $\boxed{\mathtt{ACK}}$ , even with  $\boxed{\mathtt{FIN}}$ . This may lead to erroneous results.

## Part 4

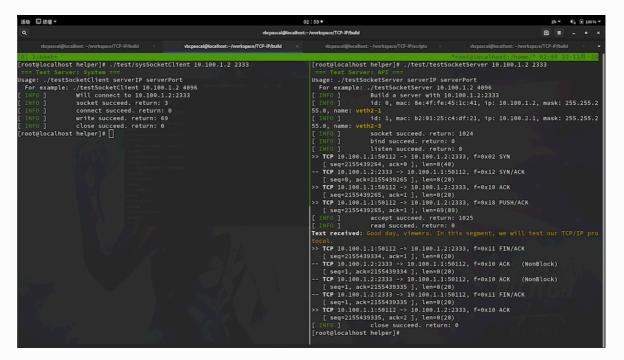
## 2. Describe your evaluating methodology, result and usage of your evaluating method.

Build a connect between **client and server**, which are implemented in tests/testSocketClient.cpp and tests/testSocketServer.cpp. Each of them are be compiled into two versions, wrapped ones and system ones. In fact, it is more convenient that using TEST\_API flag to generate them. Then, rry to connet them with each other. It can work correctly as following pictures shows.

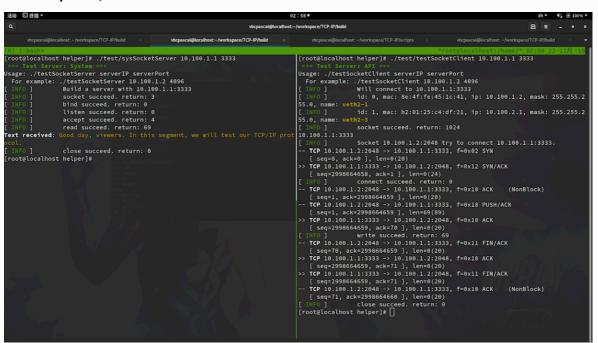
Server: System, Client: System



Server: API, Client: System



## Client: System, Server: API



Client: API, Server: API