

$C \rightarrow Nhc$

$(\alpha: \Delta^1 \rightarrow C) \mapsto (\alpha(0) \rightarrow \alpha(1) \rightarrow \dots \rightarrow \alpha(n))$ in hc

$(f: \Delta^1 \rightarrow C) \mapsto (x \xrightarrow{f} y)$ in hc .

$\exists f: \Delta^1 \text{ an iso} \dots \Rightarrow [f^{-1}]: y \rightarrow x$ in hc .

$x \xrightarrow{f} y \quad y \xrightarrow{f^{-1}}$ commutes.

$\text{Call this } h\text{ic} \xrightarrow{\alpha} \alpha: [2] \rightarrow hc$

\downarrow

$Nhc: \Delta^2 \rightarrow Nhc$

$C \rightarrow Nhc$ is wors. if $hc \rightarrow hNhc$ in ...

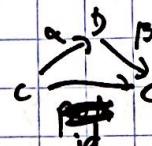
$\xrightarrow{id?} hc$

$hc \xrightarrow{h\varphi} hNhc$

$x \xrightarrow{} x$

$[f] \downarrow \xrightarrow{} h\varphi[f] = [\varphi f] \circ [f]$.

$y \xrightarrow{} y$



say $\varphi, \beta \circ \alpha$ are
ans.
 $\Rightarrow \alpha$ cons.

basepoints:

Set $_{\star}$ & Top $_{\star}$.

\star

sets: $\exists S = \{s_1, s_2, \dots, s_n\}$.

\uparrow

$\left\{ \text{pick one element} \right\}.$

Set $_{\star}$ $\ni S = \{s_1, s_2, \dots, s_n\}$.

basepoint:

Set $_{\star} \subseteq \text{Set}$

"sets w/
bpt."

Hom(Set $_{\star}$)

(X, x_0)

(Y, y_0)

X = { x_1, \dots, x_n }

Y = { y_0, \dots, y_m }

is a set-function

f: X \rightarrow Y

$x_0 \mapsto f(x_0) = y_0$.

Set $_{\star} = \text{ob: sets w/ bpt}$ maps

{+3}

$\xrightarrow{x_0} X$

singleton:



Top $_{\star} = \text{ob: ptd top-spaces}$:

X a space comes w/

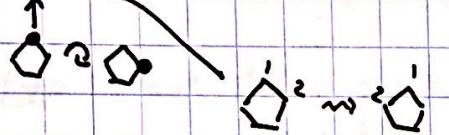
* $\xrightarrow{x_0} X$

G : a group. $\{g, g^{-1}\}$ elements.

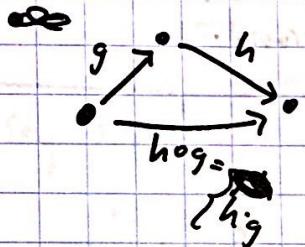
BG - ob: $\bullet \xrightarrow{id} \bullet$ an elt in G ?
ex. hom: $\{\bullet \xrightarrow{g} \bullet\}_{g \in G}$. maps.

D_5 : symmetries of 

$\{ \sigma, \tau : \sigma^5 = e, \tau^2 = e \}$.



composition in BG :



locally

small: for any objects $x, y \in \mathcal{C}$

$\text{Hom}_\mathcal{C}(x, y)$ is a set.

small: Take all morphisms in \mathcal{C}

$$\text{Hom}(\mathcal{C}) := \bigcup_{x, y \in \text{ob}(\mathcal{C})} \text{Hom}_\mathcal{C}(x, y)$$

If this guy is a set, then \mathcal{C} is "small".

Groups is locally small?

$\text{Hom}_\mathcal{C}(G, H)$ is a set. ✓.

• but not small? —

i.e. $\bigcup_{(G, H) \in \text{ob}(\mathcal{C})} \text{Hom}_\mathcal{C}(G, H)$

• $\text{ob}(\mathcal{C})$ is not a set

• $\text{ob}(\mathcal{C}) \times \text{ob}(\mathcal{C})$ is not a set

• $\text{ob}(\mathcal{C}) \times \text{ob}(\mathcal{C}) \times \text{ob}(\mathcal{C})$ is not a set

ex/1. Given a gp. hom.

$$G \xrightarrow{\varphi} H$$

$$BG \xrightarrow{\quad \quad \quad} BH$$

functor

$$\begin{array}{ccc} \bullet_G & \xrightarrow{\quad} & \bullet_H \\ g \downarrow & \longleftarrow & \downarrow \varphi(g) \\ \bullet_{G'} & & \bullet_{H'} \\ g' \downarrow & \longleftarrow & \downarrow \varphi(g') \\ \bullet_{G''} & & \bullet_{H''} \end{array}$$

(i) check $B\mathcal{C}$ is a functor

- $\text{id} \mapsto \text{id}$
- $\text{comp} \mapsto \text{comp}$

(ii) a gp hom. has some properties.

- $e \mapsto e_H$
- $g^{-1} \mapsto (\varphi(g))^{-1}$

check how these look
on the functor

bijection is an iso. in \mathcal{C} .

$\text{Hom}_\mathcal{C}(x, y) \cong \text{Hom}_\mathcal{C}(y, x)$

$\text{Hom}_\mathcal{C}(x, y) \cong \text{Hom}_\mathcal{C}(x, z) \times \text{Hom}_\mathcal{C}(z, y)$

$\text{Hom}_\mathcal{C}(x, y) \cong \text{Hom}_\mathcal{C}(x, z) \times \text{Hom}_\mathcal{C}(y, z)$

substrat.

A subgr. should have the structure
of a tf. & live inside of a
larger tf.

(1)
A substrat. should have the structure $C^0 \subseteq C$
of a cat & live inside a
layer cat.
(2)

$$(2) \cdot \text{ob}(C^0) \subseteq \text{ob}(C)$$

$$\cdot \text{mor}(C^0) \subseteq \text{mor}(C)$$

(structure of a cat.)

- (1) • every object has an id
• composition is well-defined

$\rightsquigarrow C^0$ to contain all id's of ob in C^0 .

$\rightsquigarrow C^0$ should contain compositions of
all its maps.

& they should be the same as in C .

e.g. when showing $C^{\text{mono}} \models \text{ob} = \text{ob}(C)$
 $\text{mor} = \{ \text{mono's in } C \}$.

to show $C^{\text{mono}} \subseteq C$

substrat.

WTC: (2) ✓

- (1) • identities are in C^{mono} (i.e. every id is a mono.)
• comp. of monos are monos. ex.

ox? id's \Rightarrow monos.

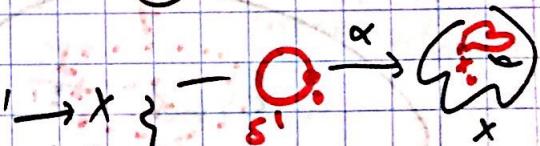
Take $(X, x_0) \in \text{Top}_*$.
Take $(S', o) \in$

circle.

~~closed~~

Look at $\text{Hom}_{\text{Top}_*}(S', o), (X, x_0) = \{ \text{cts maps } S' \xrightarrow{\alpha} X \mid \alpha|_{o} = x_0 \} \subseteq \Omega$

$$S^2 \quad S^1 = \{ |z|=1 \} \subseteq \Omega$$



forms a ret... but turns out you can give it a gp. structure.

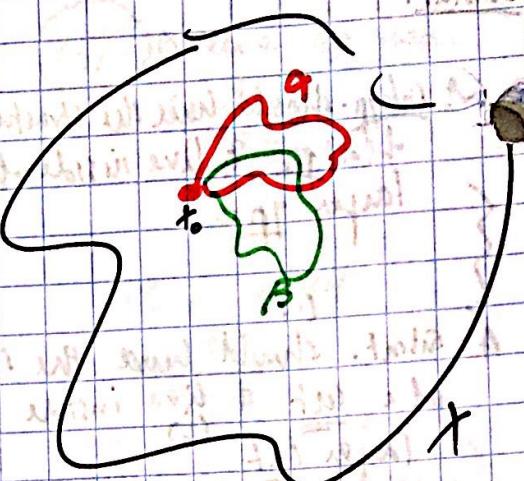
So an alt. ~~alt. $\text{Hom}(S', o)$~~ is $\alpha: (S', o) \rightarrow (X, x_0)$ is a "loop" in X .

$$(S^1, 0) \xrightarrow{\alpha} (X, x_0)$$

β

To topologish, if you can continuously deform
 $\alpha \rightsquigarrow \beta$ (in $\beta \sim \alpha$) "homotopy"

then we consider them the same.



fund. gp. of
top. space.
 (X, x_0)

$\tilde{\pi}_1(X, x_0) := \frac{\text{Hom}(S^1, X)}{\sim}$ homotopy.

$$\tilde{\pi}_1(X, x_0) = [S^1, X]_*$$

sq. brackets
= hom classes. old maps.

(another notation).

Q:

When are two categories C & D "the same"?
 i.e. what is an isomorphism in Cat?

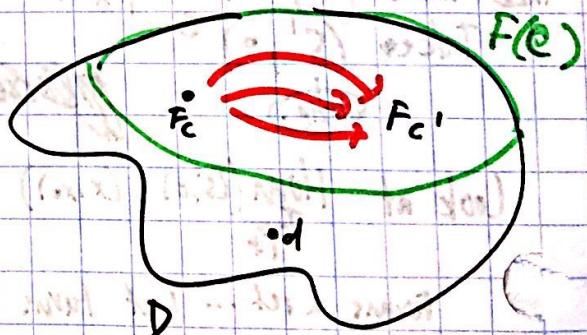
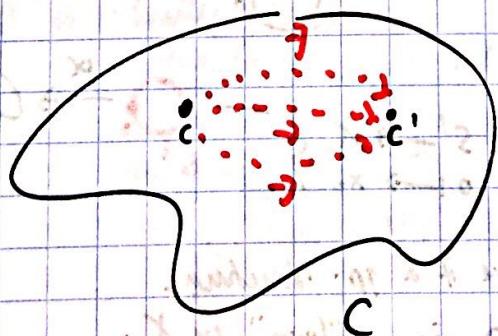
A: a functor that is fully faithful & surjective.

~~(i) full~~ Firstly a functor $F: C \rightarrow D$
 is a function $\text{ob}(C) \rightarrow \text{ob}(D)$
 - functions $\text{Hom}(c, c') \rightarrow \text{Hom}(Fc, Fc')$ Want to look at.
fully faithful.
 satisfying more things..

(i) Full if each function $\text{Hom}(c, c') \rightarrow \text{Hom}(Fc, Fc')$

is surjective.

b. for any map $Fc \rightarrow Fc'$ in D ,



(iii) faithful if $\text{Hom}(c, c') \rightarrow \text{Hom}(Fc, Fc')$ is inj.

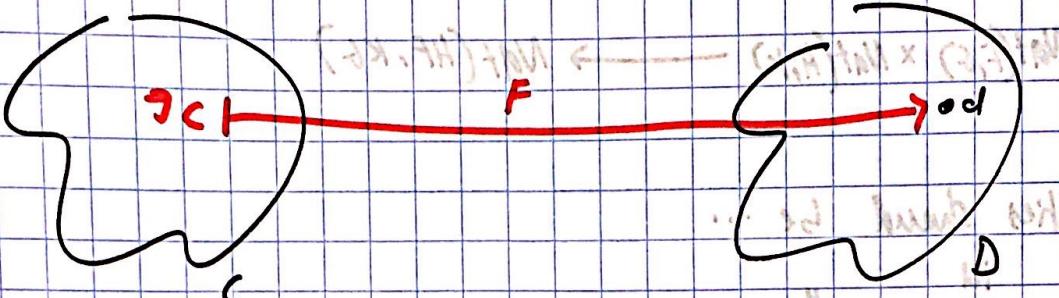
fully-faithful: $\text{Hom}_C(c, c') \xrightarrow{\sim} \text{Hom}_D(Fc, Fc')$

are bijections of sets.

essentially surjective is a condition on objects:

Guess: For any $d \in D$, there should be a $c \in C$ w/

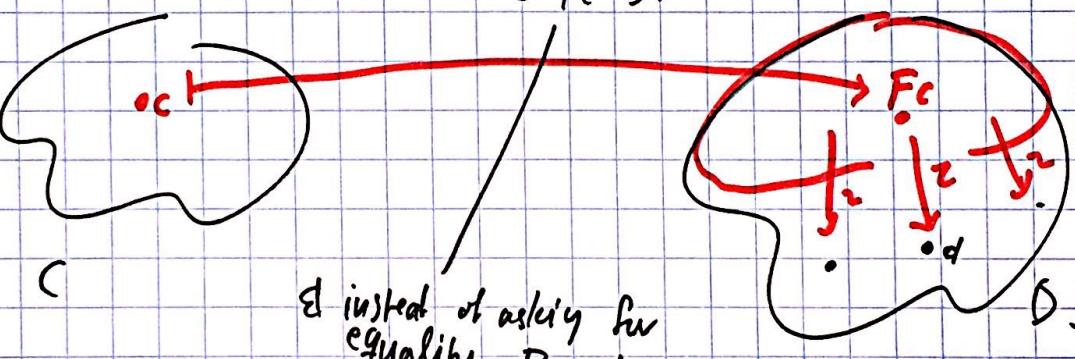
$$Fc = d$$



But this is too strong.

We take instead a slightly weaker version...

For any $d \in D$, there's a $c \in C$
 $\rightsquigarrow Fc \in D$.

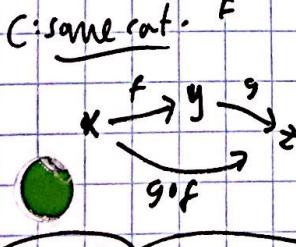


& instead of asking for
equality $Fc = d$

we ask instead for an iso. $Fc \xrightarrow{\sim} d$

injectivity on objects isn't required.

vertical ~~composition~~



$A, B \in \text{Set}$.
 $A \times B = \{(a, b) : a \in A, b \in B\}$.

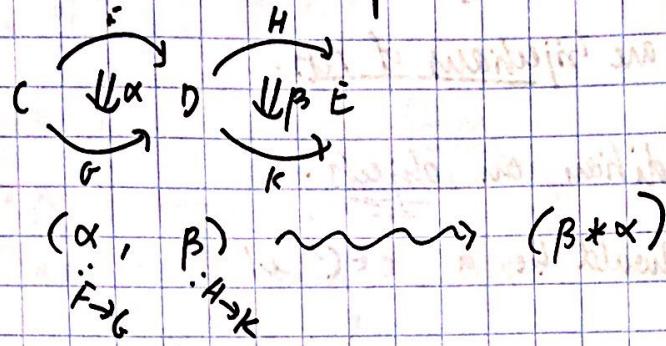
composition in ordinary cats. is a function
(set-morphism)

$$\text{Hom}(x, y) \times \text{Hom}_C(y, z) \xrightarrow{\sim} \text{Hom}_C(x, z)$$

f parallel product g $\mapsto g \circ f$

$$R^2 = R \times R.$$

in 2-cats horizontal comp.



identities should be ...

