

# Example Class Project 1

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# Content

- Implementation
  - Screenshot + Pseudocode?
- Generate randomised list
- Time complexity
  - Same  $S$  (size of smallest subarray), diff  $n$  (size of list)
  - Same  $n$ , diff  $S$
  - Using diff input datasets (random, reversed, etc)
  - Find the optimal value of  $S$
- Compare w original mergesort
  - Performance of merge sort
  - Comparison between mergesort and hybrid algorithm

# Implementation

```
HybridSort(list, start, end):
```

```
    If sizeofList < S:
```

```
        InsertionSort(list)
```

```
    Else:
```

```
        HybridSort(list, start,  
mid)
```

```
        HybridSort(list, mid, end)
```

```
        merge(list, start, mid,  
end)
```

```
MergeSort(list, start, end):
```

```
    if sizeofList <= 1:
```

```
        return
```

```
    Else:
```

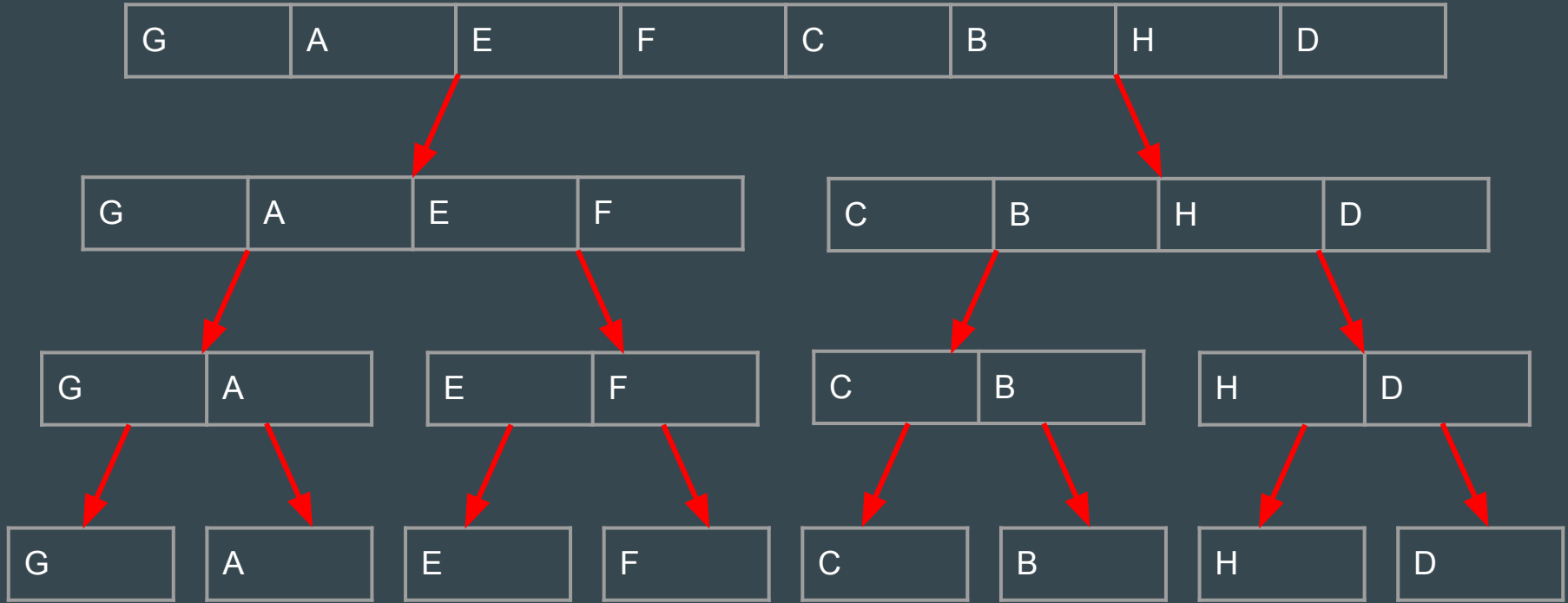
```
        MergeSort(list, start, mid)
```

```
        MergeSort(list, mid, end)
```

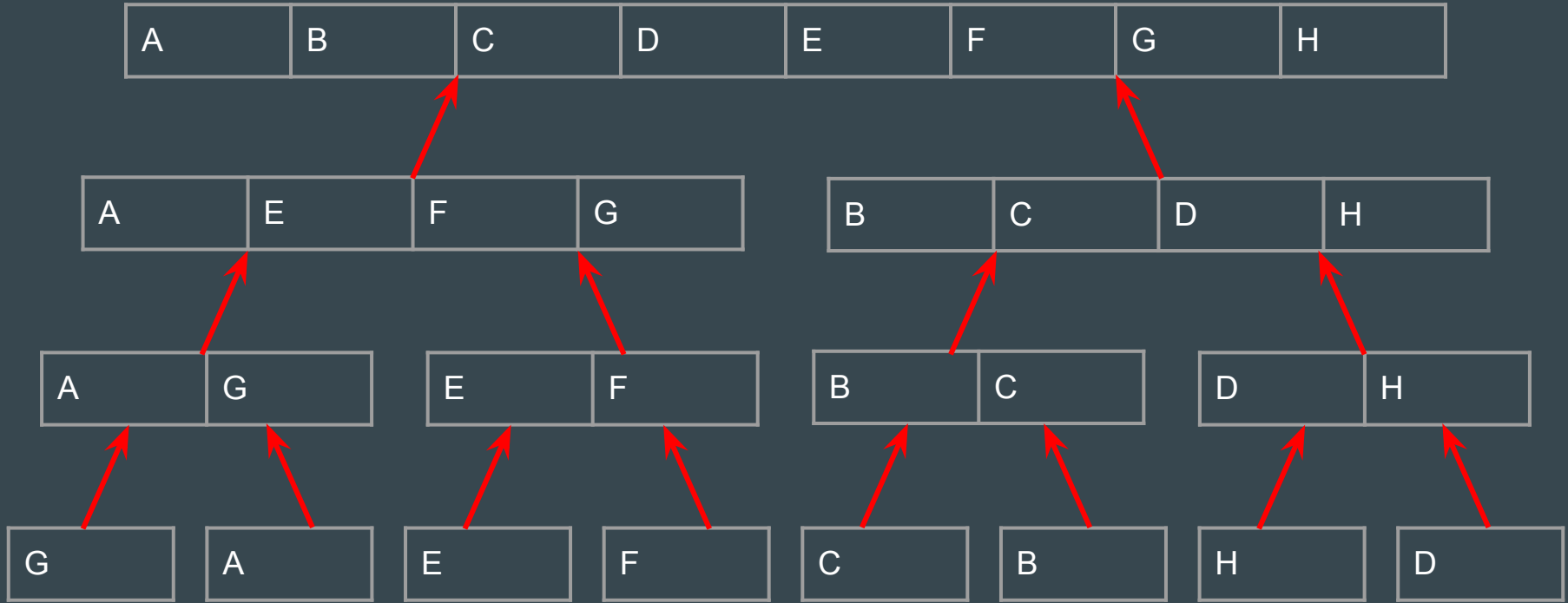
```
        merge(list, start, mid,
```

```
end)
```

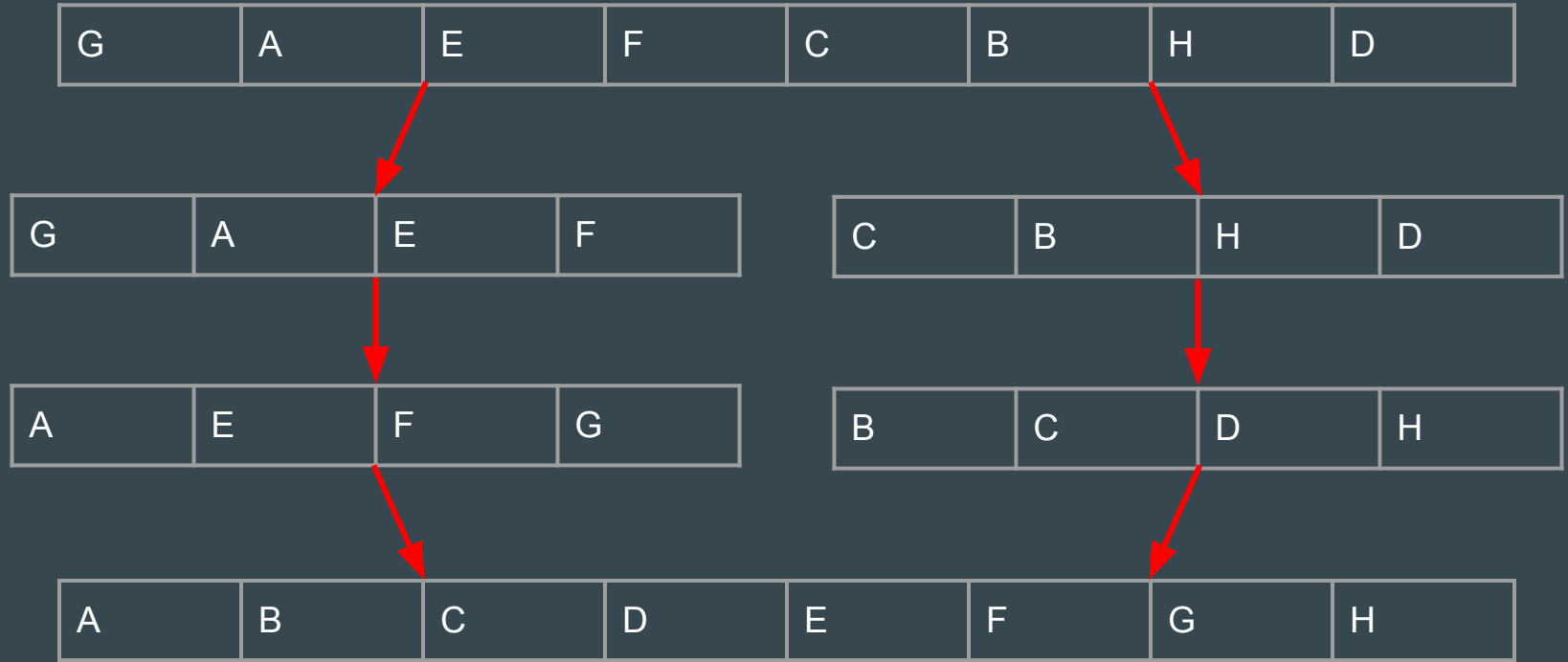
# Merge Sort



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# Implementation

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    If sizeofList < S:
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        HybridSort(list, start, mid)
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        HybridSort(list, mid, end)
```

```
        merge(list, start, mid, end)
```

```
InsertionSort(list, sizeofList):
```

```
    //Incremental Approach
```

```
    for (int i = 1; i < sizeofList; ++i):
```

```
        For (int j = i; j > 0; j--)
```

```
            if(list[j] < list[j-1])
```

```
                swap(list[j], list[j-1])
```

```
            Else break;
```

# Implementation

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HybridSort(list, start, end):
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    If sizeofList < S:
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        InsertionSort(list)
```

```
    Else:
```

```
        HybridSort(list, start, mid)
```

```
        HybridSort(list, mid, end)
```

```
        merge(list, start, mid, end)
```

```
merge(list, start, mid, end) {  
    // Split the list into left and right  
    left = list[start:mid]  
    right = list[mid:end]  
    int i, j = 0  
    int k = start;  
    //Add smaller item of the 2 sublists to the  
list  
    while(i < sizeofLeft && j < sizeofRight):  
        if left[i] < right[j]: append left to  
list  
        else: append right to list  
    // Add the remaining items to the list  
    while (i < sizeofLeft)  
        append left[i] to list  
        i++  
    while (j < sizeofRight)  
        append right[j] to list  
        j++  
}
```



# Time Complexity Theoretical Analysis

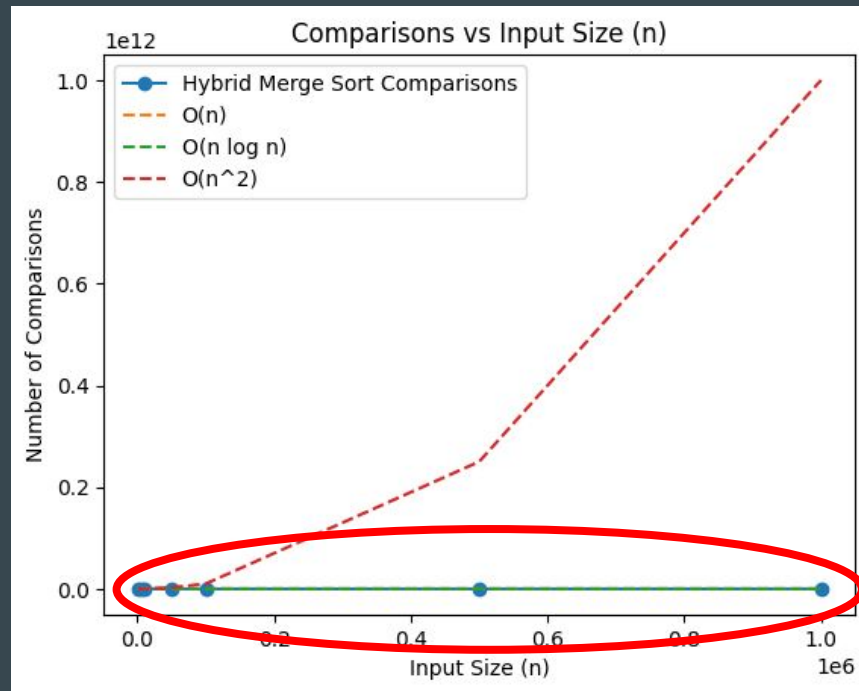
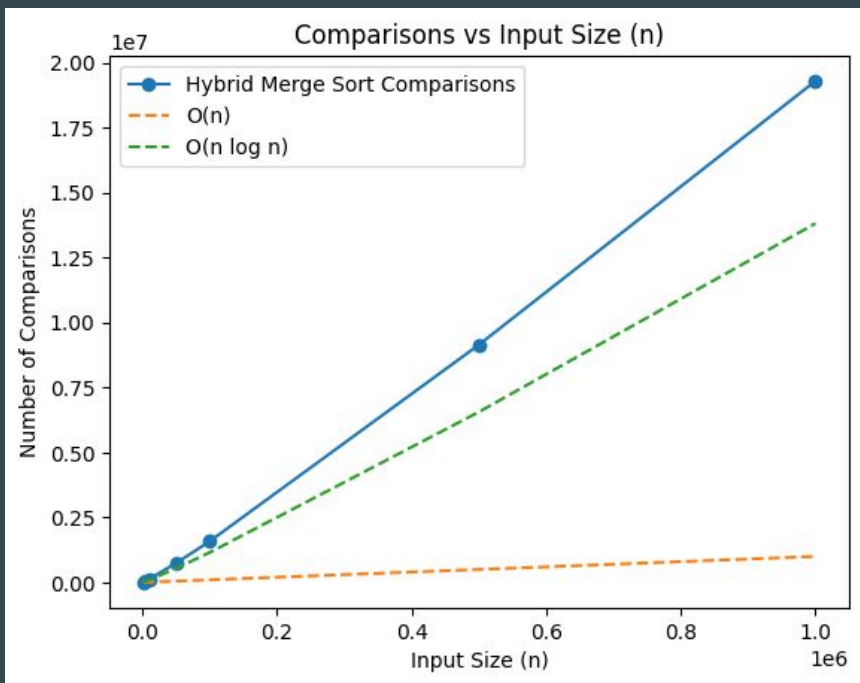
In the worst case

- For small  $S$ , behaves more like Merge Sort:  $O(n \log n)$
- For large  $S$ , behaves more like Insertion Sort:  $O(n^2)$

Assuming an optimal value of  $S$  is used,

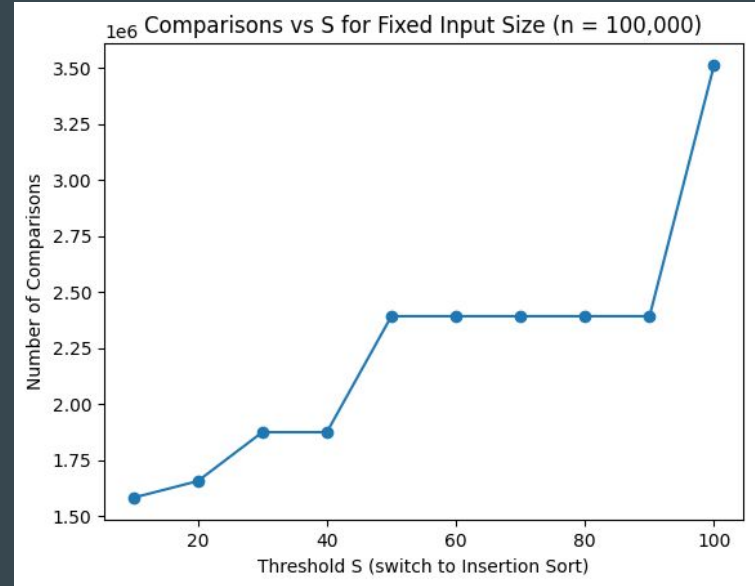
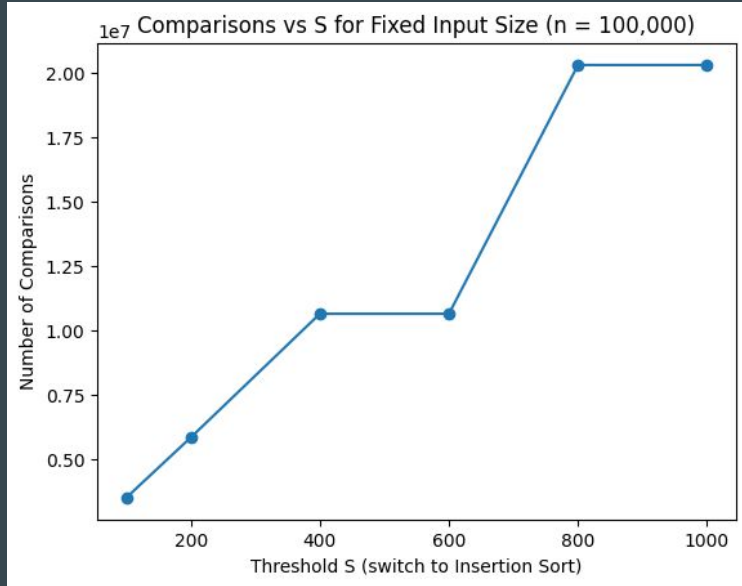
the algorithm has  $O(n \log n)$  Complexity

# Empirical Results (Fixed S)



$O(n \log$

# Empirical results (fixed n)

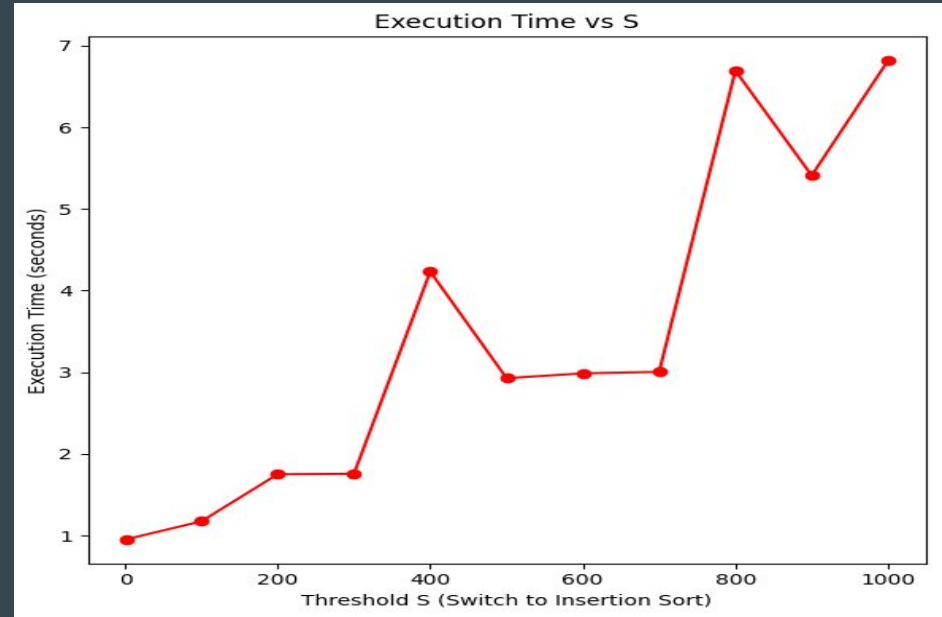


**Smaller S = less comparisons = better runtime?**

# Optimal S value

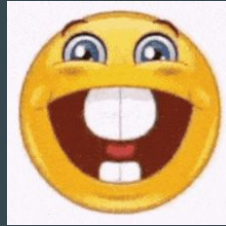
Somewhat linear relationship  
between S and runtime

Less comparisons  $\neq$  faster!!!  
Why?

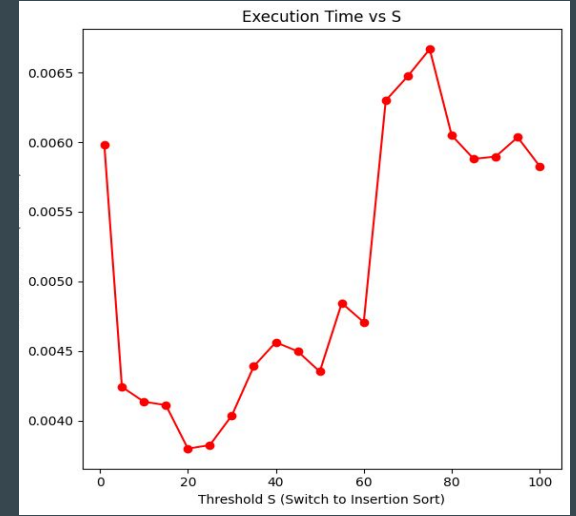
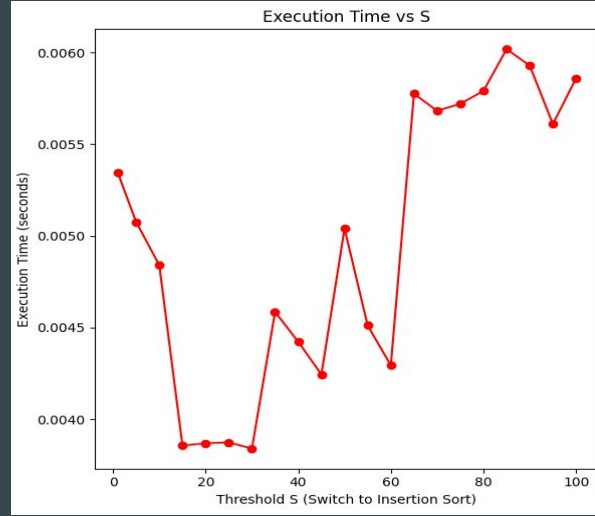
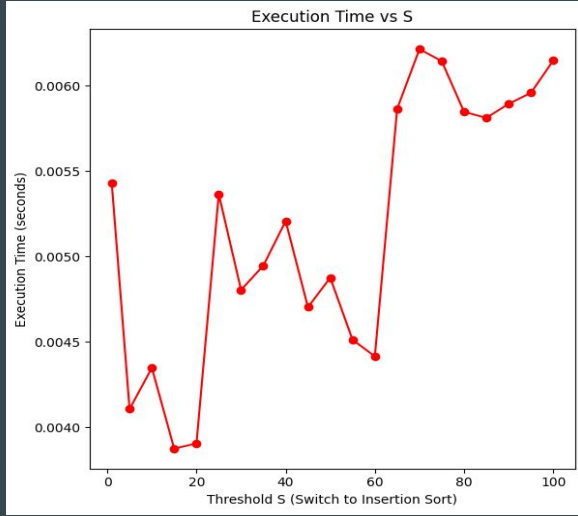


High  
recursion  
overhead

Prevalence of  
Insertion sort  
 $O(n^2)$

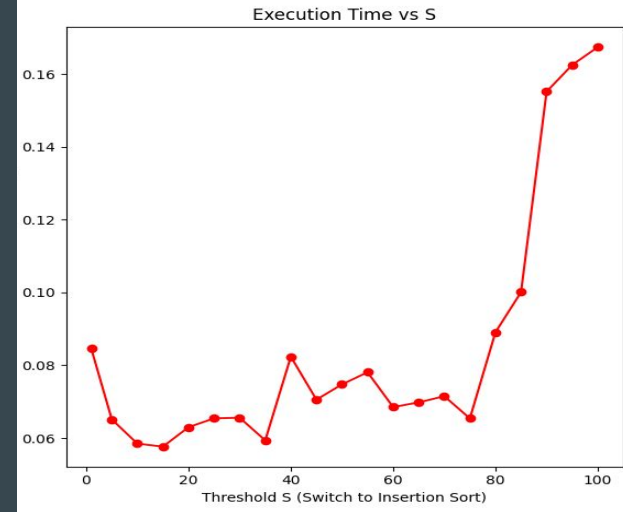
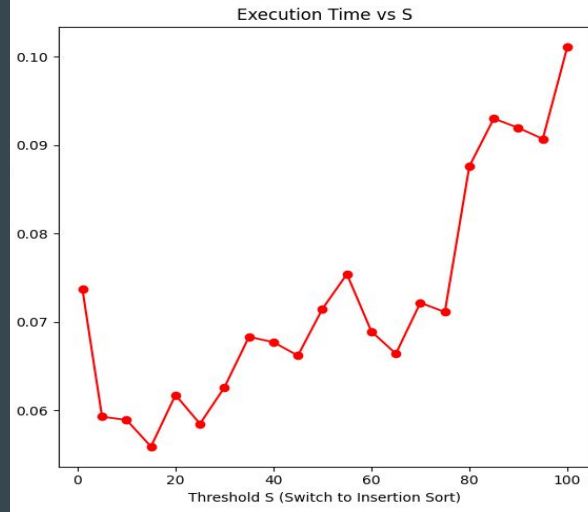
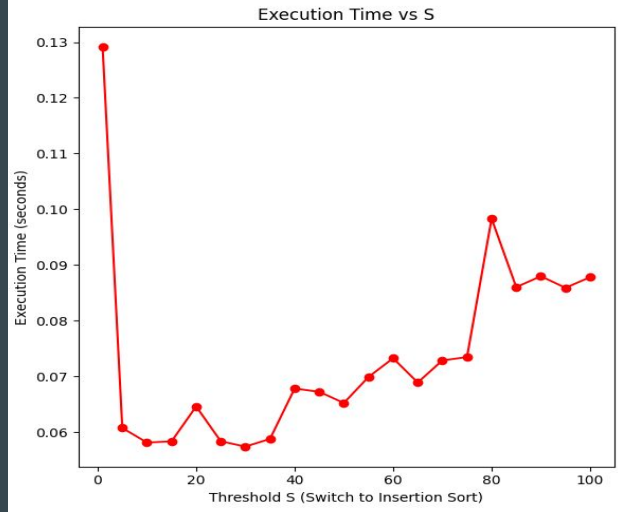


$n = 1,000$



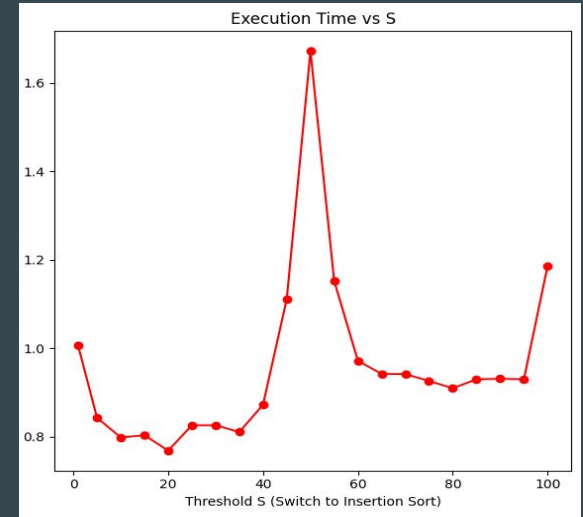
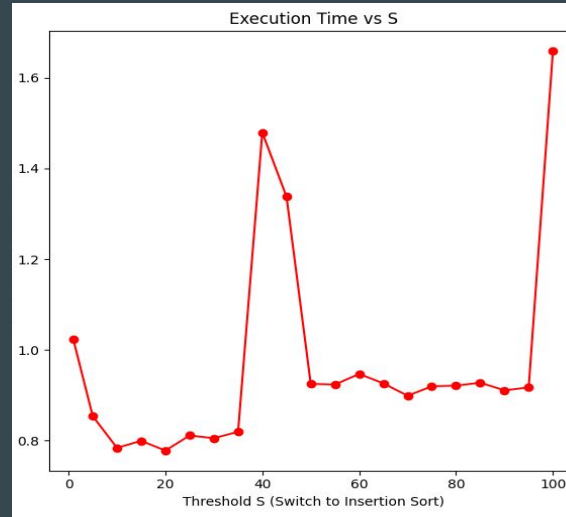
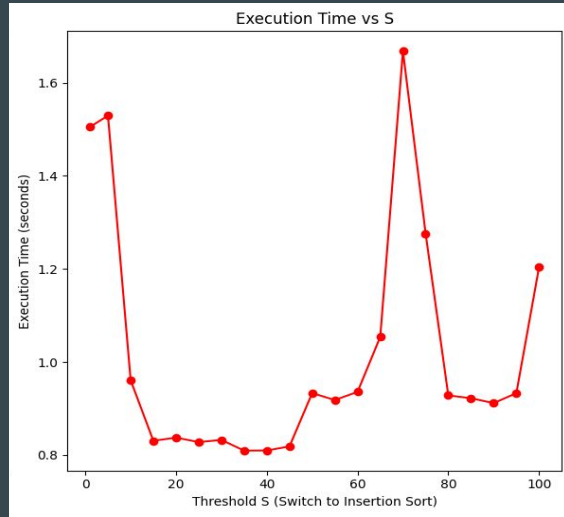
$15 < S < 30$

$n = 10,000$



$10 < S < 35$

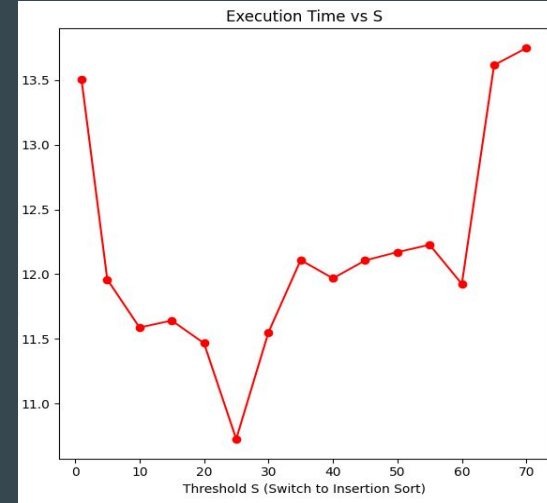
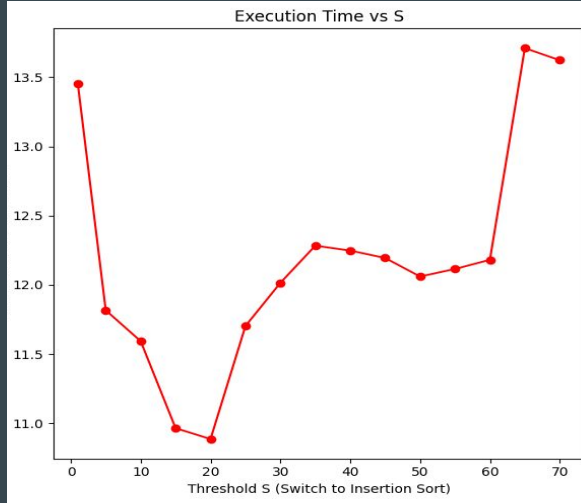
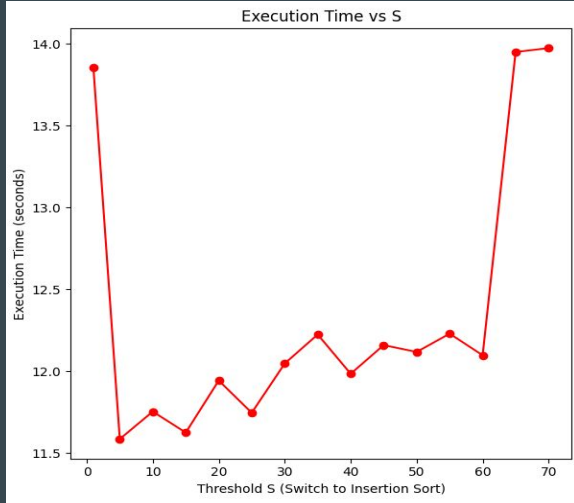
$n = 100,000$



$10 < S < 40$



$n = 1,000,000$



$10 < S < 30$

# Conclusion:

Optimal S:  $15 < S < 25$



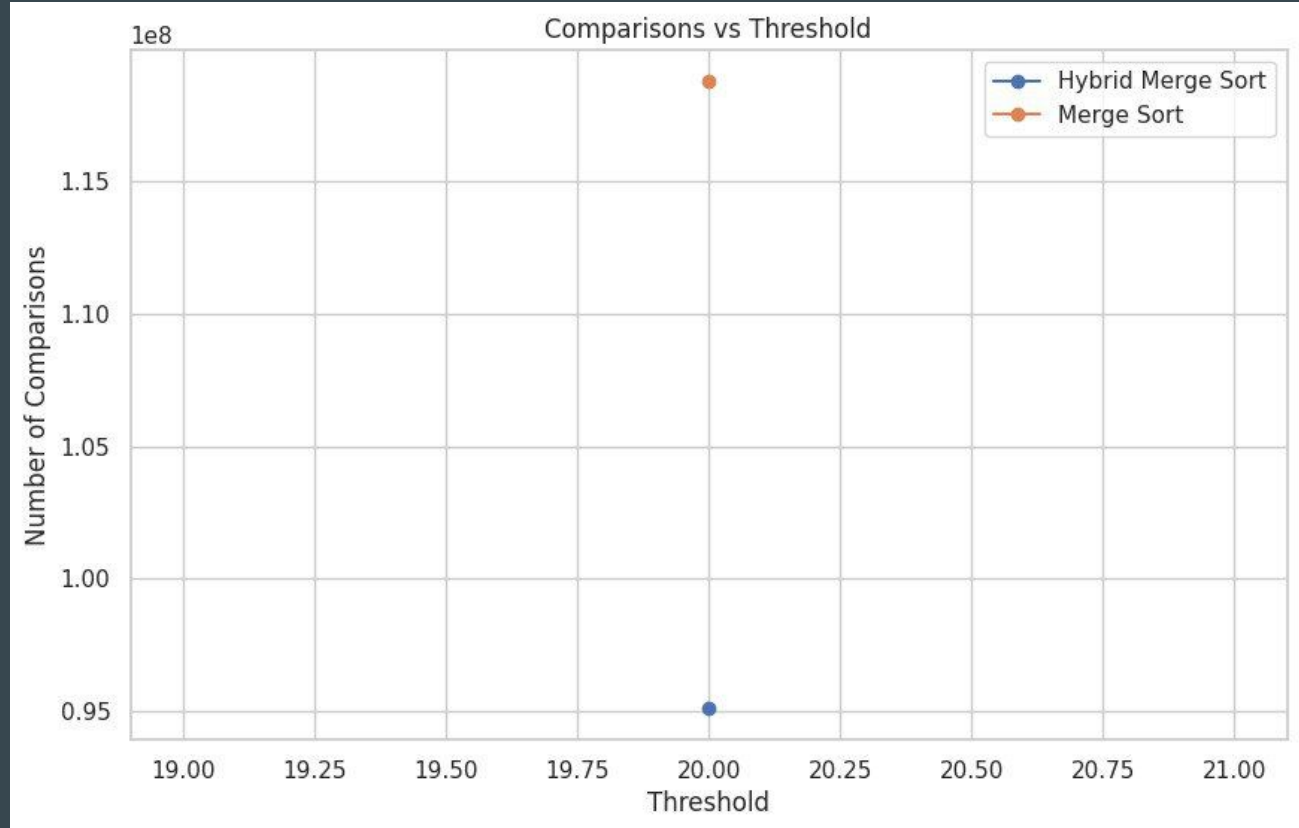
# Comparing w mergesort

S = 20, Size= 10 million

	Hybrid sort algorithm	Merge sort algorithm
No. of comparisons	95,079,616	118,788,160
Time taken(seconds)	157.6165	197.6796

# No. of Comparisons

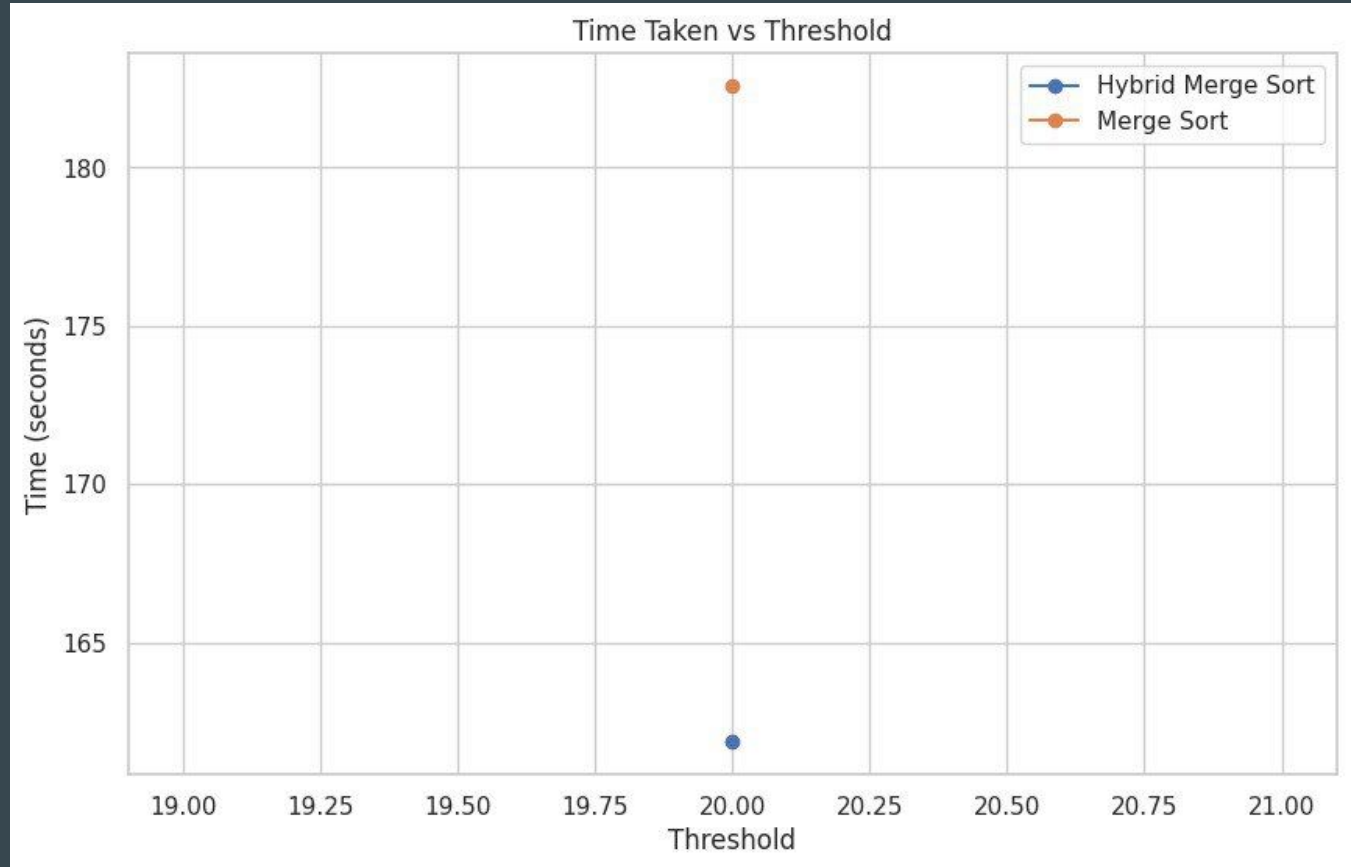
-Hybrid more efficient  
in smaller subarrays



# CPU Time

-Lesser comparison =  
faster execution time

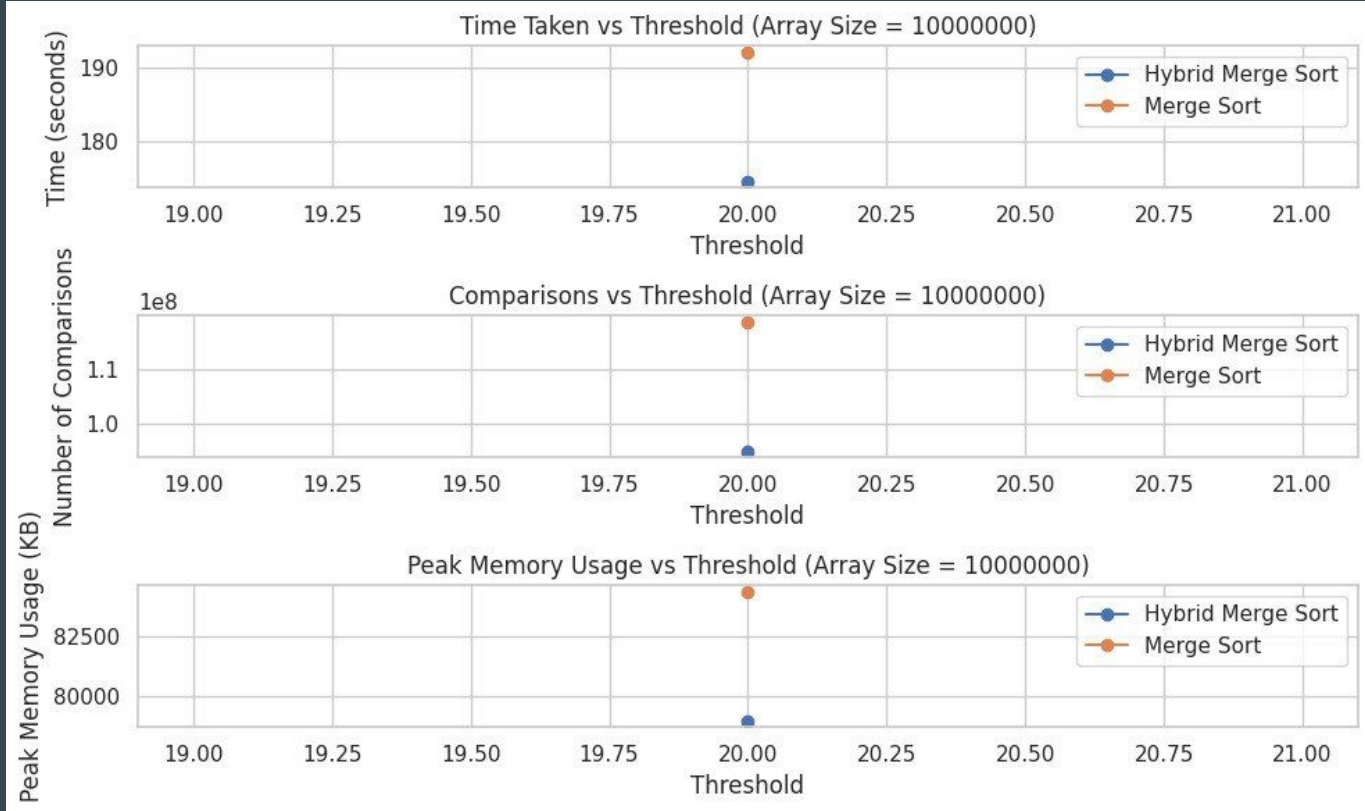
-Lesser recursive calls  
and merge operations



# Impact of optimal threshold value(S)

-Balances overhead of recursive calls and efficiency of insertion sort for small subarrays

-Optimal threshold value improved hybrid's practical performance in efficiency



**Thank you**