



AAI2007 Introduction to Algorithms

Week 1: Course Introduction

Instructor: Jinyoung Han (jinyounghan@skku.edu)



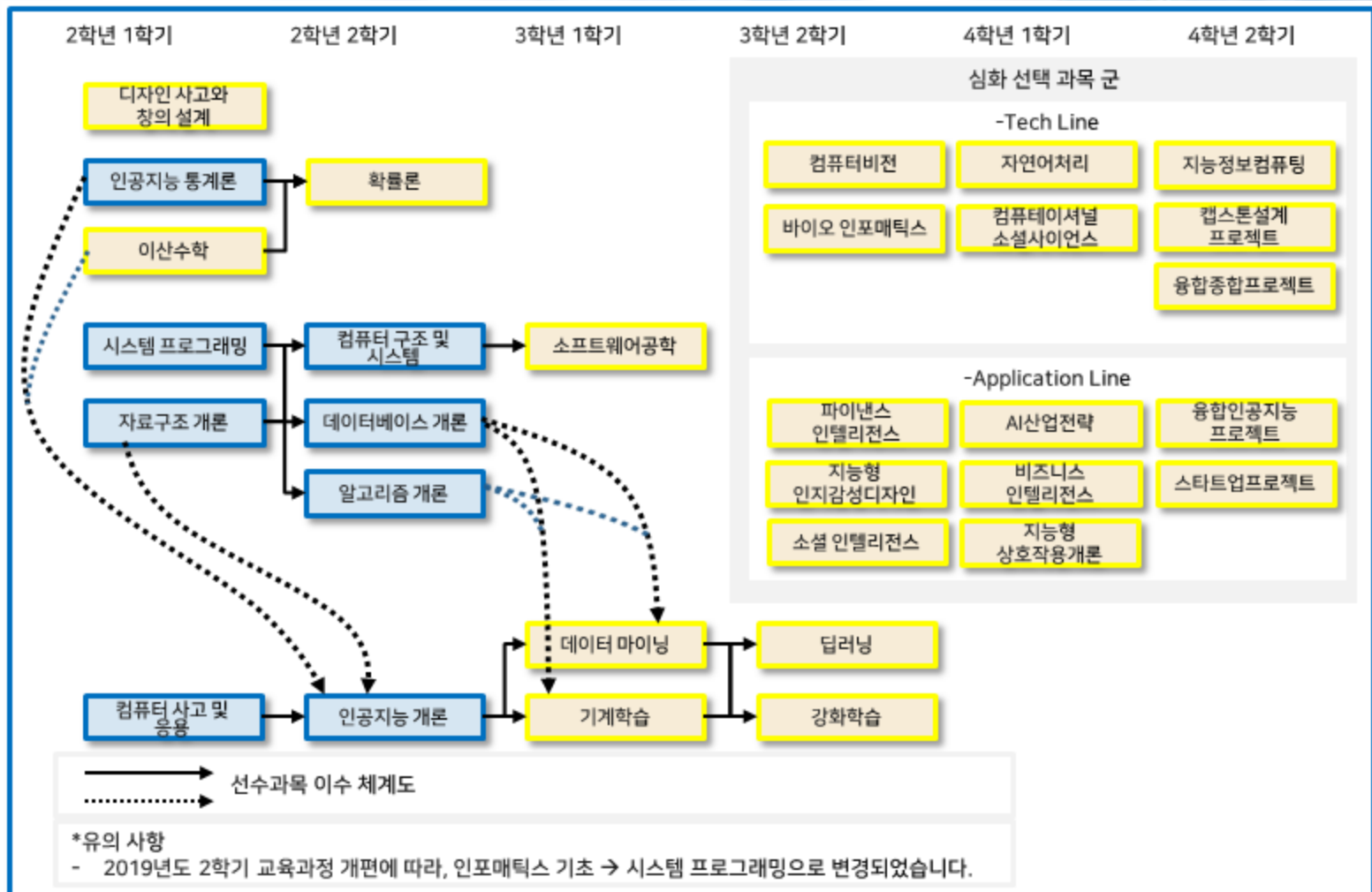
In This Lecture

Outline

1. Course Logistics
2. Algorithm Basic
3. Level Test
4. Programming Environment Setting

01. Course Logistics

Course Position



01. Course Logistics

Course Info

- Instructor: 한진영 교수
 - Office: 국제관 313호
 - Tel: 02-740-1797
 - Email: jinyoungghan@skku.edu
 - Office hour: Wed. 13:00-15:00
 - Appointment by email
- TAs
 - 김지나: jinakimcool@gmail.com
 - 강지원: kanjiwon@gmail.com
 - 윤지우: yoonjeewoo@gmail.com

01. Course Logistics

Course Overview

- Goals
 - Understanding basic computer algorithms
 - Hands-on experience on solving problems with programming
- Course Structure
 - Lecture (6:00pm~7:30pm)
 - Practice (7:30pm~9:00pm)
 - PC or Notebook
- Prerequisite
 - Data Structure, Programming
 - C or Python?

01. Course Logistics

Grading Policy

- Scores breakdown
 - Lecture attendance (10)
 - Quiz or Homework (10)
 - Practice (20)
 - Midterm/Final (50)

01. Course Logistics

Resources

- Lecture slides
 - Will be uploaded to the system
 - All the information will be given in the system, so please check the system frequently
- Reference material
 - Introduction to Algorithms
 - 쉽게 배우는 알고리즘
 - 뇌를 자극하는 알고리즘

01. Course Logistics

No Cheating

- No submission is better than cheating
 - Zero Tolerance Cheating Policy
- Definition of cheating in this class
 - Knowingly or unknowingly participating in the submission of unoriginal work for any test (e.g., lab exercise, project)
 - Answer to roll-call instead of another student is also regarded as cheating
- Penalty
 - Assign a fail grade

01. Course Logistics

Tentative Schedule

수업일	내용
9/4	Course Introduction, Algorithm Basic, Level Test
9/11	Order of Complexity, List
9/18	Stack, Queue
9/25	Tree, Binary Search Tree (BST), 건학 기념일
10/2	Priority Queue, Heap
10/9	Graph (1), 한글날
10/16	Graph (2)
10/23	Sorting
10/30	Searching
11/6	Midterm
11/13	Hash Table
11/20	Dynamic Programming (1)
11/27	Dynamic Programming (2)
12/4	Greedy Algorithms
12/11	NP-completeness
12/18	Final



Algorithm Basic

02. Algorithm Basic

Algorithm

- An algorithm: a set of instructions, typically to solve a class of problems or perform a computation
- Typical design (or problem solving methods) of algorithms
 - divide-and-conquer, dynamic programming, greedy approach, backtracking, branch-and-bound
- Algorithm Analysis
 - Efficiency analysis through computational complexity
 - Lower bounds for sorting and searching problems
 - Intractability (NP-completeness)

02. Algorithm Basic

Algorithm Description

- Natural languages
- Programming languages
- Pseudo-code
 - similar, but not identical, to C++/Java.
 - Notable exceptions: unlimited array index, variable-sized array, mathematical expressions, use of arbitrary types, convenient control structure, and etc.
 - E.g., `low <= x && x <= high -> low <= x <= high`
 - E.g., `temp = x; x = y; y = temp -> exchange x and y`

02. Algorithm Basic

An Example

- Title: Sequential Search
- Problem
 - Is the key x in the array S of n keys?
- Inputs (parameters)
 - Positive integer n , array of keys S indexed from 1 to n .
- Outputs
 - The location of x in S . (0 if x is not in S .)
- Algorithm
 - Starting with the first item in S , compare x with each item in S in sequence until x is found or S is exhausted. If x is found, answer yes; if x is not found, answer no.

02. Algorithm Basic

An Example

- Pseudo-code

```
void seqsearch(int n, // Input(1)
               const keytype S[], // (2)
               keytype x, // (3)
               index& location) // Output
{
    location = 1;
    while (location <= n && S[location] != x)
        location++;
    if (location > n)
        location = 0;
}
```

02. Algorithm Basic

An Example

- How many comparisons for searching x in S ?
 - Depends on the location of x in S
 - In worst case, we should compare (n) times.
 - In best case,
- Any better algorithm to get the same solution?
 - No, we can't, unless there is an extra information in S .

02. Algorithm Basic

An Example

- Next approach: Binary Search
- Problem
 - Is the key x in the array S of n keys?
 - Determine whether x is in the sorted array S of n keys.
- Inputs (parameters)
 - Positive integer n , **sorted** (non-decreasing order) array of keys S indexed from 1 to n
- Outputs
 - The location of x in S . (0 if x is not in S .)

02. Algorithm Basic

An Example

- Pseudo-code

```
void binarysearch(int n, // Input(1)
                  const keytype S[], // (2)
                  keytype x, // (3)
                  index& location) // Output
{
    index low, high, mid;
    low = 1; high = n;
    location = 0;
    while (low <= high && location == 0) {
        mid = (low + high) / 2; // integer div
        if (x == S[mid]) location = mid;
        else if (x < S[mid]) high = mid - 1;
        else low = mid + 1;
    }
}
```

02. Algorithm Basic

An Example

- How many comparisons for searching x in S ?
 - S is already sorted. We know it.
 - For each statement in a while-loop, the number of searching targets will be half.
 - In worst case, we should compare ($\log_2 n + 1$) times.
 - e.g. $n = 32$. $S[16]$, $S[16+8]$, $S[24+4]$, $S[28+2]$, $S[30+1]$. **$S[32]$**
 - In best case, trivial to answer.
- Any better algorithm to get the same solution?

02. Algorithm Basic

An Example

- Number of comparisons done by sequential search and binary search when x is larger than all the array items

Array Size	Number of Comparisons by Sequential Search	Number of Comparisons by Binary Search
128	128	8
1,024	1,024	11
1,048,576	1,048,576	21
4,294,967,296	4,294,967,296	33

02. Algorithm Basic

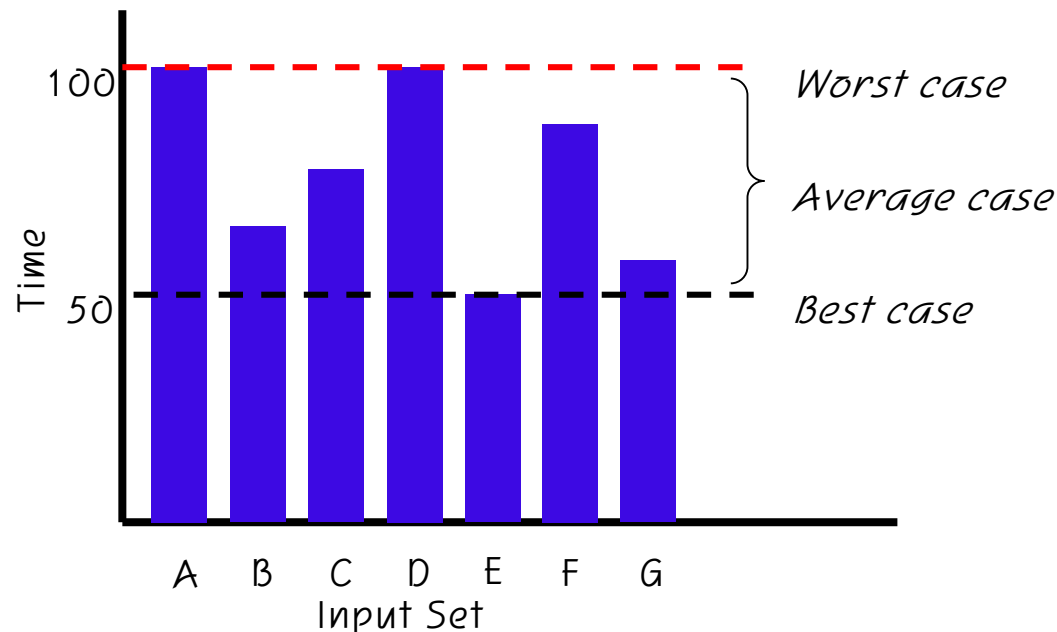
Analysis

- Time complexity analysis
 - Determination of how many times the basic operation is done for each value of the input size
 - Should be independent of CPU, OS, Programming languages...
- Metrics
 - Basic operation
 - Comparisons, assignments, etc.
 - Input size
 - The number of elements in an array
 - The length of a list
 - The number of columns and rows in a matrix
 - The number of vertices and edges in a graph

02. Algorithm Basic

Analysis

- We can consider the following three cases:
 - Best case
 - Average case
 - Worst case



02. Algorithm Basic

Analysis

- Among four possible analysis, WC, AC, and BC, which one is the right one?
- Think about ...
 - you are working for a nuclear power plant
 - what if you are working for an internet shopping mall
- Which one do you think is the most useless?
- Which one do you think is the hardest to analyze?

02. Algorithm Basic

Correctness

- Efficiency analysis vs. Correctness analysis
- In this course, when I say an algorithm analysis, I mean an efficiency analysis
- We can also analyze the correctness of an algorithm by developing a mathematical proof that the algorithm actually does what it is supposed to do
- An algorithm is incorrect...
 - if it does not stop for a given input or
 - if it gives a wrong answer for an input



Level Test



Programming Environment Setting

Thanks

Week 1: Course Introduction

Instructor: Jinyoung Han (jinyounghan@skku.edu)

