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Numerical Heat Transfer
for Applications

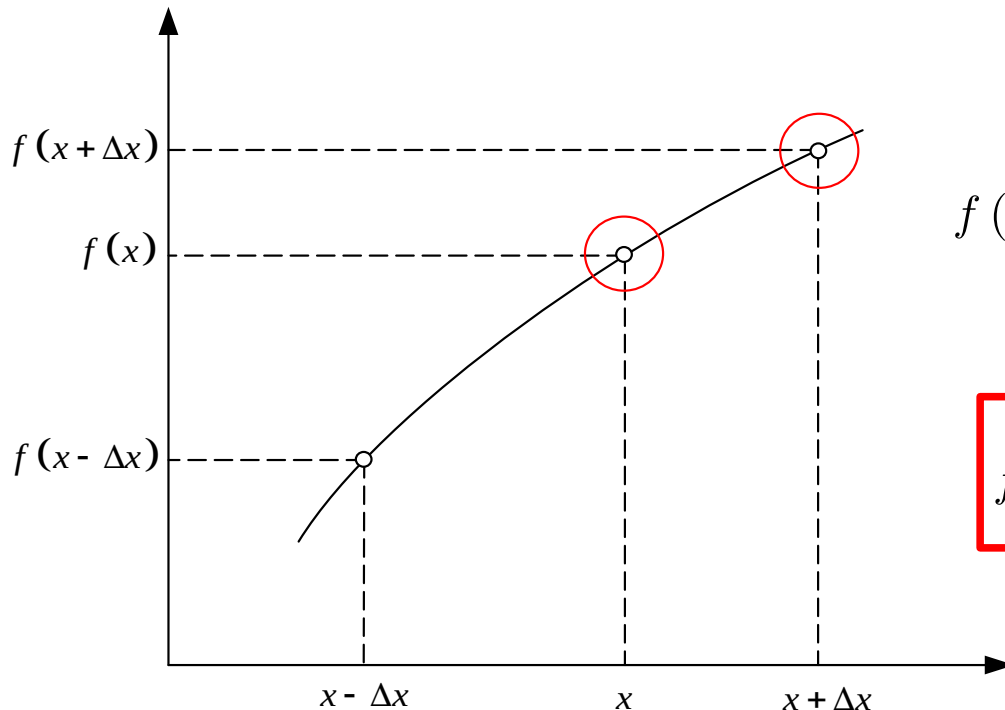
**Finite Difference Method for
Steady State Heat Conduction**

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Finite difference method concept

A formal basis for developing finite difference approximation of derivatives is through the use of Taylor series expansion. Consider Taylor series expansion of a function $f(x)$ about a given point in the forward (*i.e.*, positive x):

$$f(x + \Delta x) = f(x) + f'(x) \Delta x + f''(x) \frac{\Delta x^2}{2!} + f'''(x) \frac{\Delta x^3}{3!} + \dots$$



$$f(x + \Delta x) = f(x) + f'(x) \Delta x + O(\Delta x^2)$$



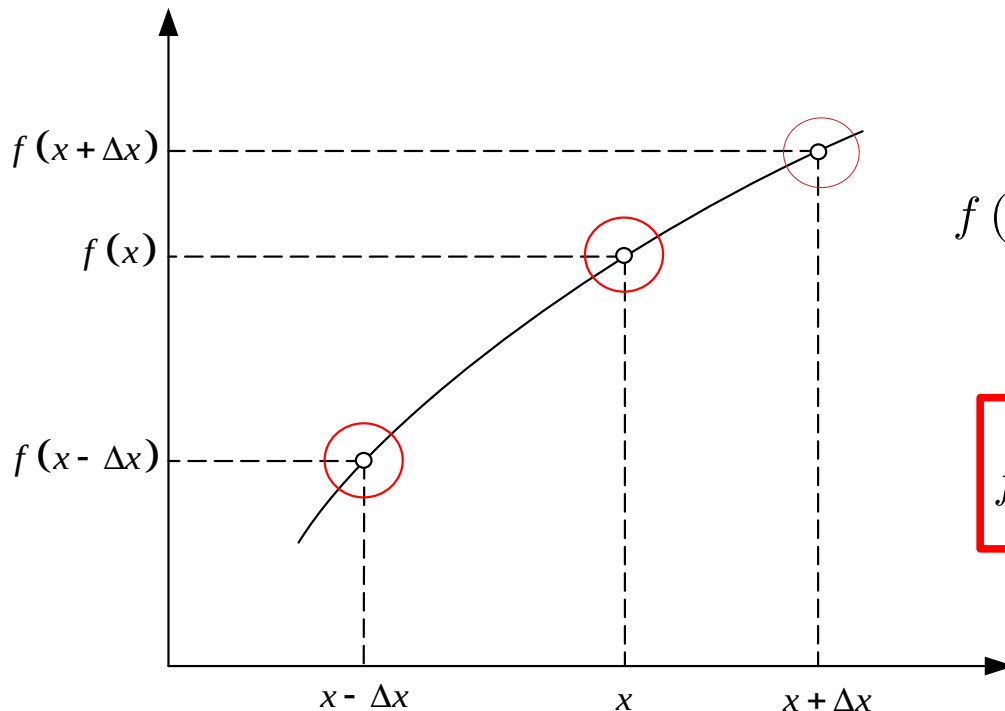
$$f'(x) = \frac{f(x + \Delta x) - f(x)}{\Delta x} + O(\Delta x)$$

Forward finite difference approx.

Finite difference method concept

Similarly, it is possible to consider Taylor series expansion of a function $f(x)$ about a given point in the backward (*i.e.*, negative x):

$$f(x - \Delta x) = f(x) - f'(x) \Delta x + f''(x) \frac{\Delta x^2}{2!} - f''(x) \frac{\Delta x^3}{3!} + \dots$$



$$f(x - \Delta x) = f(x) - f'(x) \Delta x + O(\Delta x^2)$$



$$f'(x) = \frac{f(x) - f(x - \Delta x)}{\Delta x} + O(\Delta x)$$

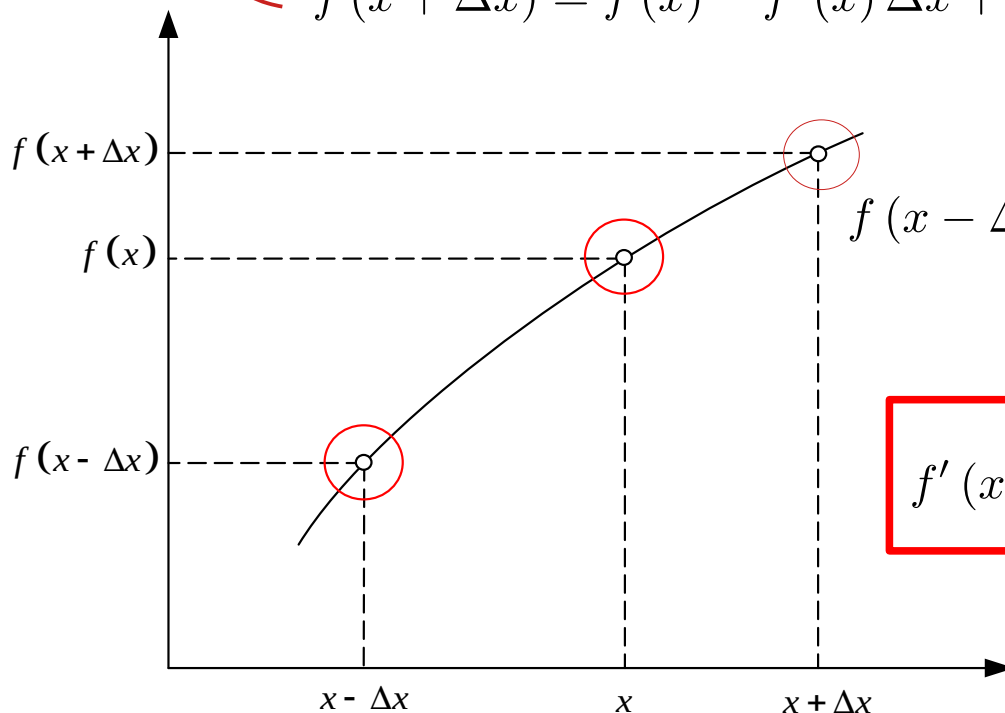
Backward finite difference approx.

Finite difference method concept

Considering both forward and backward differentiation approach is possible to obtain a third kind of differentiation: the **central approach**.

$$f(x + \Delta x) = f(x) + f'(x) \Delta x + f''(x) \frac{\Delta x^2}{2!} + O(\Delta x^3)$$

$$f(x - \Delta x) = f(x) - f'(x) \Delta x + f''(x) \frac{\Delta x^2}{2!} + O(\Delta x^3)$$



$$f(x + \Delta x) - f(x - \Delta x) = 2f'(x) \Delta x + O(\Delta x^3)$$



$$f'(x) = \frac{f(x + \Delta x) - f(x - \Delta x)}{2\Delta x} + O(\Delta x^2)$$

Central finite difference approx.

Finite difference method concept

Considering both forward and backward differentiation approach is possible to obtain also a specific approximation for second derivative:

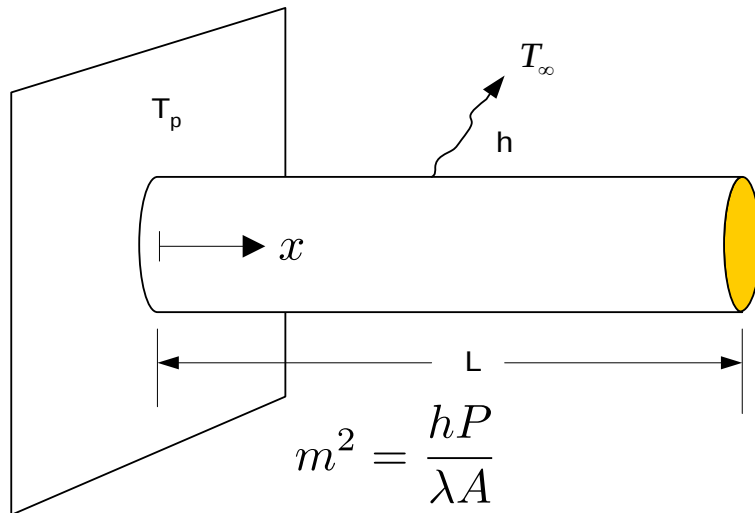
$$\begin{aligned} f(x + \Delta x) &= f(x) + f'(x) \Delta x + f''(x) \frac{\Delta x^2}{2!} + f'''(x) \frac{\Delta x^3}{3!} + O(\Delta x^4) \\ f(x - \Delta x) &= f(x) - f'(x) \Delta x + f''(x) \frac{\Delta x^2}{2!} - f'''(x) \frac{\Delta x^3}{3!} + O(\Delta x^4) \end{aligned}$$

$$f(x + \Delta x) + f(x - \Delta x) = 2f(x) + f''(x) \Delta x^2 + O(\Delta x^4)$$



$$f''(x) = \frac{f(x + \Delta x) - 2f(x) + f(x - \Delta x)}{\Delta x^2} + O(\Delta x^2)$$

Finite difference solution of fin equation



$$x^* = \frac{x}{L}$$

$$T^* = \frac{T - T_\infty}{T_b - T_\infty}$$



$$\frac{d^2 T}{dx^2} - m^2 (T - T_\infty) = 0$$

$$T(x = 0) = T_b$$

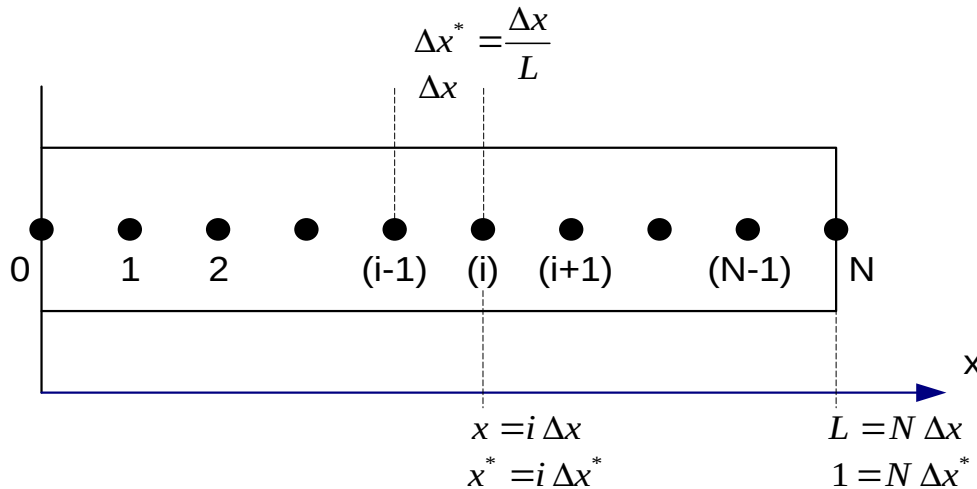
$$\left. \frac{dT}{dx} \right|_{x=L} = 0$$

$$\frac{d^2 T^*}{dx^{*2}} - (mL)^2 T^* = 0$$

$$T^*(x^* = 0) = 1$$

$$\left. \frac{dT^*}{dx^*} \right|_{x^*=1} = 0$$

Finite difference solution of fin equation



$$\frac{d^2 T^*}{dx^{*2}} - (mL)^2 T^* = 0$$

$$T^*(x^* = 0) = 1$$

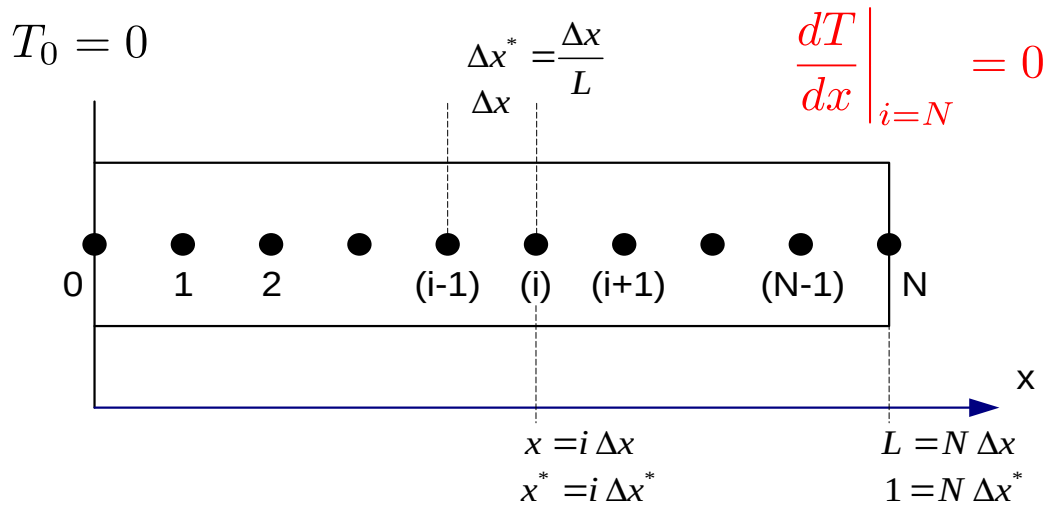
$$\left. \frac{dT^*}{dx^*} \right|_{x^*=1} = 0$$

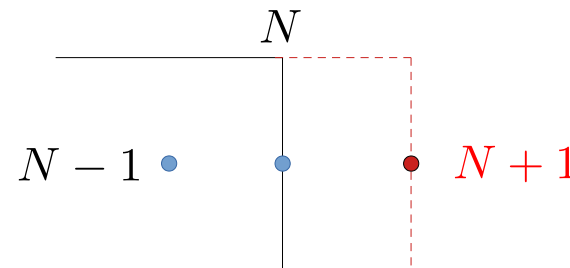
$$\frac{T_{i+1} - 2T_i + T_{i-1}}{\Delta x^2} - (mL)^2 T_i = 0 \quad 2 \leq i \leq N - 1$$



$$-T_{i-1} + [2 + (mL \Delta x)^2] T_i - T_{i+1} = 0$$

Finite difference solution of fin equation





$\left. \frac{dT}{dx} \right|_{i=N} = 0 \Rightarrow \frac{T_{N+1} - T_{N-1}}{2\Delta x} = 0$

$T_{N+1} = T_{N-1}$

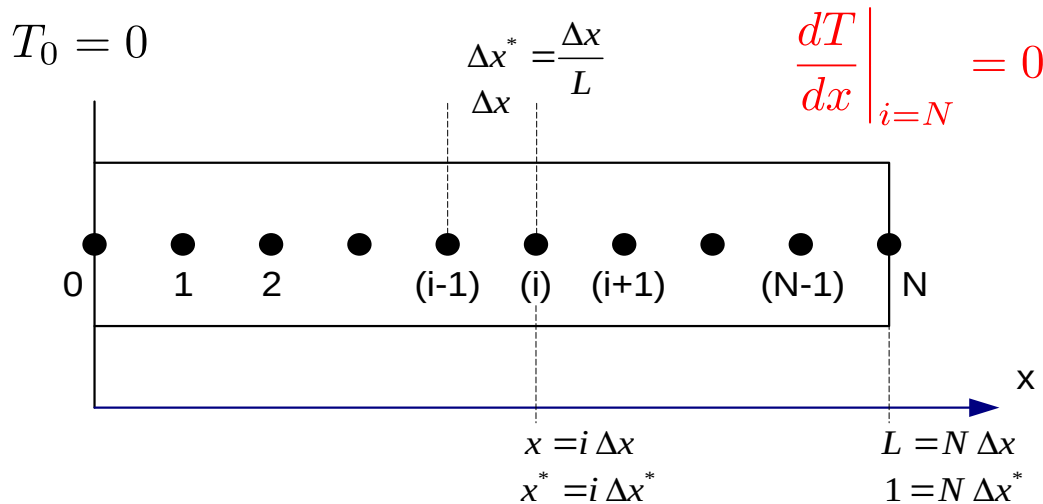
The B.C. at fin head requires a particular treatment. A possible strategy is consider a so called ghost-node which allows to use to discretized the boundary condition itself.

$$-T_{N-1} + [2 + (mL \Delta x)^2]T_N - T_{N+1} = 0 \quad i = N$$



$-2T_{N-1} + [2 + (mL \Delta x)^2]T_N = 0$

Finite difference solution of fin equation



$$\frac{d^2 T^*}{dx^{*2}} - (mL)^2 T^* = 0$$

$$T^*(x^* = 0) = 1$$

$$\left. \frac{dT^*}{dx^*} \right|_{x^*=1} = 0$$

$$[2 + (mL \Delta x)^2] T_1 - T_2 = 1 \quad i = 1$$

$$-T_{i-1} + [2 + (mL \Delta x)^2] T_i - T_{i+1} = 0 \quad 2 \leq i \leq N - 1$$

$$-2T_{N-1} + [2 + (mL \Delta x)^2] T_N = 0 \quad i = N$$

$$c = [2 + (mL \Delta x)^2]$$

Finite difference solution of fin equation

$$[2 + (mL \Delta x)^2]T_1 - T_2 = 1 \quad i = 1$$

$$-T_{i-1} + [2 + (mL \Delta x)^2]T_i - T_{i+1} = 0 \quad 2 \leq i \leq N - 1$$

$$-2T_{N-1} + [2 + (mL \Delta x)^2]T_N = 0 \quad i = N$$

$$c = [2 + (mL \Delta x)^2]$$

$N = 4$

$$cT_1 - T_2 = 1$$

$$-T_1 + cT_2 - T_3 = 0$$

$$-T_2 + cT_3 - T_4 = 0$$

$$-2T_3 + cT_4 = 0$$



$$\begin{bmatrix} c & -1 & 0 & 0 \\ -1 & c & -1 & 0 \\ 0 & -1 & c & -1 \\ 0 & 0 & -2 & c \end{bmatrix} \begin{bmatrix} T_1 \\ T_2 \\ T_3 \\ T_4 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

Tri-diagonal linear system

Ex1_FD - Finite difference solution of fin equation

Matlab code

```
clc
clear all

%Fin dimensions
%-----
%
L = 0.1112; % 111.2 mm
t = 0.0012; % 1.2 mm
b = 0.0206; % 20.6 mm
%
p = 2*(t+b);
Af = t*b;
%
h = 150; % [W/m2K]
lam = 202.6; % [W/mK ]

N= 30; % nodes' number
%-----
%
m = ((h*p)/(lam*Af))^0.5;
mL = m*L;
%
dx = L/N;
dx = dx/L; % dimensionless dx
c = 2 + (m*L*dx)^2;
%-----
```

$$c = [2 + (mL \Delta x)^2]$$

$$cT_1 - T_2 = 1 \quad i = 1$$

$$-T_{i-1} + cT_i - T_{i+1} = 0 \quad 2 \leq i \leq N - 1$$

$$-2T_{N-1} + cT_N = 0 \quad i = N$$

$$A = \begin{pmatrix} C & -1 & 0 & 0 & 0 \\ -1 & C & -1 & 0 & 0 \\ 0 & -1 & C & -1 & 0 \\ 0 & 0 & -1 & C & -1 \\ 0 & 0 & 0 & -2 & C \end{pmatrix}$$

Ex1_FD - Finite difference solution of fin equation

Matlab code

```
A = zeros(N,N);

A(1,1) = c;
A(1,2) = -1;

for i=2:(N-1)
    A(i,i-1) = -1;
    A(i,i) = c;
    A(i,i+1) = -1;
end

A(N,N-1) = -2;
A(N,N) = c;

%
%-----
%
b=zeros(N,1);
b(1) = 1;

% Linear system solution
%
x = A\b;

%-----
%Boundary value (x=0)
x = [1;x];
%
```

$$c = [2 + (mL \Delta x)^2]$$

$$cT_1 - T_2 = 1 \quad i = 1$$

$$-T_{i-1} + cT_i - T_{i+1} = 0 \quad 2 \leq i \leq N-1$$

$$-2T_{N-1} + cT_N = 0 \quad i = N$$

$$A = \begin{pmatrix} C & -1 & 0 & 0 & 0 \\ -1 & C & -1 & 0 & 0 \\ 0 & -1 & C & -1 & 0 \\ 0 & 0 & -1 & C & -1 \\ 0 & 0 & 0 & -2 & C \end{pmatrix}$$

$$b = [1, 0, 0, \dots, 0]^T$$

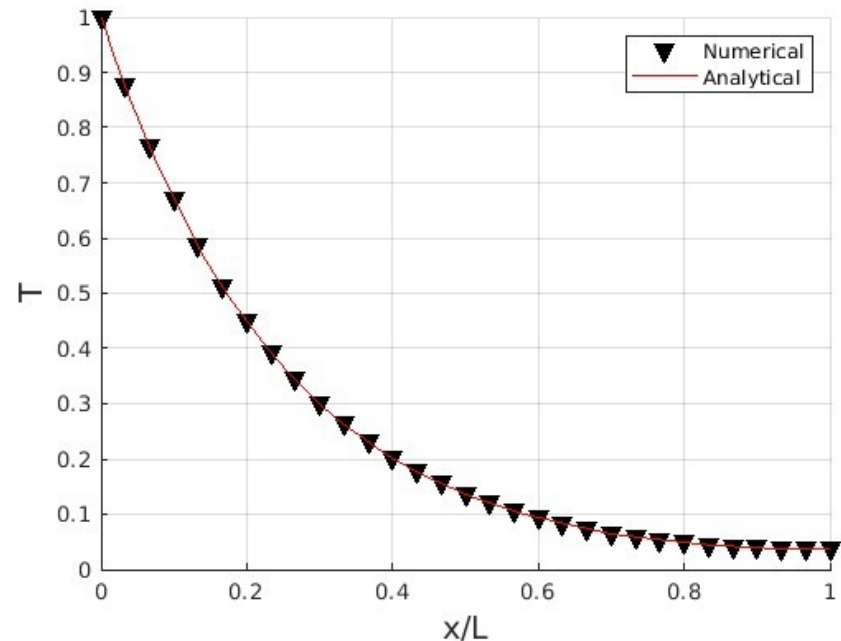
Ex1_FD - Finite difference solution of fin equation

Matlab code

```
% Analytic solution
%
%
dx1 = L/50;
dx1 = dx1/L;
lf = dx1*[0:50]';
xan = (cosh(mL*(1-lf)))/(cosh(mL));
%
% Numerical vs Analytical
%
lfn = dx*[0:N]'; %dimensionless x
grid on
%
hold on
plot(lfn, x,
'kv','markerfacecolor','k',
'markersize',8);
plot(lf, xan, 'r-');
hold off
xlim ([0 lfn(end)]);

xlabel('x/L', 'fontsize', 14);
ylabel('T', 'fontsize', 14);
legend ('Numerical', 'Analytical');
%
```

$$\frac{\theta(x)}{\theta_b} = \frac{1 - \cosh\left(mL\left(1 - \frac{x}{L}\right)\right)}{\cosh(mL)}$$



Finite difference solution of fin equation – Jacobi method

$$cT_1 - T_2 = 1$$

$$-T_1 + cT_2 - T_3 = 0$$

$$-T_2 + cT_3 - T_4 = 0$$

$$-2T_3 + cT_4 = 0$$



$$T_1^{(p+1)} = \frac{1 + T_2^{(p)}}{c}$$

$$T_2^{(p+1)} = \frac{T_1^{(p)} + T_3^{(p)}}{c}$$

$$T_3^{(p+1)} = \frac{T_2^{(p)} + T_4^{(p)}}{c}$$

$$T_4^{(p+1)} = \frac{2}{c}T_3^{(p)}$$

	T_1	T_2	T_3	T_4
0	0	0	0	0
1	$1/c$	0	0	0
2	$1/c$	$1/c^2$	0	0
3	$(c+1)/c^3$	$1/c^2$	$1/c^3$	0

The Jacobi method exhibits slow converge properties. It is very clear that in our problem are required 4 iterations to change, for the first time, T value from its initial guess in each node.

Finite difference solution of fin equation – Gauss-Siedel method

$$\begin{aligned}
 cT_1 - T_2 &= 1 \\
 -T_1 + cT_2 - T_3 &= 0 \\
 -T_2 + cT_3 - T_4 &= 0 \\
 -2T_3 + cT_4 &= 0
 \end{aligned}
 \quad \Rightarrow \quad
 \begin{aligned}
 T_1^{(p+1)} &= \frac{1 + T_2^{(p)}}{c} \\
 T_2^{(p+1)} &= \frac{T_1^{(p+1)} + T_3^{(p)}}{c} \\
 T_3^{(p+1)} &= \frac{T_2^{(p+1)} + T_4^{(p)}}{c} \\
 T_4^{(p+1)} &= \frac{2}{c} T_3^{(p+1)}
 \end{aligned}$$

	T_1	T_2	T_3	T_4
0	0	0	0	0
1	$1/c$	$1/c^2$	$1/c^3$	$1/c^4$
2
3

The Gauss-Siedel method is certainly faster than Jacobi one. After only 1 iteration, T values are different from their initial guess in each node.

Ex2_FD - Finite difference solution of fin equation - Jacobi vs. Gauss-Siedel

Matlab code

```
function x2=jacobi(A,b,N,n_max)
%
x1 = zeros(N,1); % guess vector
x2 = zeros(N,1); % updated vector

for k=1:n_max
%
app = 1/A(1,1);
x2(1) = app*(b(1) - A(1,2:N)*x1(2:N));
%
for i=2:N
app = 1/A(i,i);
ri1 = A(i,1:(i-1))*x1(1:(i-1));
ri2 = A(i,(i+1):N)*x1((i+1):N);
x2(i) = app*(b(i) - ri1 - ri2);
end
%
r=b- A*x2;
rMag = sqrt(sum(r.*r))/sqrt(sum(b.*b));
%
x1=x2;
%
if ( rMag <= 1e-12)
break
end

fprintf('%s\n',['It. ', num2str(k), ' residual
', num2str(rMag)]);
end
```

$$\mathbf{Ax} = \mathbf{b} \Rightarrow a_{ij}x_j = b_i$$

$$a_{ii}x_i + \sum_{i \neq j} a_{ij}x_j = b_i$$



$$x_i^{(k+1)} = \frac{1}{a_{ii}} \left(b_i - \sum_{i \neq j} a_{ij}x_j^{(k)} \right)$$

Ex2_FD - Finite difference solution of fin equation - Jacobi vs. Gauss-Siedel

Matlab code

```
function x2=jacobi(A,b,N,n_max)
%
x1 = zeros(N,1); % guess vector
x2 = zeros(N,1); % updated vector

for k=1:n_max
%
app = 1/A(1,1);
x2(1) = app*(b(1) - A(1,2:N)*x1(2:N));
%
for i=2:N
app = 1/A(i,i);
ri1 = A(i,1:(i-1))*x1(1:(i-1));
ri2 = A(i,(i+1):N)*x1((i+1):N);
x2(i) = app*(b(i) - ri1 - ri2);
end
%
r=b- A*x2;
rMag = sqrt(sum(r.*r))/sqrt(sum(b.*b));
%
x1=x2;
%
if ( rMag <= 1e-12)
break
end

fprintf('%s\n',['It. ', num2str(k), ' residual
', num2str(rMag)]);
end
```

$$x_i^{(k+1)} = \frac{1}{a_{ii}} \left(b_i - \sum_{i \neq j} a_{ij} x_j^{(k)} \right)$$

$$r = \frac{\| \mathbf{b} - \mathbf{A}\mathbf{x} \|}{\| \mathbf{b} \|}$$

$$\| \mathbf{x} \| = \sqrt{\sum_{k=1}^N x_k^2}$$

Ex2_FD - Finite difference solution of fin equation - Jacobi vs. Gauss-Siedel

Matlab code

```
function x2=gs(A,b,N,n_max)
%
x1 = zeros(N,1); % guess vector
x2 = zeros(N,1); % updated vector

for k=1:n_max
%
    app = 1/A(1,1);
    x2(1) = app*(b(1) - A(1,2:N)*x1(2:N));
%
    for i=2:N
        app = 1/A(i,i);
        ri1 = A(i,1:(i-1))*x2(1:(i-1));
        ri2 = A(i,(i+1):N)*x1((i+1):N);
        x2(i) = app*(b(i) - ri1 - ri2);
    end
%
r=b- A*x2;
rMag = sqrt(sum(r.*r))/sqrt(sum(b.*b));
%
x1=x2;
%
if ( rMag <= 1e-12)
    break
end

fprintf('%s\n',['It. ', num2str(k), ' residual ', num2str(rMag)]);
end
```

$$x_i^{(k+1)} = \frac{1}{a_{ii}} (b_i - r_1 - r_2)$$

$$r_1 = \sum_{i < j} a_{ij} x_j^{(k+1)}$$

$$r_2 = \sum_{i > j} a_{ij} x_j^{(k)}$$

Ex2_FD - Finite difference solution of fin equation - Jacobi vs. Gauss-Siedel

Matlab code

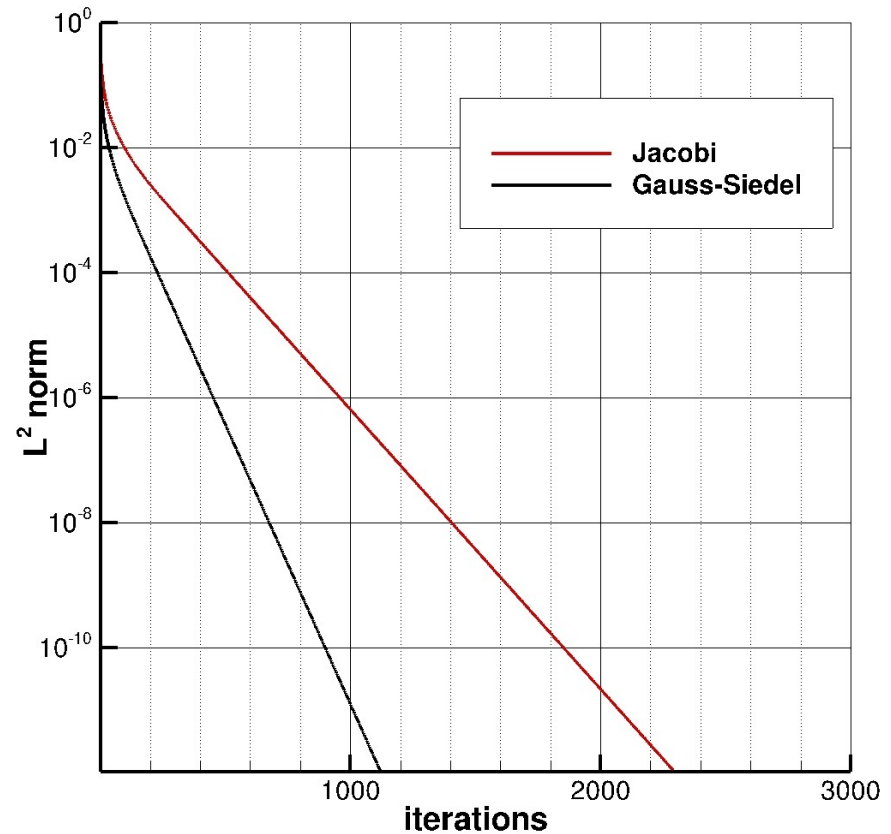
```
A      = zeros(N,N);

A(1,1)  = c;
A(1,2)  = -1;

for i=2:(N-1)
    A(i,i-1) = -1;
    A(i,i)   = c;
    A(i,i+1) = -1;
end

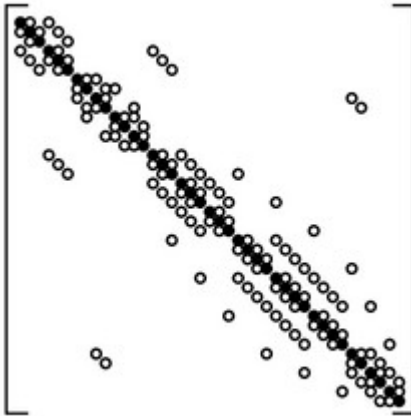
A(N,N-1) = -2;
A(N,N)   = c;
%
%-----
%
b=zeros(N,1);
b(1) = 1;
%
% Linear system solution
%
%x=A\b;
%x= gs(A,b,N,30000);
x = jacobi(A,b,N,30000);

%-----
%Boundary value (x=0)
x =[1;x];
```



Matrix sparsity

In the previous exercise we solve a linear system with a matrix having rank equal to 30. However, the matrix has only **88 non-null entries of its overall 900**. The matrix **A** is **sparse**, *i.e.* the majority of its coefficients are zero.



The sparsity is due to each node interacts only with adjacent ones. For example, also a more complex stencil with 5 nodes, produces up to 5 coefficients per matrix row, with one diagonal coefficient corresponding to a particular cell and 4 off-diagonal coefficients for the neighbour cells.

Typically current academic and industrial applications require huge meshes to obtain reliable solutions. Nowadays, $N=10^6$ is the common scale. It is easy to understand that for similar problems the non-zero terms are really lower than zeros ones. For efficiency reasons, zero coefficients are not stored in the computer's memory. Instead the storage is based on an array of non-zero coefficients and addressing arrays of the corresponding row and column indices for each coefficient.

Ex3_FD - Finite difference solution of fin equation - Thomas algorithm

The Thomas algorithm is a variant of Gauss elimination method which is specifically developed for tri-diagonal matrices. It is also known also as Tri-Diagonal Matrix Algorithm (TDMA). It reduces the computational costs from $O(N^3)$ to $O(N)$.

$$\begin{aligned} b_1 x_1 + c_1 x_2 &= d_1 \\ a_i x_{i-1} + b_i x_i + c_i x_{i+1} &= d_i, \\ a_n x_{n-1} + b_n x_n &= d_n. \end{aligned}$$

$$(\text{coeff. } x_{i-1} \text{ of eq } i-1)(\text{eq } i) - (\text{coeff. } x_i \text{ of eq } i)(\text{eq } i-1)$$

$$\begin{aligned} b_1 x_1 + c_1 x_2 &= d_1 \\ a_2 x_1 + b_2 x_2 + c_2 x_3 &= d_2 \end{aligned} \quad \Rightarrow \quad \begin{aligned} a_2 b_1 x_1 + a_2 c_1 x_2 &= a_2 d_1 \\ b_1 a_2 x_1 + b_1 b_2 x_2 + b_1 c_2 x_3 &= b_1 d_2 \end{aligned}$$

$$(b_1 b_2 - a_2 c_1) x_2 + b_1 c_2 x_3 = b_1 d_2 - a_2 d_1$$

Using this strategy (**forward substitution**) the coeffs become more complex. However, we cancel all the contributions deriving from sub-diagonals terms.

Ex3_FD - Finite difference solution of fin equation - Thomas algorithm

$$\begin{array}{rcl} a_{N-1}x_{N-2} + b_{N-1}x_{N-1} + c_{N-1}x_N = d_{N-1} & \times & a_N \\ a_Nx_{N-1} + b_Nx_N = d_N & \times & b_{N-1} \end{array} \quad \curvearrowright$$

$$(b_{N-1}b_N - c_{N-1}a_N)x_N = b_{N-1}d_N - a_Nb_{N-1}$$

The forward substitution procedure allows to obtain the **last equation** with **only one unknown**. This condition is highly appealing since in this way with a **backward substitution** is possible the finalize the solution of the overall linear system.

Basically, forward substitution is used to compute the coeffs of the new system without sub-diagonal terms. Backward substitution is used to compute the final solution.

$$c'_i = \begin{cases} \frac{c_i}{b_i}, & i = 1, \\ \frac{c_i}{b_i - a_i c'_{i-1}}, & i = 2, 3, \dots, n-1 \end{cases}$$

$$d'_i = \begin{cases} \frac{d_i}{b_i}, & i = 1, \\ \frac{d_i - a_i d'_{i-1}}{b_i - a_i c'_{i-1}}, & i = 2, 3, \dots, n. \end{cases}$$

forward
substitution

$$x_n = d'_n,$$

$$x_i = d'_i - c'_i x_{i+1}, \quad i = n-1, n-2, \dots, 1.$$

backward
substitution

Ex3_FD - Finite difference solution of fin equation - Thomas algorithm

Matlab code

```
% Matrix - vectors;
%
al      = -ones(N-1,1);
al(end) = -2;

ad      = ones(N,1)*c;
au      = -ones(N-1,1);

%-----
%
b=zeros(N,1);
b(1) = 1;

% Linear system solution
%
%x =A\b;
x = tridiagonal_vector(al,ad,au,b);

%
%-----
```

$$A = \begin{pmatrix} C & -1 & 0 & 0 & 0 \\ -1 & C & -1 & 0 & 0 \\ 0 & -1 & C & -1 & 0 \\ 0 & 0 & -1 & C & -1 \\ 0 & 0 & 0 & -2 & C \end{pmatrix} \begin{matrix} \\ \\ \\ a_u \\ a_d \end{matrix}$$

a_l a_d

The first three entries of the *tridiagonal_vector* are stored as showed in the above sketch.

In this specific case we rely on the so-called **opportunistic code-reuse strategy**. In other words, we use existing software to build new software.

Diagonal dominance

The behaviour of this problem indicates a convergence condition: the magnitude of the diagonal coefficient in each matrix row must be greater than or equal to the sum of the magnitudes of the other coefficients in the row:

$$|a_{i,i}| \geq \sum_{i \neq j} |a_{i,j}|$$

where the $>$ condition must be satisfied for at least one row. *This condition is mathematically “sufficient” which means that convergence may occur when the condition is not met.*

$$A = \begin{pmatrix} C & -1 & 0 & 0 & 0 \\ -1 & C & -1 & 0 & 0 \\ 0 & -1 & C & -1 & 0 \\ 0 & 0 & -1 & C & -1 \\ 0 & 0 & 0 & -2 & C \end{pmatrix} \quad c = [2 + (mL \Delta x)^2]$$