

Calling Convention

Abhijit A M

The need for calling convention

An essential task of the compiler

Generates object code (file) for given source code (file)

Processors provide simple features

Registers, machine instructions (add, mov, jmp, call, etc.), imp registers like stack-pointer, etc;
ability to do byte/word sized operations

No notion of data types, functions, variables, etc.

But languages like C provide high level features

Data types, variables, functions, recursion, etc

Compiler needs to map the features of C into processor's features, and then generate machine code

In reality, the language designers design a language feature only after answering the question of conversion into machine code

The need for calling convention

Examples of some of the challenges before the compiler

“call” + “ret” does not make a C-function call!

A “call” instruction in processor simply does this

Pushes IP(that is PC) on stack + Jumps to given address

This is not like calling a C-function !

Unsolved problem: How to handle parameters, return value?

Processor does not understand data types!

Although it has instructions for byte, word sized data and can differentiate between integers and reals (mov, movw, addl, addf, etc.)

Compiler and Machine code generation

Example, code inside a function

```
int a, b, c;
```

```
c = a + b;
```

What kind of code is generated by compiler for this?

```
sub 12, <esp> #normally local variables are located on stack, make space
```

```
mov <location of a in memory>, r1 #location is on stack, e.g. -4(esp)
```

```
mov <location of b in memory>, r2
```

```
add r1, r2 # result in r1
```

```
mov r1, <location of c in memory>
```

Compiler and Machine code generation

Across function calls

```
int f(int m, n) {  
    int x = m, y = n;  
    return g(x, y);  
}  
  
int x(int a) {  
    return g (a, a+ 1);  
}  
  
int g(int p, int q) {  
    p = p * q + p;  
    return p  
}
```

**g() may be called from f()
or from x()**

**Sequence of function calls
can NOT be predicted by
compiler**

**Compiler has to generate
machine code for each
function assuming
nothing about the caller**

Compiler and Machine code generation

Machine code generation for functions

Mapping C language features to existing machine code instructions.

Typical examples

`a = 100 ; ==> mov instruction`

`a = b + c; ==> mov, add instructions`

`while(a < 5) { j++; } ==> mov, cmp, jlt, add, etc. Instruction`

Where are the local variables in memory?

The only way to store them is on a stack.

Why?

Function calls

LIFO

Last in First Out

Must need a “stack” like feature to implement them

Processor Stack

Processors provide a stack pointer

%esp on x86

Instructions like push and pop are provided by hardware

They automatically increment/decrement the stack pointer. On x86 stack grows downwards (subtract from esp!)

Unlike a “stack data type” data structure, this “stack” is simply implemented with only the esp (as the “top”). The entire memory can be treated as the “array”.

Function calls

System stack, compilers, Languages

Compilers generate machine code with the 'esp'. The pointer is initialized to a proper value at the time of fork-exec by the OS for each process.

Languages like C which provide for function calls, and recursion also assume that they will run on processors with a stack support in hardware

Convention needed

How to use the stack for effective implementation of function calls ?

What goes on stack?

Local variables

Function Parameters

Return address of instruction which called a function !

Activation Record

Local Vars + parameters + return address

When functions call each other

One activation record is built on stack for each function call

On function return, the record is destroyed

On x86

ebp and esp pointers are used to denote the activation record.

How? We will see soon. You may start exploring with “gcc -S” output assembly code.

X86 instructions

leave

Equivalent to

```
mov  %ebp, %esp    # esp =  
ebp  
pop  %ebp
```

ret

Equivalent to

```
pop %ecx  
Jmp %ecx
```

call x

Equivalent to

```
push %eip  
jmp x
```

X86 instructions

endbr64

Normally a NOP

Let's see some examples now

Let's compile using

`gcc -S`

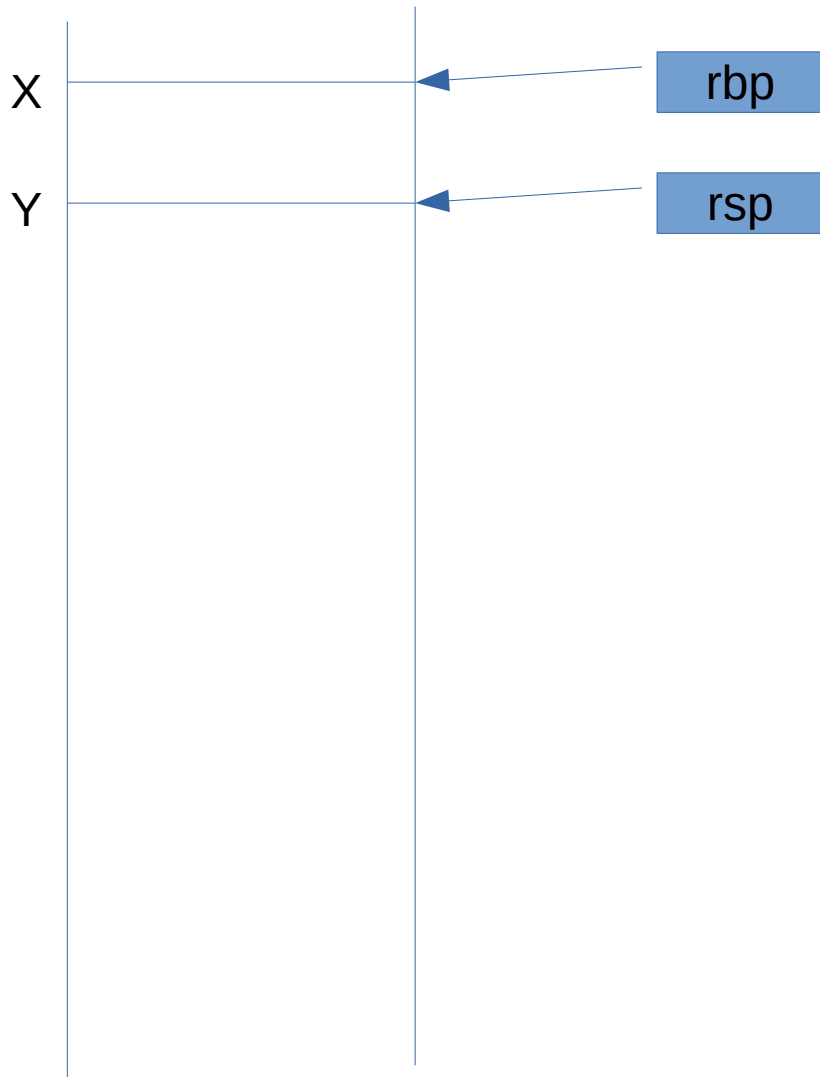
See code and understand

simple.c and simple.s

```
int f(int x) {  
    int y;  
    y = x + 3;  
    return y;  
}  
int main() {  
    int a = 20, b = 30;  
    b = f(a);  
    return b;  
}
```

```
main:  
    endbr64  
    pushq   %rbp  
    movq    %rsp, %rbp  
    subq    $16, %rsp  
    movl    $20, -8(%rbp)  
    movl    $30, -4(%rbp)  
    movl    -8(%rbp), %eax  
    movl    %eax, %edi  
    call    f  
    movl    %eax, -4(%rbp)  
    movl    -4(%rbp), %eax  
    leave  
    ret
```

```
f:  
    endbr64  
    pushq   %rbp  
    movq    %rsp, %rbp  
    movl    %edi, -20(%rbp)  
    movl    -20(%rbp), %eax  
    addl    $3, %eax  
    movl    %eax, -4(%rbp)  
    movl    -4(%rbp), %eax  
    popq    %rbp  
    ret
```



main:

endbr64

pushq %rbp

movq %rsp, %rbp

subq \$16, %rsp

movl \$20, -8(%rbp)

movl \$30, -4(%rbp)

movl -8(%rbp), %eax

movl %eax, %edi

call f

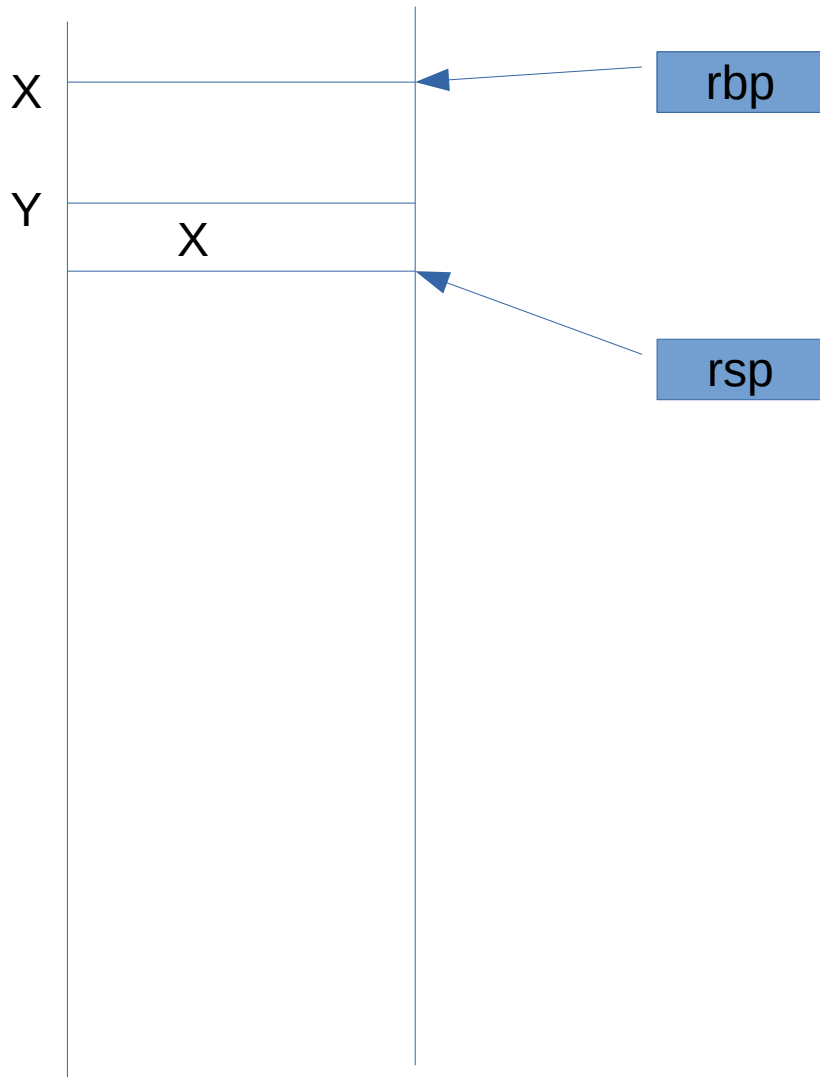
movl %eax, -4(%rbp)

movl -4(%rbp), %eax

leave

ret

```
int main() {  
    int a = 20, b = 30;  
    b = f(a);  
    return b;  
}
```



main:

endbr64

pushq %rbp

movq %rsp, %rbp

subq \$16, %rsp

movl \$20, -8(%rbp)

movl \$30, -4(%rbp)

movl -8(%rbp), %eax

movl %eax, %edi

call f

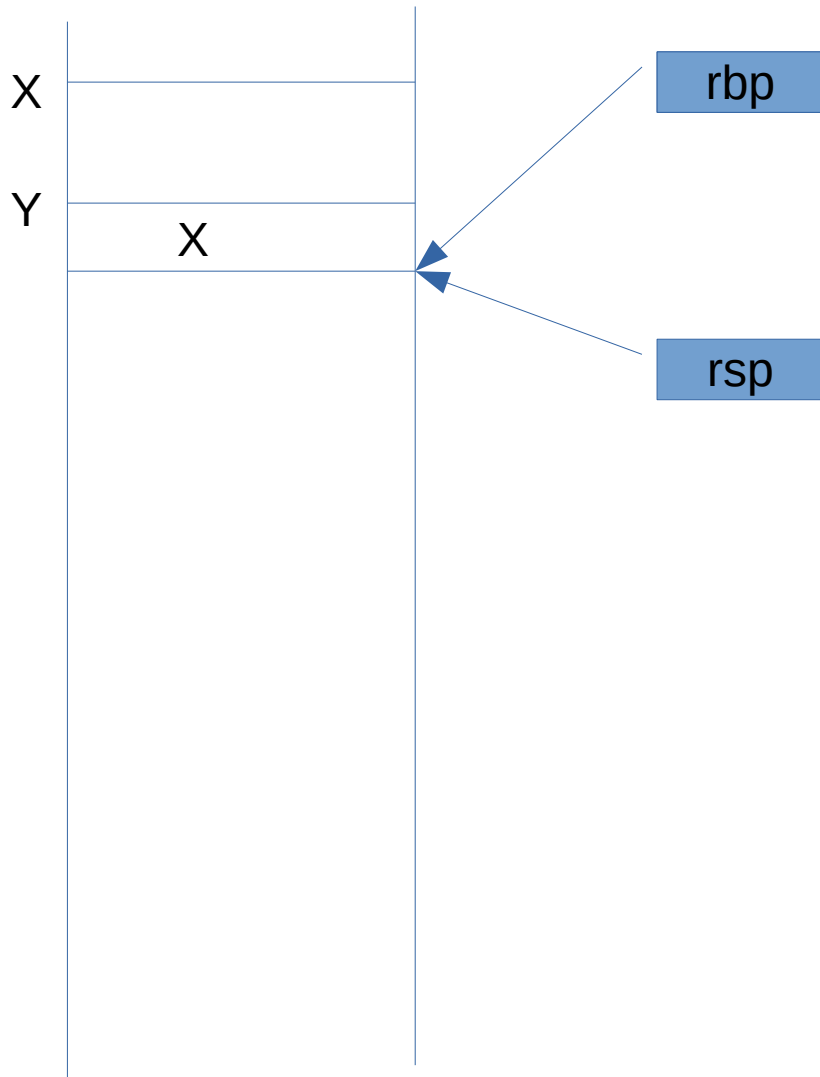
movl %eax, -4(%rbp)

movl -4(%rbp), %eax

leave

ret

```
int main() {  
    int a = 20, b = 30;  
    b = f(a);  
    return b;  
}
```



main:

endbr64

pushq %rbp

movq %rsp, %rbp

subq \$16, %rsp

movl \$20, -8(%rbp)

movl \$30, -4(%rbp)

movl -8(%rbp), %eax

movl %eax, %edi

call f

movl %eax, -4(%rbp)

movl -4(%rbp), %eax

leave

ret

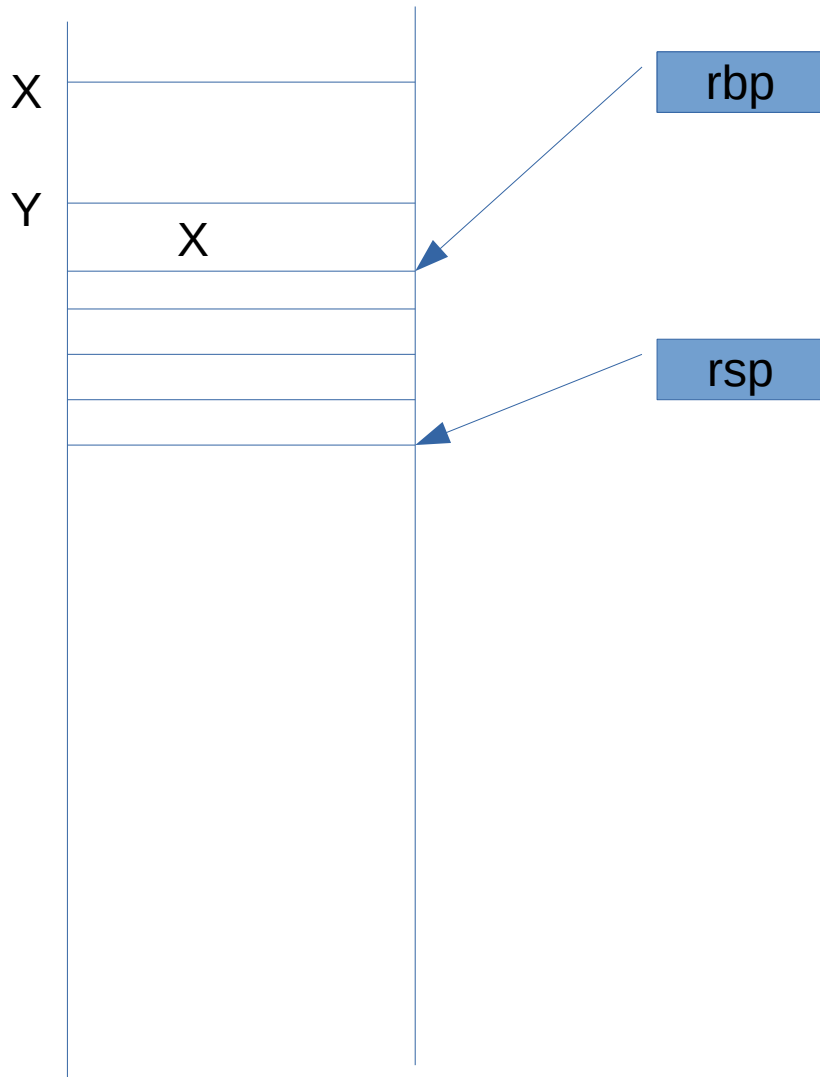
int main() {

int a = 20, b = 30;

b = f(a);

return b;

}



main:

endbr64

pushq %rbp

movq %rsp, %rbp

subq \$16, %rsp

movl \$20, -8(%rbp)

movl \$30, -4(%rbp)

movl -8(%rbp), %eax

movl %eax, %edi

call f

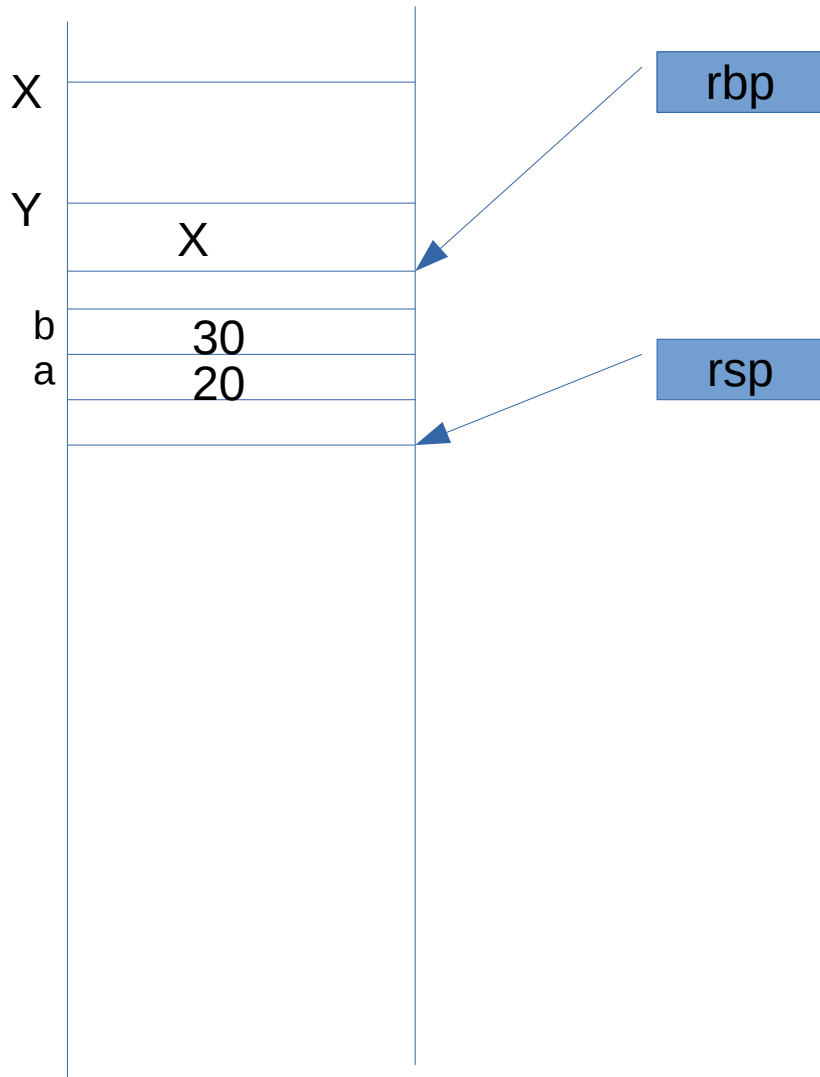
movl %eax, -4(%rbp)

movl -4(%rbp), %eax

leave

ret

```
int main() {  
    int a = 20, b = 30;  
    b = f(a);  
    return b;  
}
```



main:

endbr64

pushq %rbp

movq %rsp, %rbp

subq \$16, %rsp

movl \$20, -8(%rbp)

movl \$30, -4(%rbp)

movl -8(%rbp), %eax

movl %eax, %edi

call f

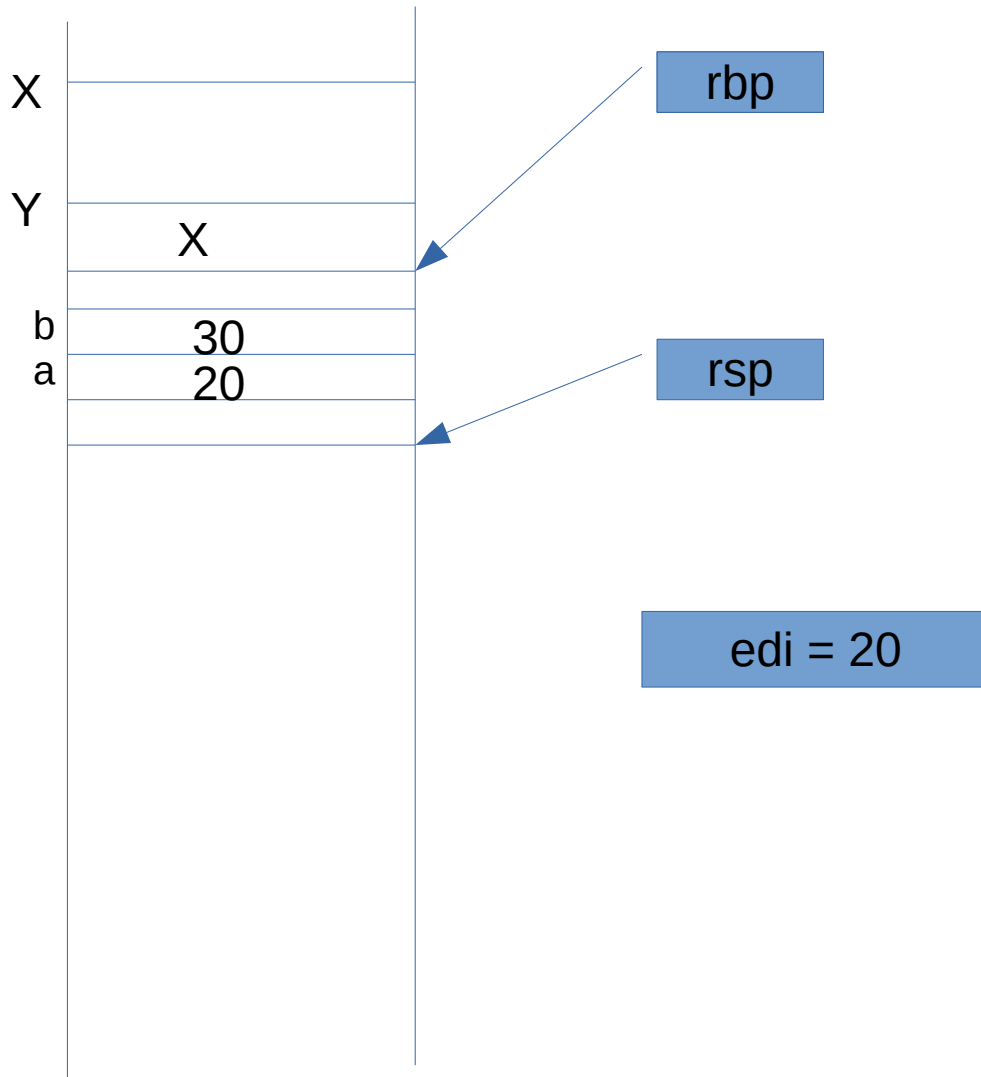
movl %eax, -4(%rbp)

movl -4(%rbp), %eax

leave

ret

```
int main() {  
    int a = 20, b = 30;  
    b = f(a);  
    return b;  
}
```



main:

endbr64

pushq %rbp

movq %rsp, %rbp

subq \$16, %rsp

movl \$20, -8(%rbp)

movl \$30, -4(%rbp)

movl -8(%rbp), %eax

movl %eax, %edi

call f

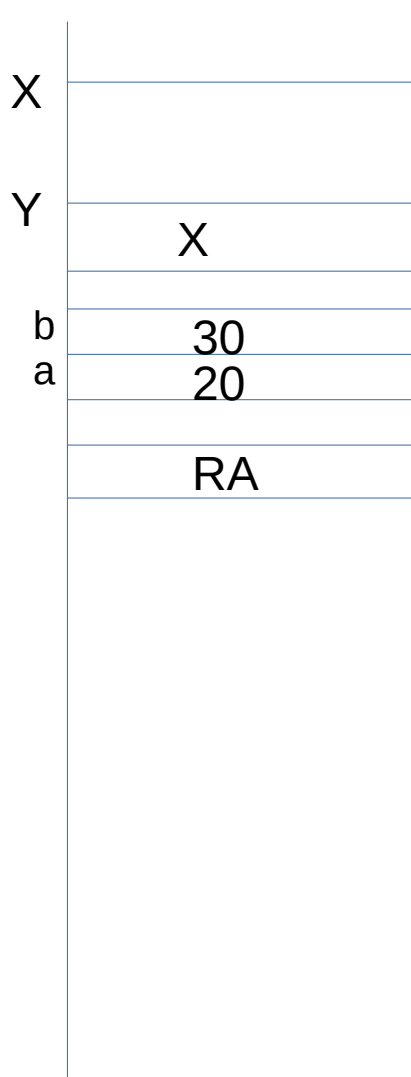
movl %eax, -4(%rbp)

movl -4(%rbp), %eax

leave

ret

```
int main() {  
    int a = 20, b = 30;  
    b = f(a);  
    return b;  
}
```



rbp

rsp

main:
endbr64

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pushq %rbp
movq  %rsp, %rbp
subq  $16, %rsp
movl  $20, -8(%rbp)
movl  $30, -4(%rbp)
movl  -8(%rbp), %eax
movl  %eax, %edi
call  f
```

RA:

```
movl  %eax, -4(%rbp)
movl  -4(%rbp), %eax
Leave
ret
```

f:

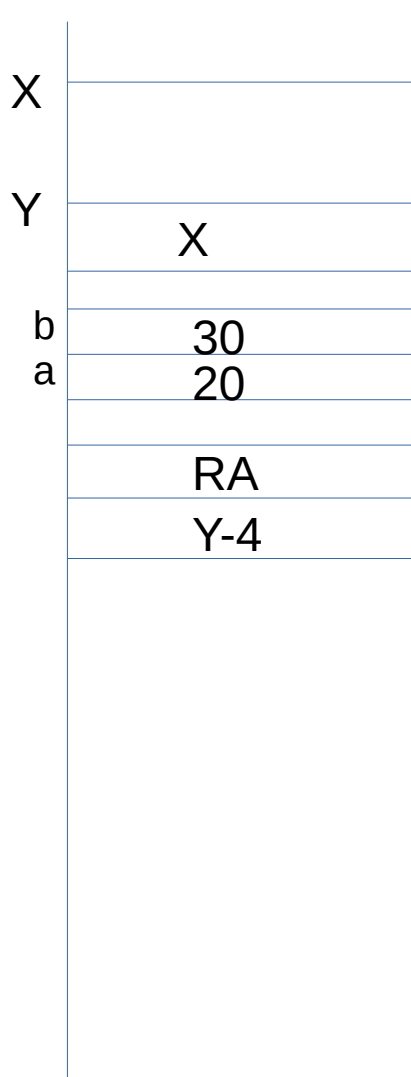
endbr64

```
pushq %rbp
movq  %rsp, %rbp
movl  %edi, -20(%rbp)
movl  -20(%rbp), %eax
addl  $3, %eax
movl  %eax, -4(%rbp)
movl  -4(%rbp), %eax
popq  %rbp
ret
```

edi = 20

```
int f(int x) {
    int y;
    y = x + 3;
    return y;
}
```

```
int main() {
    int a = 20, b = 30;
    b = f(a);
    return b;
}
```



rbp

rsp

main:
endbr64

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pushq %rbp
movq %rsp, %rbp
subq $16, %rsp
movl $20, -8(%rbp)
movl $30, -4(%rbp)
movl -8(%rbp), %eax
movl %eax, %edi
call f
```

RA:

```
movl %eax, -4(%rbp)
movl -4(%rbp), %eax
Leave
ret
```

f:

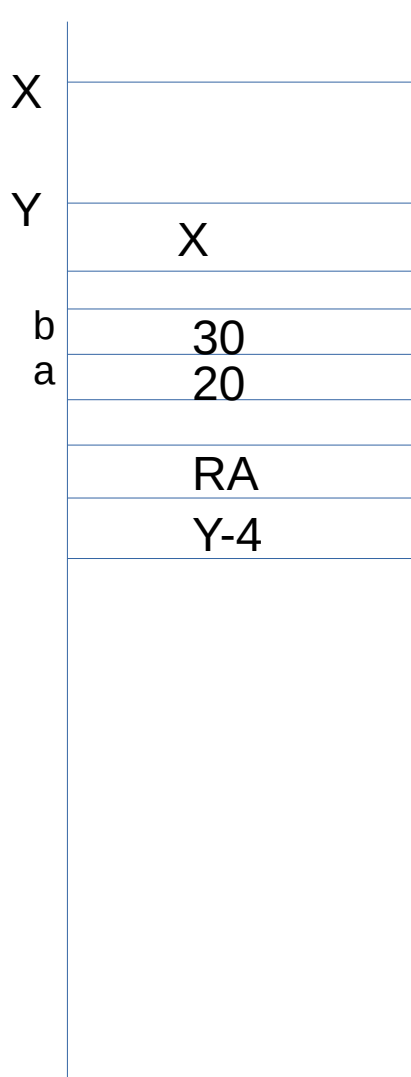
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movl %edi, -20(%rbp)
movl -20(%rbp), %eax
addl $3, %eax
movl %eax, -4(%rbp)
movl -4(%rbp), %eax
popq %rbp
ret
```

edi = 20

```
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    int y;
    y = x + 3;
    return y;
}
```

```
int main() {
    int a = 20, b = 30;
    b = f(a);
    return b;
}
```



rbp

rsp

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movl  -8(%rbp), %eax
movl  %eax, %edi
call  f
```

RA:

```
movl  %eax, -4(%rbp)
movl  -4(%rbp), %eax
Leave
ret
```

f:

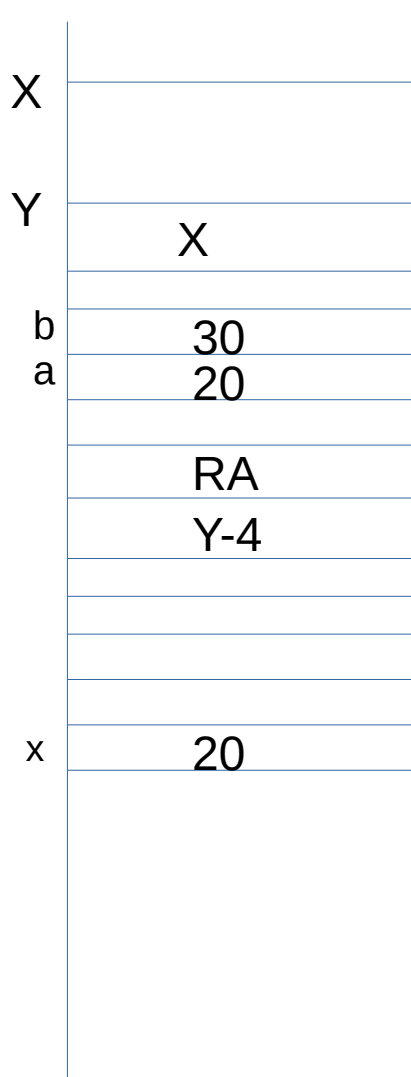
endbr64

```
pushq %rbp
movq  %rsp, %rbp
movl  %edi, -20(%rbp)
movl  -20(%rbp), %eax
addl  $3, %eax
movl  %eax, -4(%rbp)
movl  -4(%rbp), %eax
popq  %rbp
ret
```

edi = 20

```
int f(int x) {
    int y;
    y = x + 3;
    return y;
}
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```
int main() {
    int a = 20, b = 30;
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main:
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pushq %rbp
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movl $20, -8(%rbp)
movl $30, -4(%rbp)
movl -8(%rbp), %eax
movl %eax, %edi
call f
```

RA:

```
movl %eax, -4(%rbp)
movl -4(%rbp), %eax
Leave
ret
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```
int f(int x) {
    int y;
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    return y;
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```

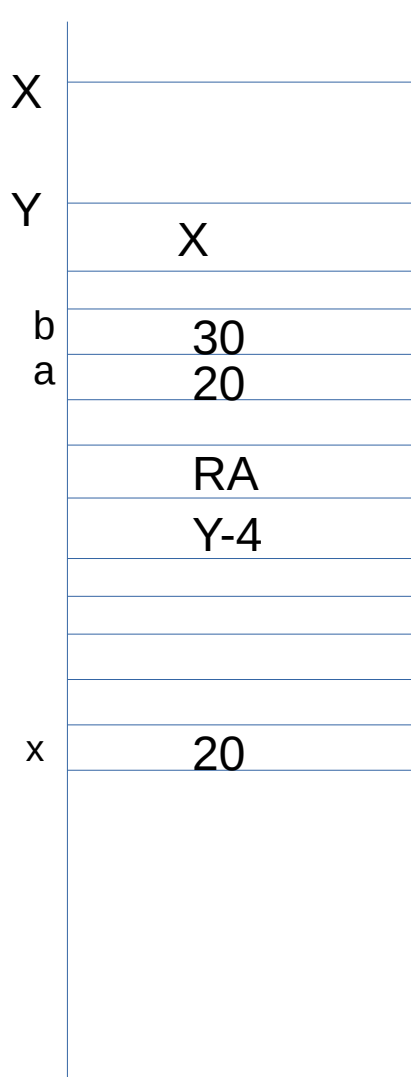
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movl %eax, -4(%rbp)
movl -4(%rbp), %eax
popq %rbp
ret
```

edi = 20

```
int main() {
    int a = 20, b = 30;
    b = f(a);
    return b;
}
```



main:
endbr64

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pushq %rbp
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movl -8(%rbp), %eax
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RA:

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movl %eax, -4(%rbp)
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Leave
ret
```

```
int f(int x) {
    int y;
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```

f:

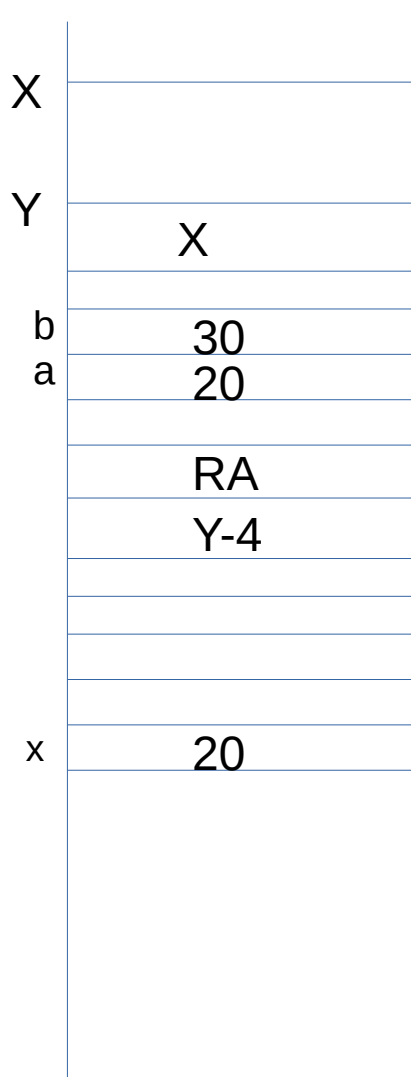
endbr64

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addl $3, %eax
movl %eax, -4(%rbp)
movl -4(%rbp), %eax
popq %rbp
ret
```

edi = 20

eax = 20

```
int main() {
    int a = 20, b = 30;
    b = f(a);
    return b;
}
```

main:
endbr64

```
pushq %rbp
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subq $16, %rsp
movl $20, -8(%rbp)
movl $30, -4(%rbp)
movl -8(%rbp), %eax
movl %eax, %edi
call f
```

RA:

```
movl %eax, -4(%rbp)
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Leave
ret
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```
int f(int x) {
    int y;
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```

f:

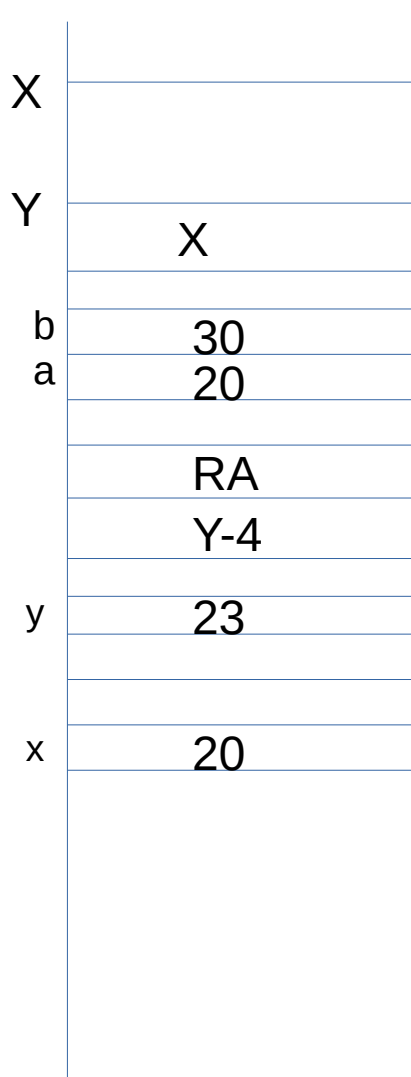
endbr64

```
pushq %rbp
movq %rsp, %rbp
movl %edi, -20(%rbp)
movl -20(%rbp), %eax
addl $3, %eax
movl %eax, -4(%rbp)
movl -4(%rbp), %eax
popq %rbp
ret
```

edi = 20

eax = 23

```
int main() {
    int a = 20, b = 30;
    b = f(a);
    return b;
}
```



rbp

rsp

main:
endbr64

```
pushq %rbp
movq %rsp, %rbp
subq $16, %rsp
movl $20, -8(%rbp)
movl $30, -4(%rbp)
movl -8(%rbp), %eax
movl %eax, %edi
call f
```

RA:

```
movl %eax, -4(%rbp)
movl -4(%rbp), %eax
Leave
ret
```

```
int f(int x) {
    int y;
    y = x + 3;
    return y;
}
```

```
int main() {
    int a = 20, b = 30;
    b = f(a);
    return b;
}
```

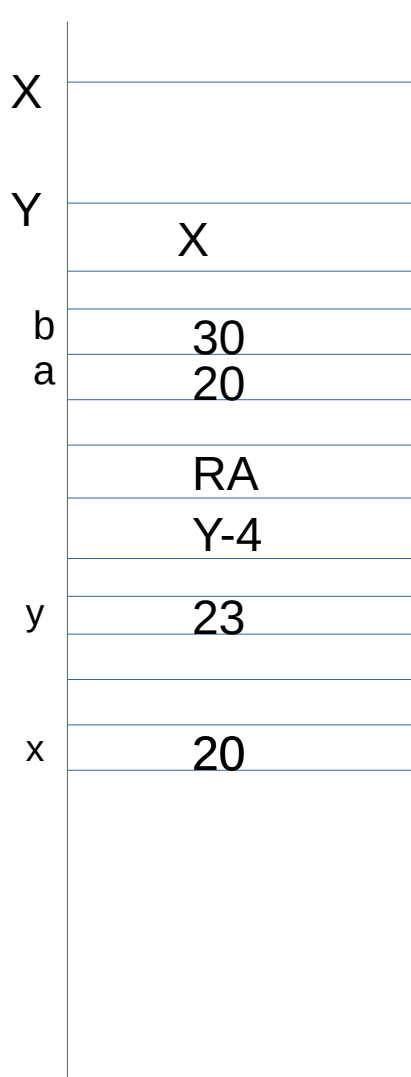
f:

endbr64

```
pushq %rbp
movq %rsp, %rbp
movl %edi, -20(%rbp)
movl -20(%rbp), %eax
addl $3, %eax
movl %eax, -4(%rbp)
movl -4(%rbp), %eax
popq %rbp
ret
```

edi = 20

eax = 23



```
main:
endbr64
```

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pushq %rbp
movq %rsp, %rbp
subq $16, %rsp
movl $20, -8(%rbp)
movl $30, -4(%rbp)
movl -8(%rbp), %eax
movl %eax, %edi
call f
```

RA:

```
movl %eax, -4(%rbp)
movl -4(%rbp), %eax
Leave
ret
```

```
int f(int x) {
    int y;
    y = x + 3;
    return y;
}
```

f:

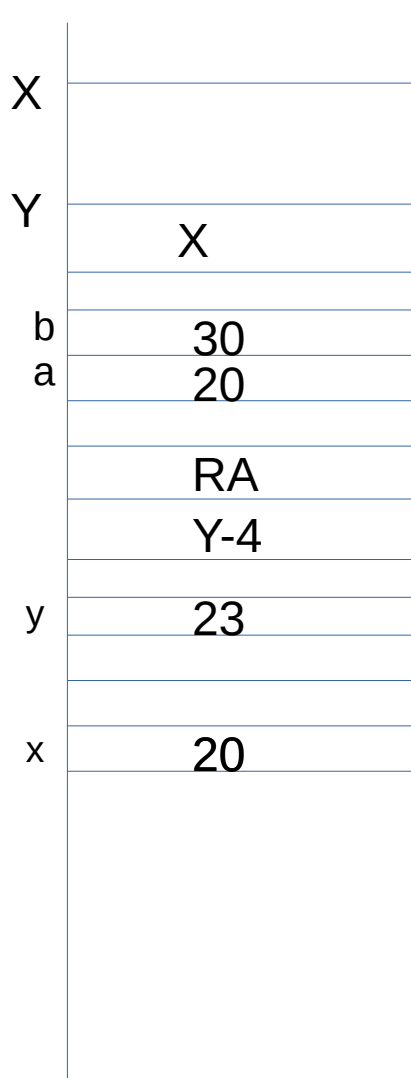
```
endbr64
```

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pushq %rbp
movq %rsp, %rbp
movl %edi, -20(%rbp)
movl -20(%rbp), %eax
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edi = 20

Eax = 23

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int main() {
    int a = 20, b = 30;
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rbp

rsp

main:
endbr64

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pushq %rbp
movq  %rsp, %rbp
subq  $16, %rsp
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movl  $30, -4(%rbp)
movl  -8(%rbp), %eax
movl  %eax, %edi
call  f

```

RA:

```

movl  %eax, -4(%rbp)
movl  -4(%rbp), %eax
Leave
ret

```

```

int f(int x) {
    int y;
    y = x + 3;
    return y;
}

```

```

int main() {
    int a = 20, b = 30;
    b = f(a);
    return b;
}

```

f:

```

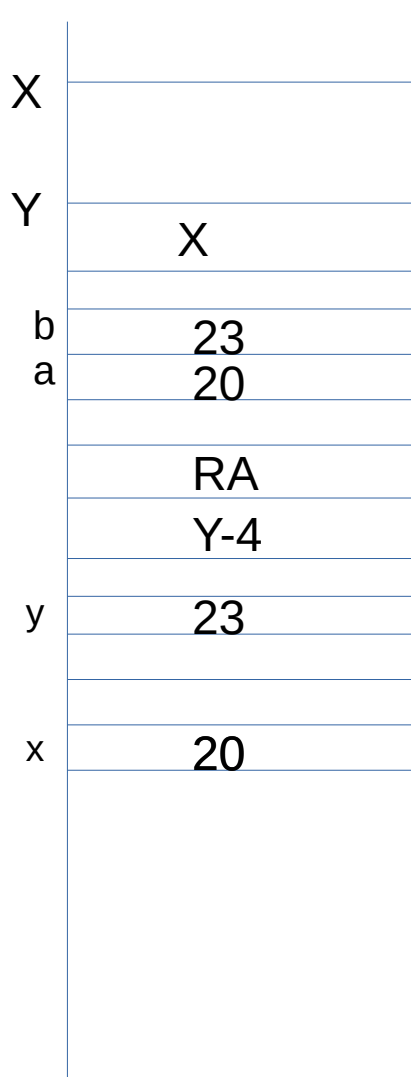
endbr64
pushq %rbp
movq  %rsp, %rbp
movl  %edi, -20(%rbp)
movl  -20(%rbp), %eax
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popq  %rbp
ret

```

edi = 20

eax = 23

eip = RA



rbp

rsp

main:
endbr64

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pushq %rbp
movq  %rsp, %rbp
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RA:
movl  %eax, -4(%rbp)
movl  -4(%rbp), %eax
leave
ret

```

```

int f(int x) {
    int y;
    y = x + 3;
    return y;
}

```

```

int main() {
    int a = 20, b = 30;
    b = f(a);
    return b;
}

```

f:

```

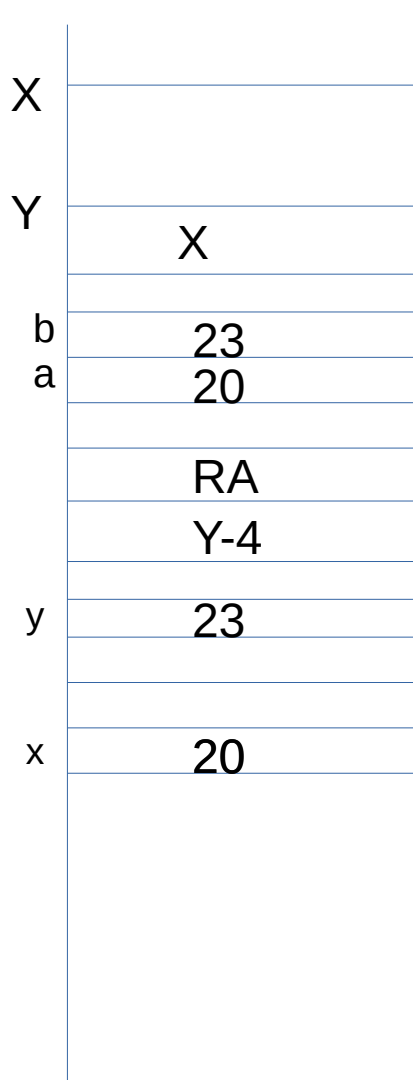
endbr64
pushq %rbp
movq  %rsp, %rbp
movl  %edi, -20(%rbp)
movl  -20(%rbp), %eax
addl  $3, %eax
movl  %eax, -4(%rbp)
movl  -4(%rbp), %eax
popq  %rbp
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```

edi = 20

eax = 23

eip = RA



rbp

rsp

main:
endbr64

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pushq %rbp
movq  %rsp, %rbp
subq  $16, %rsp
movl  $20, -8(%rbp)
movl  $30, -4(%rbp)
movl  -8(%rbp), %eax
movl  %eax, %edi
call  f

```

RA:

```

movl  %eax, -4(%rbp)
movl  -4(%rbp), %eax
leave

```

```

# mov rbp rsp; pop rbp
ret

```

f:

```

endbr64
pushq %rbp
movq  %rsp, %rbp
movl  %edi, -20(%rbp)
movl  -20(%rbp), %eax
addl  $3, %eax
movl  %eax, -4(%rbp)
movl  -4(%rbp), %eax
popq  %rbp
ret

```

edi = 20

eax = 23

eip = RA

```

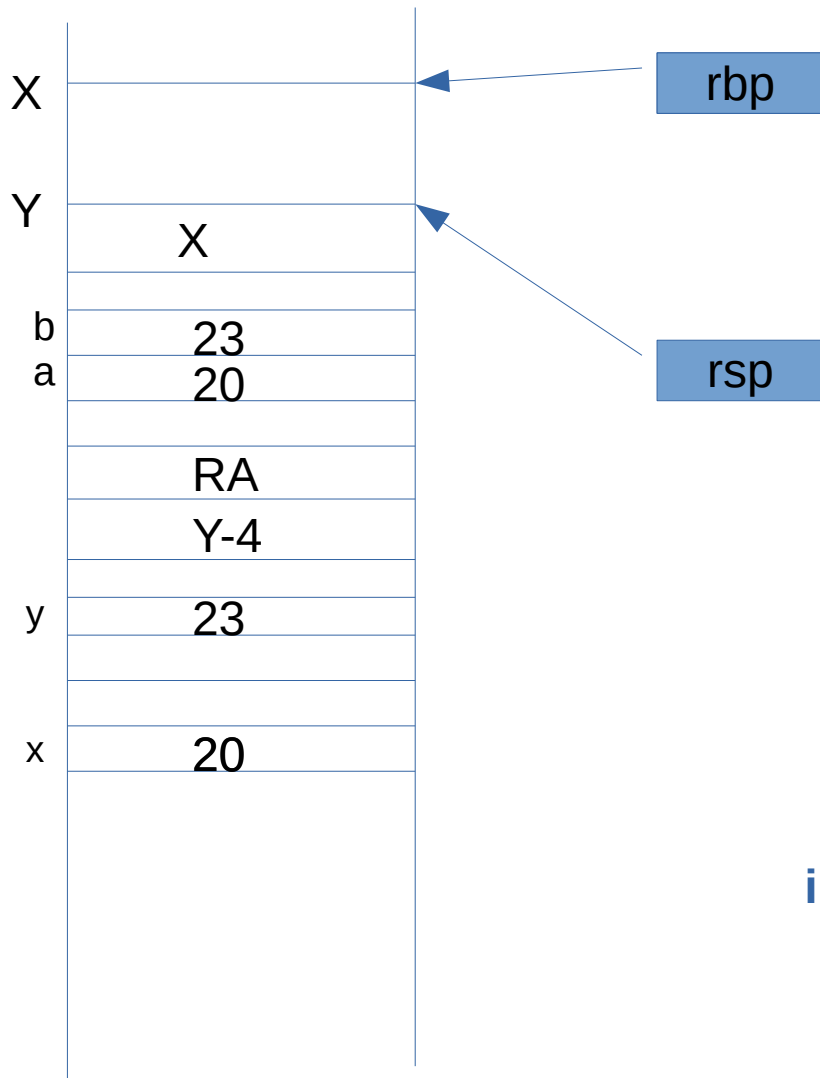
int f(int x) {
    int y;
    y = x + 3;
    return y;
}

```

```

int main() {
    int a = 20, b = 30;
    b = f(a);
    return b;
}

```



main:
endbr64

```

pushq %rbp
movq  %rsp, %rbp
subq  $16, %rsp
movl  $20, -8(%rbp)
movl  $30, -4(%rbp)
movl  -8(%rbp), %eax
movl  %eax, %edi
call  f
RA:
movl  %eax, -4(%rbp)
movl  -4(%rbp), %eax
leave
# mov ebp esp; pop ebp
ret

```

f:

```

endbr64
pushq %rbp
movq  %rsp, %rbp
movl  %edi, -20(%rbp)
movl  -20(%rbp), %eax
addl  $3, %eax
movl  %eax, -4(%rbp)
movl  -4(%rbp), %eax
popq  %rbp
ret

```

edi = 20

eax = 23

eip = RA

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int f(int x) {
    int y;
    y = x + 3;
    return y;
}

```

```

int main() {
    int a = 20, b = 30;
    b = f(a);
    return b;
}

```

Further on calling convention

This was a simple program

The parameter was passed in a register!

What if there were many parameters?

CPUs have different numbers of registers.

More parameters, more functions demand a more sophisticated convention

May be slightly different on different processors, or 32-bit, 64-bit variants also.

Caller save and Callee save registers

Local variables

- Are visible only within the function

- Recursion: different copies of variables

- Stored on “stack”

Registers

- Are only one copy

- Are within the CPU

Local Variables & Registers conflict

- Compiler's dilemma: While generating code for a function, which registers to use?

- The register might have been in use in earlier function call

Caller save and Callee save registers

Caller Save registers

Which registers need to be saved by caller function . They can be used by the callee function!

The caller function will push them (if already in use, otherwise no need) on the stack

Callee save registers

Will be pushed on to the stack by called (callee) function

How to return values?

On the stack itself – then caller will have to pop

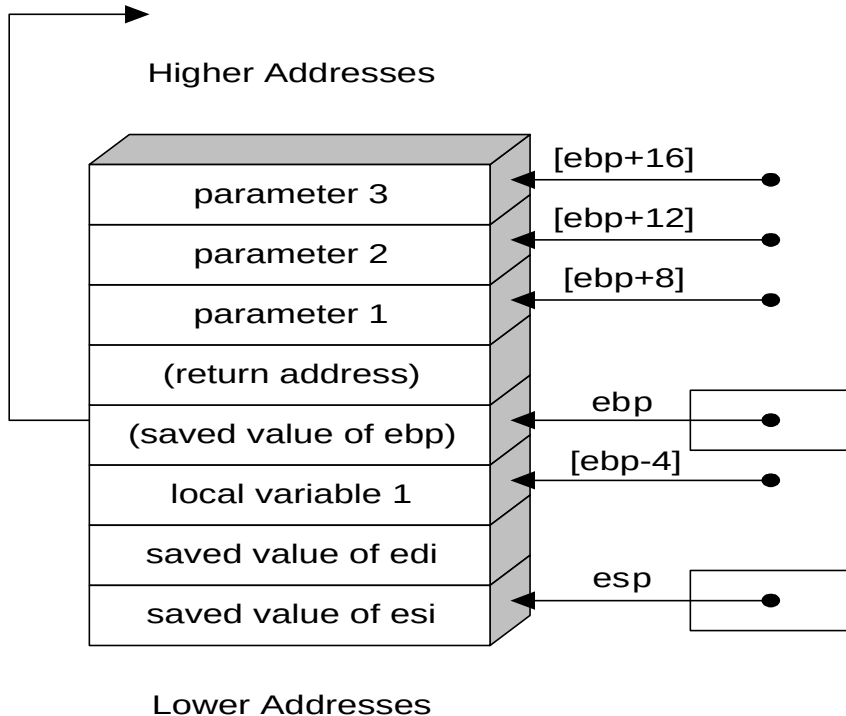
In a register, e.g. eax

X86 convention - caller, callee saved 32 bit

The caller-saved registers are EAX, ECX, EDX.

The callee-saved registers are EBX, EDI, and ESI

Activation record looks like this



F() called g()

Parameters-i refers to parameters passed by f() to g()

Local variable is a variable in g()

Return address is the location in f() where call should go back

X86 caller and callee rules(32 bit)

Caller rules on call

Push caller saved registers on stack

Push parameters on the stack – in reverse order. Why?

Subtract esp, copy data

call f() // push + jmp

Caller rules on return

return value is in eax

remove parameters from stack : Add to esp.

Restore caller saved registers (if any)

X86 caller and callee rules

Callee rules on call

1) `push ebp`

`mov ebp, esp`

ebp(+/-offset) normally used to locate local vars and parameters on stack

ebp holds a copy of esp

Ebp is pushed so that it can be later popped while returnig

2) Allocate local variables

3) Save callee-saved registers

X86 caller and callee rules

Callee rules on return

- 1) Leave return value in eax
- 2) Restore callee saved registers
- 3) Deallocate local variables
- 4) restore the ebp
- 5) return

32 bit vs 64 bit calling convention

Registers are used for passing parameters in 64 bit , to a large extent

Upto 6 parameters

More parameters pushed on stack

See

<https://aaronbloomfield.github.io/pdr/book/x86-64bit-ccc-chapter.pdf>

Beware

When you read assembly code generated using

`gcc -S`

You will find

More complex instructions

But they will essentially follow the convention mentioned

Comparison

	MIPS	x86
Arguments:	First 4 in %a0–%a3, remainder on stack	Generally all on stack
Return values:	%v0–%v1	%eax
Caller-saved registers:	%t0–%t9	%eax, %ecx, & %edx
Callee-saved registers:	%s0–%s9	Usually none

Figure 6.2: A comparison of the calling conventions of MIPS and x86

From the textbook by Misruda

simple3.c and simple3.s

```
int f(int x1, int x2, int x3, int x4, int x5, int x6, int x7, int x8, int x9, int x10) {  
    int h;  
    h = x1 + x2 + x3 + x4 + x5 + x6 + x7 + x8 + x9 + x10 + 3;  
    return h;  
}  
int main() {  
    int a1 = 10, a2 = 20, a3 = 30, a4 = 40, a5 = 50, a6 = 60, a7 = 70, a8 = 80, a9 = 90, a10 = 100;  
    int b;  
    b = f(a1, a2, a3, a4, a5, a6, a7, a8, a9, a10);  
    return b;  
}
```

simple3.c and simple3.s

main:

endbr64	movl -24(%rbp), %r9d	
pushq %rbp	movl -28(%rbp), %r8d	
movq %rsp, %rbp	movl -32(%rbp), %ecx	movl %eax, %edi
subq \$48, %rsp	movl -36(%rbp), %edx	call f
movl \$10, -44(%rbp)	movl -40(%rbp), %esi	addq \$32, %rsp
movl \$20, -40(%rbp)	movl -44(%rbp), %eax	movl %eax, -4(%rbp)
movl \$30, -36(%rbp)	movl -8(%rbp), %edi	movl -4(%rbp), %eax
movl \$40, -32(%rbp)	pushq %rdi	leave
movl \$50, -28(%rbp)	movl -12(%rbp), %edi	ret
movl \$60, -24(%rbp)	pushq %rdi	
movl \$70, -20(%rbp)	movl -16(%rbp), %edi	
movl \$80, -16(%rbp)	pushq %rdi	
movl \$90, -12(%rbp)	movl -20(%rbp), %edi	
movl \$100, -8(%rbp)	pushq %rdi	

simple3.c and simple3.s

f:

endbr64		
pushq %rbp	addl %eax, %edx	
movq %rsp, %rbp	movl -36(%rbp), %eax	
movl %edi, -20(%rbp)	addl %eax, %edx	movl 40(%rbp), %eax
movl %esi, -24(%rbp)	movl -40(%rbp), %eax	addl %edx, %eax
movl %edx, -28(%rbp)	addl %eax, %edx	addl \$3, %eax
movl %ecx, -32(%rbp)	movl 16(%rbp), %eax	movl %eax, -4(%rbp)
movl %r8d, -36(%rbp)	addl %eax, %edx	movl -4(%rbp), %eax
movl %r9d, -40(%rbp)	movl 24(%rbp), %eax	popq %rbp
movl -20(%rbp), %edx	addl %eax, %edx	ret
movl -24(%rbp), %eax	movl 32(%rbp), %eax	
addl %eax, %edx	addl %eax, %edx	
movl -28(%rbp), %eax		
addl %eax, %edx		
movl -32(%rbp), %eax		