#### System Calls, fork(), exec()

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Credits: Slides of "OS Book" ed10.

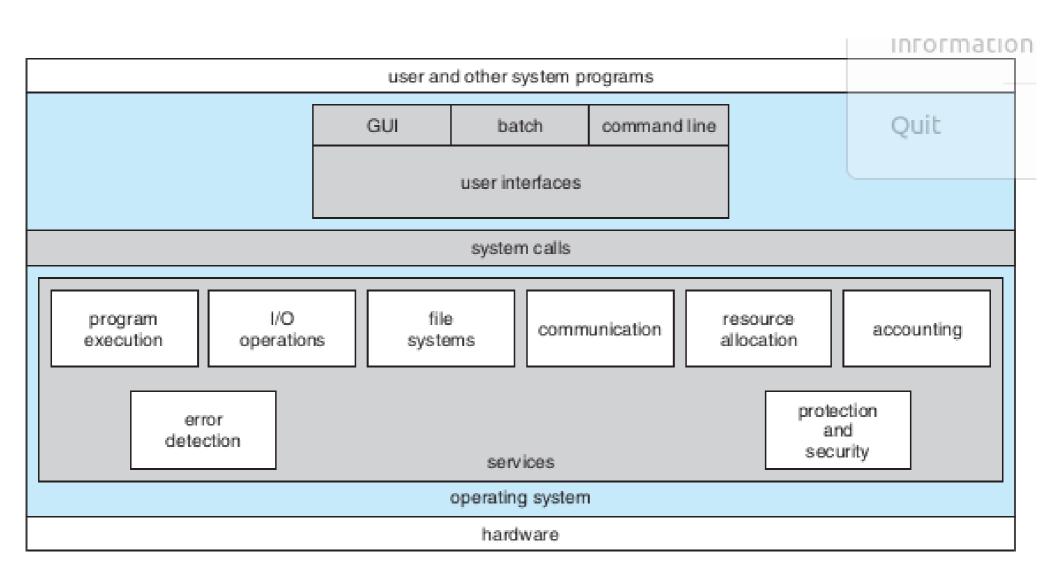


Figure 2.1 A view of operating system services.

## System Calls

- Services provided by operating system to applications
  - Essentially available to applications by calling the particular software interrupt application
    - All system calls essentially involve the "INT 0x80" on x86 processors + Linux
    - Different arguments specified in EAX register inform the kernel about different system calls
- The C library has wrapper functions for each of the system calls
  - E.g. open(), read(), write(), fork(), mmap(), etc.

## **Types of System Calls**

- File System Related
  - Open(), read(), write(), close(), etc.
- Processes Related
  - Fork(), exec(), ...
- Memory management related
  - Mmap(), shm\_open(), ...
- Device Management
- Information maintainance time,date
- Communication between processes (IPC)
- Read man syscalls

#### EXAMPLES OF WINDOWS AND UNIX SYSTEM CALLS en Launcher

	Windows	Unix Information
Process Control	<pre>CreateProcess() ExitProcess() WaitForSingleObject()</pre>	fork() exit() wait()
File Manipulation	<pre>CreateFile() ReadFile() WriteFile() CloseHandle()</pre>	open() read() write() close()
Device Manipulation	<pre>SetConsoleMode() ReadConsole() WriteConsole()</pre>	ioctl() read() write()
Information Maintenance	<pre>GetCurrentProcessID() SetTimer() Sleep()</pre>	<pre>getpid() alarm() sleep()</pre>
Communication	<pre>CreatePipe() CreateFileMapping() MapViewOfFile()</pre>	<pre>pipe() shm_open() mmap()</pre>
Protection	<pre>SetFileSecurity() InitlializeSecurityDescriptor() SetSecurityDescriptorGroup()</pre>	chmod() umask() chown()

```
int main() {
  int a = 2;
  printf("hi\n");
C Library
int printf("void *a, ...) {
  write(1, a, ...);
int write(int fd, char *, int len) {
  int ret:
  mov $5, %eax,
  mov ... %ebx,
  mov ..., %ecx
  int $0x80
  __asm__("movl %eax, -4(%ebp)");
# -4ebp is ret
  return ret;
```

### **Code schematic**

```
----user-kernel-mode-
boundary----
//OS code
int sys_write(int fd, char *, int
len) {
  figure out location on disk
  where to do the write and
  carry out the operation,
  etc.
```

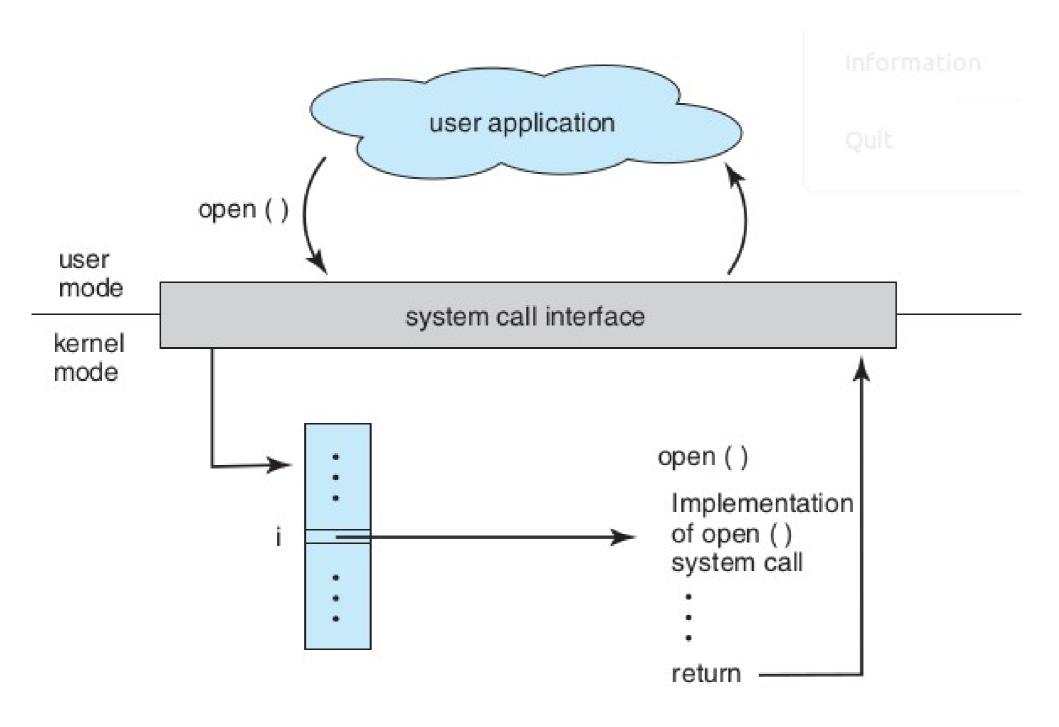


Figure 2.6 The handling of a user application invoking the open() system call.

# Two important system calls Related to processes

fork() and exec()

#### **Process**

- A program in execution
- Exists in RAM
- Scheduled by OS
  - In a timesharing system, intermittantly scheduled by allocating a time quantum, e.g. 20 microseconds
- The "ps" command on Linux

#### **Process in RAM**

- Memory is required to store the following components of a process
  - Code
  - Global variables (data)
  - Stack (stores local variables of functions)
  - Heap (stores malloced memory)
  - Shared libraries (e.g. code of printf, etc)
  - Few other things, may be

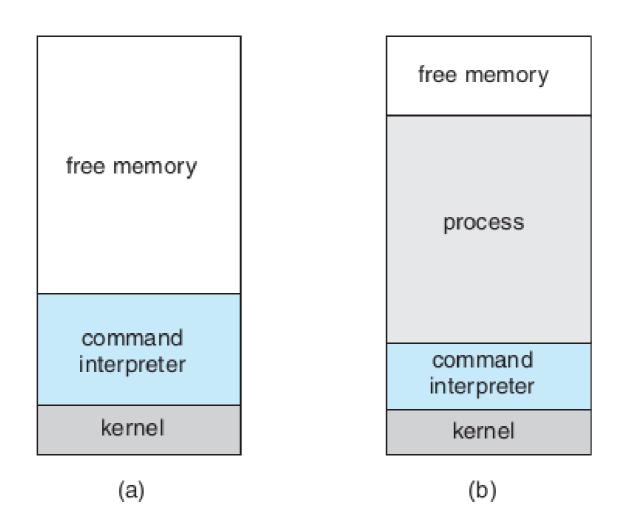


Figure 2.9 MS-DOS execution. (a) At system startup. (b) Running a program.

MS-DOS: a single tasking operating system
Only one program in RAM at a time, and only one program can run at a time

process D

free memory

process C

interpreter

process B

kernel

A multi tasking system
With multiple programs loaded in memory
Along with kernel
(A very simplified conceptual diagram. Things are more complex in reality)

## fork()

- A running process creates it's duplicate!
- After call to fork() is over
  - Two processes are running
  - Identical
  - The calling function returns in two places!
  - Caller is called parent, and the new process is called child
  - PID is returned to parent and 0 to child

## exec()

- Variants: execvp(), execl(), etc.
- Takes the path name of an executable as an argument
- Overwrites the existing process using the code provided in the executable
- The original process is OVER! Vanished!
- The new program starts running overwritting the existing process!

## Shell using fork and exec

- Demo
- The only way a process can be created on Unix/Linux is using fork() + exec()
- All processes that you see were started by some other process using fork() + exec(), except the initial "init" process
- When you click on "firefox" icon, the user-interface program does a fork() + exec() to start firefox; same with a command line shell program
- The "bash" shell you have been using is nothing but an advanced version of the shell code shown during the demo
- See the process tree starting from "init"
- Your next assignment

## The boot process, once again

- BIOS
- Boot loader
- OS kernel
- Init created by kernel by Hand(kernel mode)
- Kernel schedules init (the only process)
- Init fork-execs some programs (user mode)
  - Now these programs will be scheduled by OS
- Init -> GUI -> terminal -> shell
  - One of the typical parent-child relationships