



Course Name: **Electronic Devices Lab**  
Course Number and Section: 14:332:333:2A

Experiment: Lab Experiment 1 – Introduction to C

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1. Match the items on the left with the memory segment in which they are stored. Answers may be used more than once, and more than one answer may be required.

|                             |           |
|-----------------------------|-----------|
| 1. Static Variables – B     | A. Code   |
| 2. Local Variables – D      | B. Static |
| 3. Global Variables – B     | C. Heap   |
| 4. Constants – A B D        | D. Stack  |
| 5. Machine Instructions – A |           |
| 6. malloc() – C             |           |
| 7. String Literals – B      |           |

2. Write the code necessary to properly allocate memory (on the heap) in the following scenarios:

- a. An array `arr` of `k` integers:

```
arr = (int *)malloc(sizeof(int) * k);
```

- b. A string `str` of length `p`:

```
str = (char *)malloc(sizeof(char) * (p + 1));
```

- c. An `n x m` matrix `mat` of integers initialized to zeroes:

```
mat = (int **)calloc(n, sizeof(int *));
for(int i = 0; i < m; i++){
    mat[i] = (int *)calloc(m, sizeof(int));
}
```

3. Assume we have an array in memory that contains `int * arr = {1,2,3,4,5,6,0}`. Let the value of `arr` be a multiple of 4 and stored in register `s0`. What do the snippets of RISC-V code do?

- a. Sets register `t0` to `arr[3]`:

```
lw t0,12(s0)
```

- b. Array element with pointer `t2` incremented by 1:

```
slli t1,t0,2
add t2,s0,t1
lw t3,0(t2)
addi t3,t3,1
sw t3,0(t2)
```

- c. Multiplies `arr[0]` by -1 and sets register `t0` to its two's complement:

```
lw t0,0(s0)
xori t0,t0,0xFFF
addi t0,t0,1
```

4. What are the instructions to branch to label on each of the following conditions? The only branch instructions you may use are beq and bne.

| $s0 < s1$                                    | $s0 \leq s1$                                 | $s0 > 1$                                      |
|--|--|---|
| <i>slt t0,s0,s1</i><br><i>bne t0,0,label</i> | <i>slt t0,s1,s0</i><br><i>beq t0,0,label</i> | <i>sltiu t0,s0,2</i><br><i>beq t0,0,label</i> |

5. Open the files lab3\_ex5\_c.c and lab3\_ex5\_assembly.s. The assembly code provided (.s file) is a translation of the given C program into RISC-V. Your task is to find/explain the following components of this assembly file.

- Register representing variable k  
t0 represents variable k and is set to 0 where each iteration increments it by 1.
- Registers behaving as pointers to destination arrays and source.  
t1 is set to source array address in main function therefore having it pointed to the source array just as t2 points to the destination array.
- Assembly code for loop found in C code.

```
slli t3,t0,2
add t4,t1,t3
lw t5,0(t4)
beq t5,x0,exit
add t6,t2,t3
sw t5,0(t6)
addi t0,t0,1
jal x0,loop
```

exit: Loop checks if  $source[k] = 0$ . Source[k] set to dest[k] incrementing k by 1 per iteration. After *beq t5,x0,exit* this executes. Because we were returned the exit value the original equality must stand true.

- How pointers are manipulated in assembly.  
t1 increments by 0100<sub>2</sub> per iteration which after each the value is stored in t4 allowing t1 to retain its original value throughout the program. The value stored at the 0 offset in register t4 is stored in t5 so that the beq statement can be checked. The instruction set is similar for register t2 in which the destination address is stored. t2 is incremented by 0100<sub>2</sub> per iteration which after each it is stored in register t6 so that t2 can retain its original value throughout the program.

6. Translating between C and RISC-V Translate between the C and RISC-V code. You may want to use the RISC-V Reference Card for more information on the instruction set and syntax. In all of the C examples, we show you how the different variables map to registers – you don't have to worry about the stack or any memory-related issues. You may assume all registers are initialized to zero.

| C   | RISC-V  |
|---|---|
| <pre> //s0 → a, s1 → b //s2 → c, s3 → z int a = 4, b = 5, c = 6, z; z = a + b + c + 10; </pre>              | <pre> addi s0, x0, 4 addi s1, x0, 5 addi s2, x0, 6 add s3, s0, s1 add s3, s3, s2 addi s3, s3, 10 </pre>                             |
| <pre> // s0 → int * p = intArr; // s1 → a; *p = 0; int a = 2; p[1] = p[a] = a; </pre>                       | <pre> sw x0, 0(s0) addi s1, x0, 2 sw s1, 4(s0) slli t0, s1, 2 add t0, t0, s0 sw s1, 0(t0) </pre>                                    |
| <pre> // s0 → a, s1 → b int a = 5, b = 10; if(a + a == b) {     a = 0; }else{     b = a - 1; } </pre>       | <pre> addi s0, x0, 5 addi s1, x0, 10 add t0, s0, s0 bne t0, s1, else xor s0, x0, x0 jal x0, exit else: addi s1, s0, -1 exit: </pre> |
| <pre> // s1 = 2^30 s1 = 1; for(s0 = 0; s0 &lt; 30; s++) {     s1 *= 2; } </pre>                             | <pre> addi s0, x0, 0 addi s1, x0, 1 addi t0, x0, 30 loop: beq s0, t0, exit add s1, s1, s1 addi s0, s0, 1 jal x0, loop exit: </pre>  |
| <pre> // s0 → n, s1 → sum // assume n &gt; 0 to start int sum; for(sum = 0; n &gt; 0; sum += n - -); </pre> | <pre> addi s1, s1, 0 loop: beq s0, x0, exit add s1, s1, s0 add s0, s0, -1 jal x0, loop exit: </pre>                                 |

7. Implement a function factorial in RISC-V that has a single integer parameter  $n$  and returns  $n!$ . A stub of this function can be found in the file factorial. You will only need to add instructions under the factorial label, and the argument that is passed into the function is configured to be located at the label  $n$ . You may solve this problem using either recursion or iteration.

```
main:
    la t0,n
    lw a0,0(t0)
    jal ra,factorial
    addi a1,a0,0
    addi a0,x0,1
    ecall
    addi a0,x0,10
    ecall
factorial:
    li t1,1
    li t2,1
    bgt a0,t1,LABEL
    mul t2,t2,t1
LABEL:
    addi a0,t2,0
    jalr zero,ra,0
```