

Course Number and Section: 14:332:333:2A

Experiment: Lab Experiment 1 – Introduction to C

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1. Match the items on the left with the memory segment in which they are stored. Answers may be used more than once, and more than one answer may be required.

1. Static Variables – B	A. Code
2. Local Variables – D	B. Static
3. Global Variables – B	C. Heap
4. Constants – A B D	D. Stack
5. Machine Instructions – A	
6. malloc() – C	
7. String Literals – B	

- 2. Write the code necessary to properly allocate memory (on the heap) in the following scenarios:
  - a. An array arr of k integers:

$$arr = (int *)malloc(sizeof(int) * k);$$

b. A string str of length p:

$$str = (char *)malloc(sizeof(char) * (p + 1));$$

c. An n x m matrix mat of integers initialized to zeroes:

$$mat = (int **)calloc(n, size of (int *));$$
  

$$for(int i = 0; i < m; i + +) \{$$
  

$$mat[i] = (int *)calloc(m, size of (int));$$
  
}

- 3. Assume we have an array in memory that contains  $int * arr = \{1,2,3,4,5,6,0\}$ . Let the value of arr be a multiple of 4 and stored in register s0. What do the snippets of RISC-V code do?
  - a. Sets register t0 to arr[3]:

$$lw\ t0,12(s0)$$

b. Array element with pointer t2 incremented by 1:

c. Multiplies arr[0] by -1 and sets register t0 to its two's compliment:

4. What are the instructions to branch to label on each of the following conditions? The only branch instructions you may use are beq and bne.

s0 < s1	s0 ≤ s1	s0 > 1
slt t0, s0, s1	slt t0,s1,s0	sltiu t0, s0,2
bne t0,0, label	beq t0,0, label	beq t0,0, label

- 5. Open the files lab3\_ex5\_c.c and lab3\_ex5\_assembly.s. The assembly code provided (.s file) is a translation of the given C program into RISC-V. Your task is to find/explain the following components of this assembly file.
  - a. Register representing variable kt0 represents variable k and is set to 0 where each iteration increments it by 1.
  - b. Registers behaving as pointers to destination arrays and source.t1 is set to source array address in main function therefore having it pointed to the source array just as t2 points to the destination array.
  - c. Assembly code for loop found in C code.

slli t3, t0,2 add t4, t1, t3 lw t5,0(t4) beq t5, x0, exit add t6, t2, t3 sw t5,0(t6) addi t0, t0,1 jal x0,loop

exit: Loop checks if source[k] = 0. Source[k] set to dest[k] incrementing k by 1 per iteration. After  $beq\ t5$ , x0, exit this executes. Because we were returned the exit value the original equality must stand true.

d. How pointers are manipulated in assembly.

t1 increments by 0100<sub>2</sub> per iteration which after each the value is stored in t4 allowing t1 to retain its original value throughout the program. The value stored at the 0 offset in register t4 is stored in t5 so that the beq statement can be checked. The instruction set is similar for register t2 in which the destination address is stored. t2 is incremented by 0100<sub>2</sub> per iteration which after each it is stored in register t6 so that t2 can retain its original value throughout the program.

6. Translating between C and RISC-V Translate between the C and RISC-V code. You may want to use the RISC-V Reference Card for more information on the instruction set and syntax. In all of the C examples, we show you how the different variables map to registers – you don't have to worry about the stack or any memory-related issues. You may assume all registers are initialized to zero.

С	RISC-V
$//s0 \rightarrow a, s1 \rightarrow b$	addi s0, x0, 4
$//s2 \rightarrow c, s33 \rightarrow z$	addi s1, x0, 5
$int \ a = 4, b = 5, c = 6, z;$	addi s2,x0,6
z = a + b + c + 10;	add s3, s0, s1
	add s3, s3, s2
	addi s3, s3, 10
$//s0 \rightarrow int * p = intArr;$	sw x0,0(s0)
$//s1 \rightarrow a; *p = 0;$	addi s1,x0,2
int $a = 2$ ;	sw s1,4(s0)
p[1] = p[a] = a;	slli t0,s1,2
	add t0, t0, s0
	sw s1,0(t0)
$//s0 \rightarrow a, s1 \rightarrow b$	addi s0,x0,5
$int \ a = 5, b = 10;$	addi s1, x0, 10
$if(a+a==b)\{$	add t0, s0, s0
a=0;	bne t0, s1, else
}else{	xor s0, x0, x0
b=a-1;	jal x0, exit
}	else:
	addi s1, s0, −1
	exit:
$//s1 = 2^30$	addi s0, x0, 0
s1 = 1;	addi s1, x0, 1
$for(s0 = 0; s0 < 30; s + +) $ {	addi t0, x0, 30
s1 *= 2;	loop:
}	beq s0, t0, exit
	add s1, s1, s1
	addi s0, s0, 1
	jal x0, loop
// 0 4	exit:
$//s0 \rightarrow n, s1 \rightarrow um$	addi s1, s1, 0
// assume $n > 0$ to start	loop:
int sum;	beq s0, x0, exit
for(sum = 0; n > 0; sum += n);	add s1, s1, s0
	add s0, s0, -1
	jal x0, loop
	exit: