

Batch: D2 Roll No.:16010122323

Experiment / assignment / tutorial No. 3

Grade: AA / AB / BB / BC / CC / CD / DD

Signature of the Staff In-charge with date

TITLE : To study and implement Restoring method of division

AIM : The basis of algorithm is based on paper and pencil approach and the operation involves repetitive shifting with addition and subtraction. So the main aim is to depict the usual process in the form of an algorithm.

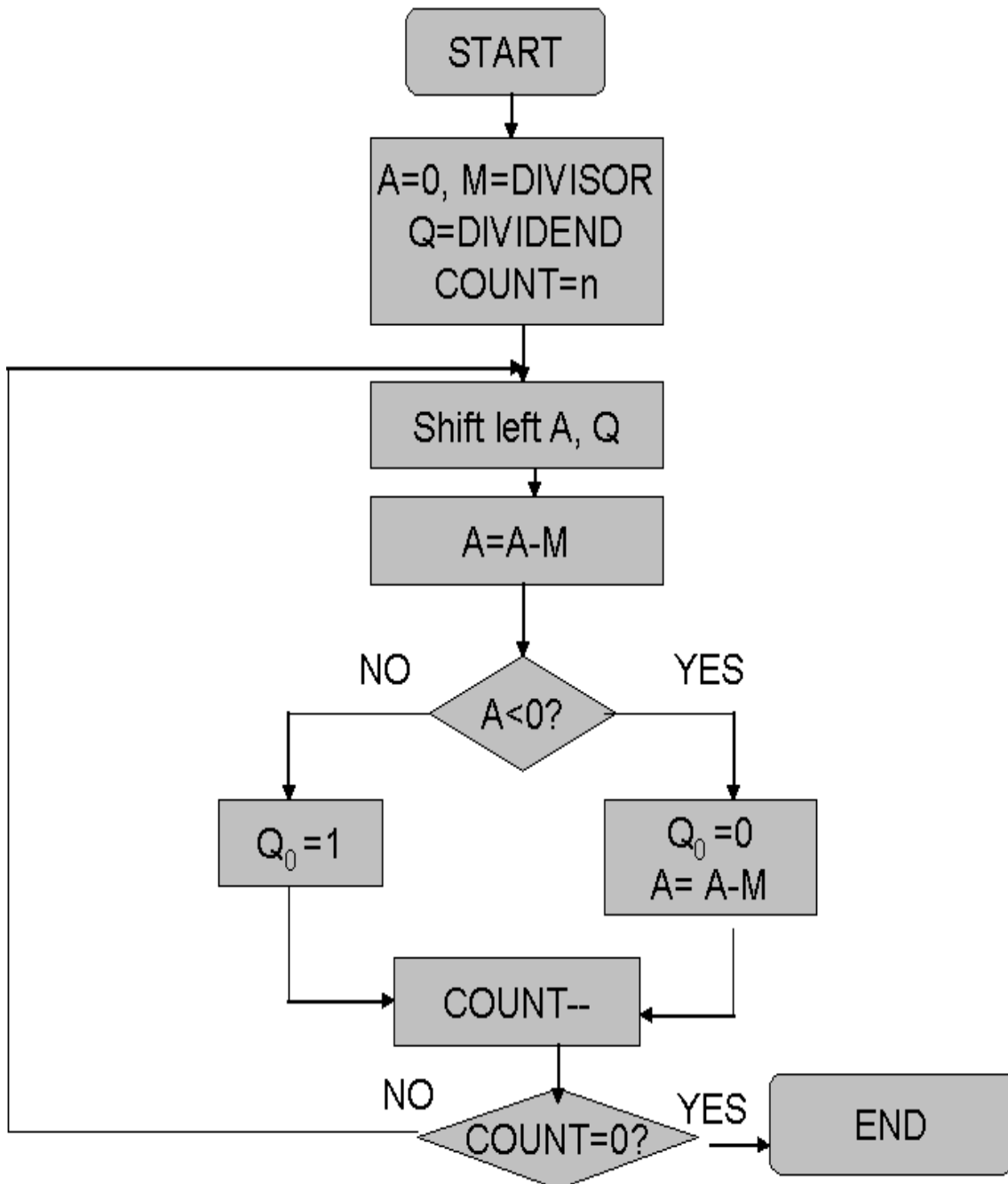
Expected OUTCOME of Experiment: (Mention CO /CO's attained here)

Books/ Journals/ Websites referred:

1. Carl Hamacher, Zvonko Vranesic and Safwat Zaky, "Computer Organization", Fifth Edition, TataMcGraw-Hill.
2. William Stallings, "Computer Organization and Architecture: Designing for Performance", Eighth Edition, Pearson.
3. Dr. M. Usha, T. S. Srikanth, "Computer System Architecture and Organization", First Edition, Wiley-India.

Pre Lab/ Prior Concepts:

The Restoring algorithm works with any combination of positive and negative numbers.

Flowchart for Restoring of Division:



Design Steps:

1. Start
2. Initialize $A=0$, $M=\text{Divisor}$, $Q=\text{Dividend}$ and $\text{count}=n$ (no of bits)
3. Left shift A, Q
4. If MSB of A and M are same
5. Then $A=A-M$
6. Else $A=A+M$
7. If MSB of previous A and present A are same
8. $Q_0=0$ & store present A
9. Else $Q_0=1$ & restore previous A
10. Decrement count.
11. If $\text{count}=0$ go to 11
12. Else go to 3
13. STOP

Example:-

Page No.
 Date: / /

Q. $M = 12$ $M = 01100$ $Q = 26$ $Q = 11010$ $M = 10100$

A	Q	Operation	Shift
00000	11010		
00001	10100	shift left	
10101	10100	$A \leftarrow A - M$	1st
10101	10100	$Q_0 \leftarrow 0$	
00001	10100	$A \leftarrow A + M$	
00011	01000	shift left	
10111	01000	$A \leftarrow A - M$	2nd
10111	01000	$Q_0 \leftarrow 0$	
00011	01000	$A \leftarrow A + M$	
00110	10000	shift left	
11010	10000	$A \leftarrow A - M$	3rd
11010	10000	$Q_0 \leftarrow 0$	
00110	10000	$A \leftarrow A + M$	
01101	00000	shift left	
00001	00000	$A \leftarrow A - M$	4th
00001	00000	$Q_0 \leftarrow 1$	
00010	00001	shift left	
10110	00010	$A \leftarrow A - M$	5th
10110	00010	$Q_0 \leftarrow 0$	
00010	00010	$A \leftarrow A + M$	

↓ ↓
2 2

A	Q		
000000	111011		
000001	11011□	shift left	1st
100001	11011□	$A \leftarrow A - M$	
100001	110110	$Q_0 \leftarrow 0$	
000001	110110	$A \leftarrow A + M$	
000011	10110□	shift left	2nd
100011	10110□	$A \leftarrow A - M$	
100011	101100	$Q_0 \leftarrow 0$	
000011	101100	$A \leftarrow A + M$	
000111	01100□	shift left	3rd
100111	01100□	$A \leftarrow A - M$	
100111	011000	$Q_0 \leftarrow 0$	
000111	011000	$A \leftarrow A + M$	
001110	110001□	shift left	4th
101110	110001□	$A \leftarrow A - M$	
101110	110000	$Q_0 \leftarrow 0$	
001110	110000	$A \leftarrow A + M$	
011101	10000□	shift left	5th
111101	10000□	$A \leftarrow A - M$	
111101	100000	$Q_0 \leftarrow 0$	
011101	100000	$A \leftarrow A + M$	
111011	00000□	shift left	6th
011011	00000□	$A \leftarrow A - M$	
011011	000001	$Q_0 \leftarrow 1$	
↓	↓		
1	27	1	
2			
18			
16			
27			

Page No.
 Date: / /

1101
 1101
 0000

1110
 0011
 0001

0001
 1101
 1110

0001
 1101
 1110

1101
 0011
 0000

1110
 0011
 0001

Eg. $M = 3$ $M = 0011$ $-M = 1101$
 $R = 7$ $R = 0111$

A	R	
0000	0111	
0000	1111	shift left
1101	1111	$A \leftarrow A - M$
1101	1110	$R_0 \leftarrow 0$
0000	1110	$A \leftarrow A + M$
0001	1101	shift left
1110	1101	$A \leftarrow A - M$
1110	1100	$R_0 \leftarrow 0$
0001	1100	$A \leftarrow A + M$
0011	1001	shift left
0000	1001	$A \leftarrow A - M$
0000	1001	$R_0 \rightarrow 1$
0001	0011	shift left
1110	0011	$A \leftarrow A - M$
1110	0010	$R_0 \leftarrow 0$
0001	0010	$A \leftarrow A + M$

1
 A

2
 R

First cycle.
 2nd cycle.
 Third cycle.
 4th cycle.

Page No.
 Date: / /

Q. $M = 5$ $m = 0101$ $-M = 1011$
 $R = 5$ $R = 0101$

A	Q	Operation	Step
0000	0101		
0000	101	shift left	1st
1011	101	$A \leftarrow A - M$	
1011	1010	$Q_0 \leftarrow 0$	
0000	1010	$A \leftarrow A + M$	
0001	010	shift left	2nd
1100	010	$A \leftarrow A - M$	
1100	0100	$Q_0 \leftarrow 0$	
0001	0100	$A \leftarrow A + M$	
0010	100	shift left	3rd
1101	100	$A \leftarrow A - M$	
1101	1000	$Q_0 \leftarrow 0$	
0010	1000	$A \leftarrow A + M$	
0101	000	shift left	4th
0000	000	$A \leftarrow A - M$	
0000	0001	$Q_0 \leftarrow 1$	
↓	↓		
0	1		

Implementation:

```
#include <stdio.h>

int a = 0, b = 0, c = 0, s = 0;
int com[] = {1, 0, 0, 0, 0};
int anum[] = {0, 0, 0, 0, 0};
int anumcp[] = {0, 0, 0, 0, 0};
int bnum[] = {0, 0, 0, 0, 0};
int acomp[] = {0, 0, 0, 0, 0};
int bcomp[] = {0, 0, 0, 0, 0};
int rem[] = {0, 0, 0, 0, 0};
int quo[] = {0, 0, 0, 0, 0};
int res[] = {0, 0, 0, 0, 0};

void binary() {
    a = abs(a);
    b = abs(b);
    int r, r2, i;
    for (i = 0; i < 5; i++) {
        r = a % 2;
        a = a / 2;
        r2 = b % 2;
        b = b / 2;
        anum[i] = r;
        anumcp[i] = r;
        bnum[i] = r2;
        if (r2 == 0) {
            bcomp[i] = 1;
        }
        if (r == 0) {
            acomp[i] = 1;
        }
    }
    c = 0;
    for (i = 0; i < 5; i++) {
        res[i] = com[i] + bcomp[i] + c;
        if (res[i] >= 2) {
            c = 1;
        } else
            c = 0;
        res[i] = res[i] % 2;
    }
    for (i = 4; i >= 0; i--) {
```



```
        bcomp[i] = res[i];
    }
}

void add(int num[]) {
    int i;
    c = 0;
    for (i = 0; i < 5; i++) {
        res[i] = rem[i] + num[i] + c;
        if (res[i] >= 2) {
            c = 1;
        } else
            c = 0;
        res[i] = res[i] % 2;
    }
    for (i = 4; i >= 0; i--) {
        rem[i] = res[i];
        printf("%d", rem[i]);
    }
    printf(":");
    for (i = 4; i >= 0; i--) {
        printf("%d", anumcp[i]);
    }
}

void shl() {
    int i;
    for (i = 4; i > 0; i--) {
        rem[i] = rem[i - 1];
    }
    rem[0] = anumcp[4];
    for (i = 4; i > 0; i--) {
        anumcp[i] = anumcp[i - 1];
    }
    anumcp[0] = 0;
    printf("\nSHIFT LEFT: ");
    for (i = 4; i >= 0; i--) {
        printf("%d", rem[i]);
    }
    printf(":");
    for (i = 4; i >= 0; i--) {
        printf("%d", anumcp[i]);
    }
}
```

```
int main() {
    printf("Name: Vedansh Savla\n");
    printf("Roll Number: 16010122323 \nDivision: D2\n-----
-----\n");
    printf("COA exp 3: To study and implement Restoring method of
division\n");
    printf("Implementation details:\n-----
-----\n");
    printf("RESTORING DIVISION ALGORITHM\n");
    printf("Enter two numbers to multiply:\n");
    printf("Both must be less than 16\n");

    do {
        printf("Enter Dividend: ");
        scanf("%d", &a);
        printf("Enter Divisor: ");
        scanf("%d", &b);
    } while (a >= 16 || b >= 16);

    printf("Expected Quotient = %d\n", a / b);
    printf("Expected Remainder = %d\n", a % b);

    if (a * b < 0) {
        s = 1;
    }

    binary();
    printf("\nUnsigned Binary Equivalents are:\n");
    printf("A = ");
    for (int i = 4; i >= 0; i--) {
        printf("%d", anum[i]);
    }
    printf("\nB = ");
    for (int i = 4; i >= 0; i--) {
        printf("%d", bnum[i]);
    }
    printf("\nB'+ 1 = ");
    for (int i = 4; i >= 0; i--) {
        printf("%d", bcomp[i]);
    }
    printf("\n\n-->\n");

    shl();
}
```

```
for (int i = 0; i < 5; i++) {
    printf("\n-->\n");
    printf("\nSUB B: ");
    add(bcomp);
    if (rem[4] == 1) {
        printf("\n-->RESTORE\n");
        printf("ADD B: ");
        anumcp[0] = 0;
        add(bnum);
    } else {
        anumcp[0] = 1;
    }
    if (i < 4)
        shl();
}
printf("\n-----\n");
printf("Sign of the result = %d\n", s);
printf("Remainder is = ");
for (int i = 4; i >= 0; i--) {
    printf("%d", rem[i]);
}
printf("\nQuotient is = ");
for (int i = 4; i >= 0; i--) {
    printf("%d", anumcp[i]);
}

return 0;
}
```

Output:



```
Name: Vedansh Savla
Roll Number: 16010122323
Division: D2
-----
COA exp 3: To study and implement Restoring method of division
Implementation details:
-----
RESTORING DIVISION ALGORITHM
Enter two numbers to multiply:
Both must be less than 16
Enter Dividend: 11
Enter Divisor: 4
Expected Quotient = 2
Expected Remainder = 3

Unsigned Binary Equivalentents are:
A = 01011
B = 00100
B'+ 1 = 11100

-->

SHIFT LEFT: 00000:10110
-->

SUB B: 11100:10110
-->RESTORE
ADD B: 00000:10110
SHIFT LEFT: 00001:01100
-->
```



```
SUB B: 11100:10110
-->RESTORE
ADD B: 00000:10110
SHIFT LEFT: 00001:01100
-->
```

```
SUB B: 11101:01100
-->RESTORE
ADD B: 00001:01100
SHIFT LEFT: 00010:11000
-->
```

```
SUB B: 11110:11000
-->RESTORE
ADD B: 00010:11000
SHIFT LEFT: 00101:10000
-->
```

```
SUB B: 00001:10000
SHIFT LEFT: 00011:00010
-->
```

```
SUB B: 11111:00010
-->RESTORE
ADD B: 00011:00010
```

```
-----
Sign of the result = 0
Remainder is = 00011
Quotient is = 00010
```


Conclusion:

Booths restoring division was implemented successfully.

Post Lab Descriptive Questions

1. **What are the advantages of restoring division over non restoring division?**

The advantage of using non - restoring arithmetic over the standard restoring division is that a test subtraction is not required; the sign bit determines whether an addition or subtraction is used. The disadvantage, though, is that an extra bit must be maintained in the partial remainder to keep track of the sign.

Date: _____

Signature of faculty in-charge

