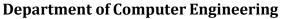


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Batch: C2 Roll No.:16010122323

Experiment / assignment / tutorial No.8

Grade: AA / AB / BB / BC / CC / CD /DD

Signature of the Staff In-charge with date

Experiment No.:8

TITLE: Study and configure RIP protocol using Cisco Packet tracer

AIM: To study and configure RIP protocol using Cisco Packet tracer

Expected Outcome of Experiment:

CO:

Books/ Journals/ Websites referred:

- 1. A. S. Tanenbaum, "Computer Networks", Pearson Education, Fourth Edition
- 2. B. A. Forouzan, "Data Communications and Networking", TMH, Fourth Edition

Pre Lab/ Prior Concepts:

IPv4 Addressing, Subnetting, Distance Vector Protocol, Router configuration Commands.

New Concepts to be learned: RIP Protocol and its configuration.

RIP (Routing Information Protocol)

RIP is a standardized Distance Vector protocol, designed for use on smaller networks. RIP was one of the first true Distance Vector routing protocols and is supported on a wide variety of systems.

RIP adheres to the following Distance Vector characteristics:

- RIP sends out periodic routing updates (every 30 seconds)
- RIP sends out the full routing table every periodic update.
- RIP uses a form of distance as its metric (in this case, hop count).
- RIP uses the Bellman-Ford Distance Vector algorithm to determine the best "path" to a particular destination

Other characteristics of RIP include:

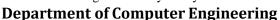
- RIP supports IP and IPX routing.
- RIP utilizes UDP port 520
- RIP routes have an administrative distance of 120.
- RIP has a maximum hop count of 15 hops.

RIP Versions

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RIP has two versions, Version 1 (RIPv1) and Version 2 (RIPv2).

RIPv1 (RFC 1058) is *classful*, and thus does not include the subnet mask with its routing table updates. Because of this, RIPv1 does not support **Variable Length Subnet Masks (VLSMs)**. When using RIPv1, networks must be contiguous, and subnets of a major network must be configured with identical subnet masks. Otherwise, route table inconsistencies (or worse) will occur.

RIPv1 sends updates as broadcasts to address 255.255.255.255.

RIPv2 (RFC 2543) is *classless*, and thus does include the subnet mask with its routing table updates. RIPv2 fully supports VLSMs, allowing discontiguous networks and varying subnet masks to exist.

Other enhancements offered by RIPv2 include:

- Routing updates are sent via multicast, using address 224.0.0.9
- Encrypted authentication can be configured between RIPv2 routers
- Route tagging is supported.

RIPv2 can interoperate with RIPv1. By default:

- RIPv1 routers will send only Version 1 packets
- RIPv1 routers will receive both Version 1 and 2 updates
- RIPv2 routers will both send and receive only Version 2 updates

We can control the version of RIP a particular interface will "send" or "receive." Unless RIPv2 is manually specified, Cisco will default to RIPv1 when configuring RIP.

RIPv1 Basic Configuration



Routing protocol configuration occurs in Global Configuration mode. On Router A, to configure RIP, we would type:

Router(config)# router rip
Router(config-router)# network 172.16.0.0
Router(config-router)# network 172.17.0.0

The first command, router rip, enables the RIP process.

The network statements tell RIP which networks you wish to advertise to other RIP routers. We simply list the networks that are directly connected to our router. Notice that we specify the networks at their classful boundaries, and we do not specify a subnet mask.

To configure Router B:

Router(config)# router rip
Router(config-router)# network 172.17.0.0
Router(config-router)# network 172.18.0.0

The routing table on Router A will look like:

RouterA# show ip route

Gateway of last resort is not set

- C 172.16.0.0 is directly connected, Ethernet0
- C 172.17.0.0 is directly connected, Serial0



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R 172.18.0.0 [120/1] via 172.17.1.2, 00:00:00, Serial0

The routing table on Router B will look like:

RouterB# show ip route

Gateway of last resort is not set

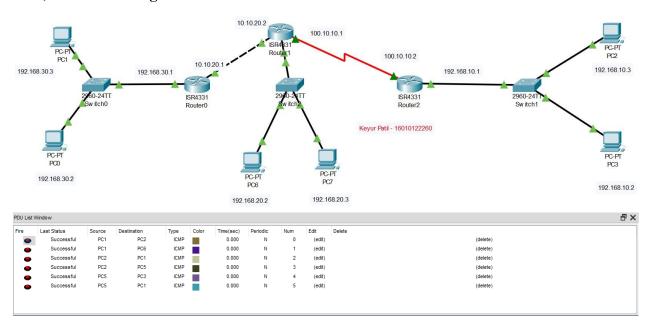
C 172.17.0.0 is directly connected, Serial0

C 172.18.0.0 is directly connected, Ethernet0

R 172.16.0.0 [120/1] via 172.17.1.1, 00:00:00, Serial0

IMPLEMENTATION: (printout of code)

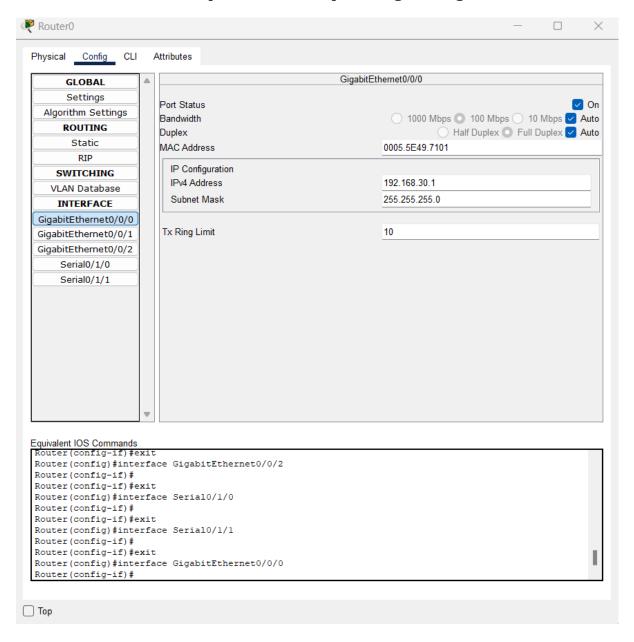
1) Static Routing





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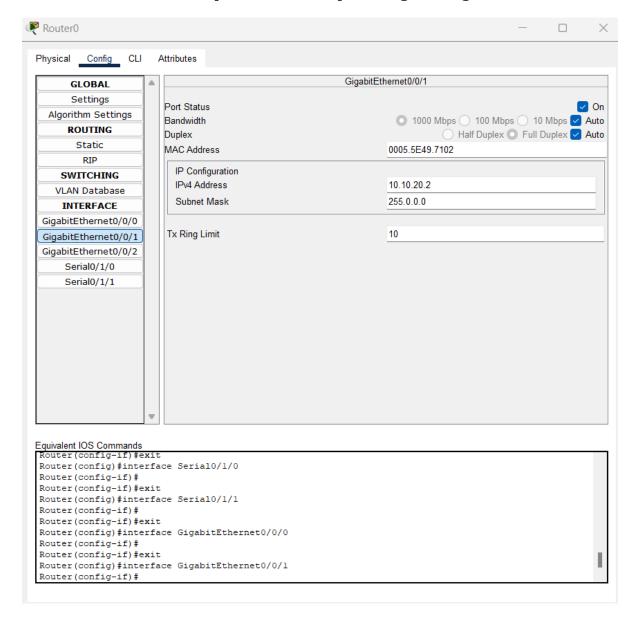






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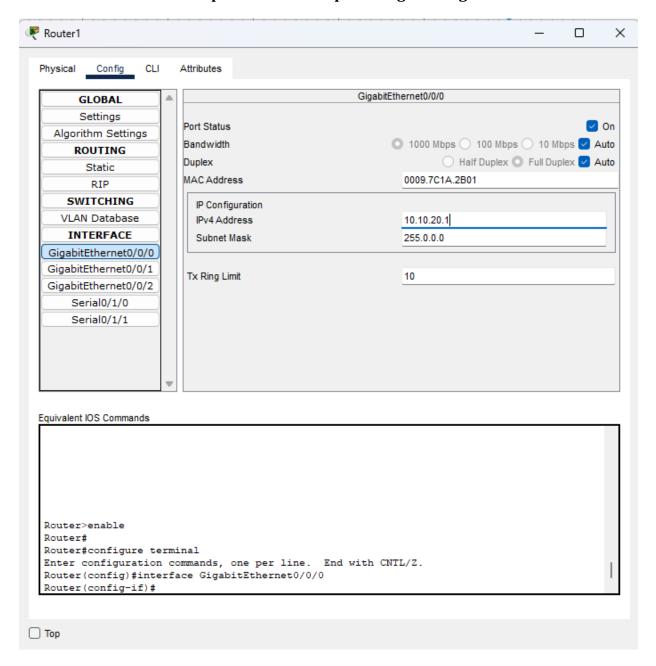


```
--- System Configuration Dialog ---
Would you like to enter the initial configuration dialog? [yes/no]: n
Press RETURN to get started!
Router>en
Router#conf t
Enter configuration commands, one per line. End with CNTL/Z.
Router(config) #interface gig 0/0/0
Router(config-if) #ip address 192.168.30.1 255.255.255.0
Router(config-if) #no shut
Router(config-if)#
%LINK-5-CHANGED: Interface GigabitEthernet0/0/0, changed state to up
%LINEPROTO-5-UPDOWN: Line protocol on Interface GigabitEthernet0/0/0, changed state to up
Router(config-if)#exit
Router(config) #interface gig 0/0/1
Router(config-if) #ip address 10.10.20.2 255.0.0.0
Router(config-if) #no shut
Router(config-if)#
%LINK-5-CHANGED: Interface GigabitEthernet0/0/1, changed state to up
Router(config-if) #exit'
% Invalid input detected at '^' marker.
Router(config-if) #exit
Router(config)#
```



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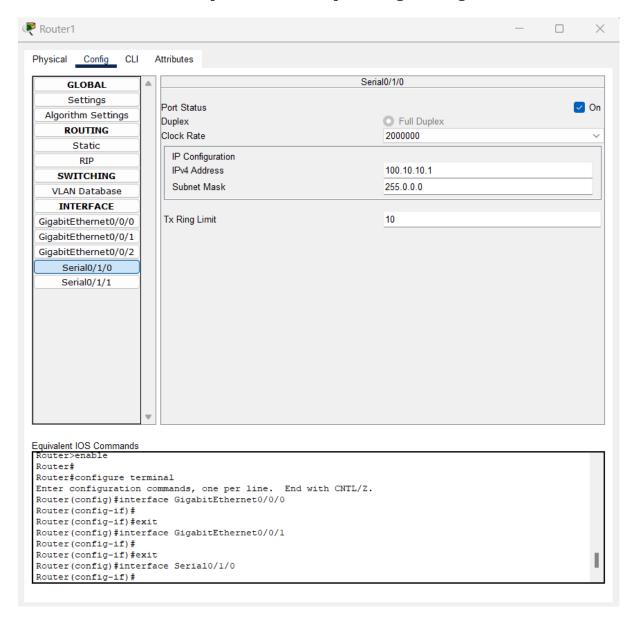






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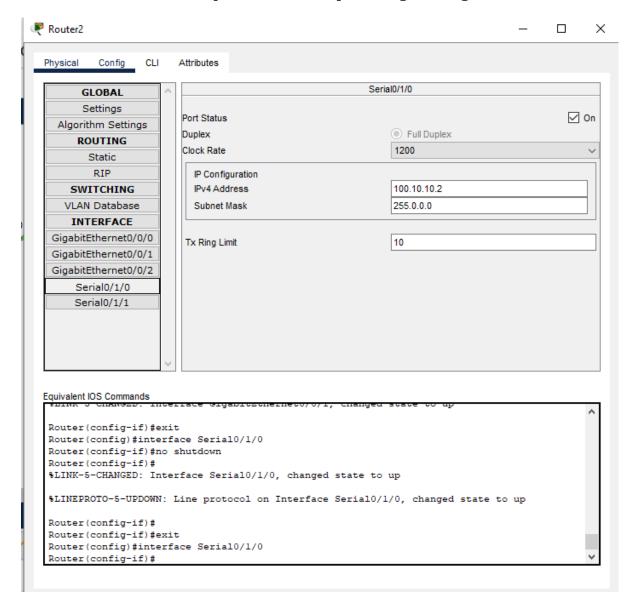






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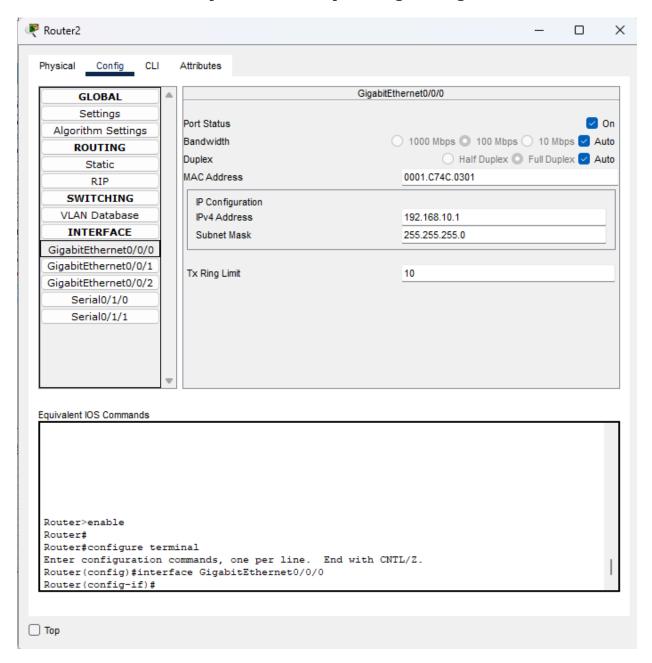






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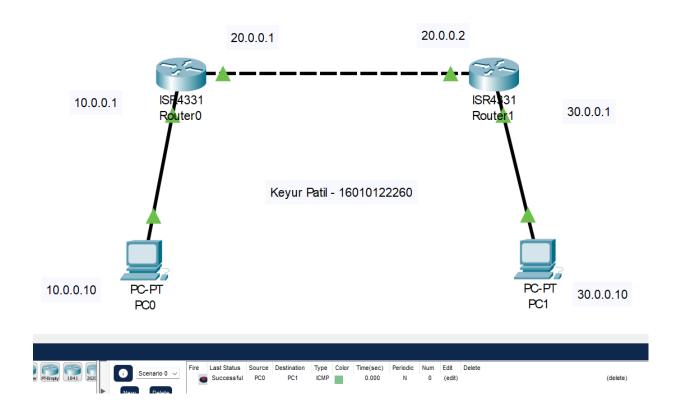


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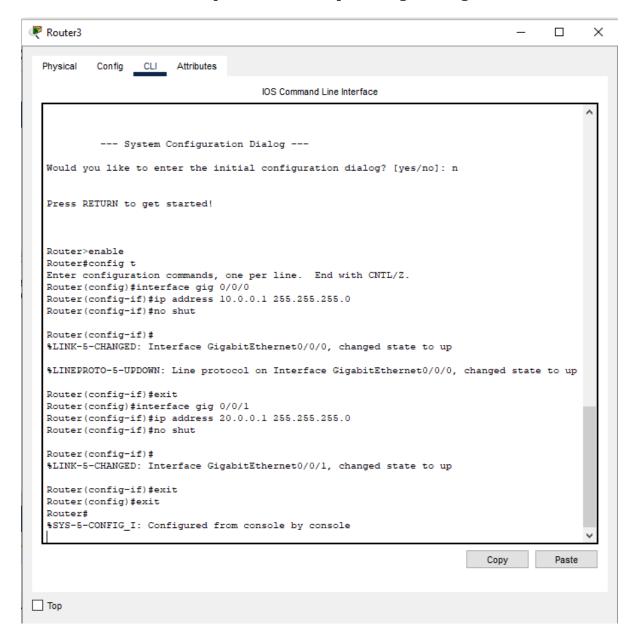
2) Dynamic Routing





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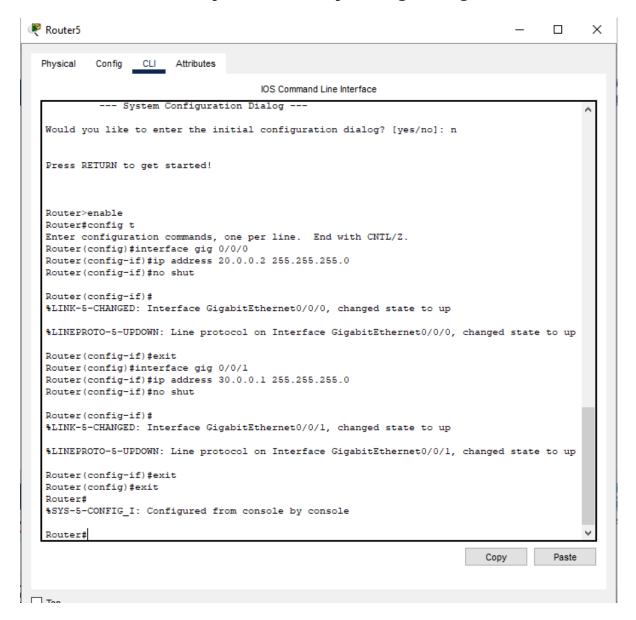






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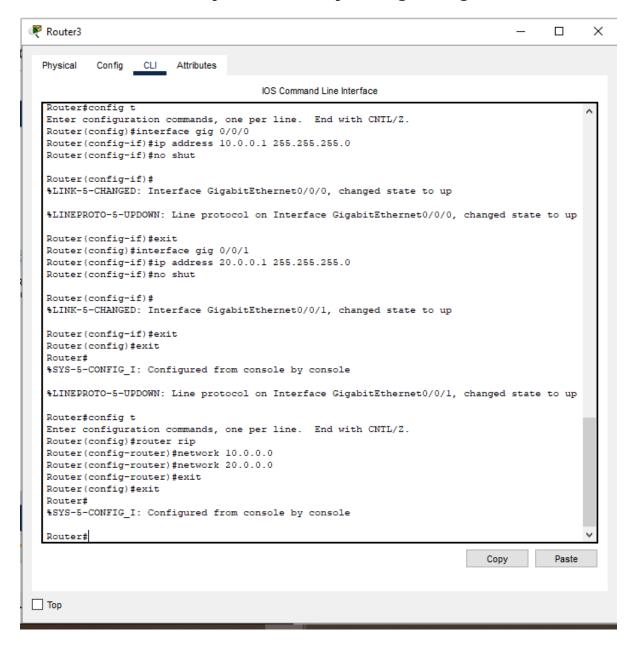






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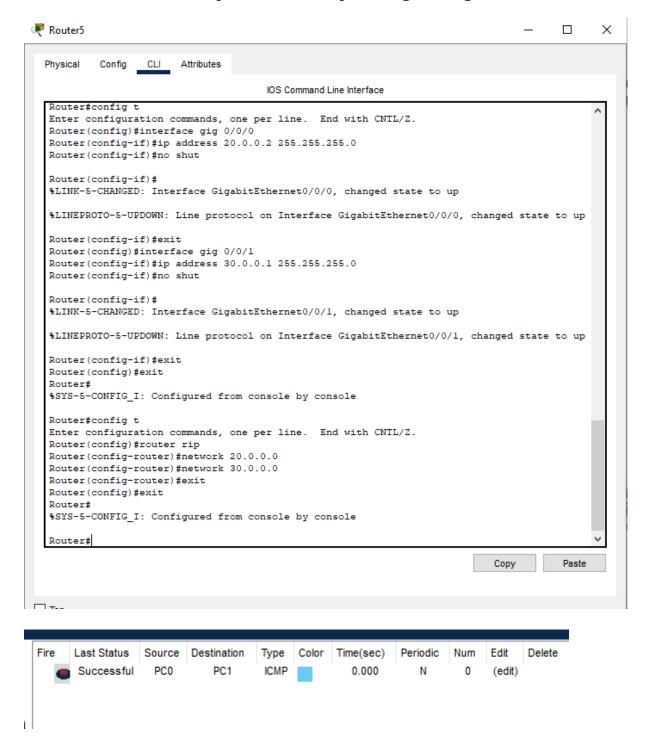




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CONCLUSION: In this experiment, we set up the RIP protocol between two routers using Cisco Packet Tracer. RIP helps routers share their routing information and find the best path to a network based on how many hops (or steps) it takes. We learned about the differences between RIPv1 and RIPv2 and saw how RIP operates in a distance vector routing system.



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Post Lab Questions

1. are two popular examples of distance vector routing protocols.

A. OSPF and RIP

B. RIP and BGP

- C. BGP and OSPF
- D. BGP and SPF
- 2. Arouting table contains information entered manually.

A. static

- B. dynamic
- C. hierarchical
- D. non static
- 3. Arouting table is updated periodically using one of the dynamic routing protocols.
- A. static

B. dynamic

- C. hierarchical
- D. non static
- 4. Which of the following is not the category of dynamic routing algorithm.
- A. Distance vector protocols
- B. Link state protocols
- C. Hybrid protocols

D. Automatic state protocols

- 5. In forwarding, the mask and destination addresses are both 0.0.0.0 in the routing table.
- A. next-hop
- B. network-specific
- C. host-specific
- D. default



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6. Differentiate between Distance Vector Routing and Link State Routing. **Ans:**

Feature	Distance Vector Routing	Link State Routing
Basic Principle	Routers share their distance to destinations with neighbors.	Routers share the state of their links with all routers.
Routing Updates	Periodic updates, often at fixed intervals.	Updates triggered by changes in the network.
Routing Information	Each router maintains a table of distances to each destination.	Each router maintains a complete map of the network topology.
Scalability	Less scalable; performance degrades in larger networks.	More scalable; can handle larger networks efficiently.
Convergence Speed	Slower convergence due to periodic updates.	Faster convergence as routers react to changes immediately.
Complexity	Simpler to implement and understand.	More complex due to the need for maintaining topology information.
Examples	RIP, IGRP	OSPF, IS-IS
Loop Prevention	Uses techniques like split horizon and poison reverse.	Uses Dijkstra's algorithm and link state advertisements.

Date:

Signature of faculty in-charge