GR NO: 21820001 ROLL NO: 223064

NAME: VEDIKA SADAVARTE.

CLASS: S.Y.B.TECH-C

BATCH: C3

ASSIGNMENT 1

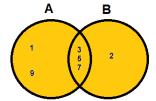
AIM: To Create ADT that implement the "set" concept.

OBJECTIVE: To Create ADT that implement the "set" concept a] Add b] Remove c]Contains d] Size e] Intersection of two set f] Union of two sets g] Difference between two sets h] subset.

THEORY:

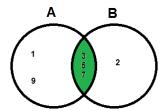
Union

The union of two sets is a new set that contains all of the elements that are in at least one of the two sets. The union is written as AUBAUB or "A or BA or B".



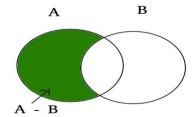
Intersection

The intersection of two sets is a new set that contains all of the elements that are in both sets. The intersection is written as $A \cap BA \cap B$ or "A and BA and B".



DIFFERENCE

Set difference is a generalization of the idea of the complement of a set and as such is sometimes called the relative complement of T with respect to S. The symmetric difference between two sets S and T is the union of S-T and T-S. "set difference." A Dictionary of Computing. . "set difference."



SUBSET

A subset is a set whose elements are all members of another set.

The symbol "⊆" means "is a subset of".

The symbol "⊂" means "is a proper subset of".

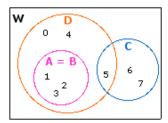
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```
W: The whole numbers: {0, 1, 2, 3, ...}
A: {1, 2, 3}
B: {1, 2, 3}
C: {5, 6, 7}
D: {0, 1, 2, 3, 4, 5}
```



PROGRAM:

```
#include<iostream>
using namespace std;
class Set
{
        int set[20];
        int size;
public:
        Set(int s)
        {
                size=s;
        }
        void create()
        {
                for(int i=0;i<size;i++)
                         cin>>set[i];
                }
        }
        int insert(int element)
        {
                set[size]=element;
                return ++size;
```

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```
}
int contains(int element)
{
        int i=0;
        while(i<size)
        {
                 if(set[i]==element)
                          return i+1;
                 i++;
        }
        return 0;
}
int set_size()
{
        return size;
}
void remove(int element)
{
        int i=0;
        while(i<size)
        {
                 if(set[i]==element)
                 {
                          for(int j=i;j<size;j++)</pre>
                          {
                                   set[j]=set[j+1];
                          }
                          cout<<"\nElement deleted...";</pre>
                          size=size-1;
                          return;
                 }
```

```
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```

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```
i++;
        }
        cout<<"\nSorry, element not found...";</pre>
}
void display()
{
        cout<<"{ ";
        for(int i=0;i<size;i++)
        {
                 cout<<set[i]<<" ";
        }
        cout<<"}";
}
void set_intesection(Set &set1,Set &set2)
{
        int i=0,j=0;
        while(i<set2.size)
        {
                 if(compare(set2.set[i],set1))
                 {
                          this->set[j]=set2.set[i];
                         j++;
                 }
                 i++;
        }
        size=j;
}
void set_union(Set &set1,Set &set2)
{
        int j;
        for(j=0;j<set1.size;j++)</pre>
```

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```
set[j]=set1.set[j];
        int i=0;
        while(i<set2.size)
        {
                 if(!compare(set2.set[i],set1))
                 {
                         this->set[j]=set2.set[i];
                         j++;
                 }
                 i++;
        }
        size=j;
}
void set_diff(Set &set1,Set &set2)
{
        int i=0,j=0;
        while(i<set1.size)
        {
                 if(!compare(set1.set[i],set2))
                 {
                         this->set[j]=set1.set[i];
                         j++;
                 }
                 i++;
        }
        size=j;
}
int subset(Set &s)
{
        int i=0,status=1;
        while(i<s.size)
```

```
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                {
                        if(!compare(s.set[i],*this))
                        {
                                 status=0;
                                 break;
                        }
                        else
                                 status=1;
                        i++;
                }
                return status;
        }
        int compare(int e,Set &s)
        {
                for(int i=0;i<s.size;i++)</pre>
                {
                        if(e==s.set[i])
                                 return 1;
                }
                return 0;
        }
};
int main()
{
        Set inter(20),BmA(20),AmB(20),un(30);
        int size1,size2,ch,status,e;
        char set_ch;
        char tryAgain;
        cout<<"Enter the size of set A: ";
        cin>>size1;
        Set s1(size1);
```

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```
cout<<"Enter the elements of set A: ";
s1.create();
cout<<"Enter the size of set B: ";
cin>>size2;
Set s2(size2);
cout<<"Enter the elements of set B: ";
s2.create();
cout<<"Set A: ";
s1.display();
cout<<"\nSet B: ";
s2.display();
do{
       cout<<"\n1.Insert";
       cout<<"\n2.Remove";
       cout<<"\n3.Contains";
       cout<<"\n4.Size";
       cout<<"\n5.Intersection";</pre>
       cout<<"\n6.Union";
       cout<<"\n7.A-B";
       cout<<"\n8.B-A";
       cout<<"\n9.Check subset";
       cout<<"\n10.Display Sets.";
       cout<<"\n11.Exit.";
       cout<<"\nEnter your choice: ";
       cin>>ch;
       switch(ch)
       {
       case 1:
              cout<<"Which set you want to enter the element in?(A/B): ";
              cin>>set_ch;
              cout<<"Enter element to be inserted: ";
               cin>>e;
```

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```
if(set_ch=='A'||set_ch=='a')
        {
                s1.insert(e);
                cout<<"\nSet A: ";
                s1.display();
        }
        else
        {
                s2.insert(e);
                cout<<"\nSet B: ";
                s2.display();
        }
break;
case 2:
        cout<<"Which set you want to delete the element from?(A/B): ";
        cin>>set_ch;
        cout<<"Enter element to be deleted: ";
        cin>>e;
        if(set_ch=='A'||set_ch=='a')
        {
                s1.remove(e);
                cout<<"\nSet A: ";
                s1.display();
        }
        else
        {
                s2.remove(e);
                cout<<"\nSet B: ";
                s2.display();
        }
break;
case 3:
```

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```
int pos;
        cout<<"In which set you want check?(A/B): ";
        cin>>set_ch;
        cout<<"Enter element to be checked: ";
        cin>>e;
        if(set_ch=='A'||set_ch=='a')
        {
                pos=s1.contains(e);
                if(pos!=0)
                {
                         cout<<"\nElement found at position: "<<pos;</pre>
                         cout<<"\nSet A: ";
                         s1.display();
                         break;
                }
        }
        else
        {
                pos=s2.contains(e);
                if(pos!=0)
                {
                         cout<<"\nElement found at position: "<<pos;</pre>
                         cout<<"\nSet B: ";
                         s2.display();
                         break;
                }
        cout<<"\nSorry,element not found...";</pre>
break;
case 4:
        cout<<"\nSet A: ";
        s1.display();
        size1=s1.set_size();
```

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```
cout<<" with Size: "<<size1;
        cout<<"\nSet B: ";
        s2.display();
        size2=s2.set_size();
        cout<<" with Size: "<<size2;
break;
case 5:
        inter.set_intesection(s1,s2);
        cout<<"\nIntersection: ";</pre>
        inter.display();
break;
case 6:
        un.set_union(s1,s2);
        cout<<"\nUnion: ";
        un.display();
break;
case 7:
        AmB.set_diff(s1,s2);
        cout<<"\nA-B: ";
        AmB.display();
break;
case 8:
        BmA.set_diff(s2,s1);
        cout<<"\nB-A: ";
        BmA.display();
break;
case 9:
        status=s1.subset(s2);
        (status==0)?cout<<"\nB Subset of A: false" :cout<<"\nB Subset of A: true";
```

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```
break;
case 10:

cout<<"Set A: ";
s1.display();
cout<<"\nSet B: ";
s2.display();
break;
case 11:
exit(1);

default: cout<<"\nInvalid Entry...";
}
cout<<"\n\nDo you want to try again? (Y/N)";
cin>>tryAgain;
}while(tryAgain=='Y'||tryAgain=='y');
```

OUTPUT:

}

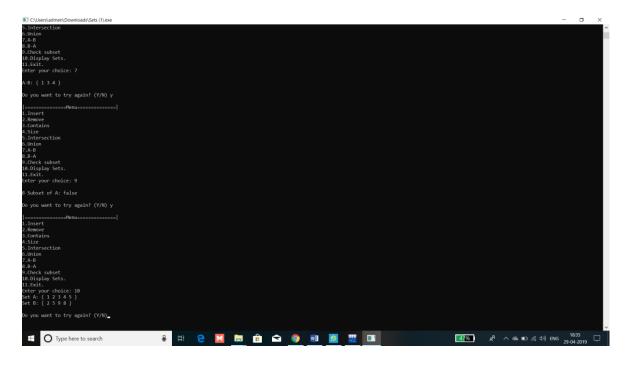
```
| Comparison | Com
```

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CLASS: S.Y.B.TECH-C

BATCH: C3



CONCLUSION:

Hence, we implemented set identity laws in C++ and learnt about its properties.

ASSIGNMENT 2

AIM: Construct a threaded binary search tree.

OBJECTIVE: Implement a threaded binary search tree by inserting values in the given order and traverse it in inorder traversal using threads.

THEORY: Insertion in Binary threaded tree is similar to insertion in binary tree but we will have to adjust the threads after insertion of each element. The idea of threaded binary trees is to make inorder traversal faster and do it without stack and without recursion. A binary tree is made threaded by making all right child pointers that would normally be NULL point to the inorder successor of the node (if it exists).

ALGORITHM:

Insertion Operation

ROLL NO: 223064 NAME: VEDIKA SADAVARTE. CLASS: S.Y.B.TECH-C BATCH: C3 Step 1: Start Step 2: If the root is null then create root node return Step 3: If the root exists then Comparer the data with root.data Step 4: While insert position is located If data is greater than node.data Go to right subtree Else Go to left subtree End while Step 5: Insert data Step 6: Stop. **Inorder Traversal** Step 1: Start **Step 2:** current=leftmost->root Step 3: While current is not null OR current is not equal to head Print current->data. Step 4: If current->rbit is true then current= current->right else current=leftmost->(current-right) Step 5: Stop. PROGRAM:

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GR NO: 21820001 **ROLL NO: 223064** NAME: VEDIKA SADAVARTE. CLASS: S.Y.B.TECH-C BATCH: C3 #include<iostream> using namespace std; struct Node { struct Node *left, *right; int info; int rf, lf=0; **}**; struct Node *insert(struct Node *root, int key) { Node *temp = root; Node *par = NULL; while (temp!= NULL) { if (key == (temp->info)) { cout<<"Duplicate Key !\n"; return root; } par = temp; if (key < temp->info) if (temp -> lf == 0)temp = temp -> left; else

break;

GR NO: 21820001 **ROLL NO: 223064** NAME: VEDIKA SADAVARTE. CLASS: S.Y.B.TECH-C BATCH: C3 } else { if (temp->rf == 0)temp =temp -> right; else break; } } Node *tmp = new Node; tmp -> info = key; tmp -> lf = 1; $tmp \rightarrow rf = 1;$ if (par == NULL) { root = tmp; tmp -> left = NULL; tmp -> right = NULL; } else if (key < (par -> info)) { tmp -> left = par -> left; tmp -> right = par; par \rightarrow If = 0; par -> left = tmp; }

```
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else
tmp -> left = par;
tmp -> right = par -> right;
par \rightarrow rf = 0;
par -> right = tmp;
}
return root;
}
struct Node *inorderSuccessor(struct Node *ptr)
{
if (ptr -> rf == 1)
return ptr->right;
ptr = ptr -> right;
while (ptr -> If== 0)
ptr = ptr -> left;
return ptr;
}
void inorder(struct Node *root)
{
if (root == NULL)
cout<<"Tree is empty";
struct Node *ptr = root;
while (ptr \rightarrow lf == 0)
ptr = ptr -> left;
```

SKILL DEVLOPMENT LAB-II-2018-19 GR NO: 21820001 **ROLL NO: 223064** NAME: VEDIKA SADAVARTE. CLASS: S.Y.B.TECH-C BATCH: C3 while (ptr != NULL) cout<<ptr -> info<<"\t"; ptr = inorderSuccessor(ptr); } } int main() struct Node *root = NULL; root = insert(root, 20); root = insert(root, 10); root = insert(root, 30); root = insert(root, 5); root = insert(root, 16); root = insert(root, 14); root = insert(root, 17); root = insert(root, 13); inorder(root); return 0; }

OUTPUT:

5 10 13 14 16 17 20 30

CONCLUSION: Successfully implemented a TBT, inserted values in given order and traversed it in inorder traversal using threads.

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Assignment No: 3

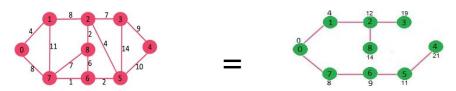
<u>Aim:</u> There are flight paths between cities. If there is a flight between city A and city B then there is an edge between the cities. The cost of the edge can be the time that flight takes to reach city B from A, or the amount of fuel used for the journey. Represent this as a graph. The node can be represented by airport name or name of the city. Use adjacency list representation of the graph or use adjacency matrix representation of the graph. Justify the storage representations used.

Objectives: To understand the various operations on graphs.

Theory:

Dijkstra's algorithm is very similar to Prim's algorithm for minimum spanning tree. Like Prim's MST, we generate a shortest path tree with given source as root. We maintain two sets, one set contains vertices included in shortest path tree, and other set includes vertices not yet included in shortest path tree.

For Example:



At every step of the algorithm, we find a vertex which is in the other set and has a minimum distance from the source. Below are the detailed steps used in Dijkstra's algorithm to find the shortest path from a single source vertex to all other vertices in the given graph.

Algorithm:

- 1. Create priority queue pq
- 2. Enqueue(pq,s)
- 3. For(i=1; i <= g > v; i++)

Distance[i]=-1

4. Distance[s]=0

```
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    5. while(!isemptyqueue(pq))
    {
       5.1v=deletemin(pq);
       5.2 for all adjacent vertices w to v
  {
  Compute new distance d=distance[v]+weight[v][w];
               If(Distance[w]==-1)
               {
                       Distance[w]=new distance d;
Insert w in priority queue with priority d
Path[w]=v}
If(Distance[w]>newdistance d)
       distance[w]=new disance d;
Update priority Of vertex w to be d;
Path[w]=v;
}}}
Program:
#include<iostream>
#define MAX 20
using namespace std;
class dijkstra
       int city;
       int distance[MAX][MAX];
```

{

}

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int d[MAX];

```
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        int visited[MAX];
public:
        void city_no();
        int minvertex();
        void matrix_fill();
        void dijkstra_code();
        void display();
};
void dijkstra::city_no()
{
        cout<<"\n enter the number of cities (including cities A and B): ";
        cin>>city;
}
int dijkstra::minvertex()
{
        int mvertex=-1;
        for(int i=0;i<city;i++)
                 if(visited[i]==0 && (mvertex==-1 || d[i]<d[mvertex]))
                         mvertex=i;
        return mvertex;
}
void dijkstra::matrix_fill()
        cout<<"\n enter the distances between the cities : ";</pre>
        for(int i=0;i<city;i++)
        {
                 cout<<"\n For city "<<i<endl;
                 for(int j=0;j<city;j++)</pre>
                {
                         if(i==j)
                                  distance[i][j]=0;
                         cin>>distance[i][j];
                }
                 d[i]=INT_MAX;
                 visited[i]=0;
        }
}
```

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void dijkstra::dijkstra_code()

```
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{
        d[0]=0;
        for(int i=0;i<city-1;i++)
                 int mvertex=minvertex();
                 visited[mvertex]=1;
                 for(int j=0;j<city;j++)</pre>
                 {
                          if((distance[mvertex][j]!=0)&&(visited[j]==0))
                          {
                                   int dist=d[mvertex]+distance[mvertex][j];
                                   if(dist<d[j])
                                            d[j]=dist;
                          }
                 }
        }
}
void dijkstra::display()
  cout<<"\n distance of cities from city 0 \n";
  cout<<"city Distance\n";</pre>
        for(int i=0;i<city;i++)</pre>
                 cout << i << "\backslash t" << d[i] << endl;
}
int main()
{
        dijkstra sp;
        sp.city no();
        sp.matrix_fill();
        sp.dijkstra_code();
        sp.display();
        return 0;
}
```

Output:

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BATCH: C3



Conclusion: From above experiment we learnt how to use shortest path algorithm using graph operation.

ASSIGNMENT 4

AIM: To implement prims algorithm for minimum spanning tree.

OBJECTIVE: For a weighted graph G, find the minimum spanning tree using prims algorithm.

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BATCH: C3

THEORY:

In computer science, Prim's (also known as Jarník's) algorithm is a greedy algorithm that finds a minimum spanning tree for a weighted undirected graph. This means it finds a subset of the edges that forms a tree that includes every vertex, where the total weight of all the edges in the tree is minimized. The algorithm operates by building this tree one vertex at a time, from an arbitrary starting vertex, at each step adding the cheapest possible connection from the tree to another vertex.

The algorithm was developed in 1930 by Czech mathematician Vojtěch Jarník and later rediscovered and republished by computer scientists Robert C. Prim in 1957 and Edsger W. Dijkstra in 1959. Therefore, it is also sometimes called the Jarník's algorithm, Prim—Jarník algorithm, Prim—Dijkstra algorithm or the DJP algorithm.

ALGORITHM:

- 1) Create a set *mstSet* that keeps track of vertices already included in MST.
- **2)** Assign a key value to all vertices in the input graph. Initialize all key values as INFINITE. Assign key value as 0 for the first vertex so that it is picked first.
- 3) While mstSet doesn't include all vertices
-a) Pick a vertex u which is not there in mstSet and has minimum key value.
-**b)** Include *u* to mstSet.
-c) Update key value of all adjacent vertices of u. To update the key values, iterate through all adjacent vertices. For every adjacent vertex v, if weight of edge u-v is less than the previous key value of v, update the key value as weight of u-v

The idea of using key values is to pick the minimum weight edge from cut. The key values are used only for vertices which are not yet included in MST, the key value for these vertices indicate the minimum weight edges connecting them to the set of vertices included in MST.

PROGRAM:

```
#include <iostream>
using namespace std;

class graph
{
   int a[100][100];
   int v;

public:
   void insert_edge(int n1,int n2,int wt)
   {
     if(n1-1>=v||n2-1>=v)
```

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```
cout<<"Vertex request out of range\n";</pre>
  else
  {
     a[n1-1][n2-1]=wt;
    a[n2-1][n1-1]=wt;
  }
}
void display()
{
  for(int i=0;i<v;i++)
  {
     for(int j=0;j< v;j++)
    {
       cout << a[i][j] << "\t";
     }
    cout<<endl;
  }
}
void update_v(int n)
{
  v=n;
}
void prims(int src)
```

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```
{
  int sp[v],dist[v],visited[v],parent[v],c=0;
  for(int i=0;i<v;i++)
  {
     visited[i]=0;
     dist[i]=9999;
  }
  dist[src-1]=0;
  parent[src-1]=-1;
  for(int i=0;i<v;i++)
  {
     int min=9999,min_ind;
     for(int j=0;j< v;j++)
    {
       if(!visited[j] && dist[j]<min )</pre>
       {
         min=dist[j];
         min_ind=j;
       }
     }
     int U=min_ind;
     visited[U]=1;
```

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};

{

```
sp[c]=U;
      C++;
      for(int V=0;V<v;V++)
      {
         if(!visited[V] && a[U][V] && a[U][V]<dist[V] && dist[U]!=9999)
          {
           parent[V]=U;
           dist[V]=a[U][V];
          }
      }
    }
    for(int i=0;i<c;i++)
    {
      cout<<sp[i]+1<<" link from "<<parent[i]+1<<endl;</pre>
    }
    cout<<endl;
  }
int main()
  char r;
  do
```

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CLASS: S.Y.B.TECH-C

```
graph g;
char op;
int v;
cout<<"Enter number of vertices: ";
cin>>v;
g.update_v(v);
do
{
 int c;
 cout<<"\n=======\n";
 cout<<"1] Insert edge\n2] Increase number of vertices\n3] Display matrix\n4] Find shortest path\n";
 cout<<"_____\n";
 cout<<"Enter your choice: ";
 cin>>c;
 switch(c)
 {
   case 1: {
        int n1,n2,wt;
        cout<<"Enter the nodes between which there is an edge\n";
        cin>>n1>>n2;
        cout<<"Enter weight: ";
        cin>>wt;
        g.insert_edge(n1,n2,wt);
```

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CLASS: S.Y.B.TECH-C

BATCH: C3

```
}
    break;
case 2: {
      int n;
      cout<<"Enter the number by which you wish to increase the vertices: ";
      cin>>n;
      v+=n;
      g.update_v(v);
    }
    break;
case 3: {
      g.display();
    }
    break;
case 4: {
      int src,dst;
      cout<<"Source: ";
      cin>>src;
      g.prims(src);
    }
    break;
default:cout<<"Error 404.....page not found\n";
```

}

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```
cout<<"Do you wish to continue(y/n): ";
    cin>>op;
}while(op=='y' || op=='Y');
cout<<"Test pass(y/n): ";
    cin>>r;
}while(r=='n' || r=='N');
cout<<"***********************\n";
cout<<"* Thank You! *\n";
cout<<"*********************\n";
return 0;</pre>
```

OUTPUT:

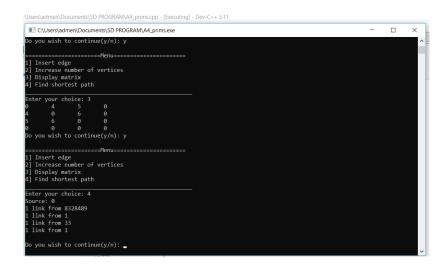
}

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CONCLUSION:

Time Complexity of the above program is $O(V^2)$. If the input graph is represented using adjacency list, then the time complexity of Prim's algorithm can be reduced to $O(E \log V)$ with the help of binary heap.

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ASSIGNMENT 5

Aim: You have a business with several offices; you want to lease phone lines to connect them up with each other; and the phone company charges different amounts of money to connect different pairs of cities. You want a set of lines that connects all your offices with a minimum total cost. Solve the problem by suggesting appropriate data structures.

Objective: Understand the problem statement, determine and implement the data structure suitable for solving above real time example.

Theory:

Kruskal's Minimum Spanning Tree:

Kruskal's algorithm is a minimum-spanning-tree **algorithm** which finds an edge of the least possible weight that connects any two trees in the forest. ... This means it finds a subset of the edges that forms a tree that includes every vertex, where the total weight of all the edges in the tree is minimized.

Algorithm:

- KRUSKAL(G):
- A = \emptyset For each vertex $v \in G.V$:
- MAKE-SET(v)
- For each edge $(u, v) \in G.E$ ordered by increasing order by weight(u, v):
- if FIND-SET(u) ≠ FIND-SET(v):
- $A = A \cup \{(u, v)\}$
- UNION(u, v)
- return A

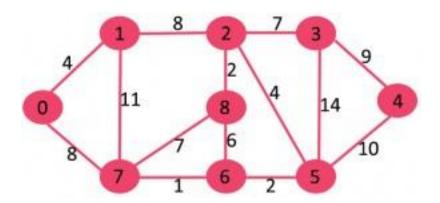
Example:

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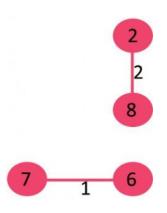


The graph contains 9 vertices and 14 edges. So, the minimum spanning tree formed will be having (9-1) = 8 edges.

Now pick all edges one by one from sorted list of edges **1.** *Pick edge 7-6:* No cycle is formed, include it.



2. Pick edge 8-2: No cycle is formed, include it.



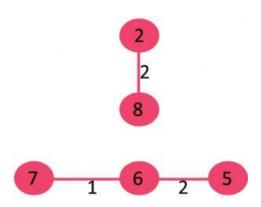
3. Pick edge 6-5: No cycle is formed, include it.

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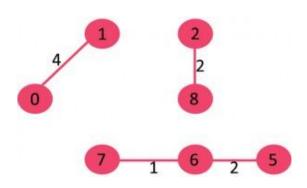
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CLASS: S.Y.B.TECH-C

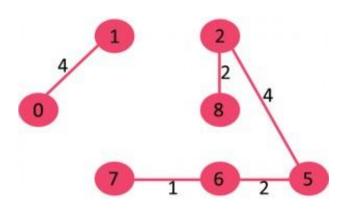
BATCH: C3



4. Pick edge 0-1: No cycle is formed, include it.



5. Pick edge 2-5: No cycle is formed, include it.



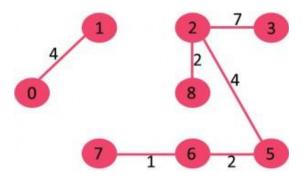
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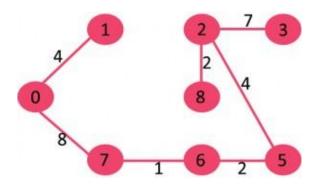
CLASS: S.Y.B.TECH-C

BATCH: C3

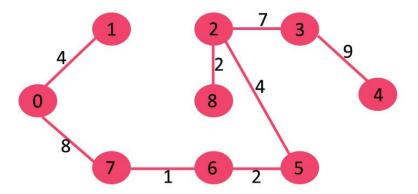
- **6.** Pick edge 8-6: Since including this edge results in cycle, discard it.
- 7. Pick edge 2-3: No cycle is formed, include it.



- 8. Pick edge 7-8: Since including this edge results in cycle, discard it.
- 9. Pick edge 0-7: No cycle is formed, include it.



- 10. Pick edge 1-2: Since including this edge results in cycle, discard it.
- **11.** *Pick edge 3-4:* No cycle is formed, include it.



Since the number of edges included equals (V-1), the algorithm stops here.

SKILL DEVLOPMENT LAB-II-2018-19 GR NO: 21820001 **ROLL NO: 223064** NAME: VEDIKA SADAVARTE. CLASS: S.Y.B.TECH-C BATCH: C3 Program: include<iostream> using namespace std; #define MAX 30 typedef struct edge { int u,v,w; }edge; typedef struct edgelist { edge data[MAX]; int count; }edgelist; edgelist elist; int G[MAX][MAX],n; edgelist spanlist;

void kruskal();

```
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NAME: VEDIKA SADAVARTE.
CLASS: S.Y.B.TECH-C
BATCH: C3
int find(int belongs[],int vertexno);
void union1(int belongs[],int c1,int c2);
void sort();
void print();
int main()
{
  int i,j;
  cout<<"\nEnter number of city's:";
  cin>>n;
cout<<"\nEnter the adjacency matrix of city ID's:\n";</pre>
for(i=0;i<n;i++)
    for(j=0;j<n;j++)
       cin>>G[i][j];
  kruskal();
  print();
}
void kruskal()
{
  int belongs[MAX],i,j,cno1,cno2;
```

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```
elist.count=0;
for(i=1;i<n;i++)
  for(j=0;j<i;j++)
  {
     if(G[i][j]!=0)
     {
       elist.data[elist.count].u=i;
       elist.data[elist.count].v=j;
       elist.data[elist.count].w=G[i][j];
       elist.count++;
     }
  }
sort();
for(i=0;i<n;i++)
  belongs[i]=i;
spanlist.count=0;
for(i=0;i<elist.count;i++)</pre>
  cno1=find(belongs,elist.data[i].u);
```

GR NO: 21820001 **ROLL NO: 223064** NAME: VEDIKA SADAVARTE. CLASS: S.Y.B.TECH-C BATCH: C3 cno2=find(belongs,elist.data[i].v); if(cno1!=cno2) { spanlist.data[spanlist.count]=elist.data[i]; spanlist.count=spanlist.count+1; union1(belongs,cno1,cno2); } } } int find(int belongs[],int vertexno) { return(belongs[vertexno]); } void union1(int belongs[],int c1,int c2) { int i; for(i=0;i<n;i++) if(belongs[i]==c2) belongs[i]=c1;

}

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CLASS: S.Y.B.TECH-C

```
void sort()
{
  int i,j;
  edge temp;
  for(i=1;i<elist.count;i++)</pre>
    for(j=0;j<elist.count-1;j++)</pre>
       if(elist.data[j].w>elist.data[j+1].w)
       {
         temp=elist.data[j];
         elist.data[j]=elist.data[j+1];
         elist.data[j+1]=temp;
       }
}
void print()
{
  int i,cost=0;
  for(i=0;i<spanlist.count;i++)</pre>
  {
    cout<<"\n"<<spanlist.data[i].v<<" "<<spanlist.data[i].v;
    cost=cost+spanlist.data[i].w;
```

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NAME: VEDIKA SADAVARTE.

CLASS: S.Y.B.TECH-C

BATCH: C3

}

cout<<"\n\nMinimum cost of the telephone lines between the cities:"<<cost<<"\n";

}

Output:

Enter number of city's:6

Enter the adjacency matrix of city ID's:

031600

305030

150564

605002

036006

 $0\,0\,4\,2\,6\,0$

2 0 1

5 3 2

1 0 3

4 1 3

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BATCH: C3

5 2 4

Minimum cost of the telephone lines between the cities:13

CONCLUSION:

Thus, we learnt about **Kruskal's algorithm** is a minimum-spanning-tree **algorithm** which finds an edge of the least possible weight that connects any two trees in the forest. And its application for traversing each node with the shortest distance to reach it.

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NAME: VEDIKA SADAVARTE.

CLASS: S.Y.B.TECH-C

BATCH: C3

Assignment No. 6

Aim:

Read the marks obtained by students of second year in an online examination of particular subject. Find out maximum and minimum marks obtained in that subject using heap data structure.

Objective:

Understand the heap data structure for heap sorting.

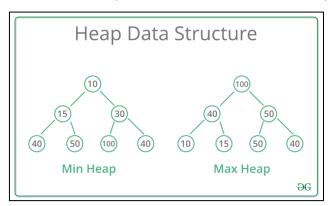
Theory:

Heap:

A Heap is a special Tree-based data structure in which the tree is a complete binary tree. Generally, Heaps can be of two types:

Max-Heap: In a Max-Heap the key present at the root node must be greatest among the keys present at all of its children. The same property must be recursively true for all sub-trees in that Binary Tree.

Min-Heap: In a Min-Heap the key present at the root node must be minimum among the keys present at all of its children. The same property must be recursively true for all sub-trees in that Binary Tree.



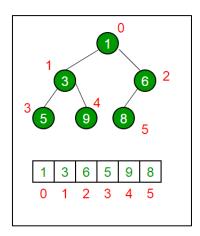
GR NO: 21820001 ROLL NO: 223064

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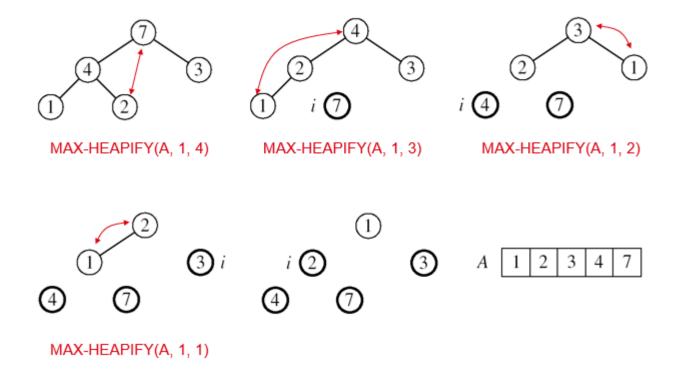
CLASS: S.Y.B.TECH-C

BATCH: C3

Since a Binary Heap is a Complete Binary Tree, it can be easily represented as array and array based representation is space efficient. If the parent node is stored at index I, the left child can be calculated by 2 * I + 1 and right child by 2 * I + 2 (assuming the indexing starts at 0).



EXAMPLE: A = [7, 4, 3, 1, 2]



Algorithm:

1. MAX-HEAPIFY(A, i, n)

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NAME: VEDIKA SADAVARTE.

CLASS: S.Y.B.TECH-C

BATCH: C3

```
I \leftarrow LEFT(i)

r \leftarrow RIGHT(i)

if I \le n and A[I] > A[i]

then largest \leftarrow I

else largest \leftarrow i

if r \le n and A[r] > A[largest]

then largest \leftarrow r

if largest != i

then exchange A[i] \leftrightarrow A[largest]

MAX-HEAPIFY(A, largest, n)
```

2. BUILD-MAX-HEAP(A)

```
n = length[A]
for i \leftarrow floor(n/2) downto 1
do MAX-HEAPIFY(A, i, n)
```

3. HEAPSORT(A)

```
BUILD-MAX-HEAP(A)
for i \leftarrow length[A] downto 2
do exchange A[1] \leftrightarrow A[i]
MAX-HEAPIFY(A, 1, i - 1)
```

Code:

```
#include <iostream>
#include <math.h>
using namespace std;

void max_heapify(int *arr, int i, int n) {
    int left_num = (2*i);
    int right_num = (2*i) + 1;
    int largest = i;

if(left_num <= n && arr[left_num] > arr[i]) {
```

```
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CLASS: S.Y.B.TECH-C
BATCH: C3
                largest = left_num;
        } else {
                largest = i;
        if(right_num <= n && arr[right_num] > arr[largest]) {
                largest = right num;
        }
        if(largest != i) {
                int temp = arr[i];
                arr[i] = arr[largest];
                arr[largest] = temp;
                max_heapify(arr, largest, n);
        }
}
void build_heap(int *arr, int n) {
        for(int i=floor(n/2); i>=0; i--) {
                max_heapify(arr,i,n);
        }
}
int main() {
        cout << "Enter no. of students: ";
        cin >> n;
        int *arr = new int[n];
        cout << "Enter Marks of M3: \n";
        for(int i=1;i<=n;i++) {
                cin >> arr[i];
        }
        //HEAP SORT
        build_heap(arr,n);
        for(int i=n;i>=1;i--) {
                int temp = arr[1];
                arr[1] = arr[i];
                arr[i] = temp;
                max_heapify(arr,1,i-1);
        }
        cout << "Maximum Marks Obtained = " << arr[n] << endl;</pre>
```

```
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BATCH: C3

cout << "Minimum Marks Obtained = " << arr[1] << endl;
return 0;
}
```

Output:

```
Enter no. of students : 6
Enter Marks of M3 :
20
28
29
17
15
22
Maximum Marks Obtained = 29
Minimum Marks Obtained = 15
```

Conclusion:

In this assignment, we have learnt the usage of heap data structure for heap sorting.

ASSIGNMENT NO 7

AIM: Insert the keys into a hash table of length m using open addressing using double hashing with $h(k)=1+(k \mod(m-1))$.

OBJECTIVE: To understand the concept of double hashing and insert the keys into hash table.

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BATCH: C3

THEORY: Double hashing is a computer programming technique used in hash tables to resolve hash collisions, in cases when two different values to be searched for produce the same hash key. It is a popular collision-resolution technique in open-addressed hash tables. Double hashing is implemented in many popular libraries.

Like linear probing, it uses one hash value as a starting point and then repeatedly steps forward an interval until the desired value is located, an empty location is reached, or the entire table has been searched; but this interval is decided using a second, independent hash function (hence the name double hashing). Unlike linear probing and quadratic probing, the interval depends on the data, so that even values mapping to the same location have different bucket sequences; this minimizes repeated collisions and the effects of clustering.

Given two randomly, uniformly, and independently selected hash functions h_1 and h_2 , the ith location in the bucket sequence for value k in a hash table T is: $h(i,k)=(h_1(k)+i.h_2(k))\mod[T]$. Generally, h_1 and h_2 are selected from a set of universal hash functions

Double Hashing Example:

Insert 89, 18, 49, 58, 69 in the table

Hash1(key) = key mod 10

Hash2(key) = 7 - (key mod 7)H(89) = $89 \mod 10 = 9$ H(18) = $18 \mod 10 = 8$ H(49) = $49 \mod 10 = 9 \rightarrow \text{collision}$ = $7 - (49 \mod 7) = 7$ H(58) = $58 \mod 10 = 8 \rightarrow \text{collision}$ = $7 - (58 \mod 7) = 5$ H(69) = $69 \mod 10 = 9 \rightarrow \text{collision}$ = $7 - (69 \mod 7) = 1$

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CLASS: S.Y.B.TECH-C

BATCH: C3

[0]	
[1]	69
[2]	
[3]	
[4]	
[5]	58
[6]	
[7]	49
[8]	18
[9]	89

ALGORITHM:

- 1. Set indx = H(K); offset = $H_2(K)$
- 2. If table location indx already contains the key, no need to insert it. Done!
- 3. Else if table location indx is empty, insert key there. Done!
- 4. Else collision. Set indx = (indx + offset) mod M.
- 5. If indx == H(K), table is full! (Throw an exception, or enlarge table.) Else go to 2.

CODE:

```
#include <iostream>
using namespace std;
int hashing(int a[],int value,int n)
{
   int i=1;
   int key = value%n;
   if (a[key]==-1)
       a[key]=value;
   else
      {
       while(a[key]!=-1)
```

```
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   {
     key=i*(1+value%(n-1));
     i+=1;
   }
    if (a[key]==-1)
    a[key]=value;
    else
     cout<<" no index to store value "<<endl;
    }
  return 0;
void print(int a[],int n)
{
  cout<<"Index "<<"
                       "<<" Value "<<endl;
  for (int i=0;i<n;i++)
  {
    if (a[i]!=-1)
                    "<<a[i]<<endl;
     cout<<i<
  }
}
```

}

```
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CLASS: S.Y.B.TECH-C
BATCH: C3
int sizeofset(int a[],int n)
{
 int cnt=0;
  for(int i=0;i<n;i++)
  {
    if (a[i]!=-1)
      cnt+=1;
  }
  return cnt;
}
int searchs(int a[],int svalue,int n)
{
  int i=1;
int key=svalue%n;
if (a[key]==svalue)
  return key;
  else
  {
   while(a[key]!=svalue)
   {
     key=i*(1+svalue%(n-1));
     i+=1;
   }
    if (a[key]== svalue)
```

```
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CLASS: S.Y.B.TECH-C
BATCH: C3
      return key;
    else
    return -1;
    }
}
int main()
{ int value,s,n;
 char choice='y';
 cout<<" enter the size of hashtable "<<endl;
    cin>>n;
    int a[n];
    for (int i=0;i<n;i++)
      a[i]=-1;
while (choice=='y')
{
  cout<<endl<<"...."<<endl;
  cout<<" enter 1 for read number of element and factor present in hashtable "<<endl;
  cout<<" enter 2 for to add value into the hashtable "<<endl;
  cout<<" enter 3 to print the hashtable "<<endl;
  cout<<" enter 4 to search the element "<<endl;
  cout<<" enter 5 to exit "<<endl;
  cin>>s;
  switch(s)
```

```
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CLASS: S.Y.B.TECH-C
BATCH: C3
  {
case 1:
    cout<<" total element in the set are "<<sizeofset(a,n)<<endl;</pre>
    cout<<" Load factor of hashtable is "<<sizeofset(a,n)/n<<endl;
    break;
case 2:
    cout<<endl<<" enter the value to store in hashtable "<<endl;
    cin>>value;
   hashing(a,value,n);
    break;
case 3:
    print(a,n);
    break;
case 4:
  int svalue;
  cout<<" enter the value to search "<<endl;
  cin>>svalue;
   if (searchs(a,svalue,n)!=-1)
    cout<<" value "<<svalue<<" is present at the key "<<searchs(a,svalue,n)<<endl;</pre>
    else
    cout<<" element not present in the hashtable "<<endl;
```

SKILL DEVLOPMENT LAB-II-2018-19 GR NO: 21820001 ROLL NO: 223064 NAME: VEDIKA SADAVARTE. CLASS: S.Y.B.TECH-C BATCH: C3 break; case 5: choice= 'n'; break; } } cout << "Hello world!" << endl; return 0; } **OUTPUT:** /*/ enter the size of hashtable 100 MENU enter 1 for read number of element and factor present in hashtable enter 2 for to add value into the hashtable enter 3 to print the hashtable enter 4 to search the element enter 5 to exit

2

SKILL DEVLOPMENT LAB-II-2018-19 GR NO: 21820001
ROLL NO: 223064
NAME: VEDIKA SADAVARTE. CLASS: S.Y.B.TECH-C
BATCH: C3
enter the value to store in hashtable
26
MENU
enter 1 for read number of element and factor present in hashtable
enter 2 for to add value into the hashtable
enter 3 to print the hashtable
enter 4 to search the element
enter 5 to exit
2
enter the value to store in hashtable
36
MENU
enter 1 for read number of element and factor present in hashtable
enter 2 for to add value into the hashtable
enter 3 to print the hashtable
enter 4 to search the element
enter 5 to exit
2
enter the value to store in hashtable
126

ROLL NO: 223064 NAME: VEDIKA SADAVARTE. CLASS: S.Y.B.TECH-C BATCH: C3 MENU enter 1 for read number of element and factor present in hashtable enter 2 for to add value into the hashtable enter 3 to print the hashtable enter 4 to search the element enter 5 to exit 2 enter the value to store in hashtable 226 MENU enter 1 for read number of element and factor present in hashtable enter 2 for to add value into the hashtable enter 3 to print the hashtable enter 4 to search the element enter 5 to exit 2 enter the value to store in hashtable 136 MENU

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ROLL NO: 223064 NAME: VEDIKA SADAVARTE. CLASS: S.Y.B.TECH-C BATCH: C3 enter 1 for read number of element and factor present in hashtable enter 2 for to add value into the hashtable enter 3 to print the hashtable enter 4 to search the element enter 5 to exit 3 Index Value 26 26 28 126 29 226 36 36 38 136 MENU enter 1 for read number of element and factor present in hashtable enter 2 for to add value into the hashtable enter 3 to print the hashtable enter 4 to search the element enter 5 to exit 4 enter the value to search 226 value 226 is present at the key 29

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/*/

GR NO: 21820001 ROLL NO: 223064

NAME: VEDIKA SADAVARTE.

CLASS: S.Y.B.TECH-C

BATCH: C3

Assignment 8

Aim:

Store information of student in sequential file. It should contain roll no, name, division address. Allow user to add, delete, search information of student.

Objective:

Understand sequential file organization, implement it using C++.

Theory: The simplest organization for a file is sequential. A *sequential file* is a sequence of records. The records may or may not be kept in sorted order in the sequence. In standard C input/output all files are sequential files. A record of a file is not necessarily declared to be of type structure. Although file records are typically of type structure, a file record may also be declared to be of type integer, float, character, or any other C type. All records of a file need not be the same type.

Advantages of sequential file

• It is simple to program and easy to design. • Sequential file is best use if storage space.

Disadvantages of sequential file

- Sequential file is time consuming process.
- It has high data redundancy.
- Random searching is not possible.

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CLASS: S.Y.B.TECH-C

BATCH: C3

Program:

```
#include <iostream>
#include <fstream>
#include <string>
using namespace std;
int main()
{ int k=0;
string line;
string find;
char name[25];
int roll=0;
char d;
string adres;
ofstream myfile,f2;
myfile.open("data.txt",ios::app);
while(k!=4){
cout<<"press 1 for adding data"<<endl;
cout<<"press 2 for update "<<endl;</pre>
cout<<"press 3 for searching"<<endl;</pre>
cin>>k;
if(k==1)
cout<<"enter roll no. "<<endl;
```

GR NO: 21820001 ROLL NO: 223064

NAME: VEDIKA SADAVARTE.

CLASS: S.Y.B.TECH-C

```
cin>>roll;
cout<<"enter Name"<<endl;
 cin>>name;
cout<<"enter division "<<endl;
cin>>d; cout<<"enter adre."<<endl;
cin>>adres; myfile<<roll<<" ";
myfile<<name<<" ";
myfile<<d<<" ";
myfile<<adres<<endl;
}
if(k==2)
cout<<"Enter roll no, you want to update "<<endl;
cin>>find;
ifstream file,f1;
file.open("data.txt",ios::in);
file>>line;
while (!file.eof() && line!=find)
{
getline(file,line);
if(find!=line)
```

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{
cout<<"deaTAILS NOT FOUND";
}
else{
cout<<"enter Name"<<endl;
cin>>name;
cout<<"enter dividion "<<endl;
cin>>d;
cout<<"enter address"<<endl;
cin>>adres; myfile<<find;
myfile<<name;
myfile<<d;
myfile<<adres<<endl;
}}
if(k==3)
{
cout<<"Enter roll no, you want to update "<<endl;
cin>>find;
ifstream file;
file.open("data.txt",ios::in);
file>>line;
while (!file.eof() && line!=find)
{
```

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ROLL NO: 223064 NAME: VEDIKA SADAVARTE. CLASS: S.Y.B.TECH-C BATCH: C3 getline(file,line); file>>line; } if(find!=line) { cout<<"deaTAILS NOT FOUND"; } else { cout<<"roll.no"<<endl; cout<endl; file>>line; cout<endl; file>>line; cout<endl; file>>line; cout<<line<<endl; }</pre> } } return 0; }

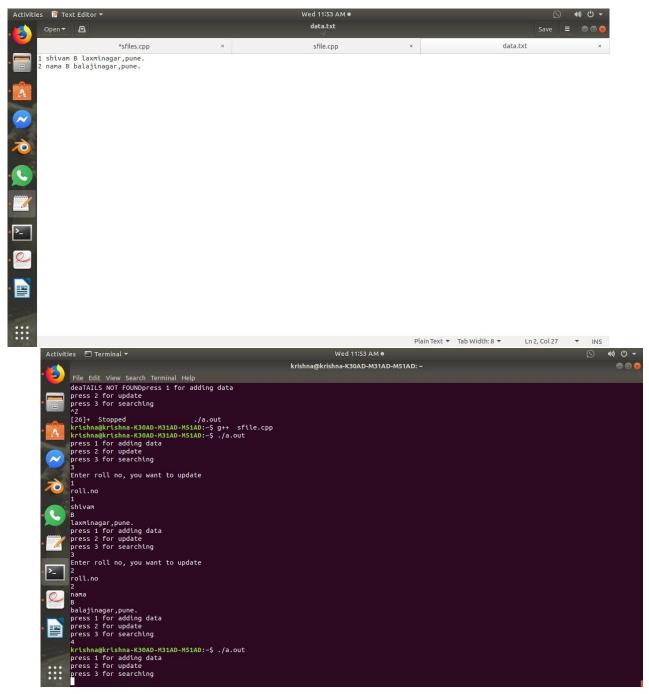
Output:

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Conclusion: In this assignment, we have learnt the sequential file organization.

ASSIGNMENT 9

GR NO: 21820001 ROLL NO: 223064

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BATCH: C3

Title:

Index Sequential File

Problem Statement:

Department maintains a employee information. The file contains employee ID, name, designation and salary. Allow user to add, delete information of employee. Display information of particular employee. If employee does not exist an appropriate message is displayed. If it is, then the system displays the employee details. Use index sequential file to main the data.

Objective:

To make use of index sequential files to maintain and operation on data.

Software And Hardware Requirement:

64-bit Open source Linux or its derivative.

Open Source C++ Programming tool like G++/GCC.

Theory:

Index Sequential File:

This is basically a mixture of sequential and indexed file organization techniques. Records are held in sequential order and can be accessed randomly through an index. Thus, these files share the merits of both systems enabling sequential or direct access to the data.

The index to these files operates by storing the highest record key in given cylinders and tracks. Note how this organization gives the index a tree structure. Obviously this type of file organization will require a direct access device, such as a hard disk.

Indexed sequential file organization is very useful where records are often retrieved randomly and are also processed in (sequential) key order. Banks may use this organization for their auto-bank machines i.e. customers randomly access their accounts throughout the day and at the end of the day the banks can update the whole file sequentially.

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Advantages of Indexed Sequential Files:

- 1. Allows records to be accessed directly or sequentially.
- 2. Direct access ability provides vastly superior (average) access times.

Disadvantages of Indexed Sequential Files:

- 1. The fact that several tables must be stored for the index makes for a considerable storage overhead.
- 2 .As the items are stored in a sequential fashion this adds complexity to the addition/deletion of records. Because frequent updating can be very inefficient, especially for large files, batch updates are often performed.

PROGRAM:

```
#include <iostream>
#include<fstream>
#include<string>
using namespace std;
typedef struct seq_file
{
   int id;
   char name[20],desg[20];
   long int sal;
}record;
typedef struct ind_file
{
   int id;
```

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```
}index;
class file
{
  record data;
  index info;
public:
  void get_data()
  {
    cout<<"Enter id: ";
    cin>>data.id;
    cout<<"Enter name: ";
    cin>>data.name;
    cout<<"Enter designation: ";
    cin>>data.desg;
    cout<<"Enter salary: ";
    cin>>data.sal;
    info.id=data.id;
  }
  void add()
  {
    fstream out1;
    fstream out2;
    out1.open("pos.txt",ios::app);
    out2.open("rec.txt",ios::app);
```

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```
get_data();
  out2.write((char*)&data,sizeof(data));
  out1.write((char*)&info,sizeof(info));
  out1.close();
  out2.close();
}
void search_rec(int id)
{
  int pos=0,loc=-1;
  fstream out1;
  fstream out2;
  out1.open("pos.txt");
  out2.open("rec.txt");
  loc=sizeof(info)*pos;
  out2.seekg(loc,ios::beg);
  for(pos=0;out2.read((char*)&info,sizeof(info));pos++)
  {
    loc=sizeof(info)*pos;
    out2.seekg(loc,ios::beg);
    out2.read((char*)&info,sizeof(info));
    if(info.id==id)
      break;
    }
```

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};

```
}
    if(loc==-1)
    {
      cout<<"Record not found\n";</pre>
    }
    else
    {
      pos--;
      pos=sizeof(data)*pos;
      out1.seekg(pos,ios::beg);
      out1.read((char*)&data,sizeof(data));
      cout<<"Record found\n";</pre>
      cout<<data.id<<"\t"<<data.name<<"\t"<<data.desg<<"\t"<<data.sal<<endl;
    }
    out1.close();
    out2.close();
  }
int main()
  char r;
  do
  {
    char op;
```

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```
file f;
do
{
 int c;
 cout<<"\n========\n";
 cout<<"1] Add record\n2] Search record\n3] Delete record\n";</pre>
 cout<<"____
                         \n";
 cout<<"Enter your choice: ";
 cin>>c;
 switch(c)
 {
   case 1: {
        f.add();
      }
      break;
   case 2: {
        int id;
        cout<<"Enter id to search: ";
        cin>>id;
        f.search_rec(id);
      }
      break;
   case 3: {
```

```
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```

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```
}
           break;
      case 4: {
           }
           break;
      default:cout<<"Error 404.....page not found\n";
    }
    cout<<"Do you wish to continue(y/n): ";</pre>
    cin>>op;
  }while(op=='y' | | op=='Y');
  cout<<"Test pass(y/n): ";</pre>
  cin>>r;
}while(r=='n' || r=='N');
cout<<"***************\n";
cout<<"* Thank You! *\n";</pre>
cout<<"****************\n";
return 0;
```

OUTPUT:

}

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■ Select C:\Users\admen\Documents\SD PROGRAM\A9_empsq.exe X] Add record Search record 3] Delete record Enter your choice: 1 Enter id: 12345 nter name: vedika Enter designation: ceo Enter salary: 250000 Do you wish to continue(y/n): y -----Menu-----1] Add record Search record B] Delete record Enter your choice: 1 Enter id: 2354 Enter name: pooja Enter designation: vp Enter salary: 260000 Do you wish to continue(y/n): y -----Menu----l] Add record Search record Delete record

Conclusion:

In above assignment, we made the use of index sequential files to operate on employee data.

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