

VIII®

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OPERATING SYSTEMS

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Processing large amounts of data in an operating system involves several techniques and algorithms. Here are some key approaches and well-known algorithms:

- Memory Management: The OS manages memory efficiently to ensure that data can be loaded and processed. This includes techniques like paging and segmentation.
- File Systems: The OS uses file systems to organise and store large volumes of data. Popular file systems include NTFS, ext4, and HFS+.
- I/O Scheduling: Algorithms like the C-SCAN, FCFS, and SSTF are used to schedule disk I/O requests, optimising data access.
- Buffering: Buffering involves using a buffer to temporarily hold data, reducing the frequency of I/O operations. Techniques like double buffering and circular buffering are used.
- Caching: Caching algorithms like LRU (Least Recently Used) and LFU (Least Frequently Used) help improve data access times by keeping frequently used data in memory.
- Data Compression: OS can use compression algorithms like Lempel-Ziv-Welch (LZW) or Deflate to reduce the size of data for storage and transmission.
- Parallel Processing: Multithreading and multiprocessing techniques allow the OS to distribute data processing tasks across multiple cores or processors.
- Virtual Memory: Virtual memory techniques, including demand paging and page replacement algorithms like LRU and FIFO, are used to efficiently manage memory for large datasets.
- Data Indexing: For searching and retrieval of data, indexing algorithms like B-trees, Hashing, and various search algorithms are employed.
- Distributed File Systems: In a networked environment, distributed file systems like Hadoop HDFS and Google File System (GFS) are used to process and store large datasets across multiple machines.
- Data Partitioning: Large datasets can be divided into smaller partitions, and MapReduce-style algorithms are used for parallel processing, with frameworks like Hadoop.
- Database Management Systems: When dealing with large datasets in a database, DBMS systems like MySQL, PostgreSQL, and NoSQL databases are employed for efficient data retrieval and management.

• Streaming Algorithms: For processing data streams, algorithms like Count-Min Sketch and HyperLogLog are used to estimate statistics from large data streams in realtime.

Operating systems use scheduling algorithms to manage the execution of multiple processes or threads efficiently. The scheduling algorithm's primary goal is to allocate CPU time to different tasks or processes while maximising system throughput, fairness, and responsiveness. Here are some common scheduling algorithms used by operating systems:

• First-Come, First-Served (FCFS):

- It schedules processes based on their arrival time.
 The process that arrives first gets executed first.
- It is simple but can lead to poor utilisation of CPU time, especially if long processes arrive early (known as the "convoy effect").

• Shortest Job First (SJF):

- Prioritises processes with the smallest CPU burst or execution time.
- It minimises average waiting time but requires knowledge of each process's execution time, which may not always be available.

• Round Robin (RR):

- - Allocates fixed time slices (time quantum) to each process in a circular manner.
- Provides fairness and prevents a single long-running process from hogging the CPU.
- Shorter time slices lead to more frequent context switches.

• Priority Scheduling:

- Assigns a priority to each process, and the CPU is allocated to the highest priority process.
- - Can be either preemptive (higher priority process can interrupt a lower one) or non-preemptive.

• Multilevel Queue Scheduling:

- Divides the ready queue into multiple priority queues, each with its own scheduling algorithm.
- Higher-priority queues receive more CPU time, and lower-priority queues are serviced only when the higherpriority queues are empty.

• Multilevel Feedback Queue Scheduling:

- A variation of multilevel queue scheduling that allows processes to move between queues based on their behaviour (e.g., aging, I/O-bound vs. CPU-bound).
- Provides adaptability and dynamic priority adjustments.

• Lottery Scheduling:

- Allocates "lottery tickets" to processes, and a random drawing determines the process to execute.
- Provides a way to implement probabilistic scheduling and fairness.

• Completely Fair Scheduler (CFS):

- Used in Linux, it aims to divide CPU time fairly among processes, using a red-black tree to maintain the execution order.
- Achieves fairness and responsiveness in dynamic workloads.

• Earliest Deadline First (EDF):

- - Used in real-time systems, processes are scheduled based on their deadlines.
- Ensures that processes with the earliest deadlines get CPU time, critical for real-time applications.

Ubuntu, like many other Linux distributions, uses the Completely Fair Scheduler (CFS) as its default process scheduler. The CFS is part of the Linux kernel's scheduler and aims to provide a fair and efficient allocation of CPU time to processes. It uses a red-black tree to maintain the execution order and strives to divide CPU time fairly among processes, ensuring that no process hogs the CPU.

Simulation for CFS:

Code:

```
1. #include <stdio.h>
2. #include <stdlib.h>
3. #include <pthread.h>
4. #include <unistd.h>
5. #include <time.h>
6. #include <sched.h>
7.
8. #define NUM_THREADS 40
9. #define DATA_FILE "data.txt"
10. #define RESULT_FILE "results.txt"
11.
12. void *thread_function(void *arg) {
13.  int thread_id = *(int *)arg;
```

```
int sum = 0;
14.
15.
       // Open the data file for reading
16.
       FILE *data file = fopen(DATA FILE, "r");
17.
       if (data file == NULL) {
18.
           perror ("Failed to open data file");
19.
           pthread exit(NULL);
20.
       }
21.
22.
23.
       // Read integers from the data file and calculate the sum
       int num;
24.
25.
       while (fscanf(data file, "%d", &num) == 1) {
26.
            sum += num;
27.
28.
29.
       fclose(data file);
30.
       // Get the core number where the thread is running
31.
       int core num = sched getcpu();
32.
33.
34.
       // Open the results file for writing
       FILE *results file = fopen(RESULT FILE, "a");
35.
       if (results file == NULL) {
36.
37.
           perror ("Failed to open results file");
           pthread exit(NULL);
38.
39.
        }
40.
41
       // Write the result, thread ID, and core number to the
   results file
       fprintf(results file, "Thread %d (Core %d): Sum = %d\n",
   thread id, core num, sum);
       fclose(results file);
43.
44.
       pthread exit(NULL);
45.
46. }
47.
48. int main() {
49.
      pthread t threads[NUM THREADS];
50.
       int thread ids[NUM THREADS];
51.
       clock t start, end;
52.
       double cpu time used;
53.
54.
       start = clock();
55.
       // Open the results file and clear its contents
57.
       FILE *results file = fopen(RESULT FILE, "w");
       if (results file == NULL) {
58.
           perror ("Failed to open results file");
59.
            return 1;
60.
61.
        }
       fclose(results file);
62.
63.
```

```
for (int i = 0; i < NUM THREADS; i++) {
64.
            thread ids[i] = i;
65.
            int result = pthread create(&threads[i], NULL,
   thread function, &thread ids[i]);
            if (result) {
67.
                perror("Thread creation failed");
68.
                return 1;
69.
70.
71.
72.
       for (int i = 0; i < NUM THREADS; i++) {
73.
74.
            pthread join(threads[i], NULL);
75.
76.
       end = clock();
77.
78.
       cpu time used = ((double)(end - start)) / CLOCKS PER SEC;
79.
       printf("All threads have finished. Total execution time: %f
80.
   seconds.\n", cpu time used);
81.
      return 0;
82.
83. }
```

Terminal Window After Execution:

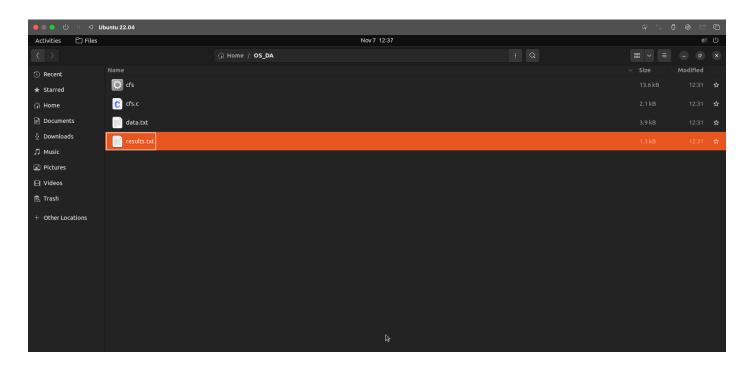
```
ashish_22bai1062@ubuntu:~$ cd OS_DA
ashish_22bai1062@ubuntu:~/OS_DA$ touch cfs.c
ashish_22bai1062@ubuntu:~/OS_DA$ touch data.c
ashish_22bai1062@ubuntu:~/OS_DA$ touch data.txt
ashish_22bai1062@ubuntu:~/OS_DA$ touch results.txt
ashish_22bai1062@ubuntu:~/OS_DA$ gcc cfs.c -o cfs -lpthread -lrt
cfs.c: In function 'thread_function':
cfs.c:32:20: warning: implicit declaration of function 'sched_getcpu'; did you mean
 'sched_getparam'? [-Wimplicit-function-declaration]
            int core_num = sched_getcpu();
   32 I
                           sched getparam
ashish_22bai1062@ubuntu:~/OS_DA$ ./cfs
All threads have finished. Total execution time: 0.011944 seconds.
ashish_22bai1062@ubuntu:~/OS_DA$
                                                 I
```

Data File:

Results File:

```
Thread 0 (Core 2): Sum = 499500
Thread 1 (Core 0): Sum = 499500
Thread 3 (Core 3): Sum = 499500
Thread 5 (Core 3): Sum = 499500
Thread 4 (Core 0): Sum = 499500
Thread 7 (Core 1): Sum = 499500
Thread 8 (Core 2): Sum = 499500
Thread 6 (Core 3): Sum = 499500
Thread 9 (Core 2): Sum = 499500
Thread 12 (Core 3): Sum = 499500
Thread 11 (Core 2): Sum = 499500
Thread 16 (Core 3): Sum = 499500
Thread 13 (Core 3): Sum = 499500
Thread 15 (Core 2): Sum = 499500
Thread 19 (Core 3): Sum = 499500
Thread 18 (Core 2): Sum = 499500
Thread 14 (Core 0): Sum = 499500
Thread 20 (Core 2): Sum = 499500
Thread 10 (Core 0): Sum = 499500
Thread 17 (Core 0): Sum = 499500
Thread 22
          (Core 2): Sum = 499500
Thread 21 (Core 0): Sum = 499500
Thread 23 (Core 3): Sum = 499500
Thread 24 (Core 0): Sum = 499500
Thread 26 (Core 1): Sum = 499500
Thread 29 (Core 0): Sum = 499500
Thread 27 (Core 2): Sum = 499500
Thread 25 (Core 1): Sum = 499500
Thread 28 (Core 0): Sum = 499500
Thread 30 (Core 0): Sum = 499500
Thread 32 (Core 1): Sum = 499500
Thread 33 (Core 0): Sum = 499500
Thread 38 (Core 0): Sum = 499500
Thread 37 (Core 1): Sum = 499500
Thread 31 (Core 3): Sum = 499500
Thread 35 (Core 3): Sum = 499500
Thread 39 (Core 2): Sum = 499500
Thread 34 (Core 3): Sum = 499500
Thread 36 (Core 2): Sum = 499500
```

File System View:



A Real-Time Operating System (RTOS) is a specialized operating system designed for systems and applications that require precise and deterministic timing and rapid response to external events. Unlike general-purpose operating systems (such as Windows, Linux, or macOS), which are optimized for multitasking and may prioritize tasks based on various factors, RTOS is focused on meeting stringent timing requirements and ensuring that tasks are executed within specific time constraints. Here are some key characteristics and features of real-time operating systems:

- 1. **Deterministic Behavior: ** RTOS systems are known for their deterministic behavior. They guarantee that tasks will be executed within predetermined time frames, making them suitable for applications where timing is critical. This predictability is crucial in fields like aerospace, automotive, industrial automation, and medical devices.
- 2. **Prioritization:** Real-time tasks are assigned priorities, and the scheduler ensures that higher-priority tasks are executed before lower-priority tasks. This priority-based scheduling helps meet timing deadlines.

- 3. **Hard Real-Time vs. Soft Real-Time: ** RTOS can be categorized into hard real-time and soft real-time systems. In a hard real-time system, missing a deadline can have catastrophic consequences, while in a soft real-time system, occasional missed deadlines might be acceptable. Hard real-time systems demand more stringent timing guarantees.
- 4. **Task Synchronization and Communication: ** RTOS provides mechanisms for tasks to synchronize and communicate with each other. Common synchronization tools include semaphores, mutexes, and message queues, which ensure that tasks do not interfere with each other.
- 5. **Interrupt Handling:** RTOS is equipped to efficiently handle hardware interrupts and manage real-time tasks in response to these events. This is crucial for systems where immediate response to external events is vital.
- 6. **Minimal Overhead: ** RTOS is designed to have minimal overhead. This means that the operating system itself is lightweight and doesn't introduce significant delays or uncertainty into task execution.
- 7. **Resource Management:** Real-time operating systems manage resources efficiently. This includes memory allocation, CPU utilization, and other system resources to ensure that they are allocated optimally.
- 8. **Embedded Systems: ** RTOS is commonly used in embedded systems, such as microcontrollers, where it is crucial to meet strict timing requirements in applications like automotive control systems, medical devices, robotics, and communication equipment.
- 9. **Task Constraints:** RTOS allows you to define timing constraints for tasks, such as maximum execution time, periodicity, and deadlines, to ensure that the system operates within the required parameters.

10. **Safety-Critical Applications: ** Real-time operating systems are often used in safety-critical applications where reliability and predictability are paramount. This includes systems in the aerospace industry, nuclear reactors, and medical devices.

Some well-known real-time operating systems include:

- **FreeRTOS: ** An open-source RTOS with a large user base and a wide range of supported microcontrollers.
- **VxWorks: ** A commercial real-time operating system used in various critical applications, including aerospace and automotive.
- **QNX: ** Known for its reliability and used in applications like automotive infotainment systems and medical devices.
- **RTOS-32:** An RTOS designed for Windows systems, providing real-time capabilities on Windows platforms.

RTOS plays a crucial role in ensuring that critical systems operate with high reliability and predictability. When choosing an RTOS for a specific application, factors like timing requirements, hardware compatibility, and licensing costs should be carefully considered.

CODE:

```
1. #include <stdio.h>
2. #include "FreeRTOS.h"
3. #include "task.h"
4. #include "semphr.h"
5. #include "esp_timer.h"
6.
7. #define NUM_TASKS 200
8. #define NUM_CORES 4
9. #define DATA_FILE "data.txt"
10. #define RESULT_FILE "results.txt"
11.
12. // Define a mutex to protect the file access
13. SemaphoreHandle_t fileMutex;
14.
```

```
15. void thread function(void *pvParameters) {
       int task id = (int)pvParameters;
16.
       int sum = 0;
17.
18.
19.
       // Record the start time
20.
       int64 t start time = esp timer get time();
21.
       // Open the data file for reading
22.
       FILE *data file = fopen(DATA FILE, "r");
23.
       if (data file == NULL) {
24.
           perror ("Failed to open data file");
25.
26.
           vTaskDelete(NULL);
27.
       }
28.
29.
       // Read integers from the data file and calculate
  the sum
       int num;
30.
      while (fscanf(data_file, "%d", &num) == 1) {
31.
32.
           sum += num;
33.
       }
34.
35.
      fclose(data file);
36.
       // Record the end time
37.
       int64 t end time = esp timer get time();
38.
39.
      // Get the core number where the task is running
40.
41.
       int core num = xPortGetCoreID();
42.
43.
       // Take the mutex to protect file access
       if (xSemaphoreTake(fileMutex, portMAX DELAY) ==
  pdTRUE) {
45.
           // Open the results file for writing
46.
           FILE *results file = fopen(RESULT FILE, "a");
           if (results file == NULL) {
               perror ("Failed to open results file");
48.
               vTaskDelete(NULL);
49.
50.
           }
51.
           // Write the result, task ID, core number, and
52.
  execution time to the results file
           fprintf(results file, "Task %d (Core %d): Sum =
   %d, Start Time = %lld, End Time = %lld\n", task id,
  core num, sum, start time, end time);
           fclose(results file);
```

```
55.
          // Release the mutex
56.
           xSemaphoreGive(fileMutex);
57.
58.
      }
59.
60. vTaskDelete(NULL);
61.}
62.
63. int main() {
      // Initialize the mutex for file access
       fileMutex = xSemaphoreCreateMutex();
65.
66.
67.
      if (fileMutex == NULL) {
68.
           perror("Failed to create mutex");
69.
           return 1;
70.
      }
71.
72.
      // Create FreeRTOS tasks on multiple cores
73.
      TaskHandle t tasks[NUM TASKS];
74.
     for (int i = 0; i < NUM TASKS; i++) {
           int task id = i;
75.
           // Create tasks on different cores (cyclic
76.
  assignment)
           int core num = task id % NUM CORES;
77.
           xTaskCreatePinnedToCore(thread function,
   "Task", configMINIMAL STACK SIZE, (void *) task id,
   tskIDLE PRIORITY + 1, &tasks[i], core num);
79.
       }
80.
81.
      // Start the FreeRTOS scheduler
82.
      vTaskStartScheduler();
83.
84. return 0; // Should never reach here
85.}
```

Thank You!!