

Project 1.

1. GTThreads package

The GTThreads package consists of the following components:

- **gt_bitops** - describes the functions for assigning and resetting the bitmask, as well as the function for determining the lowest bit set (needed to define the highest priority for uthread);
- **gt_signal** - describes the functions for sending signals (SIGVTALRM, etc.), which are used to interrupt the execution of the current thread after a certain time interval;
- **gt_a_spinlock** - describes the function of the spinlock, which allows threads to “wait” while the current thread is executing;
- **gt_tailq** - describes various algorithms for adding (at the head of the queue, at the end of the queue, etc.) and removing uthread from kthread ;
- **gt_pq** - describes functions for finding a suitable uthread to add to the runqueue. It also includes queue and priority structures;
- **gt_kthread** - describes the functions of kthread’s operations (such as creation, initialization, etc.) and also gthread application’s operations over kthreads and uthreads (starting, initialization and termination);
- **gt_uthread** - describes the functions of uthread’s operations (such as creation, initialization, etc.) ;
- **gt_matrix** - the main program for creating and multiplying matrices :
 - gthread_app_init - initializes, allocates memory and creates kthread ;
 - init_matrices - initializes and creates 3 matrices;
 - uthread_create (& utids [inx], uthread _ mulmat , uarg , uarg -> gid) - allocates memory and creates a new uthread , then adds it to the active runqueue for the selected kthread;
 - gthread_app_exit - while there are active uthreads, it executes the uthread_schedule (& sched_find_best_uthread) function . The function's main task is to move the uthread between runqueues and shift active and expires runqueues.

1.1 Implementation of the O(1) priority scheduler:

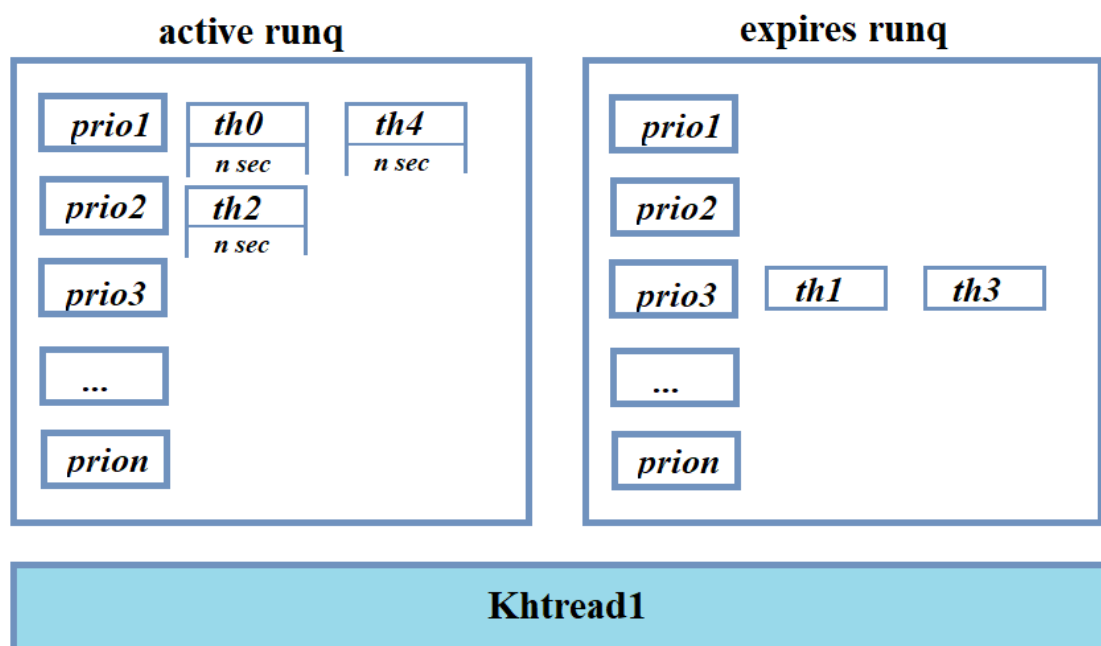


Figure 1. The algorithm of O(1)

Initially, the program creates kthread on the first logical processor using the “kthread_init” function in “gtthread_app_init”. All subsequent kthreads (depending on the number of logical processors) will be its “clones” according to the “kthread_create” function. Each kthread has two runqueues active and expires (Figure 1).

At the time of creation, a uthread is assigned to a specific kthread using the “ksched_find_target” function. Then the uthread is added to the kthread’s active runqueue by the priority (the least bit set corresponds to the highest priority).

Following the function “uthread _ schedule (& sched _ find _ best _ uthread)”, each uthread in the “*Running*” state will be executed until it receives a signal “SIGVTALRM” (timeslice is set in gt _ signal.h). At this time, all other uthreads will receive the SIGUSR1 signal. After receiving a signal, if the uthread is not finished, its state will be changed to “*Runnable*”, and it will be transferred to the expires queue. If the work is completed, then the uthread's state will become “*Done*”, and it will be moved to the zombie runqueue with decreasing the value of “kthread_cur_uthreads”. The “uthread_context_func” function is responsible for changing the uthread state.

If there is no uthread left in the active queue (“uthread_tot” = 0), the active and expires runqueues will be swapped. The “sched_find_best_uthread” function is responsible for changing runqueues. After that, the execution of uthread will start again. The algorithm will repeat until all uthreads have finished executing. A simplified algorithm is shown below (Figure 2).

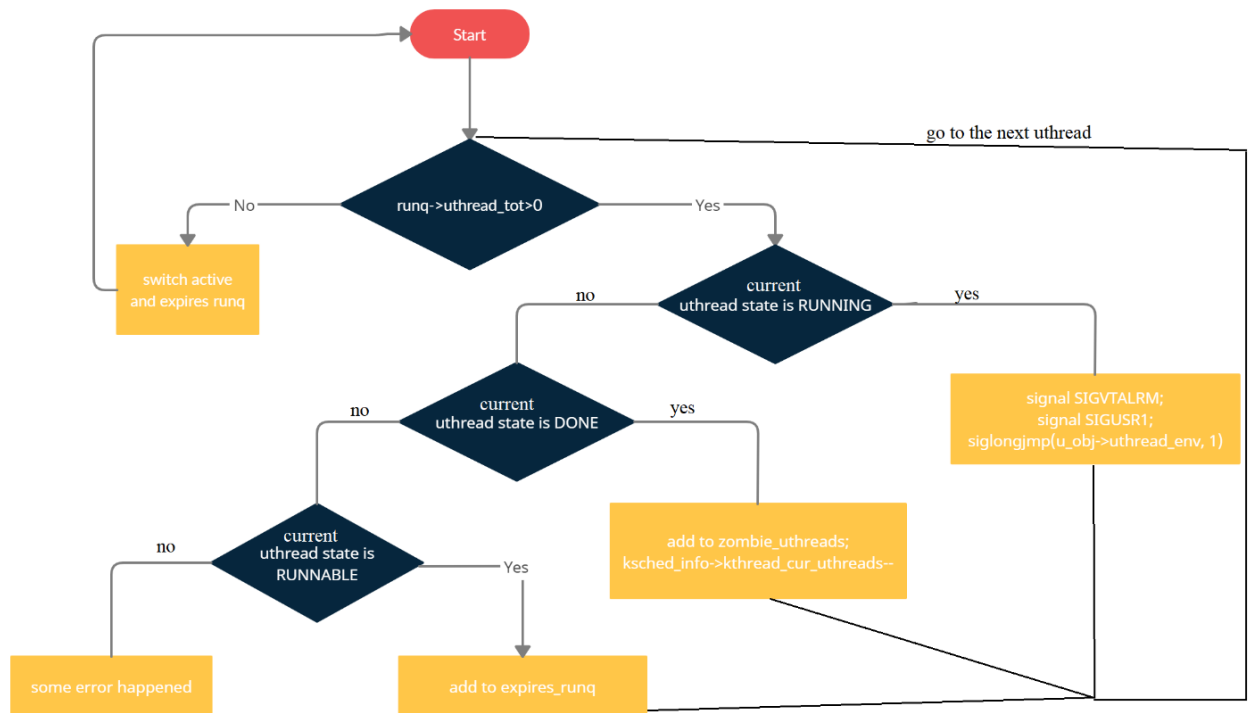


Figure 2. Also the algorithm of O(1)

2. The main ideas of Credit scheduler

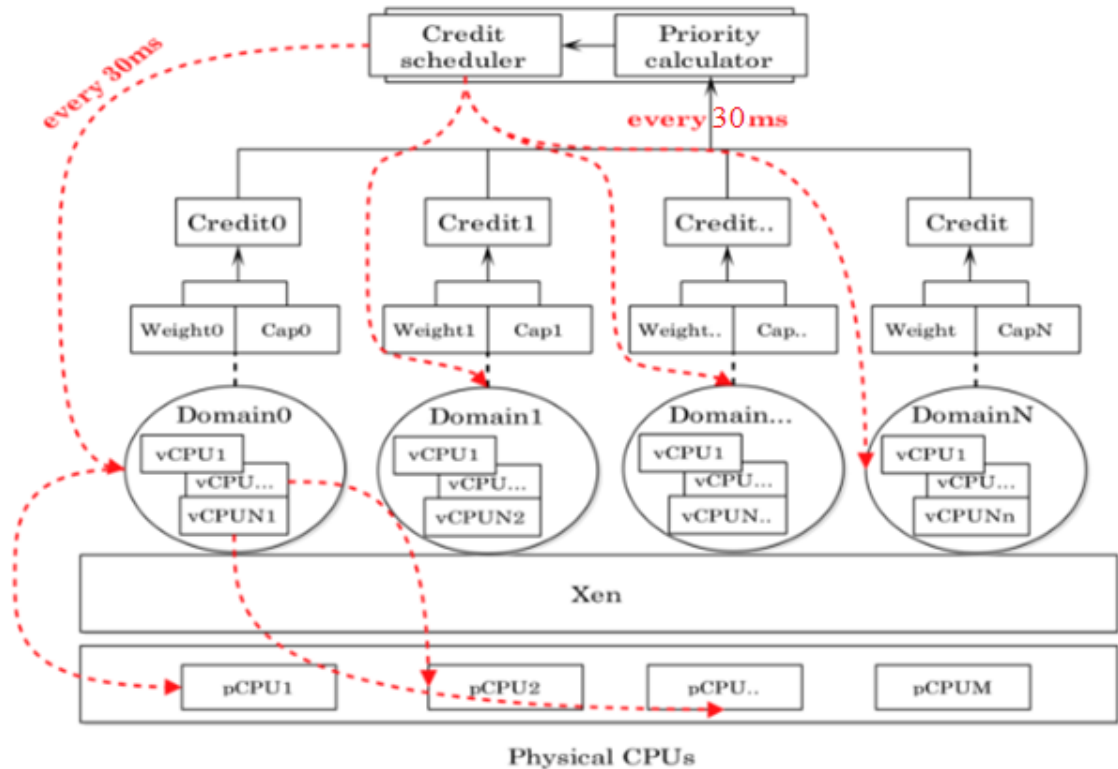


Figure 3. Algorithm of credit scheduler[1]

Credit scheduler is a fair share virtual CPU scheduler. The scheduler awards credit to each vCPU periodically and charges each vCPU to run on the pCPU[1].

Each **credit** contains two key **parameters**:

- *Weight* is proportional to vCPU value. That is, the higher a vCPU's weight, the more pCPU it can use;
- *Cap* is expressed as a percentage of the maximum number of pCPUs that a domain can utilize[1].

Each vCPU can have two **priorities**:

- “*Over*”, if the number of credits is less than 0;
- “*Under*”, if the number of credits is more than 0.

Time in Credit scheduler:

- *Time slice* (scheduling quantum) - uses 30 ms by default, after which it interrupts the vCPU and recalculates its priority;
- *Context-switch rate limit* - use 1ms by default. It is defining the minimum number of microseconds a vCPU would be allowed to run uninterrupted before being context switched[1].

Algorithm.

Each physical processor manages a queue of the vCPU. vCPUs are sorted based on priority (“Over” or “Under”). After a certain amount of time (30 ms), the vCPU will be interrupted, and the execution of the next vCPU will start. While the vCPU's priority is “Under”, it will be placed at the end of the queue. The vCPU with priority “Over” will be able to run before any that are out of credit. Credit priorities are recounted every 30 ms [2].

Thus, we can conclude that threads with more credits will have more time to execute.

3. Implementation of the Credit based scheduler

3.1 Algorithm

Comparing the $O(1)$ and the Credit schedulers' work, we can conclude that their principles are similar. In the Credit scheduler a kthread has a runqueue of runnable uthreads (for this project, we can assume that each kthread has 2 runqueues Under and Over) (Figure 4). A uthread will work a specific time before being interrupted (as in the $O(1)$ scheduler). However, in a Credit scheduler, "more important" threads with more credits must execute longer before being preempted by another thread.

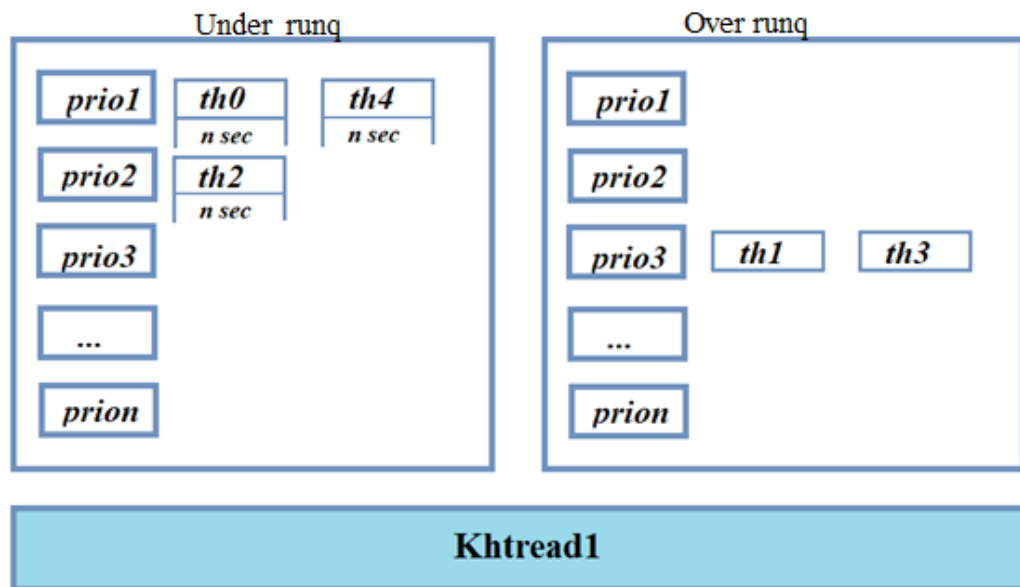


Figure 4. The credit scheduler algorithm

We can implement it that way: after each performance, the uthread will lose a certain number of credits. Before the next execution starts, the number of the thread's credits must be checked. If the number of credits is greater than 0, then the uthread is added to the end of Under queue (active); if the number of credits is lower than 0, then the uthread will be transferred to Over queue (expires). However, we should keep in mind, that when uthreads get into Over queue, they should get their initial credits back. When there are no active uthreads left in Under queue, then queues will change as in the $O(1)$ scheduler. The algorithm is shown below (Figure 5).

In that way the main idea of the credit scheduler will be saved. For example, in our project we can assume, that uthread loses 25 credits when executed, so a uthread with 100 credits will be executed four times in Under queue, while uthread with 25 credits will only be executed once. Only after that, they will be transferred to the Over queue.

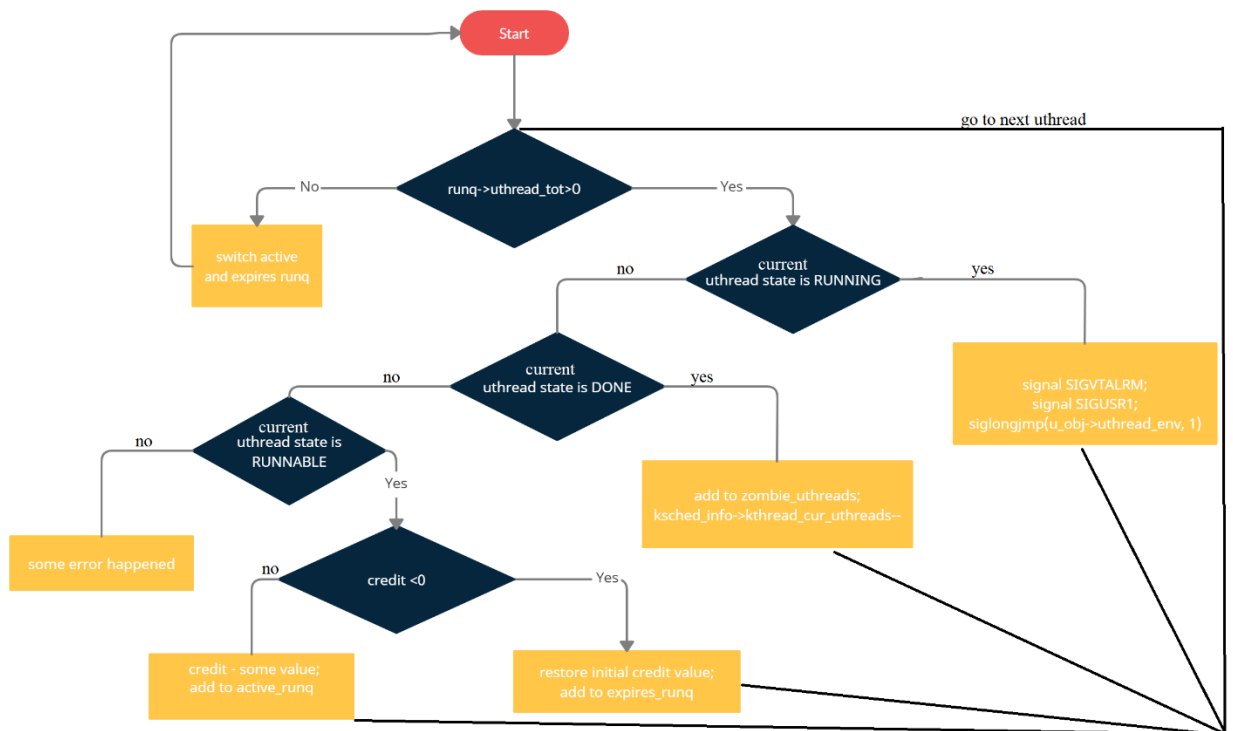


Figure 5. Also the credit scheduler algorithm

3.2 gt_yield

In this project the `gt_yield` function has been added to complete the execution voluntarily. When it is called, the function changes the corresponding flag (`gt_yield_flag`) from 0 to 1 (added to `uthread_struct`). The “`uthread_schedule`” function checks the value of this flag, and if the flag is equal to one, then the credit value of the yielding uthread is changed (in our case, it adds 25 credits extra), and the uthread is moved to the expires runqueue. In real life, it is usually used with I/O operations.

During testing, the “`gt_yield`” function was called for an arbitrary uthread. The result is shown in Figure 6. The function is disabled by default. To active it, the below code should be uncommented (Figure 7).

```

Thread(id:4, group:4, cpu:1) started
Thread(id:36, group:4, cpu:3) finished (TIME : 3 s and 18446744073709548987 us) completed its execution with total time= 0 s and 52198 us and run time =0 s and 5560 us
Thread(id:100, group:4, cpu:3) started1Thread id= 100 yielded
Thread yield id: 100, credit:125
Thread(id:37, group:5, cpu:3) started
Thread(id:4, group:4, cpu:1) finished (TIME : 3 s and 18446744073709549256 us) completed its execution with total time= 0 s and 53026 us and run time =0 s and 5665 us

```

Figure 6. Gt _ yield ()

```

{ //if (ptr->tid == 100 && sh_type==1 && c<2)
// {
//     gt_yield();
//     c++;
// }

```

Figure 7. Code of `gt_yield()`

3.3 Load balancing

The balancing algorithm was also added in the Credit scheduler. It is used when a certain kthread is idle, but there are still uthreads to execute. So a free kthread can intercept them. To do this, we introduce a new runqueue (`kthread_runqueue_t * other_kthread_runq; runqueue_t * other_runq`). If other runqueue's `uthread_mask` value is not zero, then a free kthread can intercept some uthreads. The result is shown below in Figure 8.

```

Thread running id: 62, credit:0)The thread with id: 25 is in expired runqueue with replinshed credits= 25
The thread with id: 61 is in expired runqueue with replinshed credits= 50
Thread: 126, group: 14 removed from CPU:0
..... uthread_context_func .....

```

Figure 8. Load balancing

3.4 Matrix

Each utread is executed on a matrix of its own. We create 128 uthreads and divide them into 16 groups (according to NUM_GROUPS). In each group, we have eight uthreads with a different number of credits and matrix size.

The code is provided below in Figure 9.

```

else if (sh_type == 1)
{for(credit = 25; credit <= 100; credit += 25)
{for(size = 32; size <= 256; size *= 2)
{for(inx=0; inx < 8; inx++)
{
uarg = &uargs[th_num];
uarg->_A = &A;
uarg->_B = &B;
uarg->_C = &C;
uarg->tid = th_num;
uarg->gid = (th_num % NUM_GROUPS);
uarg->m_dim = size;
uthread_create(&utids[th_num], uthread_mulmat, uarg, uarg->gid, credit);
th_num+=1;}
gr++;}}}

```

Figure 9. Matrix for the Credit scheduler

For the O(1) scheduler, we also create 128 uthreads. However, since there are no credits, there are 32 the same uthreads for each matrix size. The code is also provided below in Figure 10.

```

if (sh_type == 0)
{
for(size = 32; size <= 256; size *= 2)
{
for(inx=0; inx < 32; inx++)
{
uarg = &uargs[th_num];
uarg->_A = &A;
uarg->_B = &B;
uarg->_C = &C;
uarg->tid = th_num;
uarg->gid = (th_num % NUM_GROUPS);
uarg->m_dim = size;
uthread_create(&utids[th_num], uthread_mulmat, uarg, uarg->gid, 0);
th_num+=1;}
gr++;}}

```

Figure 10. Matrix for the O (1) scheduler

3.5 Time calculating

To calculate the time, we add new elements to the uthread_struct structure:

```

struct timeval time_start; /*time when uthread is created */
struct timeval time_end; /* time when uthread's state is done*/
struct timeval run_time; /* total time of uthread's execution*/
struct timeval run_time_t; /* uthread's executing time*/
struct timeval time_pr; /*time when uthread is preempted by another uthread*/
struct timeval time_sh; /*time when uthread adds to runq */.

```

Also, a new structure “uthread_time_t” is added:
struct timeval runtime; /*total time of executing*/
struct timeval totaltime; /* total time*/

The calculations:

totaltime = time_end – start_time;
run_time_t = time_pr - time_sh;
run_time = cyMMA bcex run_time_t.

The standard deviation and the mean value of both, the individual thread run times and the total execution times, should be identified.

To find the mean value, we added the runtime and totaltime of all threads in one group and divided it by 8 (the number of threads in one group).

In order to find the value of the standard deviation, we used the formula (Figure 11), where m is the average value, x is the time value, N = 8:

$$SD = \sqrt{\frac{\sum |x - \mu|^2}{N}}$$

Figure 11. Formula of standard deviation

4. Conclusion

During the experiments, where the NUM_CPUS = 4 was assumed, the following results were obtained :

```
-----Timing-----
SD(rt)185.893784      SD(tt)119934.234375      M(rt)5219.000000      M(tt)195980.125000
SD(rt)320.661987      SD(tt)125269.953125      M(rt)5299.125000      M(tt)607353.375000
      SD(rt)189.068115      SD(tt)119572.250000      M(rt)5208.500000      M(tt)195648.250000
      SD(rt)271.341766      SD(tt)125098.882812      M(rt)5228.125000      M(tt)606855.250000
SD(rt)324.692505      SD(tt)235201.921875      M(rt)19994.375000      M(tt)373200.625000
SD(rt)414.674683      SD(tt)229858.000000      M(rt)20743.875000      M(tt)1189793.375000
      SD(rt)208.356934      SD(tt)235140.937500      M(rt)20768.875000      M(tt)373955.750000
      SD(rt)137.454025      SD(tt)229796.234375      M(rt)20357.875000      M(tt)1189249.875000
SD(rt)923.244995      SD(tt)462984.906250      M(rt)80631.500000      M(tt)671614.250000
SD(rt)719.586548      SD(tt)326601.812500      M(rt)80856.375000      M(tt)948242.750000
      SD(rt)699.270691      SD(tt)461338.437500      M(rt)79737.625000      M(tt)669946.125000
      SD(rt)657.438843      SD(tt)323408.468750      M(rt)79878.375000      M(tt)945091.500000
SD(rt)2665.223633      SD(tt)197520.218750      M(rt)323741.000000      M(tt)5108550.500000
SD(rt)5254.513184      SD(tt)108887.468750      M(rt)325581.375000      M(tt)5362605.000000
      SD(rt)8170.186035      SD(tt)215822.500000      M(rt)327248.250000      M(tt)5128719.500000
      SD(rt)3266.476807      SD(tt)114484.281250      M(rt)322742.125000      M(tt)5363630.500000
~/gtthreads$
```

O(1) scheduler ts=100ms

For the O1 scheduler, just kept in mind, that every 4 groups have a matrix of the same size (1-4 line: 32; 5-8 line: 64; 9-12 line: 128; 13-16 line: 256). Because printing operation was done for credit scheduler, the first two columns are not valid for O(1).

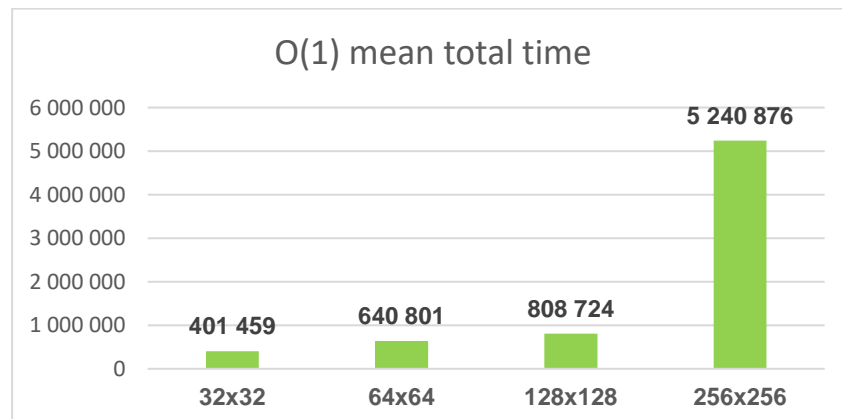
```
-----Timing-----
Cr:25 Mat:32      SD(rt)717.163940      SD(tt)23079.810547      M(rt)5266.375000      M(tt)42699.250000
Cr:25 Mat:64      SD(rt)439.593292      SD(tt)91338.601562      M(rt)19964.500000      M(tt)244682.875000
Cr:25 Mat:128      SD(rt)668.197754      SD(tt)2381580.750000      M(rt)79920.500000      M(tt)2914278.750000
Cr:25 Mat:256      SD(rt)1606.848022      SD(tt)331948.218750      M(rt)317058.750000      M(tt)6033226.500000
Cr:50 Mat:32      SD(rt)357.985687      SD(tt)23423.427734      M(rt)5207.000000      M(tt)41878.500000
Cr:50 Mat:64      SD(rt)121.242271      SD(tt)93218.062500      M(rt)20274.750000      M(tt)248359.500000
Cr:50 Mat:128      SD(rt)1673.513062      SD(tt)390344.593750      M(rt)83115.000000      M(tt)731749.375000
Cr:50 Mat:256      SD(rt)2298.714844      SD(tt)103053.398438      M(rt)323608.875000      M(tt)6354675.500000
Cr:75 Mat:32      SD(rt)113.458969      SD(tt)23053.615234      M(rt)5017.750000      M(tt)46176.375000
Cr:75 Mat:64      SD(rt)445.480347      SD(tt)92517.156250      M(rt)19993.500000      M(tt)259179.125000
Cr:75 Mat:128      SD(rt)672.440186      SD(tt)349069.375000      M(rt)79418.875000      M(tt)686104.125000
Cr:75 Mat:256      SD(rt)1121.707275      SD(tt)1147924.000000      M(rt)316729.750000      M(tt)3497764.500000
Cr:100 Mat:32      SD(rt)96.938126      SD(tt)23424.921875      M(rt)5090.500000      M(tt)46287.875000
Cr:100 Mat:64      SD(rt)290.651306      SD(tt)94211.234375      M(rt)20317.250000      M(tt)263168.000000
Cr:100 Mat:128      SD(rt)1455.297607      SD(tt)375604.406250      M(rt)82239.625000      M(tt)723278.625000
Cr:100 Mat:256      SD(rt)1081.555908      SD(tt)1381595.375000      M(rt)321742.750000      M(tt)4056619.500000
vmakhambayeva3@advos-05:~/gtthreads$
```

Credit scheduler ts=100ms

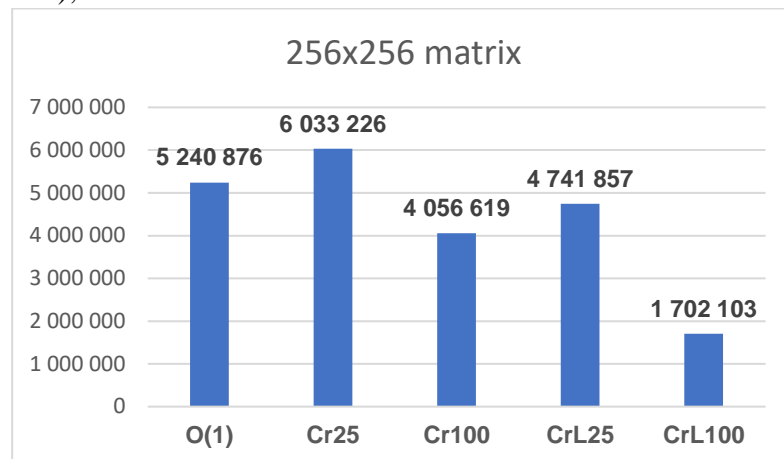
```
-----Timing-----
Cr:25 Mat:32      SD(rt)649.804016      SD(tt)23610.681641      M(rt)5337.500000      M(tt)43882.750000
Cr:25 Mat:64      SD(rt)1039.998169      SD(tt)92122.835938      M(rt)20185.000000      M(tt)248025.125000
Cr:25 Mat:128      SD(rt)603.219666      SD(tt)2853903.500000      M(rt)79671.375000      M(tt)3382433.500000
Cr:25 Mat:256      SD(rt)19256.984375      SD(tt)1386176.000000      M(rt)329330.750000      M(tt)4741857.000000
Cr:50 Mat:32      SD(rt)557.469482      SD(tt)23786.156250      M(rt)5349.500000      M(tt)43759.625000
Cr:50 Mat:64      SD(rt)310.655762      SD(tt)98073.273438      M(rt)20418.000000      M(tt)253927.875000
Cr:50 Mat:128      SD(rt)3727.824463      SD(tt)189377.031250      M(rt)82943.000000      M(tt)494600.125000
Cr:50 Mat:256      SD(rt)2714.906494      SD(tt)72737.484375      M(rt)321010.500000      M(tt)2514754.250000
Cr:75 Mat:32      SD(rt)283.316956      SD(tt)23434.041016      M(rt)5108.000000      M(tt)48429.625000
Cr:75 Mat:64      SD(rt)243.548248      SD(tt)93275.390625      M(rt)19923.000000      M(tt)262343.125000
Cr:75 Mat:128      SD(rt)1737.435059      SD(tt)344631.281250      M(rt)80450.375000      M(tt)680065.125000
Cr:75 Mat:256      SD(rt)7823.870605      SD(tt)1470689.125000      M(rt)323534.625000      M(tt)3871140.250000
Cr:100 Mat:32      SD(rt)271.312361      SD(tt)23649.660136      M(rt)5183.000000      M(tt)48205.875000
Cr:100 Mat:64      SD(rt)2269.215820      SD(tt)95432.320312      M(rt)21022.250000      M(tt)270987.750000
Cr:100 Mat:128      SD(rt)8617.101562      SD(tt)178164.687500      M(rt)85075.875000      M(tt)526434.875000
Cr:100 Mat:256      SD(rt)2636.642822      SD(tt)450801.093750      M(rt)321593.250000      M(tt)1702103.625000
vmakhambayeva3@advos-05:~/gtthreads$
```

Credit scheduler with load balance ts=100ms

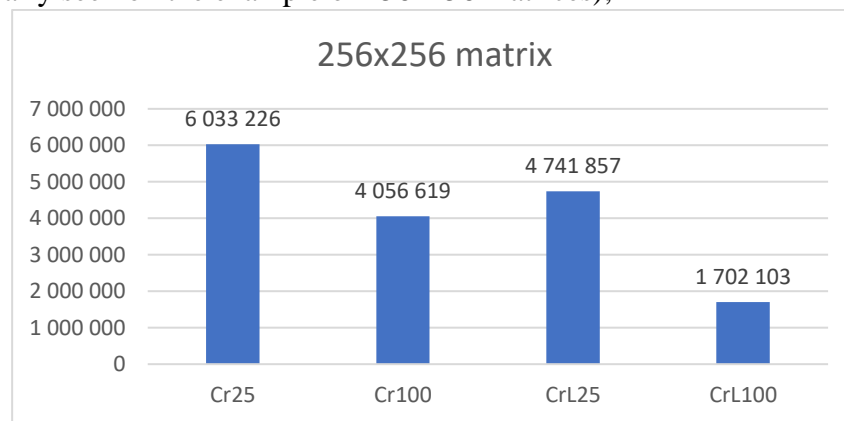
- 1) O(1) works as planned, run time and total time for similar processes (thread with one matrix size) are identical. You can observe a uniform increase in run time (approximately four times) depending on the matrix's size. The total execution time also increases depending on the size of the matrix;



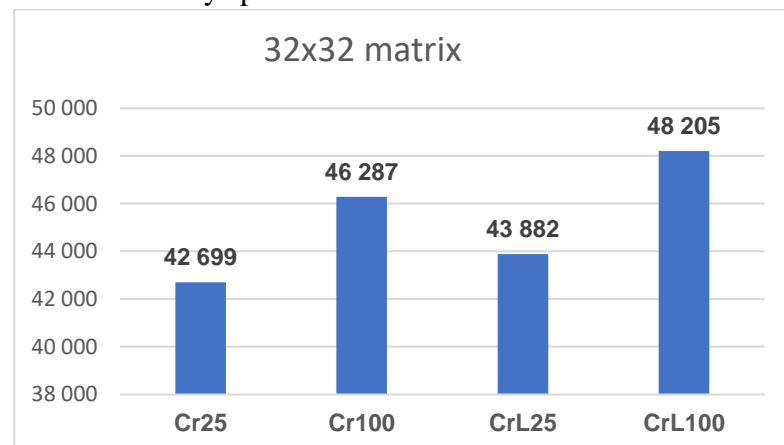
- 2) The number of credits affects the thread's total execution time in the Credit scheduler, as expected. The more credits, the less total execution time (it is more clearly seen on the example of 256x256 matrices);



- 3) Credit scheduler with balancing can significantly reduce the total execution time of the uthread (it is more clearly seen on the example of 256x256 matrices);



- 4) For thread with small matrix computation, the results are not indicative since they manage to finish their execution before they spend all their credit.



Thus, we have proved that O(1) is a fair share scheduler. However, using the credit scheduler, it is possible to execute "more important" tasks faster since they will wait less time until their turn comes. Since there were eight groups and only 4 CPUs, some experimental values may not be accurate. Similar results were obtained with $t_s = 30\text{ms}$.

```

0-----Timing-----
Cr:25 Mat:32 SD(rt)294.344177 SD(tt)70641.851562 M(rt)5323.000000 M(tt)110473.250000
Cr:25 Mat:64 SD(rt)187.473953 SD(tt)69233.476562 M(rt)5134.625000 M(tt)357302.500000
Cr:25 Mat:128 SD(rt)237.283279 SD(tt)70604.187500 M(rt)5439.125000 M(tt)110448.250000
Cr:25 Mat:256 SD(rt)176.545624 SD(tt)69239.296875 M(rt)5284.875000 M(tt)357216.750000
Cr:50 Mat:32 SD(rt)120.740425 SD(tt)70638.789062 M(rt)20015.500000 M(tt)124861.250000
Cr:50 Mat:64 SD(rt)1164.389648 SD(tt)70096.328125 M(rt)20716.625000 M(tt)372560.875000
Cr:50 Mat:128 SD(rt)402.585388 SD(tt)70383.750000 M(rt)21090.500000 M(tt)125958.500000
Cr:50 Mat:256 SD(rt)887.155396 SD(tt)69853.882812 M(rt)21217.750000 M(tt)372917.500000
Cr:75 Mat:32 SD(rt)573.609131 SD(tt)139465.515625 M(rt)79040.750000 M(tt)1174189.625000
Cr:75 Mat:64 SD(rt)1172.882446 SD(tt)79017.101562 M(rt)79480.375000 M(tt)1248292.625000
Cr:75 Mat:128 SD(rt)622.885132 SD(tt)126989.656250 M(rt)81224.750000 M(tt)1227843.000000
Cr:75 Mat:256 SD(rt)948.828247 SD(tt)74895.679688 M(rt)81993.500000 M(tt)1281111.500000
Cr:100 Mat:32 SD(rt)5045.372559 SD(tt)212190.218750 M(rt)323034.750000 M(tt)5219296.000000
Cr:100 Mat:64 SD(rt)4832.487793 SD(tt)177664.562500 M(rt)324247.375000 M(tt)5314508.000000
Cr:100 Mat:128 SD(rt)6426.998047 SD(tt)232640.593750 M(rt)324530.500000 M(tt)5241245.000000
Cr:100 Mat:256 SD(rt)3439.080322 SD(tt)185550.671875 M(rt)323942.875000 M(tt)5324476.000000
vmakhambayeva3@advos-05:~/gtthreads$

```

O(1) scheduler ts=30ms

```

-----Timing-----
Cr:25 Mat:32 SD(rt)793.148254 SD(tt)22405.917969 M(rt)5205.625000 M(tt)43261.750000
Cr:25 Mat:64 SD(rt)444.370056 SD(tt)96067.835938 M(rt)19513.500000 M(tt)282339.250000
Cr:25 Mat:128 SD(rt)633.036316 SD(tt)682333.625000 M(rt)80645.375000 M(tt)2968268.250000
Cr:25 Mat:256 SD(rt)2142.972656 SD(tt)97642.507812 M(rt)317645.750000 M(tt)632092.500000
Cr:50 Mat:32 SD(rt)555.335815 SD(tt)24031.353516 M(rt)5329.375000 M(tt)44819.250000
Cr:50 Mat:64 SD(rt)540.091675 SD(tt)96827.359375 M(rt)20415.375000 M(tt)260951.750000
Cr:50 Mat:128 SD(rt)1614.640991 SD(tt)1250655.125000 M(rt)80758.125000 M(tt)1946913.625000
Cr:50 Mat:256 SD(rt)8356.992188 SD(tt)178274.218750 M(rt)330228.125000 M(tt)6394836.000000
Cr:75 Mat:32 SD(rt)134.587738 SD(tt)22717.478516 M(rt)4914.875000 M(tt)46502.500000
Cr:75 Mat:64 SD(rt)621.112183 SD(tt)89774.804688 M(rt)19718.125000 M(tt)247981.250000
Cr:75 Mat:128 SD(rt)497.201416 SD(tt)280262.625000 M(rt)79979.375000 M(tt)573964.875000
Cr:75 Mat:256 SD(rt)2361.929199 SD(tt)274133.250000 M(rt)319306.250000 M(tt)4700123.000000
Cr:100 Mat:32 SD(rt)89.667351 SD(tt)23918.945312 M(rt)5157.625000 M(tt)47739.250000
Cr:100 Mat:64 SD(rt)243.403351 SD(tt)93310.007812 M(rt)20222.250000 M(tt)260950.500000
Cr:100 Mat:128 SD(rt)1249.771118 SD(tt)349912.906250 M(rt)80102.000000 M(tt)710425.875000
Cr:100 Mat:256 SD(rt)3720.808594 SD(tt)208287.234375 M(rt)323591.750000 M(tt)5465387.500000
vmakhambayeva3@advos-05:~/gtthreads$

```

Credit scheduler ts=30ms

```

-----Timing-----
Cr:25 Mat:32 SD(rt)483.149292 SD(tt)23674.050781 M(rt)5234.000000 M(tt)42703.250000
Cr:25 Mat:64 SD(rt)1524.558594 SD(tt)660375.062500 M(rt)20425.625000 M(tt)624172.500000
Cr:25 Mat:128 SD(rt)579.104675 SD(tt)87210.492188 M(rt)79282.500000 M(tt)4387070.000000
Cr:25 Mat:256 SD(rt)3337.857666 SD(tt)1403660.500000 M(rt)319345.250000 M(tt)4235343.500000
Cr:50 Mat:32 SD(rt)206.840134 SD(tt)25680.548828 M(rt)5389.125000 M(tt)44978.750000
Cr:50 Mat:64 SD(rt)264.355225 SD(tt)100693.375000 M(rt)20557.250000 M(tt)266474.000000
Cr:50 Mat:128 SD(rt)1288.162354 SD(tt)506465.593750 M(rt)81292.250000 M(tt)1158247.125000
Cr:50 Mat:256 SD(rt)4476.158691 SD(tt)433272.656250 M(rt)324741.875000 M(tt)2353822.750000
Cr:75 Mat:32 SD(rt)36.204281 SD(tt)23190.296875 M(rt)4994.000000 M(tt)45750.250000
Cr:75 Mat:64 SD(rt)337.834504 SD(tt)87498.414062 M(rt)19863.750000 M(tt)246936.625000
Cr:75 Mat:128 SD(rt)617.854431 SD(tt)278652.562500 M(rt)79267.875000 M(tt)565261.750000
Cr:75 Mat:256 SD(rt)4769.881836 SD(tt)1356477.125000 M(rt)319867.125000 M(tt)4488480.500000
Cr:100 Mat:32 SD(rt)191.627899 SD(tt)25663.453125 M(rt)5388.500000 M(tt)50067.750000
Cr:100 Mat:64 SD(rt)127.338135 SD(tt)94504.460938 M(rt)20486.000000 M(tt)271911.625000
Cr:100 Mat:128 SD(rt)1088.655029 SD(tt)170570.312500 M(rt)81017.875000 M(tt)495027.750000
Cr:100 Mat:256 SD(rt)8408.640625 SD(tt)462537.781250 M(rt)325231.750000 M(tt)2279037.000000
vmakhambayeva3@advos-05:~/gtthreads$

```

Credit scheduler with load balance ts=30ms

Reference:

- 1) “An introduction to the credit scheduler in Xen,” *Welcome to xfhelen's Homepage*, 10-Feb-2017. [Online]. Available: <https://xfhelen.github.io/2016/06/01/2016-06-01/>. [Accessed: 1-Feb-2021].
- 2) “Credit Scheduler,” *Xen*. [Online]. Available: https://wiki.xenproject.org/wiki/Credit_Scheduler. [Accessed: 1-Feb-2021].