



Basic's

Rust Memory Diagram



# Basic's

## ✓ Memory Diagram (4 Bytes, Decimal Values)

lua

Variable: i = 2 (i32)

Address	Byte (Decimal)	Comment
0x1000	2	Least significant byte (LSB)
0x1001	0	
0x1002	0	
0x1003	0	Most significant byte (MSB)



# Basic's

## ✓ Memory Diagram for i8

sql

Variable: x = 5 (i8)

Address	Byte (Decimal)	(Only 1 byte)
0x3000	5	



# Rust Movement of Address

```
let x: i32 = 42;  
let ptr: *const i32 = &x;
```

Here, `ptr` is a **raw pointer** pointing to variable `x`.

Variable: `x = 42 (i32)`

`x (Data):`

Address	Byte (Decimal)	(Stored in 4 bytes as i32)
0x2000	42	
0x2001	0	
0x2002	0	
0x2003	0	

Pointer Variable: `ptr`

`ptr (Pointer):`

Address	Pointer Value	(Address of x)
0x3000	0x2000	



# Rust Movement of Address



## Concept: Two Pointers, Single Data

- The **data** stays in **one place in memory**.
- Pointers simply **store the address** of that data.
- If a pointer is reassigned or moved, only the **pointer address** changes.



# Rust Movement of Address

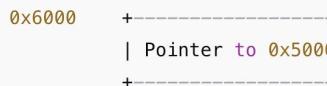
## ✓ Memory Diagram – Two Pointers, One Data

pgsql

Data Variable: x = 55 (Stored in Stack)



Pointer 1 (ptr1):



Pointer 2 (ptr2):



```
fn main() {
    let x: i32 = 55;

    let ptr1: *const i32 = &x;
    let ptr2: *const i32 = ptr1; // Copying/moving pointer (not data)

    println!("{:p}", ptr1);      // Same address
    println!("{:p}", ptr2);      // Same address
}
```



# Rust Movement of Address

## ✓ Memory Diagram

lua

```
DATA SEGMENT (read-only):
```



STACK:



```
let names: &str = "coderrange";
```



# Rust Movement of Address

## 📦 Diagram – Movement of Box Pointer

lua

HEAP:



STACK (before move):



STACK (after move):

b1 = INVALID

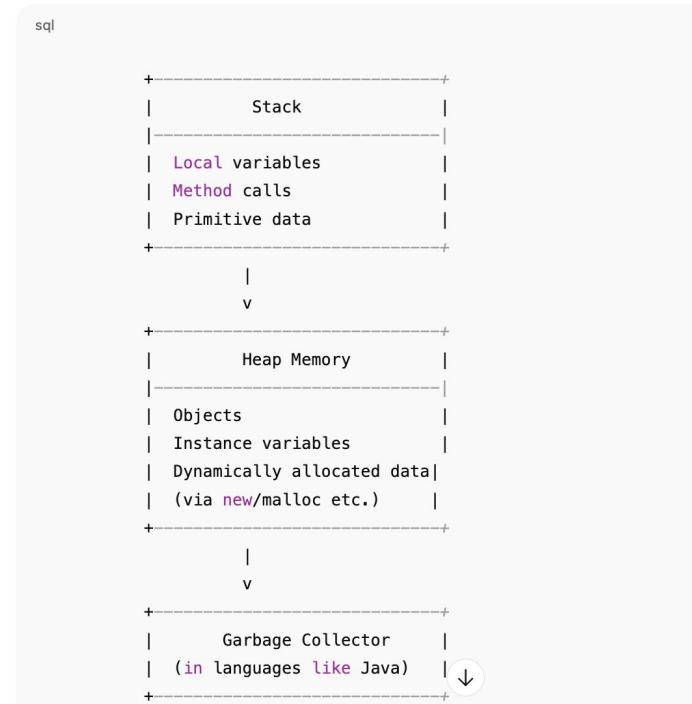


```
let b1 = Box::new(10); // Box owns data on HEAP  
  
let b2 = b1;           // Ownership of pointer (Box) MOVES to b2  
  
// println!("{}", b1); // ✗ Error: b1 is invalid after move  
  
println!("{}", b2);   // ✓ b2 is now the owner
```



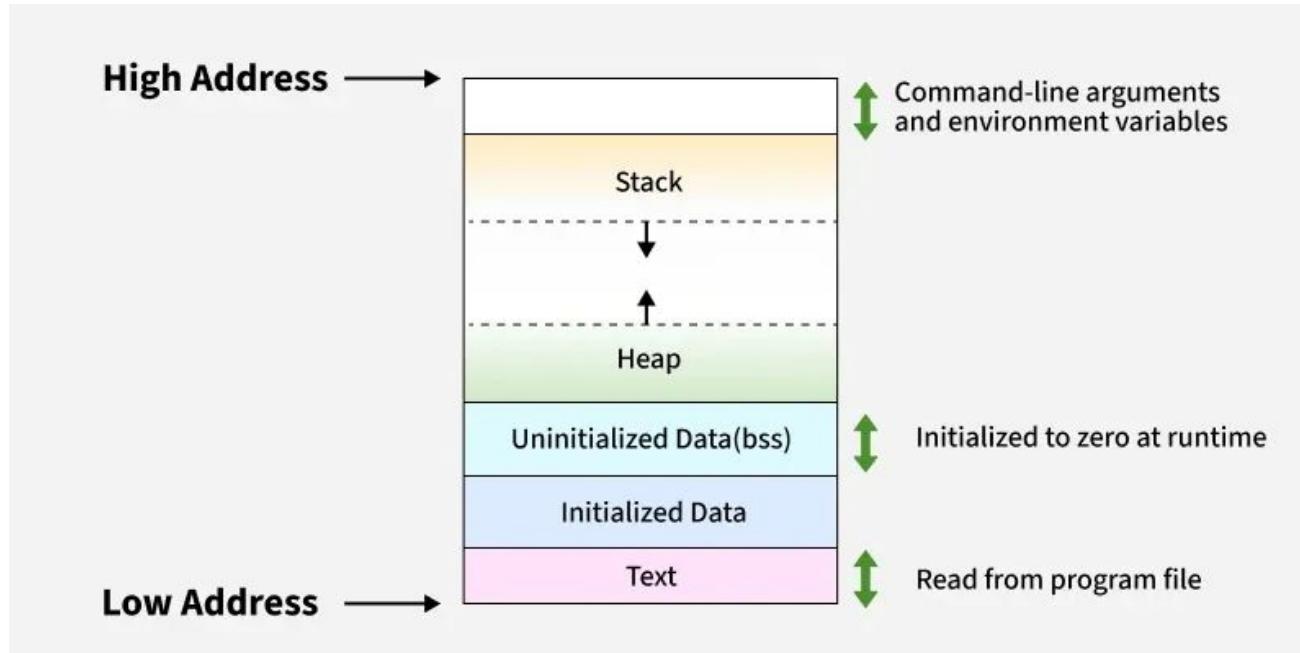
# Memory Diagram

## 🧠 Heap Memory – Conceptual Diagram





# Memory Diagram





# Memory Diagram

## Visual Memory Diagram

vbnet

```
let name = String::from("Rust");
```

STACK:

```
+-----+  
| Pointer to Heap (name) |  
| Length = 4 |  
| Capacity >= 4 |  
+-----+
```

HEAP:

```
+---+---+---+---+  
| R | u | s | t |  
+---+---+---+---+
```

When you do `name.push_str(" Language");`, r



# Memory Diagram

## 1 Pointers vs. References

pgsql

speed	100

ref: &speed

Points to 'speed'

pgsql

speed	100

raw_ptr: *const T

Unsafe pointer to 'speed'



# Memory Diagram

## 3 Ownership and Moves

rust

```
let a = String::from("Hello");
```



```
let b = a; // Ownership moved to b
```



a is now **invalid** (moved).



# Memory Diagram

## 4 Borrowing (&T and &mut T)

Immutable Borrow:

kotlin

```
let data = 42;
let ref1 = &data; // Immutable borrow
```



Points to 'data'



# Memory Diagram

Mutable Borrow (Only One Allowed):

wasm

```
let mut data = 42;  
let ref_mut = &mut data; // Mutable borrow
```



Points to 'data'



# Memory Diagram

## 5 Lifetimes

rust

```
fn longest<'a>(x: &'a str, y: &'a str) -> &'a str
```

Diagram:

rust

'a lifetime scope	
x: &'a str	Both references valid
y: &'a str	for same lifetime 'a
return: &'a str	



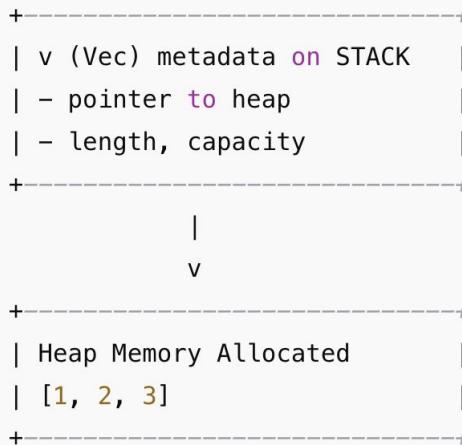


# Memory Diagram

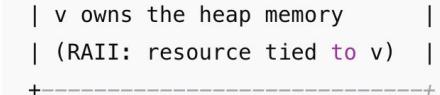
## ✓ Best RAII + Heap Memory Diagram

sql

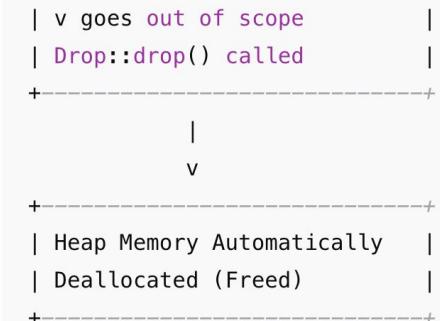
### 1 Resource Allocation (Vec Created)



### 2 Ownership Scope



### 3 Scope Ends (Drop Triggered)





# Memory Diagram

## 1 `mut` vs. `&mut` – (Most Common Confusion)

Situation	Code Snippet	Meaning
<code>mut</code> before variable	<code>let mut x = 10;</code>	<code>x</code> itself is mutable.
<code>mut</code> in reference	<code>let y = &amp;mut x;</code>	Reference to mutable <code>x</code> . You can mutate <b>what y points to</b> , not y itself.



# Memory Diagram

## 2 Raw Pointers: `*const` vs. `*mut`

Keyword	Example	Meaning
<code>*const T</code>	<code>let ptr: *const i32</code>	Immutable raw pointer.
<code>*mut T</code>	<code>let ptr: *mut i32</code>	Mutable raw pointer.
Dereferencing	<code>unsafe { *ptr }</code>	Always inside <code>unsafe</code> block.



# Memory Diagram

3 ref vs. &

Keyword	Use Case	Example
ref	Pattern matching	let ref y = x;
&	Reference in binding	let y = &x;

ref is rarely used in modern Rust; using & is clearer for most cases.



# Memory Diagram

## 4 static vs. const

Keyword	Example	Stored In	Lifespan	Mutable?
const	const MAX: u32 = 100;	DATA/TEXT	Compile-time	Never
static	static COUNT: u32 = 10;	DATA	Entire program	Yes (unsafe with static mut )



# Memory Diagram

## 5 Shadowing with let

### Confusion Point

---

Variable can be redefined using `let` again in the same scope.

rust

```
let x = 5;  
let x = x + 1;    // Shadowing, not mutation
```



# Memory Diagram

## 6 Ownership / Borrowing Errors

Confusion Point	Mistake Example
Trying to use moved value	After <code>let b = a;</code> , <code>a</code> becomes invalid if non-Copy.
Borrow checker complaints	Holding both <code>&amp;mut</code> and <code>&amp;</code> simultaneously causes compiler errors.



# Memory Diagram

## 7 Lifetime Annotations ( 'a )

- Even experienced devs struggle with:
  - Missing lifetime annotations.
  - Incorrect 'a placement.
  - Functions returning references to local variables.



## Memory Diagram

### 8 Interior Mutability Confusion

- Use of:
  - `RefCell<T>` – run-time borrow checking.
  - `Cell<T>` – copyable interior mutability.
  - `Mutex<T>` – for multi-threaded interior mutability.

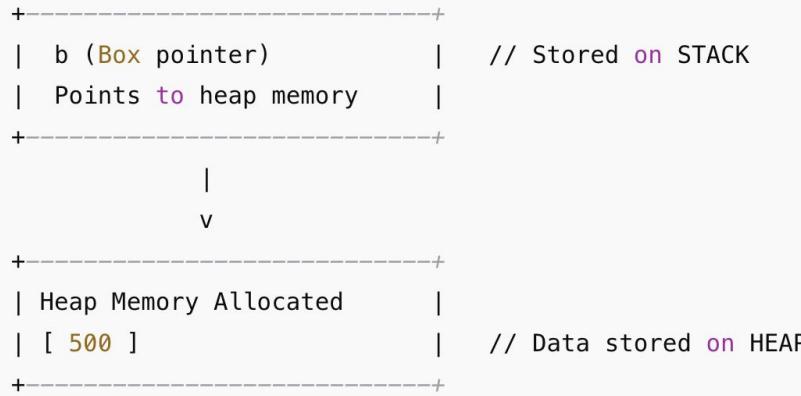


# Memory Diagram

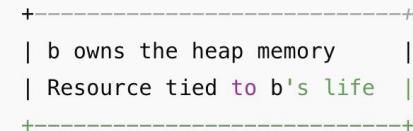
## ✓ RAI + Box – Best Box Diagram

pgsql

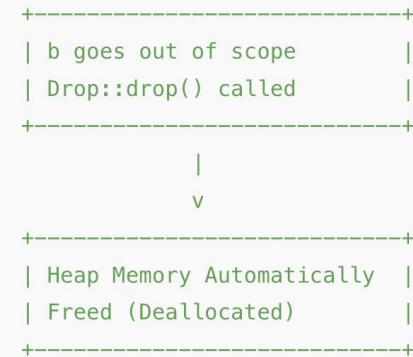
### 1 Resource Allocation (Box Created)



### 2 Ownership Scope (RAII Active)



### 3 Scope Ends (RAII Cleanup)





## Memory Diagram

### ✓ Example Code Using Box

rust

```
fn main() {
    let b = Box::new(500);      // Allocates 500 on HEAP
    println!("{}", b);          // b owns the heap memory

} // Scope ends -> Drop called -> HEAP memory freed
```

$\&b \rightarrow 7000$   
6000

6000 →  
500



## Memory Diagram

`&b → 7000`

`6000`

`let b1 =b`

`6000→`

`500`

`Let &b1 = &b`



# Memory Diagram

## ✓ Example Code Using Box

rust

```
fn main() {
    let b = Box::new(500);      // Allocates 500 on HEAP
    println!("{}", b);          // b owns the heap memory
}
// Scope ends -> Drop called -> HEAP memory freed
```

&boxing → 10004  
6007

6007 → heap  
500

&boxer → 7000  
10004



# Memory Diagram

## ✓ 1 Tuple Memory Diagram

rust

```
let t = (10, 3.14, true);
```

+	-----	-----	-----	-----		
	10 (i32)		3.14 (f64)		true (bool)	
+	-----	-----	-----	-----		



# Memory Diagram

## ✓ 2 Array Memory Diagram

rust

```
let arr = [1, 2, 3, 4];
```

+-----+	-----+	-----+	-----+	-----+	-----+	-----+
1	2	3	4		(All i32)	
+-----+	-----+	-----+	-----+	-----+	-----+	-----+



# Memory Diagram



3

## Slice Memory Diagram

rust

```
let arr = [10, 20, 30, 40];
let slice = &arr[1..3]; // [20, 30]
```

- A **pointer** to start of data.
- A **length** (number of elements in slice).

pgsql

+-----+		+-----+
Pointer to arr[1]		
Length = 2		
+-----+		+-----+
	v	
+-----+-----+		(View of arr)
20   30		+-----+



# Memory Diagram

## ✓ Slice – STACK (metadata) + Referenced Data (STACK)

- Example:

rust

```
let arr = [10, 20, 30, 40];
let slice = &arr[1..3]; // slice = [20, 30]
```



# Memory Diagram

## 🔍 When Slice Points to Heap:

- Example:

rust

```
let v = vec![10, 20, 30, 40]; // Data on HEAP
let slice = &v[1..3];           // Slice view over HEAP data
```



# Memory Diagram

Tuple:

+-----+	-----+	-----+	-----+	-----+
Element 1	Element 2	Element 3	(All on STACK)	
+-----+	-----+	-----+	-----+	

Array:

+-----+	-----+	-----+	-----+	-----+
1	2	3	4	(All on STACK)
+-----+	-----+	-----+	-----+	

Slice (of an array):

+-----+	-----+
Pointer + Length (STACK)	
+-----+	
v	
+-----+	
20	30   (Data in STACK or HEAP dep  ing on source)
+-----+	



# Memory Diagram

## 1 Tuple + mut — Challenges

rust

```
let mut t = (1, 2.0, true); // Entire tuple is mutable  
  
t.0 = 5;          // ✓ Can modify individual elements  
t.1 = 3.14;       // ✓  
t.2 = false;      // ✓
```

## 2 Array + mut — Challenges

rust

```
let mut arr = [1, 2, 3, 4];  
  
arr[0] = 99;      // ✓ Modify elements directly
```

## 3 Slice + mut — Challenges

rust

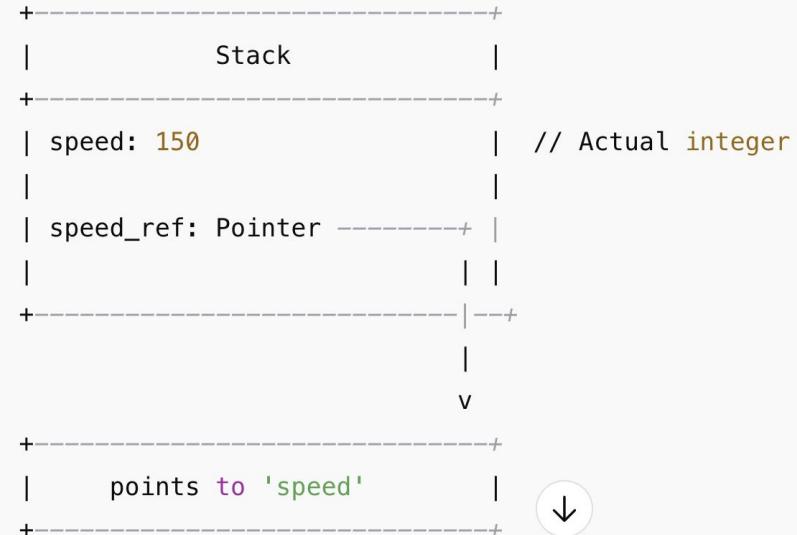
```
let mut arr = [10, 20, 30, 40];  
let slice = &mut arr[1..3];    // Mutable slice of arr  
  
slice[0] = 99;      // ✓ Modifies arr[1]
```



# Memory Diagram

✓ Best Memory Diagram

pgsql





# Memory Diagram 0

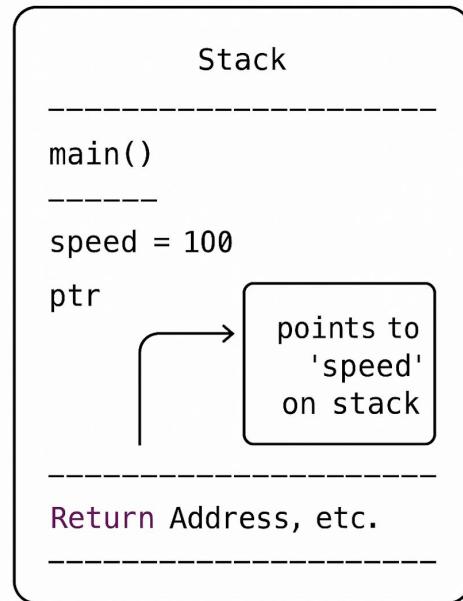
## 🦀 Rust Raw Pointer – Memory Diagram

### ✓ Rust Code (Raw Pointer Example)

```
rust

fn main() {
    let speed = 100;                      // stack allocation
    let ptr: *const i32 = &speed;          // raw pointer to speed

    unsafe {
        println!("Value via ptr: {}", *ptr);
    }
}
```





# Coding Mistakes

Box	:: new
String	:: string push_str



# Memory Diagram 1

## ✓ Rust Code (Stack Object)

```
rust

struct Car {
    speed: u32,
}

fn main() {
    let car1 = Car { speed: 100 }; // stack allocation
    println!("Speed: {}", car1.speed);
}
```

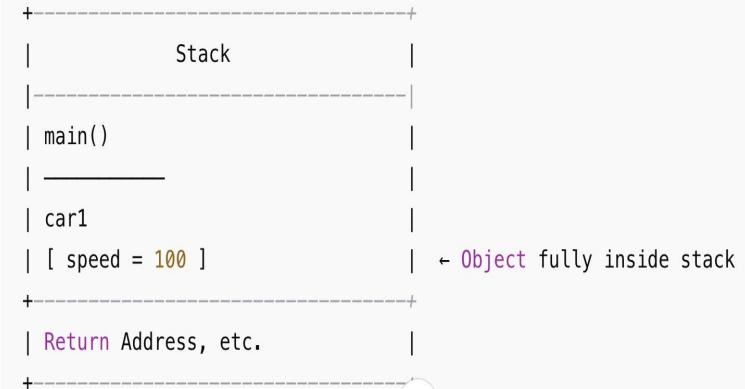


## 🧠 Memory Diagram – Rust Stack Object

pgsql

Copy

### 🦀 Rust Stack Object – Memory Diagram





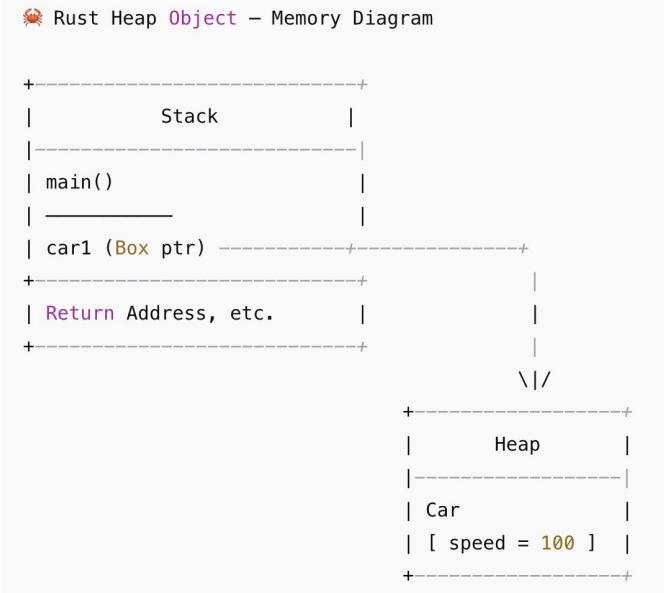
# Memory Diagram 1

## ✓ Rust Code (Heap Allocation)

```
rust                                ⚡ C

    struct Car {
        speed: u32,
    }

    fn main() {
        let car1 = Box::new(Car { speed: 100 }); // heap allocation
        println!("Speed: {}", car1.speed);
    }
}
```

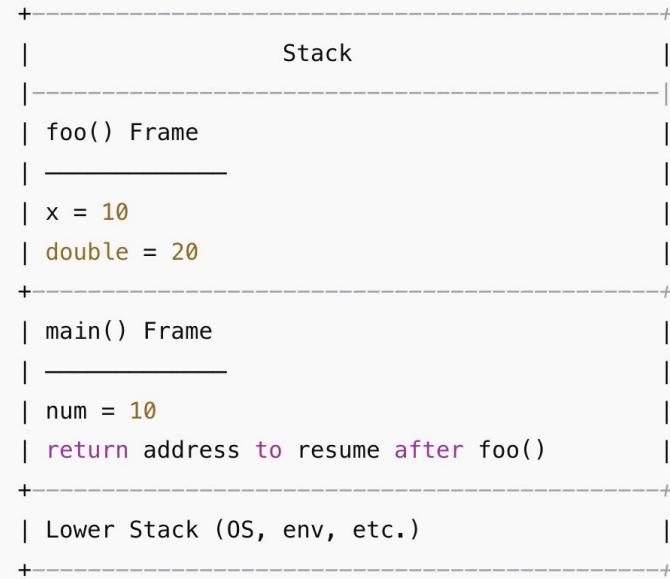




# Memory Diagram 2

```
fn foo(x: i32) {  
    let double = x * 2;  
    println!("Double: {}", double);  
}  
  
fn main() {  
    let num = 10;  
    foo(num); // function call  
}
```

🦀 Rust Function Call – Memory Diagram





## Memory Diagram 2

HEAP:

```
+-----+  
| "coderrange" |   <- vec[0]  
+-----+
```

STACK:

```
vec (pointer, len, cap)  
+-----+  
| vec ptr    | ---> points to HEAP  
| len         |  
| capacity    |  
+-----+
```

ref\_to\_first (borrows from HEAP)

```
+-----+  
| &vec[0]          |  
| Points to HEAP data |  
+-----+
```

Lifetime ties:

ref\_to\_first cannot outlive vec



## Memory Diagram 2

HEAP:

```
+-----+  
| "rustacean" | <- Stored in HEAP (owned by String s)  
+-----+
```

STACK:

s (String): owns pointer to heap

r (&String): immutable borrow, points to heap data  
Lifetime 'a: r cannot outlive s



# Memory Diagram 2

## Quick Code Summary

Element	Memory Location	Role	(
s (String)	Stack (pointer + len + cap) + Heap	Owns heap data ("rustacean")	)
r (&String)	Stack	Borrowed reference, points to heap	
Lifetime	Logical rule	Ensures r doesn't outlive s	



## Memory Diagram 2

```
fn main() {  
    let s = String::from("rustacean");  
  
    let r = borrow_string(&s);  
  
    println!("{}", r); // ✅ Safe: r cannot outlive s  
}
```



# Memory Diagram 2

## Return Ownership:

rust

```
fn get_ref() -> String {  
    let s = String::from("coderrange");  
    s // Returning ownership  
}
```

Usage:

rust

```
fn main() {  
    let result = get_ref();  
    println!("{}", result); // Safe: result now owns the String  
}
```





## Memory Diagram 2

```
fn get_ref<'a>(input: &'a String) -> &'a String {  
    input      // Borrowing from input safely  
}
```

Usage:

rust

```
fn main() {  
    let s = String::from("coderrange");  
    let r = get_ref(&s);  
  
    println!("{}", r);  // ✅ Safe: r cannot outlive s  
}
```

& 'a





## Memory Diagram 2

&'a String

&'a Str

**implicit lifetimes,**

```
fn news(args: &String) -> &String {  
    args  
}
```

```
fn news<'a>(args: &'a String) -> &'a String { ... }
```



# Memory Diagram 3

## ✓ `{}` vs `{:?}` in Rust – Internally Handled by Traits

Syntax	Trait Used	Purpose
<code>{}</code>	<code>std::fmt::Display</code>	Human-readable output
<code>{:?}</code>	<code>std::fmt::Debug</code>	Developer/debug output



# Memory Diagram 3



1

## Debug – Auto-Generated Example:

Rust creates an implementation that enables formatted printing using `{:?}`.

```
rust

#[derive(Debug, Clone, PartialEq)]
struct User {
    name: String,
    age: u32,
}
```

rust

```
impl std::fmt::Debug for User {
    fn fmt(&self, f: &mut std::fmt::Formatter) -> std::fmt::Result {
        f.debug_struct("User")
            .field("name", &self.name)
            .field("age", &self.age)
            .finish()
    }
}
```



# Memory Diagram 3

rust

```
println!("{:?}", my_data);
```

➡ The compiler looks for an implementation of the `Debug` trait:

rust

```
impl std::fmt::Debug for MyType {
    fn fmt(&self, f: &mut std::fmt::Formatter<'_>) -> std::fmt::Result {
        f.debug_struct("MyType")
            .field("some_field", &self.some_field)
            .finish()
    }
}
```



# Memory Diagram 3

Method	Behavior
<code>.unwrap()</code>	Panic if None/Err
<code>.expect()</code>	Panic with custom message
<code>.unwrap_or()</code>	Use fallback value if None/Err
<code>.unwrap_or_else()</code>	Use closure to handle error



# Memory Diagram 3

## 📦 ·unwrap() – Source Code in Option<T>

From Rust standard library:

rust



```
impl<T> Option<T> {
    pub fn unwrap(self) -> T {
        match self {
            Some(val) => val,
            None => panic!("called `Option::unwrap()` on a `None` value"),
        }
    }
}
```



# Memory Diagram 3

## 🧠 C++ Function Arguments – Memory Diagram

Assume:

cpp

```
void example(int x, int& y, int* z);
```

Then we call:

cpp

```
int a = 10, b = 20, c = 30;
example(a, b, &c);
```

### 📦 Stack Memory Layout During Function Call

pgsql

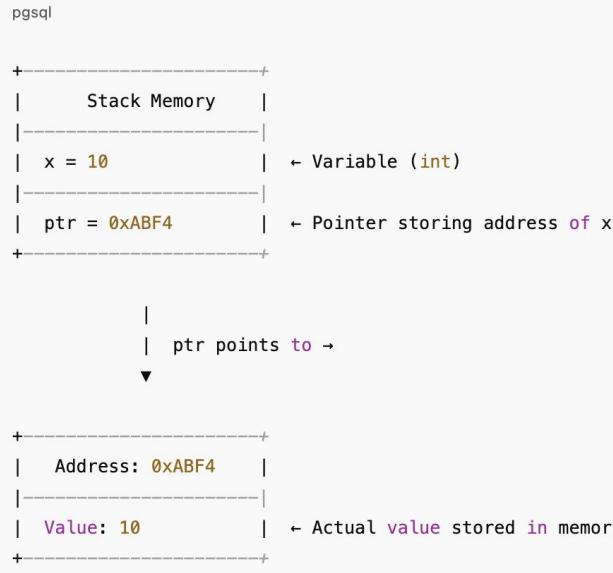
Stack Frame: example()	
Return Address	
x = 10	← Pass-by-Value (copy of a)
y → refers to b	← Pass-by-Reference (alias)
z → points to &c	← Pass-by-Pointer (explicit addr)



# Memory Diagram 4

## Pointer

### Diagram:

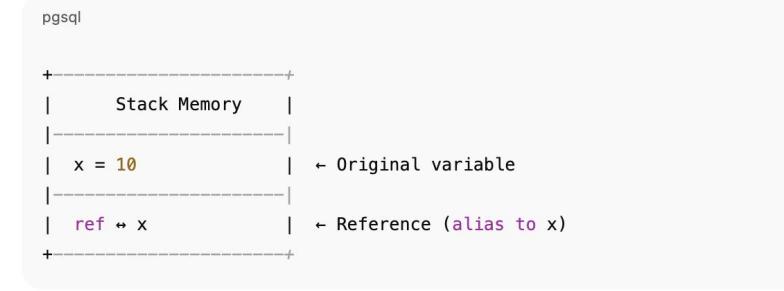


### C++ Reference Variable – Memory Diagram

### Scenario:

You have an `int x = 10;` and `int& ref = x;`

### Memory Diagram:

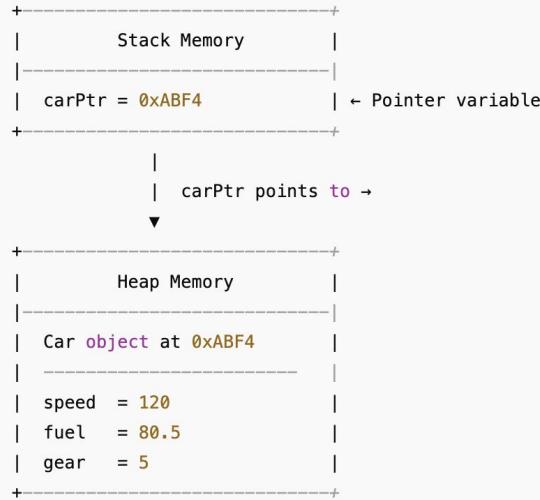




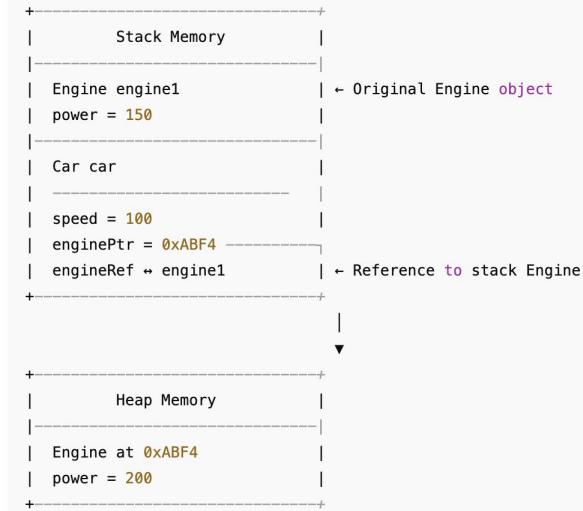
# Memory Diagram 5



**pgsql**



pgs





# Memory Diagram 6

## 📦 Shallow Copy – Memory Diagram

Let's assume:

```
cpp

class Car {
public:
    int speed;
    int* fuel;
};
```

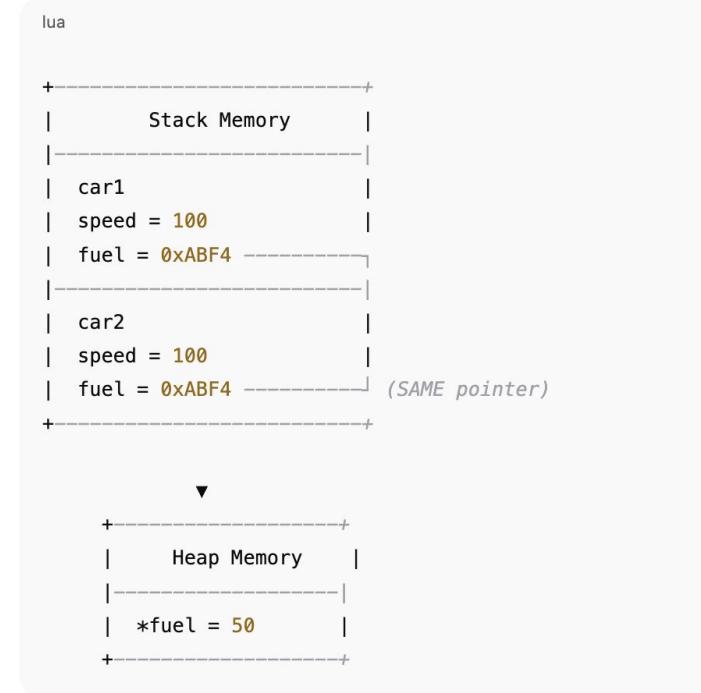
We do:

```
cpp

Car car1;
car1.speed = 100;
car1.fuel = new int(50);

Car car2 = car1; // shallow copy
```

### 🧠 Memory Layout:





# Memory Diagram 7

## 🧠 C++ Deep Copy – Memory Diagram

### 🎯 Scenario:

You have a class like:

```
cpp

class Car {
public:
    int speed;
    int* fuel;
};

You want to deep copy car1 to car2, so that:
```

- A new memory location is created for `fuel`.
- Both objects are fully independent.

### 📦 Deep Copy – Memory Layout

lua

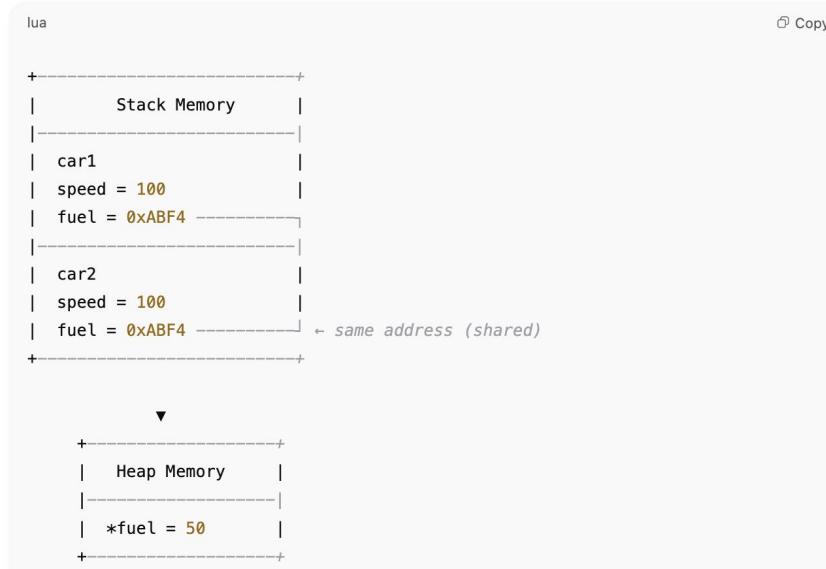
```
+-----+
|       Stack Memory      |
+-----+
| car1
| speed = 100
| fuel = 0xABF4
+-----+
| car2
| speed = 100
| fuel = 0xAC20
+-----+
```

```
▼           ▼
+-----+     +-----+
|   Heap Memory   |     |   Heap Memory   |
+-----+     +-----+
| *fuel = 50     |     | *fuel = 50
| (car1)        |     | (car2)        |
+-----+     +-----+
```



# Memory Diagram 8

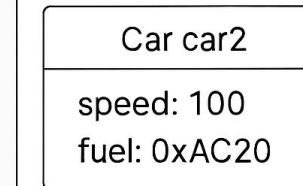
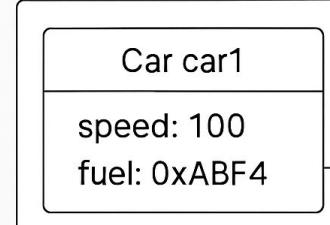
## A. Default = Operator (Shallow Copy) – Memory Diagram



! Risk: Double delete, shared mutation, crash

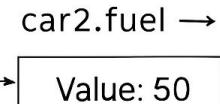
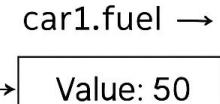
## DEEP COPY

### STACK MEMORY



### STACK MEMORY

### HEAP MEMORY





# Memory Diagram 9

## ⌚ Object Memory Layout on Stack

pgsql

STACK MEMORY (No vptr, No dynamic dispatch)

```
+-----+  
| Dog d; |  
+-----+  
|-----|  
| Base class part → Animal:  
|   - speak() function (statically bound)  
|-----|  
| Derived class part → Dog:  
|   - speak() function (separate, hides base)  
+-----+  
  
+-----+  
| Animal* ptr = &d; |  
+-----+  
|-----|  
| Points to base part of Dog object |  
| ptr->speak() → Calls Animal::speak() ONLY |  
+-----+
```

- 🚫 No vptr
- 🚫 No vtable
- 🚫 No runtime dispatch



## 📦 Memory Layout (NO virtual, STATIC binding)

pgsql

STACK MEMORY

```
+-----+  
| Dog d; |  
+-----+  
|-----|  
| Base class: Animal  
|   - age      = 5 |  
|-----|  
| Derived class: Dog  
|   - tailLength = 20 |  
+-----+  
  
+-----+  
| Animal* ptr = &d; |  
+-----+  
|-----|  
| Points to base part of Dog |  
| ptr->speak() → Calls Animal::speak() |  
+-----+
```

- 🚫 No vptr
- 🚫 No vtable
- ✅ Static memory layout





# Memory Diagram 10

## Memory Layout Diagram — No virtual

### ✓ STACK MEMORY

pgsql

Copy Edit

#### STACK MEMORY

```
+-----+
| Vehicle v;           |
+-----+
| start() -> Vehicle::start() |
+-----+
```

```
+-----+
| Car c;           |
+-----+
| start() -> Car::start() |
+-----+
```

```
+-----+
| Vehicle* ptr = &c;           |
+-----+
| Points to Car object       |
| Call ptr->start()         |
| ⚡ But it calls Vehicle::start() |
+-----+
```

#### NO VTABLE EXISTS

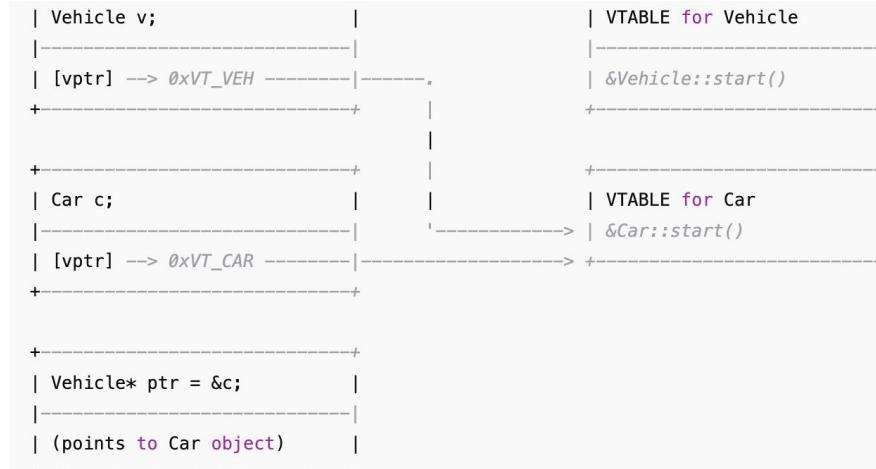
```
+-----+
| ⓘ No vtable generated    |
| ⓘ No vptr in object      |
| ⓘ No dynamic dispatch    |
+-----+
```



# Memory Diagram 11



## Correct Diagram – Stack Object & Vtable with Overriding

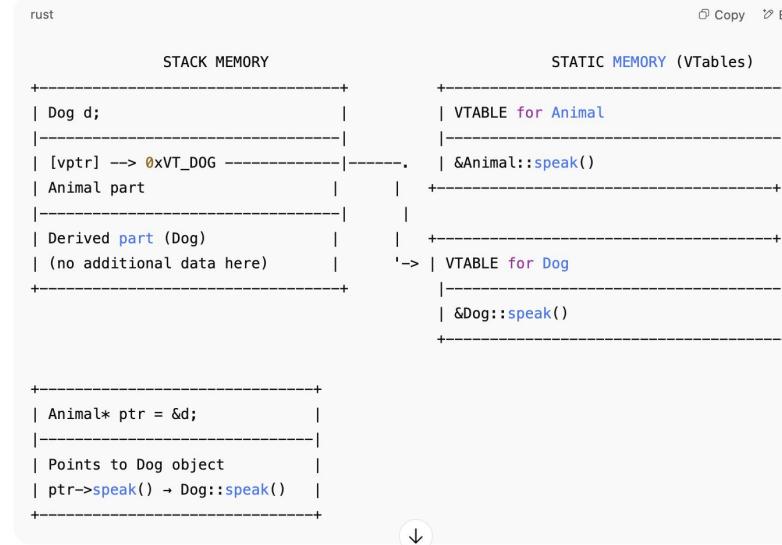




# Memory Diagram 11

## 🧠 Correct Diagram — Stack Object & Vtable with Overriding

### ✓ INHERITANCE OBJECT DIAGRAM (STACK + VTABLE)





# Memory Diagram 12

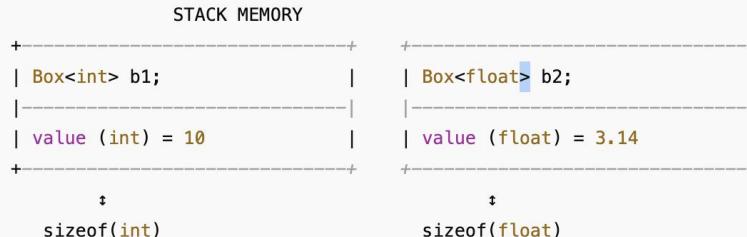
## 📌 Code

```
cpp

Box<int> b1;
Box<float> b2;
```

## ⌚ Memory Layout

pgsql



- ✓ No vptr
- ✓ No shared code
- ✓ No dynamic dispatch



## Template Function Example

cpp

```
template <typename T>
T add(T a, T b) {
    return a + b;
}
```



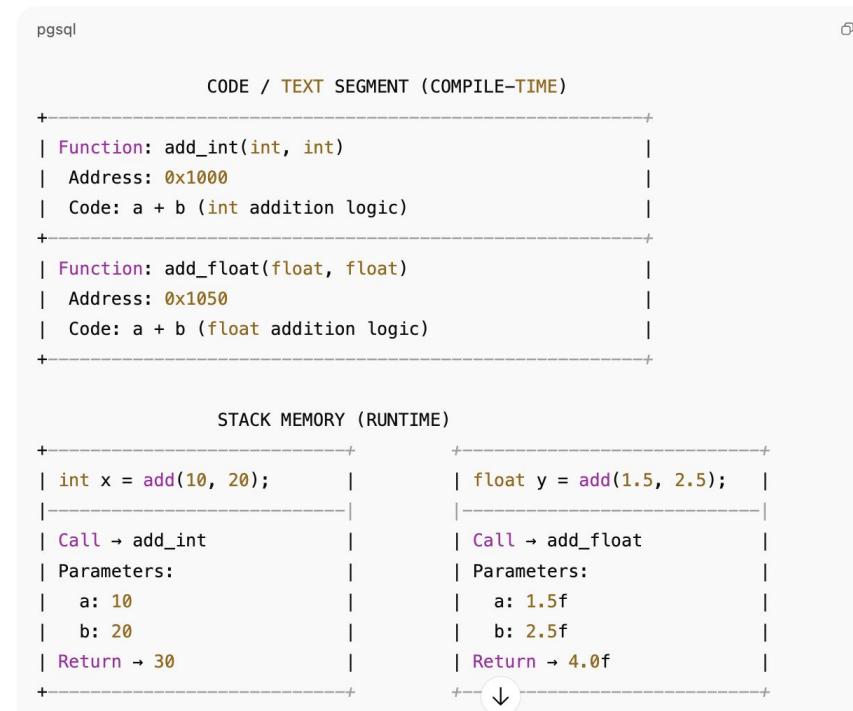
# Memory Diagram 12

## Template Function Example

cpp

```
template <typename T>
T add(T a, T b) {
    return a + b;
}
```

### STATIC + STACK MEMORY LAYOUT





## Reference

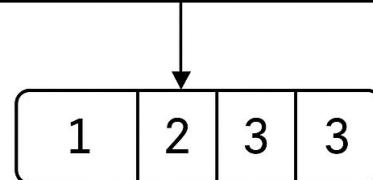
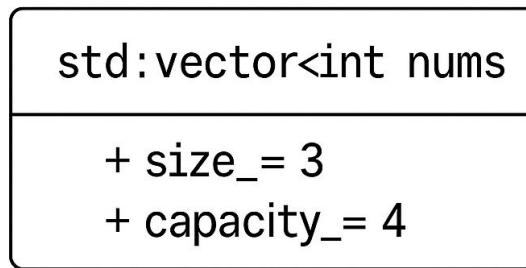
<https://chatgpt.com/share/686bce8a-e55c-8002-b96d-a0b3eaf1fe08>

<https://chatgpt.com/share/686bcea6-a310-8002-933a-f2489dd24615>



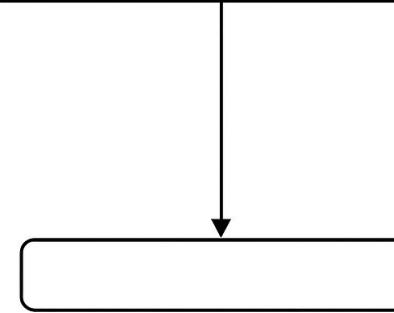
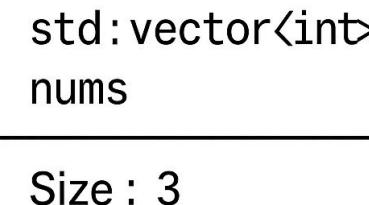
## Memory Diagram 13

**CODE / TEXT SEGMENT**  
(COMPILE-TIME)



Size 3

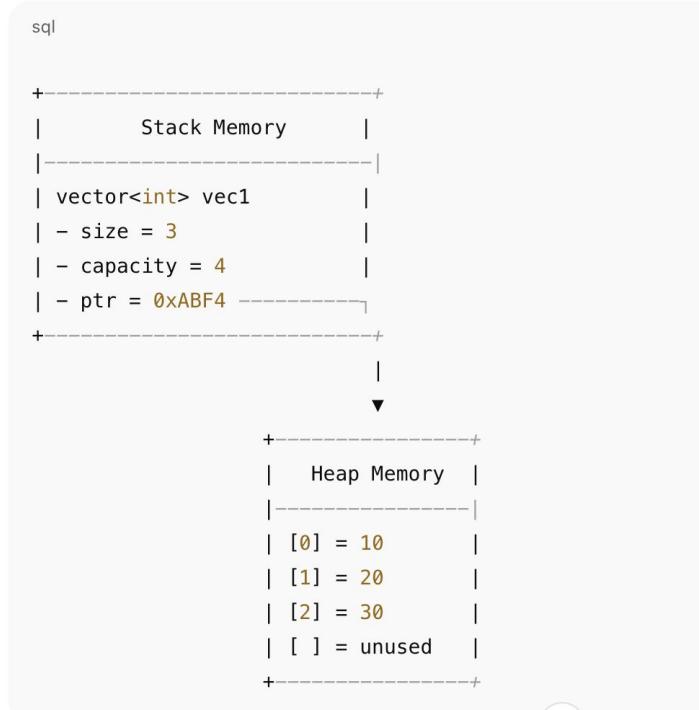
**STACK MEMORY**  
(RUNTIME)





# Memory Diagram 13

## A. std::vector – Memory Diagram



## B. std::list – Memory Diagram (Doubly Linked List)





# Memory Diagram 13

```
vector<int> vec = {10, 20, 30};
```



Heap Memory

[0] = 10

[1] = 20 —→ iter

[2] = 30

Heap Memory

## C. std::map – Memory Diagram (Red-Black Tree)

sql

```
+-----+  
| Stack Memory |  
+-----+  
| map<int, string> m1 |  
| - root = 0x1500 |  
+-----+
```

```
+-----+  
| Heap Memory |  
+-----+  
| [Node] |  
| key = 2 |  
| value = "B" |  
| left = 0x1400 |  
| right = 0x1600 |  
| parent = nullptr |  
+-----+
```

```
+-----+ +-----+  
| [Node] | | [Node] |  
| key = 1 | | key = 3 |  
| value = "A" | | value = "C" |  
| left = nullptr | | right = nullptr |  
| right = nullptr | | left = nullptr |  
| parent = 0x1500 | | parent = 0x1500 |  
+-----+ +-----+
```

Copy Edit



# Memory Diagram 14 - Iterator

```
+-----+  
|      Stack Memory      |  
+-----+  
| vec  
|   ptr = 0xC0DE  
|   size = 3  
|   capacity = 3  
|  
| it1  
|   current = 0xC0DE  
|  
+-----+  
| it2  
|   current = 0xC0DE      | ← same address (shared start)  
+-----+  
  
▼  
  
+-----+  
|      Heap Memory      |  
+-----+  
| *vec (0xC0DE)  
|   [10] [20] [30]  
+-----+  
  
! Risk: If `vec` is modified (e.g., resized), iterators may become **dangling** .
```





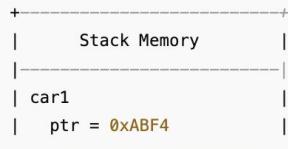
# Memory Diagram 15

## 1. std::unique\_ptr<T> — Owns Object Exclusively

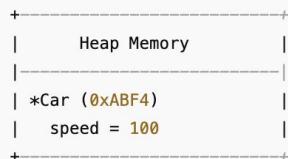
cpp

```
std::unique_ptr<Car> car1 = std::make_unique<Car>(100);
```

pgsql



▼



✓ Ownership: Single owner (car1)

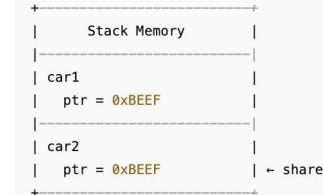
✗ No copies allowed

## 2. std::shared\_ptr<T> — Shared Ownership (Ref Counted)

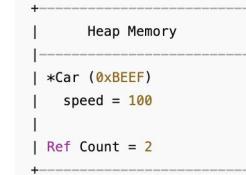
cpp

```
std::shared_ptr<Car> car1 = std::make_shared<Car>(100);
std::shared_ptr<Car> car2 = car1;
```

sql



▼



✓ Shared ownership

✗ Automatically deleted when Ref Count = ↓



# Memory Diagram 15

## 3. std::weak\_ptr<T> — Non-owning Observer

cpp

```
std::shared_ptr<Car> car1 = std::make_shared<Car>(100);
std::weak_ptr<Car> weak_car = car1;
```

pgsql

Stack Memory
car1
ptr = 0xDEAD
weak_car
ptr = 0xDEAD (weak) ← no ref count

▼

Heap Memory
*Car (0xDEAD)
speed = 100
Ref Count = 1



# Memory Diagram 16

pgsql

```
+-----+  
|      Stack Memory      |  
+-----+  
| Main Thread           |  
|   counter = 0xBEEF    |  
+-----+  
| Thread t1             |  
|   access → 0xBEEF     |  
+-----+  
| Thread t2             |  
|   access → 0xBEEF     |  
+-----+
```

▼

```
+-----+  
|      Heap Memory       |  
+-----+  
| *counter (0xBEEF)     |  
|   value = 0 initially  |  
|   value = 1 or 2 (RACE CONDITION) |  
+-----+
```

```
+-----+  
|      Stack Memory      |  
+-----+  
| Main Thread           |  
|   counter = 0xBEEF    |  
|   mtx = std::mutex    |  
+-----+  
| Thread t1             |  
|   lock_guard(mtx)     |  
|   access → 0xBEEF     |  
+-----+  
| Thread t2             |  
|   lock_guard(mtx)     |  
|   access → 0xBEEF     |  
+-----+
```

▼

```
+-----+  
|      Heap Memory       |  
+-----+  
| *counter (0xBEEF)     |  
|   value = 0 → 2 safely |  
+-----+
```

✓ Thread-safe via `lock`  
! Slower due to locking



# Memory Diagram 17

## Correct Singleton Memory Diagram (Single Thread)

Code:

```
cpp

class Singleton {
public:
    static Singleton* getInstance() {
        if (instance == nullptr)
            instance = new Singleton();
        return instance;
    }

private:
    static Singleton* instance;
    Singleton() {} // private constructor
};

// Define static member
Singleton* Singleton::instance = nullptr;
```

When you call:

```
cpp

Singleton* s1 = Singleton::getInstance();
Singleton* s2 = Singleton::getInstance();
```

## Simple Stack + Heap Diagram:

sql



▼



Static variable:

Singleton::instance → 0xBEEF



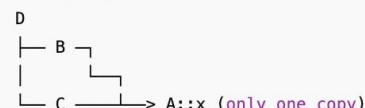
# Memory Diagram 18

## Virtual Inheritance — Memory Diagram

sql

```
+-----  
|           Stack  
|-----  
| obj (of class D)  
|-----  
| B → [vptr to A]  
| C → [vptr to A]  
|           ↑ shared  
| A::x = 10  
+-----
```

## Memory layout:



✓ One shared A instance → unambiguous



# Memory Diagram 18

## ✓ When VTables Are Used

Usage	VTable Exists?	Dispatch Type
<code>&amp;dyn Trait</code>	<input checked="" type="checkbox"/> Yes	Dynamic
<code>Box&lt;dyn Trait&gt;</code>	<input checked="" type="checkbox"/> Yes	Dynamic
<code>Arc&lt;dyn Trait&gt;</code>	<input checked="" type="checkbox"/> Yes	Dynamic
<code>&amp;impl Trait</code>	<input type="checkbox"/> No	Static
<code>struct T; impl Trait for T</code>	<input type="checkbox"/> No	Static



# Memory Diagram 18

```
trait Pilot {
    fn fly(&self);
}

trait Engineer {
    fn work(&self);
}

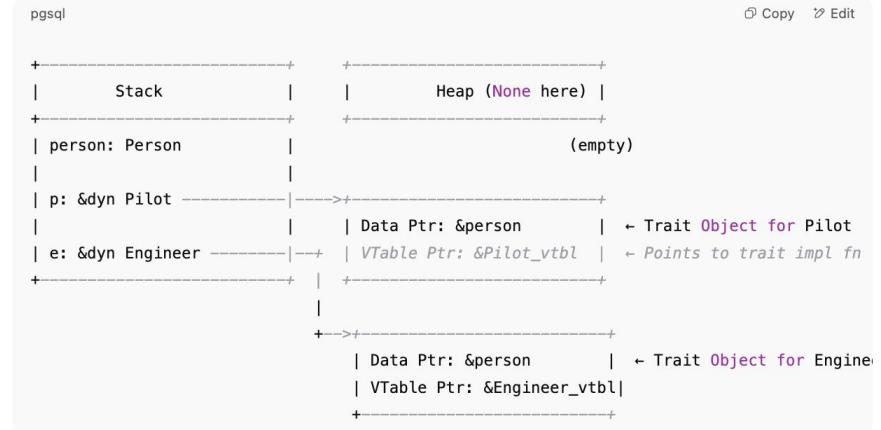
struct Person;

impl Pilot for Person {
    fn fly(&self) {
        println!("Pilot flying");
    }
}

impl Engineer for Person {
    fn work(&self) {
        println!("Engineer working");
    }
}

fn main() {
    let person = Person;
    let p: &dyn Pilot = &person;      // Trait object
    let e: &dyn Engineer = &person;   // Trait object
}
```

## 🧠 Memory Diagram (Simplified)





# Best Real Time Formula

## 📊 Refined Formula for Rust Development Thinking:

#	Concept	Is it Core to Rust?	Notes
1	Memory Allocation (Stack / Heap)	✓ Core	Ownership, borrowing, move semantics are directly tied to memory model.
2	Memory Initialization	✓ Core	Rust enforces mandatory initialization before use.
3	Condition Handling (if, match)	✓ Core	Pattern matching is a powerful feature in Rust.
4	Communication (Function Calls, Return)	✓ Core	Functions, methods, trait methods are essential, with explicit data movement (ownership).
5	Loops (for, while)	✓ Core	Rust also promotes iterator-based loops (preferred idiom).
6	Exception Handling	⚠ Partial	Rust uses <code>Result</code> and <code>Option</code> for error handling instead of exceptions. <code>panic!</code> is for unrecoverable errors.
7	Business Logic	✓ Domain-specific	Core to solving actual problems. Rust doesn't enforce domain patterns but is capable.
8	Transformation (Data Processing)	✓ Strong	Iterators, functional programming patterns (map, filter, fold) make Rust great for transformation logic.
9	Resources (Files, DB, Threads, OS)	✓ Strong	<code>std::fs</code> , <code>tokio</code> , <code>threads</code> , <code>async</code> handle resource interaction safely.
10	Domain Knowledge	✓ Important	Language-agnostic but essential for applying Rust to real-world problems.
11	Functionality (Deliverables)	✓ Result ↓	This is your outcome. All above points lead to implementing meaningful features.



# Best Real Time Formula

## Suggested Rust Mindmap / Formula:

plaintext

1. Ownership / Borrowing / Lifetimes
2. Memory Management (Stack / Heap / Move / Copy)
3. Pattern Matching / Enums
4. Error Handling (Result, Option)
5. Traits / Trait Objects (Abstraction)
6. Concurrency / Async (Safe Parallelism)
7. Functional Data Transformations (Iterators)
8. Resource Safety (Files, Threads, DB)
9. Business Logic & Domain Modeling
10. Deliver Features Safely & Performantly