

# ASSIGNMENT 4

VENKATESH  
RAIZADA7

1.

$$(a) \forall x \forall y (x+y = y+x) \wedge (x \times y = y \times x)$$

$$(b) \forall x \forall y \exists z (x+y=z) \wedge (x \times y=z) \wedge \text{Integer}(z)$$

$$(c) \forall x \forall y \forall z (x+y) \times z = (x \times z) + (y \times z)$$

$$(d) \forall x \forall y \forall z ((x+y)+z = x+(y+z)) \wedge ((x \times y) \times z = x \times (y \times z))$$

$$(e) \forall x (x+0=x) \wedge (x \times 1=x)$$

e

$$\alpha = \forall x (P(x) \vee Q(x))$$

$$\beta = \forall x P(x) \vee \forall x Q(x)$$

Let us create a truth table

$P(x)$	$Q(x)$	$\alpha$	$\beta$
T	T	T	T
T	F	T	T
F	T	T	T
F	F	F	F

$\therefore$  If  $\alpha$  is true,  $\beta$  is also true

$\therefore \alpha \models \beta$

3. Following are abbreviations for various words used in the question.

HS : Han Solo.

PL : Princess Leia

OW : Obi-Wan Kenobi

MF : Millennium Falcon

RA : Rebel Alliance

LS : Light Saber

a. Owns (HS, MF)

b. Unhappy (PL)

c. Loves (PL, HS)

d.  $\forall x (Owns(x, MF) \vee Unhappy(x) \Rightarrow Visits(x, OW))$

e.  $\forall x Visits(x, OW) \Rightarrow Wise(x)$

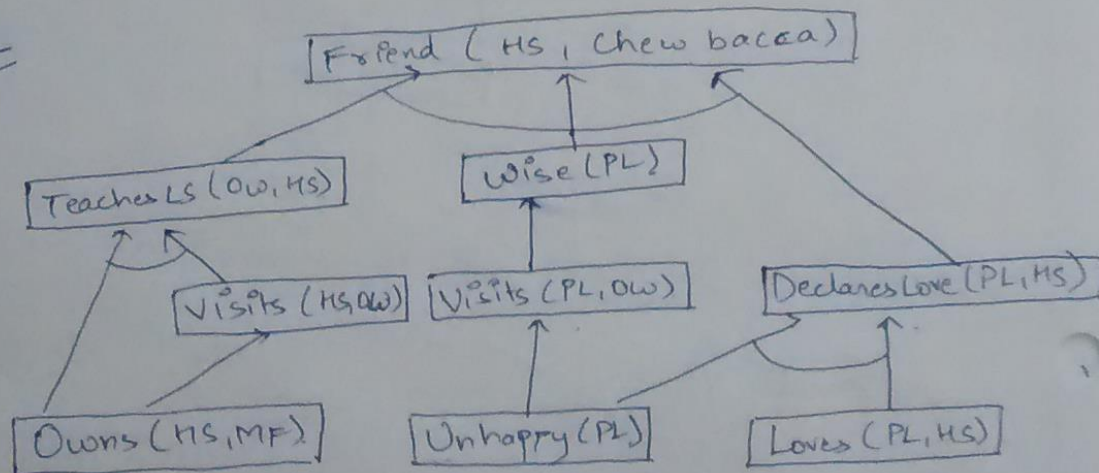
f.  $\forall x (Owns(x, MF) \wedge Visits(x, OW) \Rightarrow Teaches LS(OW, x))$

g.  $\forall x ((Unhappy(x) \vee Owns(x, MF)) \wedge Teaches LS(OW, x) \Rightarrow Joins(x, RA))$

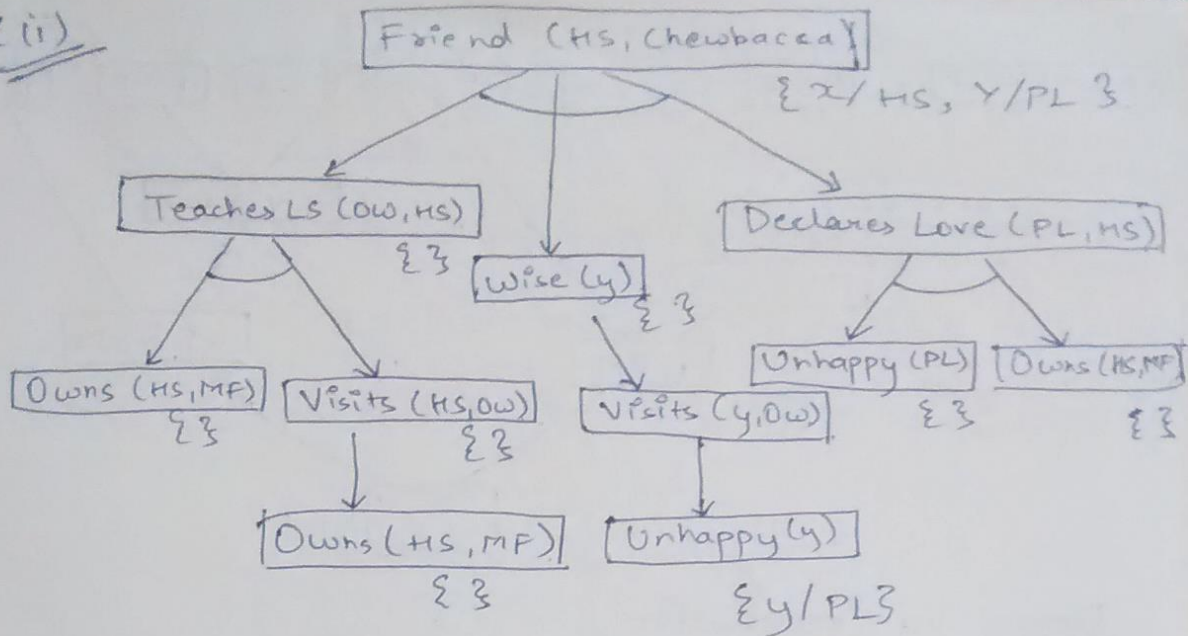
h.  $\forall x \forall y Unhappy(x) \wedge Loves(x, y) \Rightarrow Declares Love(x, y)$

i.  $\forall x \forall y (Teaches LS(OW, x) \wedge Declares Love(y, x) \wedge Wise(y) \Rightarrow Fend(x, Chewbacca))$

(i)



(ii)



4.

- A : Zeus is able to prevent evil  
 W : Zeus is willing  
 P : Zeus Prevents evil.  
 I : Zeus is Impotent.  
 M : Zeus is Malevolent.  
 E : Zeus exists

Convert all statements to propositional logic:-

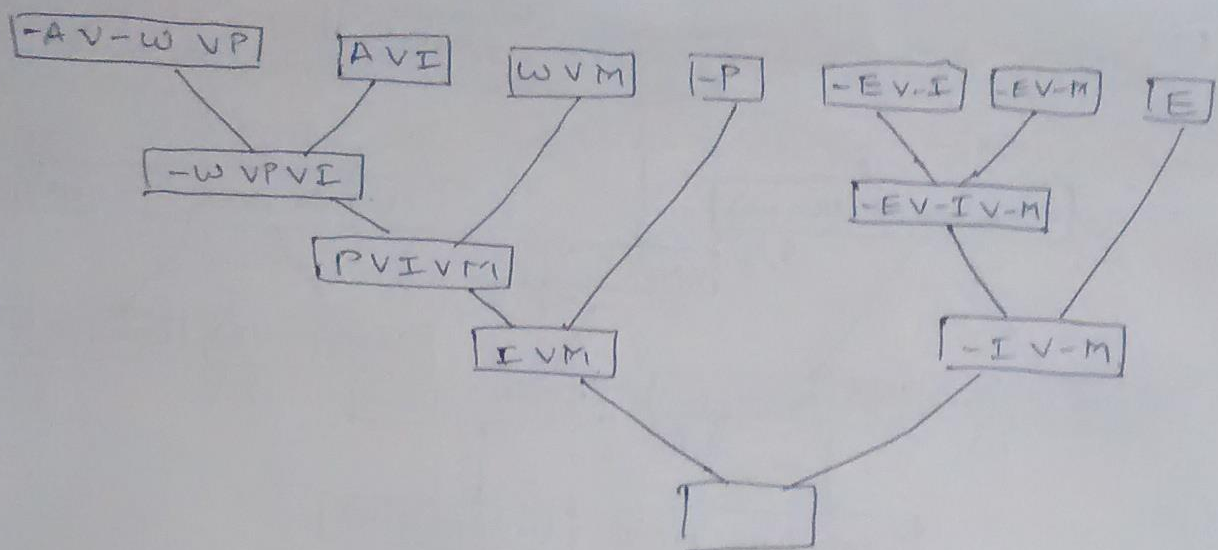
a.  $A \wedge W \Rightarrow P$   
 $\neg(A \wedge W) \vee P$   
 $\neg A \vee \neg W \vee P$

b.  $\neg A \Rightarrow I$   
 $\neg(\neg A) \vee I$   
 $A \vee I$

c.  $\neg W \Rightarrow M$   
 $\neg(\neg W) \vee M$   
 $W \vee M$

d.  $\neg P$

e.  $E \Rightarrow \neg M \wedge \neg I$   
 $\neg E \vee (\neg M \wedge \neg I)$   
 $(\neg E \vee \neg M) \wedge (\neg E \vee \neg I)$



Hence by resolution.

$\bar{E}$  : Zeus does not exist.

5 Using resolution to show negation leads to contradiction.

(a)  $S: \forall x (P(x) \Rightarrow P(x))$

$\bar{S}: \neg \forall x (P(x) \Rightarrow P(x))$

(Removing biconditionals)

$\neg \forall x (\neg P(x) \vee P(x))$

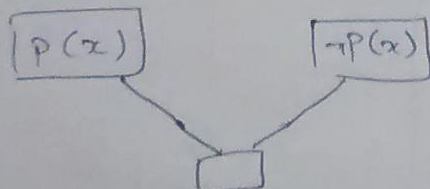
$\exists x \neg (\neg P(x) \vee P(x))$

(Applying De Morgan's Law)

$\exists x (P(x) \wedge \neg P(x))$

(Drop universal identifiers)

$P(x) \wedge \neg P(x)$



$\therefore \bar{S}$  is a contradiction

$\therefore S$  is valid.



(b)  $S: \neg(\neg \exists x P(x)) \Rightarrow (\forall x \neg P(x))$

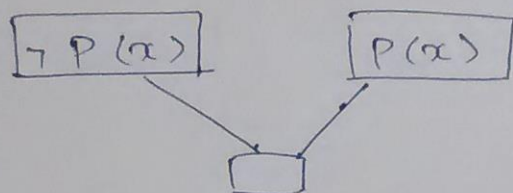
$\bar{S}: \neg((\neg \exists x P(x)) \Rightarrow (\forall x \neg P(x)))$  (Remove Biconditionals)

$\neg(\exists x P(x) \vee \forall x \neg P(x))$  (Applying De Morgan's)

$\neg \exists x P(x) \wedge \neg \forall x \neg P(x)$

$\forall x \neg P(x) \wedge \exists x P(x)$  (Drop universal quantifiers)

$\neg P(x) \wedge P(x)$



$\therefore \bar{S}$  is a contradiction.  
 $\therefore S$  is valid.

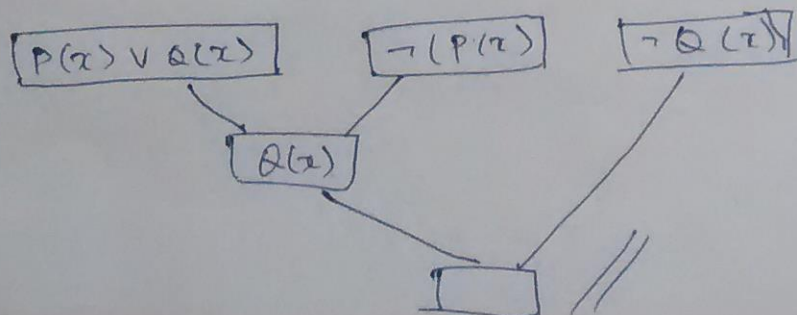
(c)  $S: (\forall x(P(x) \vee Q(x))) \Rightarrow ((\forall x P(x)) \vee (\exists x Q(x)))$

$\bar{S}: \neg((\forall x(P(x) \vee Q(x))) \Rightarrow ((\forall x P(x)) \vee (\exists x Q(x))))$

$\neg(\neg(\forall x(P(x) \vee Q(x))) \vee ((\forall x P(x)) \vee (\exists x Q(x))))$

$\forall x(P(x) \vee Q(x)) \wedge (\exists x \neg P(x)) \wedge (\forall x \neg Q(x))$

$(P(x) \vee Q(x)) \wedge (\neg P(x)) \wedge (\neg Q(x))$



$\therefore \bar{S}$  is a contradiction.

$\therefore S$  is valid.