# Dijkstra's Algorithm & Poseds code

Functions:

Distance (upt 1, upt 2) - cak. distance from one walkfupt to another

Xy locations of this is cost in our case

Still build be birm, energy, de)

Collision Check (obstacle list, gold bounds, current location)

Return True/ False depending on if current location

is a valid location

Calculate Noch Index (corrent work, gold information)

Previously defined numbering (conting gold bounds, step size

Dijkstra (Start docation, goal location, gold into, obstacle 1854)

+ Roas the Dijkstra algorithm, returns the dwired path/wayspoints

\* Optional Node Neighbur Cost (whent made, grid into, obstade list)

I Calculates cost of moving to all neighbor newly I updates unulsited with distibution as appropriate

### New Python Items:

Classes - many use, but in our case its an easy way to create a costom variable structure for storing our nodes

Example:

what ever variable name you would like

class data:

ruariables you want to pass into class

def -- init -- (self 2, input 1, input 2, cheese):

self. x1 = input 1 Name you use to reference

self. y1 = input 2

self. quality = cheese

then usage of class: temp = data (1.0, 3.0, 4) Print temp-quality + 4

#### Dictionaries.

Lists/arrays you are used to are indexed by numbers. And you if you want to stone doming [3] = 5, then doming[0] - [2] must be created in memory (even of o vila).

A dictionary uses keys for indexing. Keys can be strings, numbers, ... And importantly, because the indexing is with keys, its not dependent on having a monotonially thereasing inches.

We will use a det. to store our visited and unvisited wales then we can search through the det. as needed:

### U Sage

Unvisited nodes = dret () + initialize the dictronomy corrent\_worde = worde (0,0,0,-1) ~ a worde class to store Node-index = calc\_index (connect worle, -.) y counds, cost, ppanent Unvisited\_ modes { node\_index] = corrent\_ mode is get work index of convent work 4 store correct mode at Key index vole-Index

Minimum: Want to search through unvisited worlds to find Next work with lowest cost to visit

set search of keys min (example, key = lambda x: x['key'] Edictionary to search

For our use;

index reade to visit = mm (unuisited modes, key = lambda, X: unuisited modes [x]. cost

what And trying to find min of.

## Dijkstra Forction "Bred code"

- Initialize visited and unvisited modes dictionaries

- Current Node = Storting point, no travel cost, Parent work index = 1

- Calculate location index of corrent mode (based upon where node exists in grid, need to create numbering scheme) + calcincher function

when finding I find path, -1 is cary was to know me have found start points

Put corrent made in unvisited males dictionary

While correct wale not egold to god location

Correct make = node number of minimum cost in unvisited makes

Correct make = unvisited makes I corrected index

pot/stone correct works in visited makes

delete correct make from unvisited makes = Don't want to

visit make again

check if correct make is goal location, if so

break out of loop.

\* Can make this a function call

with cornert node, need to calculate costs to travel to all neighbors

From corrent work, of connected water. One approach is
to cycle shrough all of them. Will have to check if
any weighbor is an obstacle, out of bounds, or the nock itself
or i in range (-grid size to grid size)
for j in range (-grid size to grid size)

- Calculate cost to go from correct made to correct +(i,j)
= temp made = node (x,y, cost, correct made make) parent

- Calculate temp voile location index - used for storing/looking - Check if temp wale is valid a Not same location as corrent woole + Not in obstacle > Not out of gold boundary - if check is good, see if temp node exists in unvisited water dictionary If remp wode does exist in unvisited wides and cost from temp node is less than node in unvisited modes -> then update cost and parent wale index In unvisited modes dictionary If temp made does not exist in unvisited moder -> then work is new. Stick/stone temp wode in unvisited nodes at index => unvisited nodes[temp Nacle Loop Around + After this the goal has been found = temp noch Time to find the path from goal to start - Initialize path x ? y as a list or array with goal counds - Get location index of goal location Thirte Parent Noble Index & -1 > Remember this is the - Add > visited nodes [ferent rule Index] to path x ? y -New Parent work index = parent of visited\_note [parent work index] Loup around is = visited\_work [parent male helow] - parent Plat path (use Mat Plat Lib)