

9<sup>th</sup> Aug 2023

## TRIES - 1 (TRYs)

Tries → Tree like DS that stores data from top to bottom.

Trie of characters (a-z)

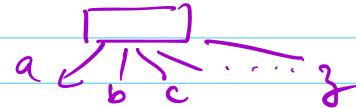
class Node {

Node[26] child // a→0  
b→1

boolean isEnd :

} initially if child[i] = null

default isEnd = false.



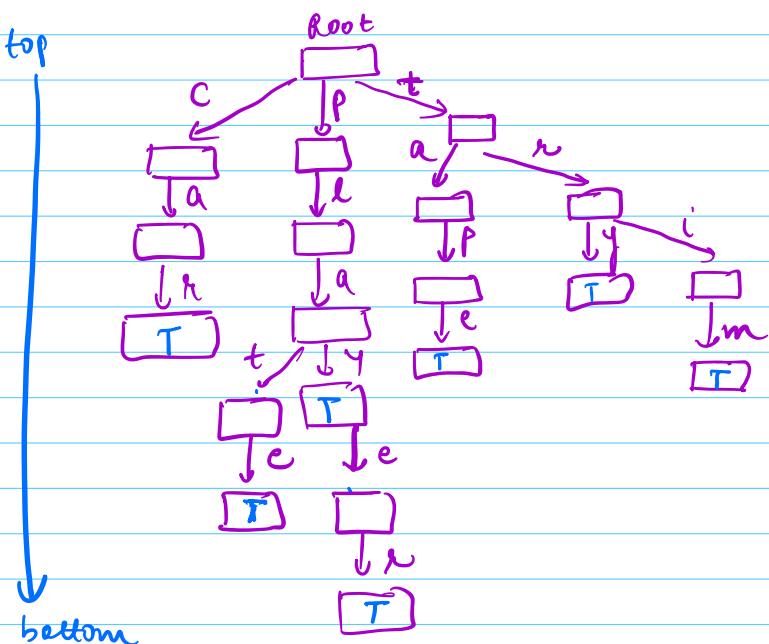
Real life use cases

- ① Auto complete
- ② Word search
- ③ Spelling checker

- ④ Internet routing  
Encryption.

List of words

- 1> tape
- 2> try
- 3> trim
- 4> play
- 5> plate
- 6> car
- 7> player.



## Search

check("cat")  $\Rightarrow$  false (not present)

If all characters of the word are not travelled  
 $\Rightarrow$  not present.

check("try")  $\Rightarrow$  true.}

check("tri")  $\Rightarrow$  false } If all characters of the word  
are travelled then, it is either present OR  
it is a prefix of any word present.

boolean search (root, word) {

temp = root;

for i  $\rightarrow$  0 to (word.length - 1) {

ch = word[i]

if (temp.child[ch - 'a'] == null)  
return false;

temp = temp.child[ch - 'a'];

}

return temp.isEnd;

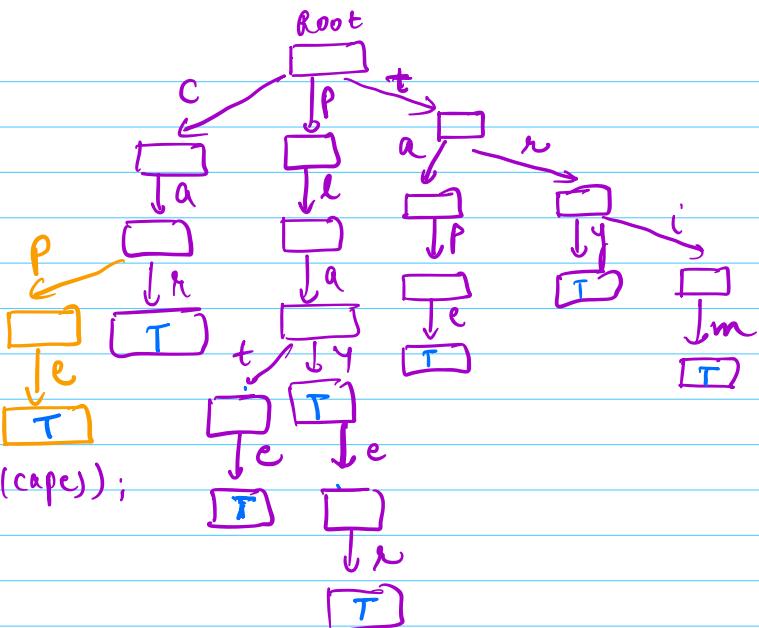
Tc: O(word length)

Hashset  $\rightarrow$  complete word is required to check.

Trie  $\rightarrow$  can check as we type characters  $\rightarrow$  faster.

### List of words

- 1> tape
- 2> try
- 3> time
- 4> play
- 5> plate
- 6> car
- 7> player.
- 8> cape (mient (cape));



```

void insert (root, word) {
    temp = root;
    for i=0 to (word.length-1) {
        ch = word[i]
        if (temp.child [ch-'a'] == null)
            temp.child [ch-'a'] = new Node()
        temp = temp.child [ch-'a'];
    }
    temp.isEnd = true; // temp.freq++;
}
  
```

Tc: O(N + word length), Sc: O(N + word length)

Keep track if a word is inserted multiple times. i.e.

freq  
~~boolean isEnd~~ → **int freq**

class Dict {

    Node (root);

    Dict () {

        root = new Node();

    }

    Dict d = new Dict();  
    d.getRoot();

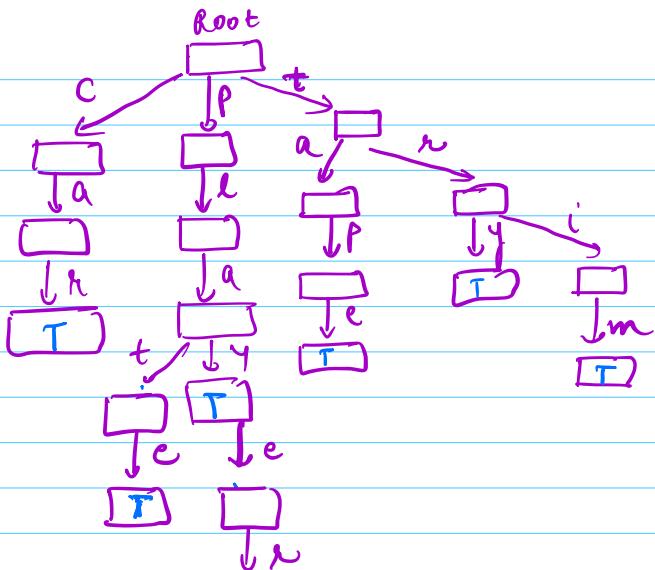
    insert (root, "x");

8:23 am IST

### Deletion

delete("trim")

Sell  $\rightarrow$  Travel to the end of the word & update isEnd to false.



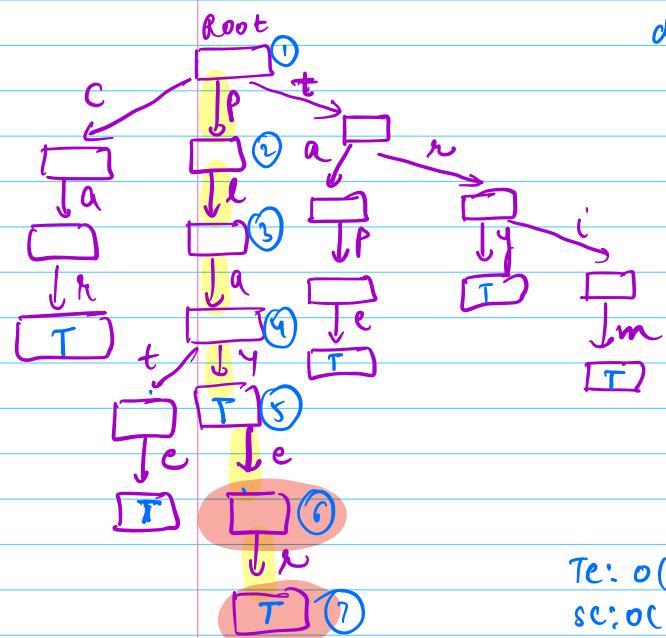
### Traversals:

a) Travel complete  $\Rightarrow$  valid prefix of any input

b) Space wastage

Sell  $\rightarrow$  delete un-utilized nodes

only leaf node.



delete(trim)

1> Travel all nodes & insert in the stack.

2> Update isEnd = false for last node.

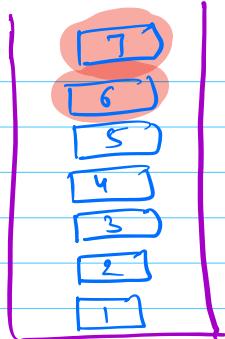
3> Pop elements from stack, if leaf node & isEnd = false  
     $\Rightarrow$  delete it.

$\Rightarrow$  if  $\text{child}[i] = \text{null}$ .

Tc: O( word length )

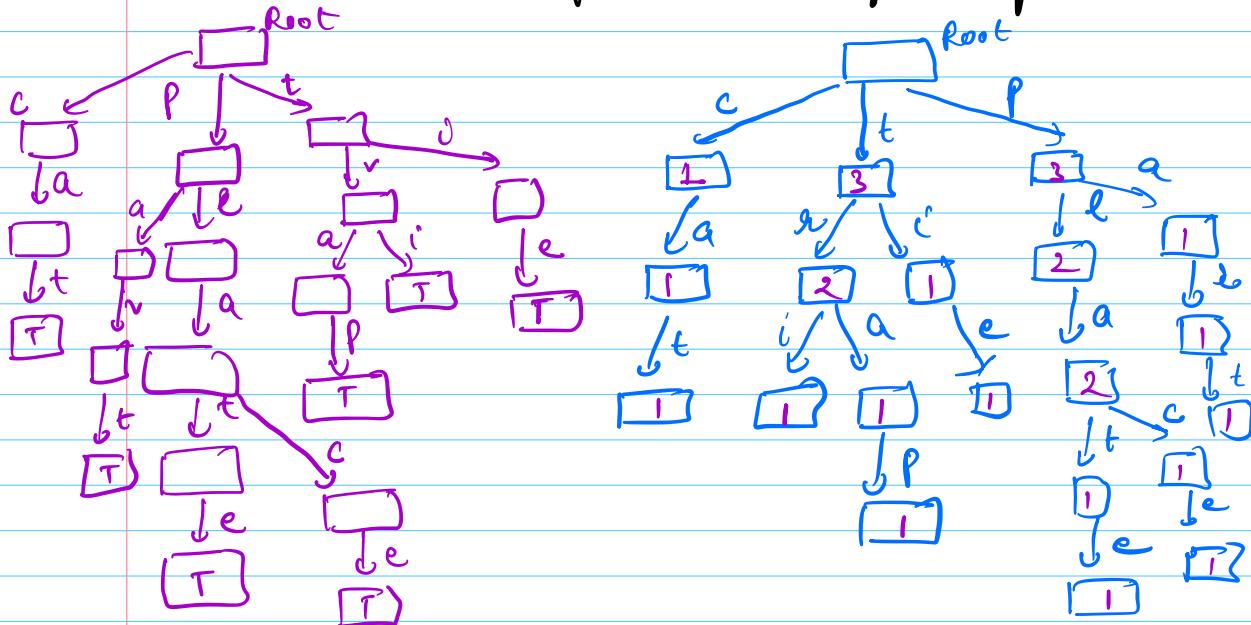
Sc: O( word length )

delete player



- Q Given a list of words, find shortest prefix of each word that uniquely defines the word.  
(no word is a prefix of another word)

Eg → ["tri", "trap", "plate", "cat", "part", "place", "tie"]  
tri      tra      plat      c      pa      place      ti



- 1)  $\dagger$  nodes  $\rightarrow$  store # times that node is visited.  
(while inserting words in the trie)
- 2) Travel the word till a node with freq = 1 is reached

Tc:  $O(N \times \text{word length})$   
Sc:  $O(N \times \text{word length})$

```

void insert (root, word) {
    temp = root;
    for i = 0 to (word.length - 1) {
        ch = word[i];
        if (temp.child [ch - 'a'] == null)
            temp.child [ch - 'a'] = new Node();
        temp = temp.child [ch - 'a'];
        temp.freq++;
    }
}

```

```
string search ( root, word ) {  
    temp = root;  
    string ans = "";  
    for ( i = 0 to ( word.length - 1 ) ) {  
        ch = word[i];  
        temp = temp. child [ ch - 'a' ];  
        ans += ch;  
        if ( temp.freq == 1 )  
            return ans;  
    }  
}
```

```
class Node {
```

```
    Node[26] child;  
    int freq;
```

```
}
```