

1/9/2023

Contest 6 discussion

Q1 Unique paths 3.

Given $n \times m$ grid

- $A[i][j]$
- 0 empty cells
 - 1 start cell (only 1)
 - 2 end cell (only 1)
 - -1 obstacles (can't walk over it)

Return the no of 4 directional walks from starting cell to ending cell, that walk over every non-obstacle cell only once.



```
int empty = 1  
int UP ( int [][] grid ) {  
    n = grid.length, m = grid[0].length.
```

```
    for ( i = 0 to n ) {  
        for ( j = 0 to m ) {  
            if ( grid[i][j] == 0 ) {  
                empty++;  
            }  
            else if ( grid[i][j] == 1 ) {  
                sx = i;  
                sy = j;  
            }  
        }  
    }  
}
```

```

    } solve( grid, sx, sy );
}

```

```

void solve( grid[ ][ ], x, y ) {

```

```

    if( x < 0 || x > grid.length || y < 0 || y > grid[0].length || grid[x][y] < 0 )
        return;

```

```

    if( grid[x][y] == 2 ) {

```

```

        if( empty == 0 )
            ans++;

```

```

        return;
    }

```

```

    grid[x][y] = -2 // Do

```

```

    empty--;
    // Do

```

```

    solve( grid, x+1, y );

```

```

    solve( grid, x-1, y );

```

```

    solve( grid, x, y+1 );

```

```

    solve( grid, x, y-1 );

```

```

    grid[x][y] = 0

```

```

    empty++;
    } Undo.

```

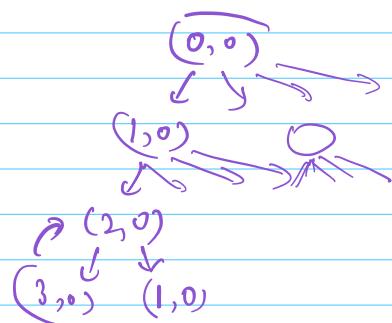
	0	1	2	3
0	█			
1			█	

①

$\text{int[] r} = \{1, -1, 0, 0\}$
 $\text{int[] c} = \{0, 0, +1, -1\}$
 $\text{new r} = x + r[i]$
 $\text{new c} = y + c[i]$

empty = 10

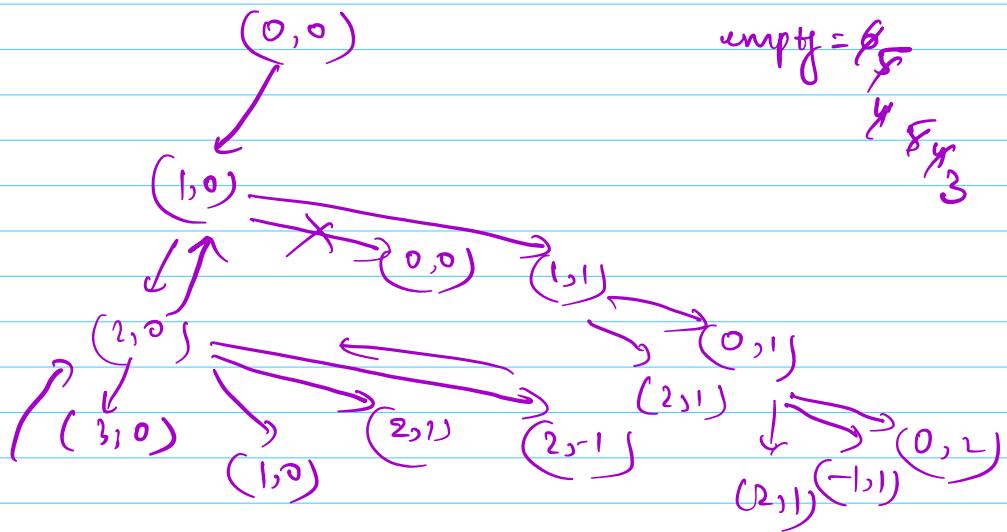
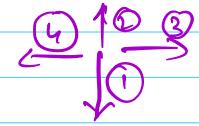
TC $\leq O(mn)$



	0	1	2
0	-2	-2	
1	-2	-2	
2	0	0	0

empty = 1

$$sx = 0, sy = 0$$



Q2

Mice & pizza

- each pizza needs some amount of time to be eaten.
- find max no. of pizzas we can eat in B minutes.

$$A = [1, 2, 3], \quad B = 3, \quad \text{ans} = 2, \quad A[i] \approx 10^9$$

$$A = [10, 2, 6], \quad B = 13, \quad \text{ans} = 2, \quad B \approx 10^9$$

- ① sort the array
- ② add & compare.

int solve (int A[], int B) {

```

sort (A);
int ans = 0;
long sum = 0;
for (int i = 0; i < N - 1; i++) {
    sum += A[i];
    if (sum > B) {
        break;
    }
    ans++;
}
return ans;
}
```

greedy

we can subtract as well.

TC: $O(N \log N + N) \approx O(n \log N)$

$$B = 2 \times 10^9$$

$$A = [2 \times 10^9, 2 \times 10^9]$$

Q3
Wood cost
Find min cost to join woods into a single piece.

Cost of joining 2 woods = sum of len of the 2 woods.

$$A[i] \approx 10^9$$

$$A = [1, 2, 3, 4, 4, 6] \rightarrow A = [\underbrace{10^9}, \underbrace{10^9}, \underbrace{10^9}, \underbrace{10^9}, \underbrace{10^9}]$$

$$\begin{matrix} & 1 \\ & | \\ [3, 3, 4, 4, 6] \end{matrix}$$

$$\begin{matrix} & 4 \\ & | \\ [4, 4, 6, 6] \end{matrix}$$

$$\begin{matrix} & 6 \\ & | \\ [6, 6, 8] \end{matrix}$$

$$\begin{matrix} & 8 \\ & | \\ [8, 12] \end{matrix}$$

$$[10^9, 10^9, 10^9, 2 \times 10^9]$$

$$[10^9, 2 \times 10^9, 2 \times 10^9]$$

$$[2 \times 10^9, 3 \times 10^9]$$

$$[5 \times 10^9]$$

$$[20]$$

long solve (init[] A) {

Priority Queue <long> pq;

for (int i = 0 to N-1)
pq.add(A[i]);

long ans = 0

while (pq.size > 1) {

long s = pq.poll() + pq.poll();

pq.add(s);
ans += s;

3

return ans;
}

for (i=0 to A.size(); i++) {

 Long x = new Long (A.get(i));
 pq.add(x);

}

int x

Integer

Object

} Autoboxing.