

9th Aug 2023

TRIES - 1 (TRYs)

Tries → Tree like DS that stores data from top to bottom.

Trie of characters (a-z)

class Node {

Node[26] child // a → 0
 b → 1

boolean isEnd ;

// initially if i child[i] = null

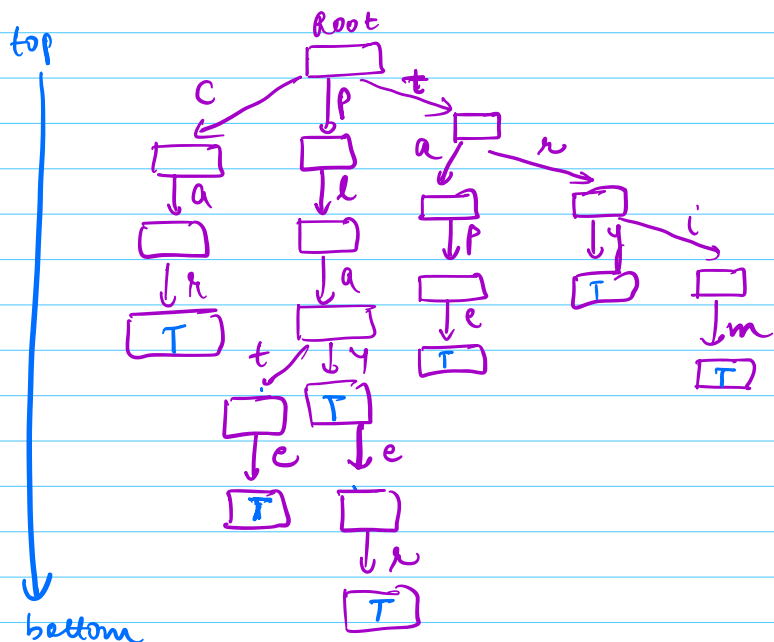
default isEnd = false.

Real life uses

- ① Auto complete
- ② Word search
- ③ Spelling checker
- ④ Internet routing
- ⑤ Encryption.

List of words

1> tape
2> try
3> trim
4> play
5> plate
6> car
7> player.



Search

check("cat") \Rightarrow false (not present)

If all characters of the word are not travelled
 \Rightarrow not present.

check("try") \Rightarrow true.

check("tri") \Rightarrow false } If all characters of the word
are travelled then, it is either present OR
it is a prefix of any word present.

```
boolean search (root, word) {  
    temp = root;  
    for i  $\rightarrow$  0 to (word.length - 1) {  
        ch = word[i]  
        if (temp.child[ch - 'a'] == null)  
            return false;  
        temp = temp.child[ch - 'a'];  
    }  
    return temp.isEnd;  
}
```

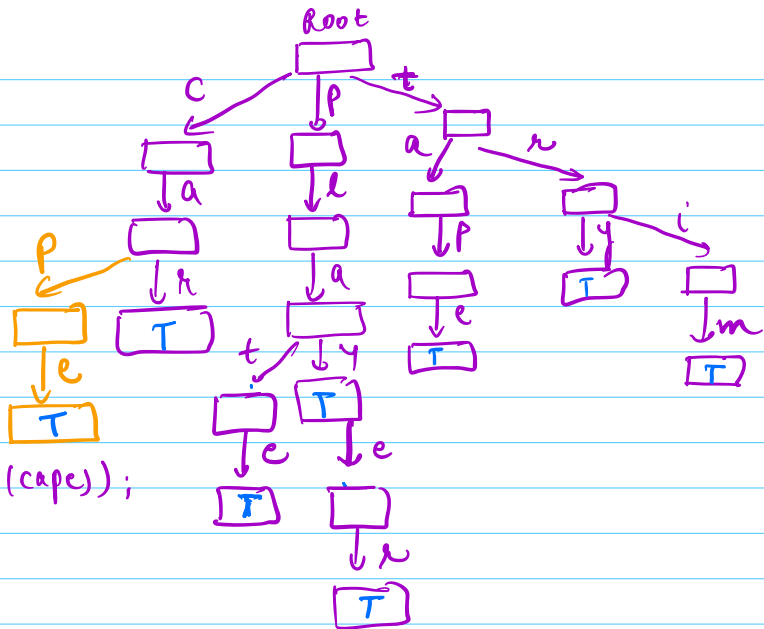
TC: $O(\text{word length})$

HashSet \rightarrow complete word is required to check.

Tree \rightarrow can check as we type characters \rightarrow faster.

List of words

- 1> tape
- 2> try
- 3> trim
- 4> play
- 5> plate
- 6> car
- 7> player
- 8> cape (insert (cape));



```

void insert (root, word) {
    temp = root;
    for i = 0 to (word.length - 1) {
        ch = word[i]
        if (temp.child[ch - 'a'] == null)
            temp.child[ch - 'a'] = new Node()
        temp = temp.child[ch - 'a'];
    }
    temp.isEnd = true; // temp.freq++;
}
    
```

TC: $O(N * \text{word length})$, SC: $O(N * \text{word length})$

Keep track if a word is inserted multiple times. i.e.

freq

~~boolean isEnd~~ → mit freq.

```
class Dict {
```

```
    Node root;
```

```
    Dict () {
```

```
        root = new Node();
```

```
    } public Node getRoot() {  
        return this.root;  
    }
```

```
    Dict d = new Dict();  
    d.getRoot();
```

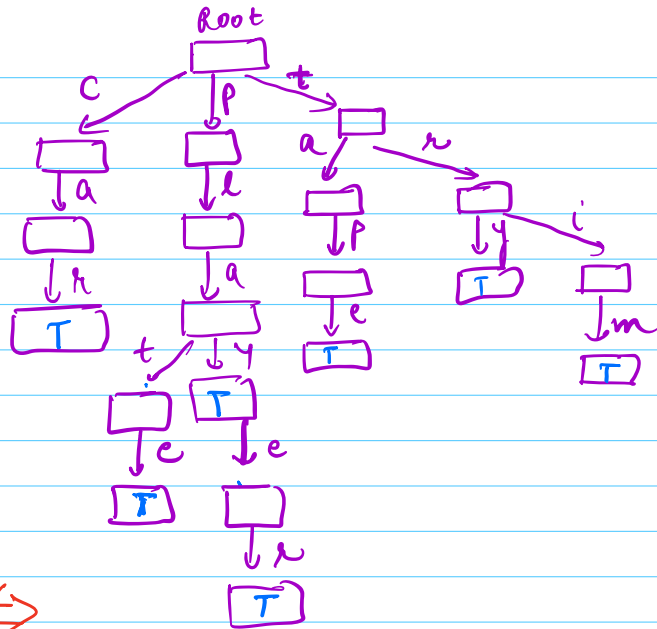
```
    insert (root, "x");
```

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Deletion

delete("trim")

Sell → Travel to the end of the word & update isEnd to false.

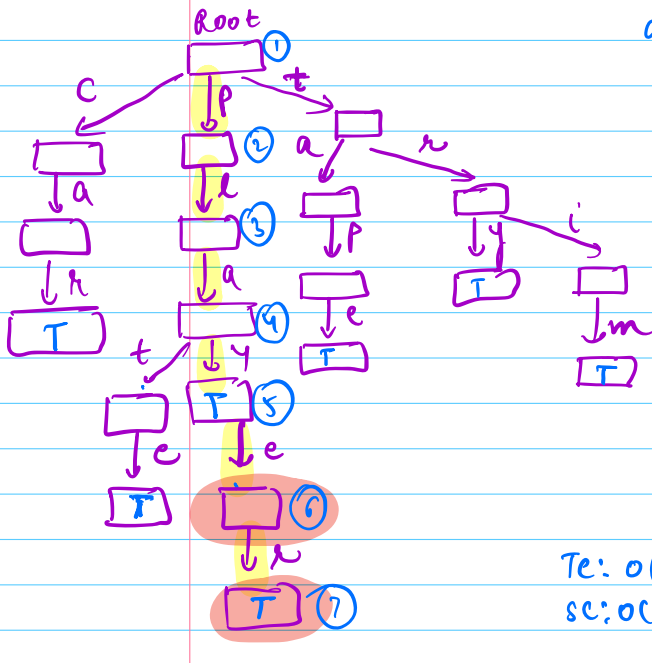


Travels:

a> Travel complete → valid prefix of any input

b> Space wastage

Sell2 → delete un-utilized nodes
↳ only leaf node.



delete(trim)

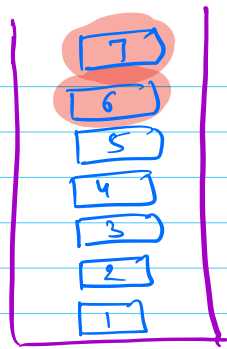
1> Travel all nodes & insert in the stack.

2> Update isEnd = false for last node.

3> Pop elements from stack, if leaf node & isEnd = false → delete it.

↳ if child[i] = null.

Te: $O(\text{word length})$
Sc: $O(\text{word length})$

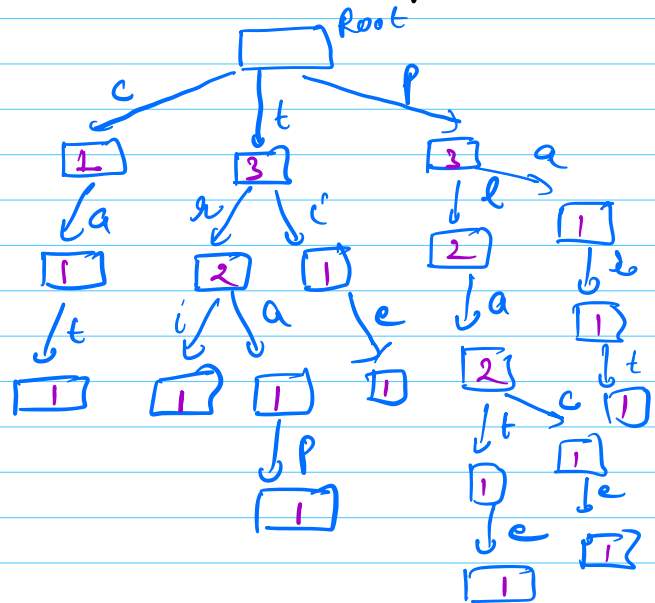
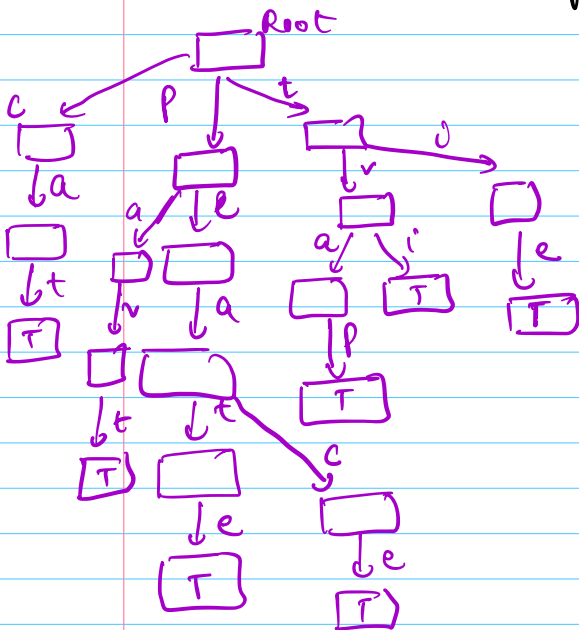


delete player

Q

Given a list of words, find shortest prefix of each word that uniquely defines the word.
(no word is a prefix of another word)

Eg → ["tri", "trap", "plate", "cat", "part", "plac", "tie"]
tri tra plat c pa plac ti



1) If nodes \rightarrow store # times that node is visited.
(while inserting words in the trie)

2) Travel the word till a node with $freq = 1$
is reached

TC: $O(N \times \text{word length})$
SE: $O(N \times \text{word length})$

```
void insert (root, word) {  
    temp = root;
```

```
    for i = 0 to (word.length - 1) {
```

```
        ch = word[i]
```

```
        if (temp.child[ch - 'a'] == null)
```

```
            temp.child[ch - 'a'] = new Node()
```

```
        temp = temp.child[ch - 'a'];
```

```
    }  
    temp.freq++;
```

```
}
```

```
string search (root, word) {
```

```
    temp = root;
```

```
    string ans = "";
```

```
    for (i = 0 to (word.length - 1) {
```

```
        ch = word[i]
```

```
        temp = temp.child[ch - 'a'];
```

```
        ans += ch;
```

```
        if (temp.freq == 1)
```

```
            return ans;
```

```
    }
```

```
}
```

```
class Node {
```

```
    Node[26] child;
```

```
    int freq;
```

```
}
```