

Subset Sum problem

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Exhaustive Algorithm

The exhaustive search algorithm checks all possible multisets $T \subseteq S$ until finding a T (if any) with the appropriate sum. It should return True (it found such a multiset T) or False (it didn't). It will be extremely slow. For maximum credit, come up with pruning techniques to shorten the search.

There are three types:

- 1.Bit wise
- 2.Recursion
- 3.Evaluation

Here we am using the recursion method because we keen of it:

Approach: For the recursive approach we will consider two cases.

- 1. Inclusion
- 2. Exclusion
- 1.Consider the element and now the next tempsum = tempsum of present + value of the element and index = index+1
- 2.Leave the element and now the next **temp sum = tempsum of present** and **index = index+1**

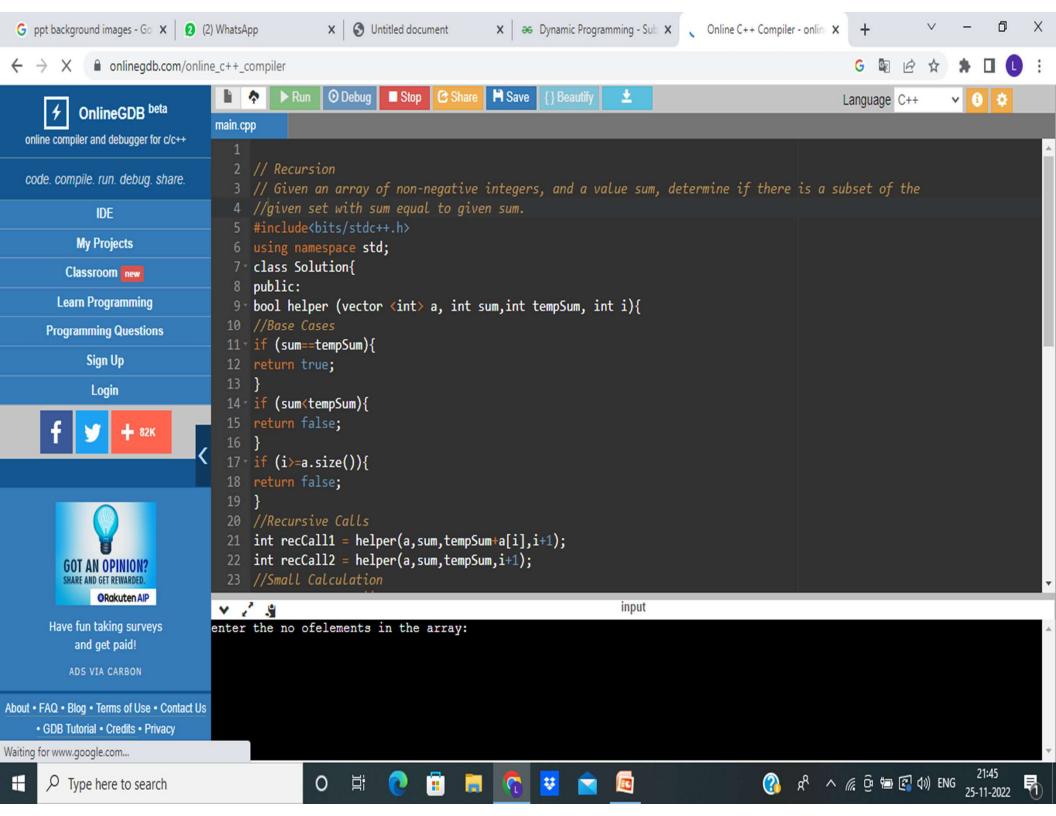
ALGORITHM:

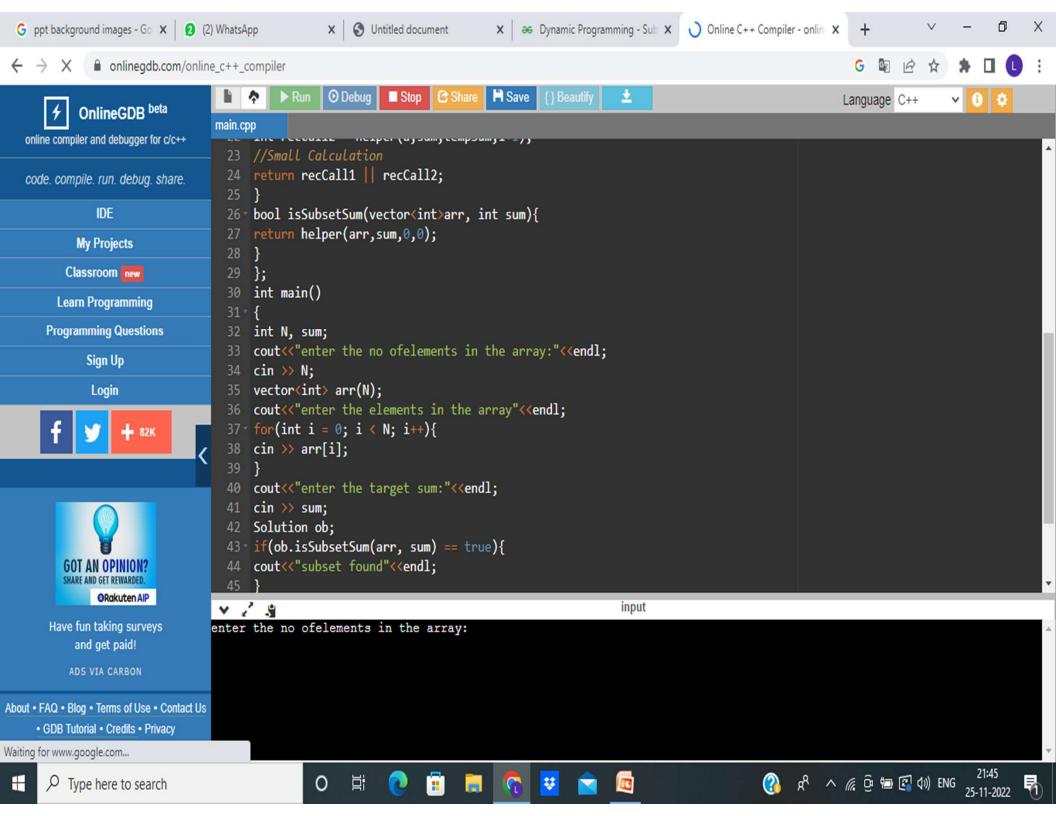
Using vector header file and bool:

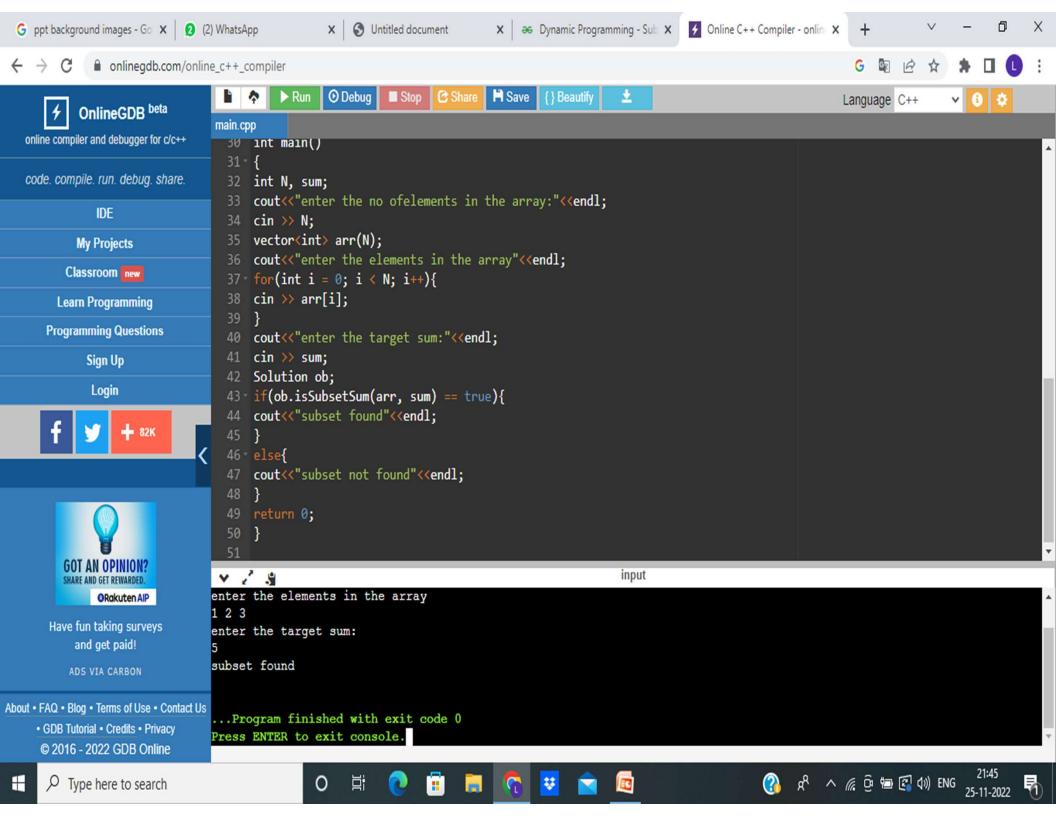
helper(a,sum,tempsum,i)-----///here a = array; sum = target value; tempsum== depends on inclusion and exclusion; i= index value in array

helper(a,sum,tempSum+a[i],i+1) || helper(a,sum,tempSum,i+1); [inclusion or exclusion] // base case:(using bool)

if (sum==tempSum) then return true
if (sum=a.size()) then return false







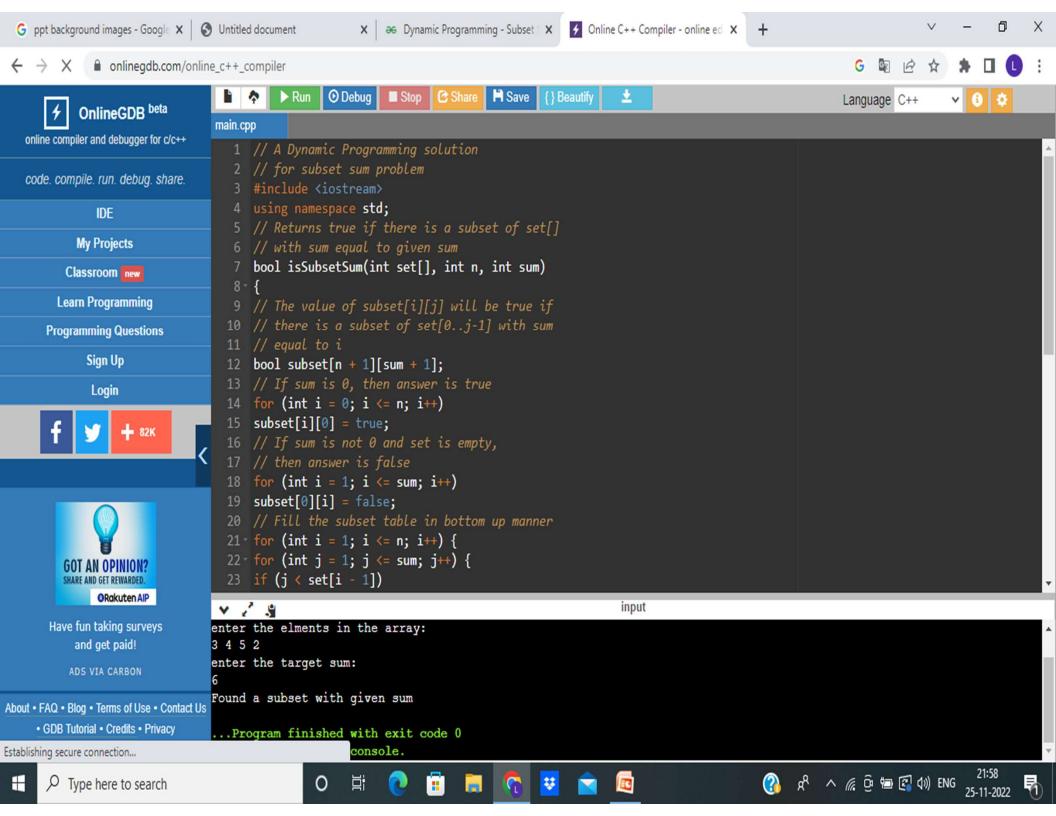
Dynamic Programming

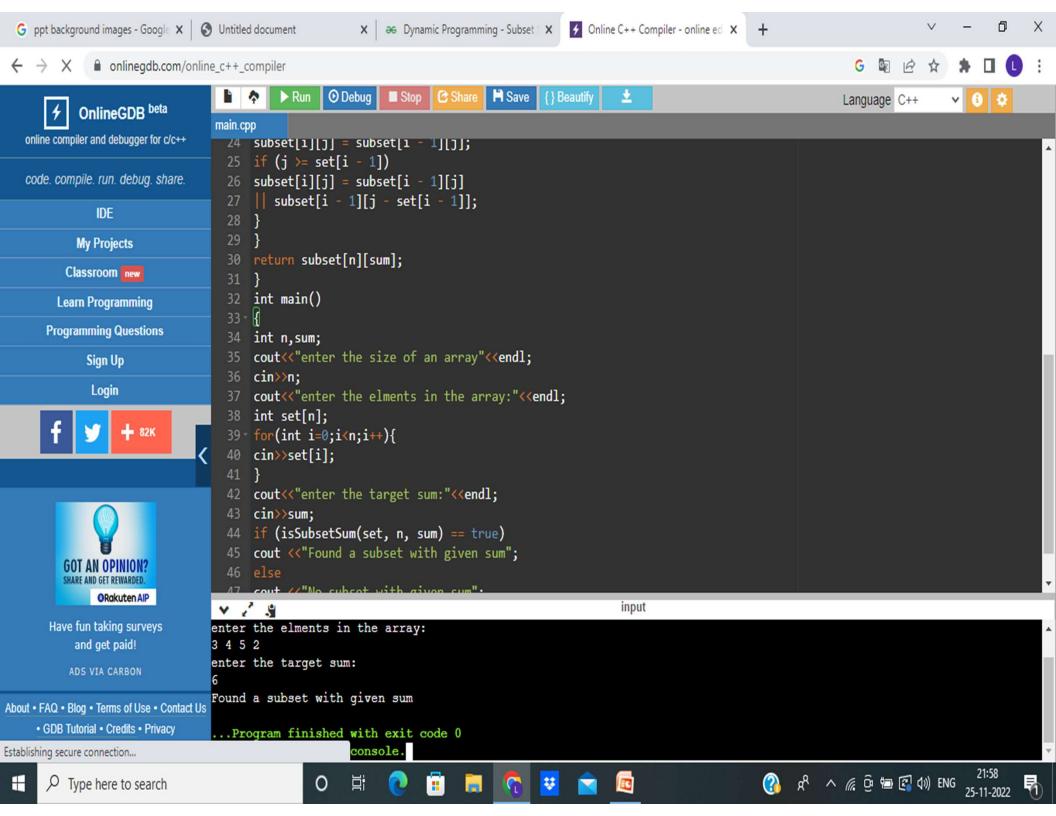
So we will create a 2D array of size (arr.size() + 1) * (target + 1) of type boolean. The state DP[i][j] will be true if there exists a subset of elements from A[0....i] with sum value = 'j'. The approach for the problem is:

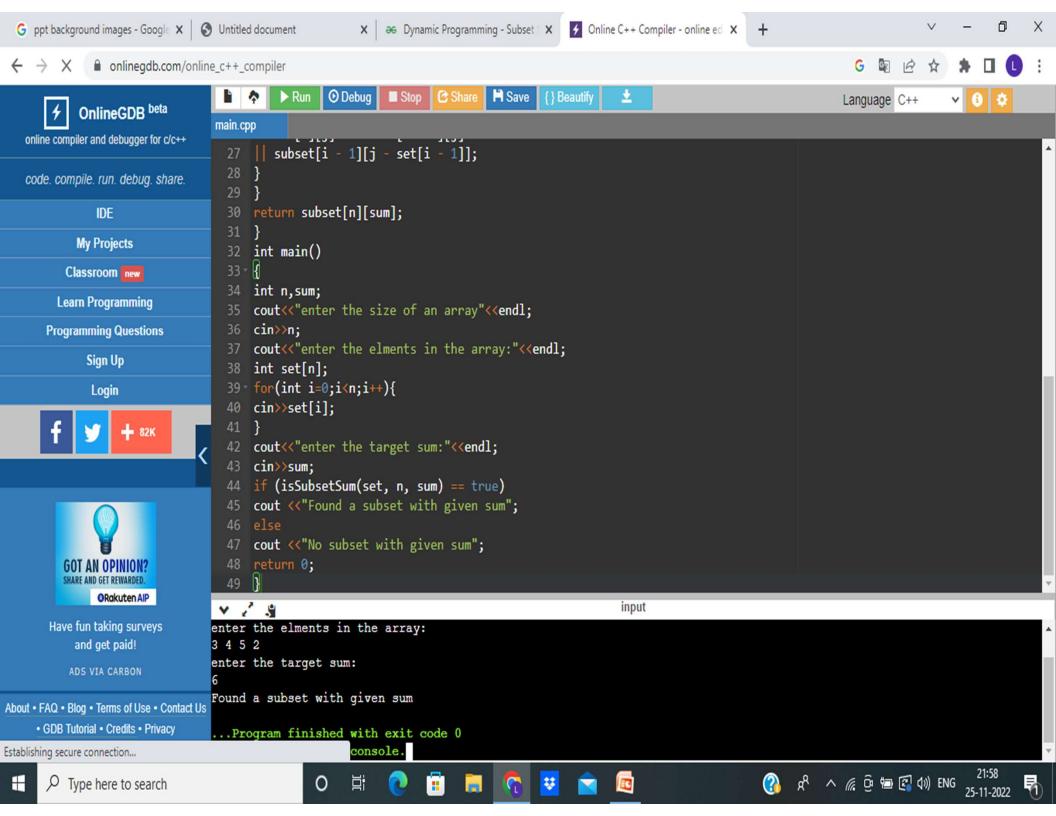
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if (A[i-1] > j)
     DP[i][j] = DP[i-1][j]
else
     DP[i][j] = DP[i-1][j] OR DP[i-1][j-A[i-1]]
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- 1. This means that if current element has value greater than 'current sum value' we will copy the answer for previous cases
- 2. And if the current sum value is greater than the 'ith' element we will see if any of previous states have already experienced the sum='j' OR any previous states experienced a value 'j A[i]' which will solve our purpose.

The below simulation will clarify the above approach: set[]={3, 4, 5, 2} target=6 0 1 2 3 4 5 6 OTFFFFFF 3 T F F T F F 4 T F F T T F F







FOR RECURSION COMPLEXITY IS:

The time complexity will be exponential because it will check for all the subset\

FOR DYNAMIC PROGRAMMING:

Time Complexity: O(sum*n), where sum is the 'target sum' and 'n' is the size of array.

Auxiliary Space: O(sum*n), as the size of 2-D array is sum*n. + O(n) for recursive stack space

