

Contest

- 1.5 hr contest with 3 questions.
- Online Assessment works for companies.
- Test Yourself
- If you are doing regular assignments then a single revision.

Q: Given a string s of lowercase characters, return the count of pairs (i, j) such that

- $i < j$

(2) $s[i] == 'a'$ and $s[j] == 'g'$

Ex: $s =$ "a b e g a g"
 ^{0 1 2 3 4 5}

Ans: 3

(0 3)

(0 5)

(4 5)

Quiz 1: Count in "a c g d g a g"
 ^{0 1 2 3 4 5 6}

Ans: 4

0 2

0 4

0 6

5 6

Quiz 2: Count in "bcaggag" ^{0 1 2 3 4 5 6 7}

Ans: 5

2 3
2 4
2 7
5 7
6 7

Brute Force:

For every 'a', we need to find the count of g on the right side of a.

=> Nested Loops

```
int countAg (String s) {
```

```
    int ans = 0
```

```
    for (int i = 0; i < N; i++) {
```

```
        if (s[i] == 'a') {
```

```
            for (int j = i+1; j < N; j++) {
```

```
                if (s[j] == 'g')
```

```
                    ans++
```

```
            }
```

```
        }
```

```
    return ans;
```

```
}
```

TC: $O(N^2)$

SC: $O(1)$

$\begin{matrix} & \nearrow^{N-1} & & \nearrow^{N-2} & & \nearrow^{N-3} & & \dots & & \nearrow^0 \\ a & a & a & a & a & \dots & a & \end{matrix}$
N 'a'

Optimised Approach: Carry Forward

		1			2			3		3
		a	c	b	a	g	k	a	g	g
count A	0	1	1	1	2	2	2	3	3	3
ans	0	0	0	0	0	↓ ⁺ 2	2	2	↓ ⁺ 5	↓ ⁺ 8

ans = 8

```

int countAg (String Δ) {
    int ans = 0
    int countA = 0
    for (int i = 0; i < N; i++) {
        if (Δ[i] == 'a') {
            countA ++;
        }
        else if (Δ[i] == 'g') {
            ans += countA;
        }
    }
    return ans
}

```

TC: $O(N)$

SC: $O(1)$

Intro to Subarrays

- A subarray is a contiguous part of an array
- It is formed by selecting a range of elements from the array.
- It can have one or more elements but it must be contiguous part of original array.

4	1	2	3	-1	6	9	8	12
---	---	---	---	----	---	---	---	----

2 3 -1 6 is a subarray of size 4

9 is a subarray of size 1

4 1 2 3 -1 6 9 8 12 is a subarray with all elements

4 12 is not a subarray
as not contiguous

1 2 6 is not a subarray
as not contiguous

3 2 1 4 is not a subarray
as order is different

Quiz 3: $A[] = \{ 2, 4, 1, 6, -3, 7, 8, 4 \}$

(A) 1 6 8 X

(B) 1 4 X

(C) 6 1 4 2 X

(D) 7 8 4 YES

Representation of Subarray

0	1	2	3	4	5	6	7	8
4	1	2	3	-1	6	9	8	12

The subarray 2, 3, -1, 6 can be represented by

1) Start 2 to 5 (Start and end index of subarray)
End Index

2) Start and Size
Start = 2
Size = 4

Q How many subarrays starting from index 0?

0 1 2 3 4 5 6
4 2 10 3 12 -2 15

Start End

(0, 0) (0, 4)

(0, 1) (0, 5)

(0, 2) (0, 6)

(0, 3)

ans = 7

Q How many subarrays starting from index 1?

⁰ 4 ¹ 2 ² 10 ³ 3 ⁴ 12 ⁵ -2 ⁶ 15

(1, 1) (1, 3) (1, 5)

(1, 2) (1, 4) (1, 6)

ans = 6

Formula to count total No. of subarrays

No. of subarrays starting from index 0 = N

" " " " " " 1 = $N-1$

" " " " " " 2 = $N-2$

⋮

No. of subarray starting from index $N-2 = 2$

" " " " " " $N-1 = 1$

Sum of N natural numbers = $\frac{N * (N+1)}{2}$

Q Print the subarray of the array that starts from the given start index and end at the given end index.

0	1	2	3	4	5
1	2	-1	2	5	6

start = 1
end = 3
output
2 -1 2

```
void printSubarray(int arr[], int start, int end) {
    for (i = start; i <= end; i++) {
        printf("%d ", arr[i]);
    }
    printf("\n");
}
```

TC: $O(N)$

SC: $O(1)$

Q Print all possible subarrays of the array.

[1 2 3]

output

[1]

[1 2]

[1 2 3]

[2]

[2 3]

[3]

$$\frac{3 \times (3+1)}{2} = \frac{3 \times 4}{2} = 6$$

```
void printAllSubarrays ( int arr[], int n ) {
```

```
    for (int i = 0; i < n; i++) {      i is start
```

```
        for (int j = i; j < n; j++) {  j is end
```

Print the subarray $\rightarrow (i, j)$

```
            for (int k = i; k <= j; k++) {
```

```
                print (arr[k]);
```

```
            }
```

```
            print ('\n');
```

```
        }
```

```
    }
```

TC: $O(N^3)$ SC: $O(1)$

printSubarray (arr, i, j)

	1	2	3	4	5	
					⋮	
			2			1
		2	3			1 2
3		2	3	4		1 2 3
3 4		2	3	4 5		1 2 3 4
3 4 5						1 2 3 4 5
	4			5		
	4 5					

$$\frac{5 \times 6}{2} = 15$$

Q Given an array of N integers, return the len of smallest subarray which contains both maximum and minimum element of the array.

0	1	2	3	4	5	6	7	8	9	10
2	2	6	4	5	1	5	2	6	4	1

max = 6

min = 1

Brute Force: 1) Find Max and Min TC: $O(N)$

Naive

Explore all

2) Explore all subarray and for subarray which has both max and min.

TC: $O(N^3)$

Consider them for ans.

TC: $O(N^3)$

SC: $O(1)$

Observations :

1) Max and Min can come in any order.

2)

MAX MIN

MIN MAX

=> Iterate if found :

MAX \Rightarrow then go ahead and find next min.

MIN \Rightarrow then go ahead and find next max.

0 1 2 3 4 5 6 7 8 9 10
2 2 6 4 5 1 5 2 6 4 1

\uparrow j max = 6
min = 1

ans = ~~1~~ 3

TC: $O(N + N^2)$ SC: $O(1)$
 $\approx O(N^2)$

Can't Forward:

0 1 2 3 4 5 6 7 8 9 10
2 2 6 4 5 1 5 2 6 4 1

\Rightarrow Calculate MAX MIN beforehand

i	A[i]	last_min_index	last_max_index	ans
-	-	-1	-1	11 / INT_MAX
0	2	-1	-1	11
1	2	-1	-1	11
2	6	-1	2	11
3	4	-1	2	11
4	5	-1	2	11
5	1	5	2	$\frac{5-2+1}{2} = 4$ 4

6	5	5	2	4
7	2	5	2	4
8	6	5	8	4
9	4	5	8	
10	1	10		

$(8-5)$
 $+1$
 $= 4$

$(10-8)$
 $+1$
 $= 3$

```
int minSubarray(int A[], int n) {
```

```
    int min = minOfArray(A);
```

```
    int max = maxOfArray(A);
```

```
    int last_min_index = -1
```

```
    int last_max_index = -1
```

```
    int ans = N
```

```
    for (i = 0; i < N; i++) {
```

```
        if (A[i] == min_value) {
```

```
            last_min_index = i
```

```
            if (last_max_index != -1) {
```

```
                ans = min(ans, i - last_max_index + 1);
```

```
            }
```

```
    }
```

```
else if(A[i] == maxVal) {
```

```
    last_max_index = i
```

```
    if (last_min_index != -1) {
```

```
        ans = min(ans, i - last_min_index + 1);
```

```
    }
```

```
}
```

```
}
```

```
return ans;
```

```
}
```

TC: $O(N)$

SC: $O(1)$

Next Class:

→ Subarrays

→ Printing all subarrays

→ Contribution Technique

Doubts

0 based Index

0	1	2	3	4	5	6	7
1	2	3	1	2	1	2	3

Prefix Array

0	1	2	3	4	5	6	7
1	3	6	7	9	10	12	13

$$\text{sum} \Rightarrow 2 \quad 5 \Rightarrow \text{ps}[5] - \text{ps}[2] = 10 - 3 = 7$$

$$\text{sum} \Rightarrow 0 \quad 3 \Rightarrow \text{ps}[3] = 7$$

To avoid if else on the $l == 0$

1 based Index

0	1	2	3	4	5	6	7
1	2	3	1	2	1	2	3

Prefix Array

0	1	2	3	4	5	6	7	8
0	1	3	6	7	9	10	12	13

$$\text{sum} \Rightarrow 2 \quad 5 \Rightarrow 3 \quad 6 \Rightarrow \text{ps}[6] - \text{ps}[2] = 10 - 3 = 7$$

$$\text{sum} \Rightarrow 0 \quad 3 \Rightarrow 1 \quad 4 \Rightarrow \text{ps}[4] - \text{ps}[0] = 7 - 0 = 7$$

