,
Quantum Annealing for Air Traffic Management

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I. INTRODUCTION

II. PROBLEM DESCRIPTION

The problem at hand is the deconflicting of transatlantic wind-optimal trajectories. As it was done in [1] we are using the same wind-optimal trajectories of a single day, July 29 2012. These wind-optimal trajectories are given as $(\mathbf{x}_i)_{i=1}^n$, where $\mathbf{x}_i = (x_{i,t})_{t=\tau_{i,0}}^{\tau_{i,1}}$ and $x_{i,t}$ is the location (as latitude, longitude, and altitude) of the ith flight at time t. The times $\tau_{i,0}$ and $\tau_{i,1}$ are the times at which the wind-optimal trajectory for the ith flight begins and ends, respectively. Furthermore, the times are given in units of one minutes $T_i = (\tau_{i,0}, \tau_{i,0} + 1, \ldots, \tau_{i,1})$. Each flight i is at a constant speed v_i, to within (classical) machine precision.

A conflict between two flights is defined as a pair of trajectory points which are too close to each other in space and time.

$$\{(x_{i,t}, x_{j,t'}) \mid \mathcal{D}(x_{i,t}, x_{j,t'}) < \Delta_x, |t - t'| < \Delta_t\} ,$$

where $\mathcal{D}(x, y)$ is the spatial distance between two points x and y given as latitude, longitude and altitude. Following [1], the space threshold is $\Delta_x = 3$ nautical miles and the time threshold is $\Delta_t = 3$ minutes. In this paper, we consider the following means to deconflict the trajectories: First, we can delay each flight i at departure time by a departure delay d_i

$$x_{i,t} \to x_{i,t+d_i} \quad \forall\, t \in T_i$$

Second, we can avoid a conflict by maneuvers of both involved flights. We assume, however, that the maneuvers will not introduce new conflicts. In doing so, these maneuvers can be view as resulting in time shifts only.

III. CLASSICAL PREPROCESSING

It is beneficial to reduce the data to conflicting regions in space and decoupling the spacial and temporal components of the problem. As a first step, we detect all pairs of trajectory points which are separated by a spacial distance below Δ_x

$$\{(x_{i,t}, x_{j,t'}) \mid \mathcal{D}(x_{i,t}, x_{j,t'}) < \Delta_x\} ,$$

Two spatially conflicting trajectory points might never become conflicting in time if the corresponding times are far apart. By introducing a constant maximum delay D_{max} we can dismiss all spatial conflicts which can never become conflicting in time

$$\{(x_{i,t}, x_{j,t'}) \mid \mathcal{D}(x_{i,t}, x_{j,t'}) < \Delta_x, |t - t'| \geq \Delta_t + D_{\max}\} .$$

With this, we are left with a set of potentially conflicting pairs of trajectory points

$$C_0^{ij} = \{(x_{i,t}, x_{j,t'}) \mid \mathcal{D}(x_{i,t}, x_{j,t'}) < \Delta_x, |t - t'| < \Delta_t + D_{\max}\} .$$

A. Conflict Clustering

As a next step, we group together conflicting trajectory point pairs which are subsequent in time

$$\begin{aligned}
C_\parallel^{ij} = \{((x_{i,t}, x_{j,t'}), (x_{i,s}, x_{j,s'})) \mid\ & (x_{i,t}, x_{j,t'}) \in C_0^{ij}, \\
& (x_{i,s}, x_{j,s'}) \in C_0^{ij}, \\
& |t - s| < \Delta_t', \\
& |t' - s'| < \Delta_t'\}
\end{aligned}$$

where we set $\Delta_t' = 2$ minutes. For a given pair of flights (i, j) there might be multiple "disjunct" subsets in $C_\parallel^{ij}$

$$\bigcup_n C_{\parallel n}^{ij} = C_\parallel^{ij}$$

where

$$\begin{aligned}
& |t - s| \geq \Delta_t' \wedge |t' - s'| \geq \Delta_t' \\
& \forall (x_{i,t}, x_{j,t'}) \in C_{\parallel n}^{ij}, \\
& \forall (x_{i,s}, x_{j,s'}) \in C_{\parallel n'}^{ij}, \\
& n \neq n' .
\end{aligned}$$

Together with the remaining, spatially isolated, conflicting trajectory points

$$C_\times^{ij} = C_0^{ij} \setminus C_\parallel^{ij} ,$$

these subsets of trajectory point clusters are called *potential conflicts*.

$$C_k \in C = \{C_{\parallel n}^{ij} | \forall i, j, n\} \cup \{C_\times^{ij} | \forall i, j\}$$

Here, we introduced an conflict index $k \in 1, \ldots, K$, with $K = |C|$. For each conflict index k, we will denote the pair of involved flights by $I_k = (i_k, j_k)$.

IV. MAPPING TO QUBO

[1] O. Rodionova, D. Delahaye, B. Sridhar, and H. Ng., Proceedings of Advanced Aircraft Efficiency in a Global Air Transport System (AEGATS'16) Conference (2016).