Introduction to Database Systems

2023-Fall

Database & DBMS

- A very large, integrated collection of data.

- Models real-world *enterprise.*
  - Entities (e.g., students, courses)
  - Relationships (e.g., electives)

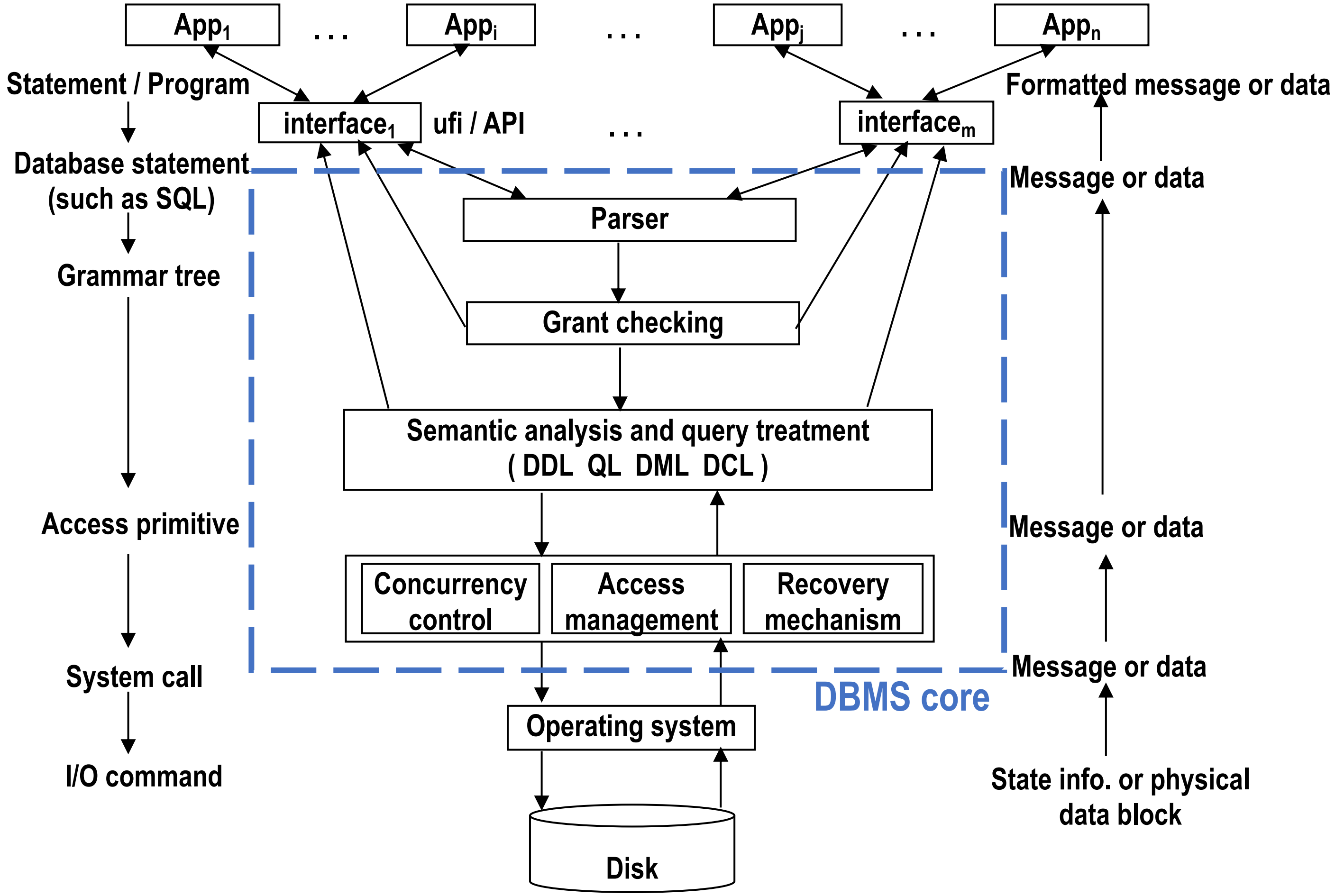
- A *Database Management System (DBMS)* is a software package designed to store and manage databases.

4. Database Management Systems （1/5）

Contents

4.1 The Architecture of DBMS

- ***The Components of DBMS Core***

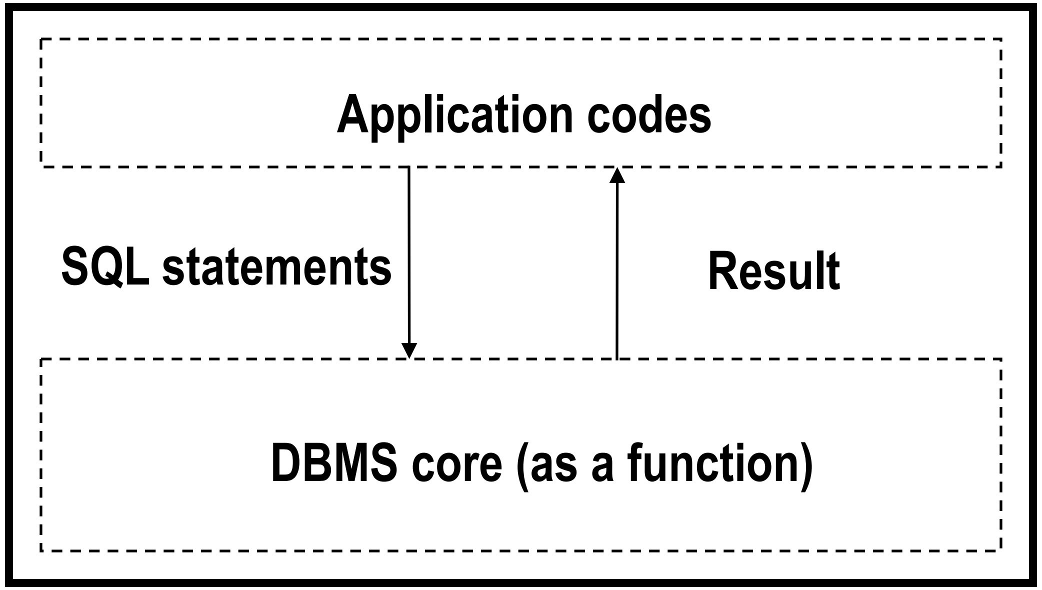


4.2 The Process Structure of DBMS

- Single process structure
- Multi processes structure
- Multi threads structure
- Communication protocols between processes / threads

Single process structure

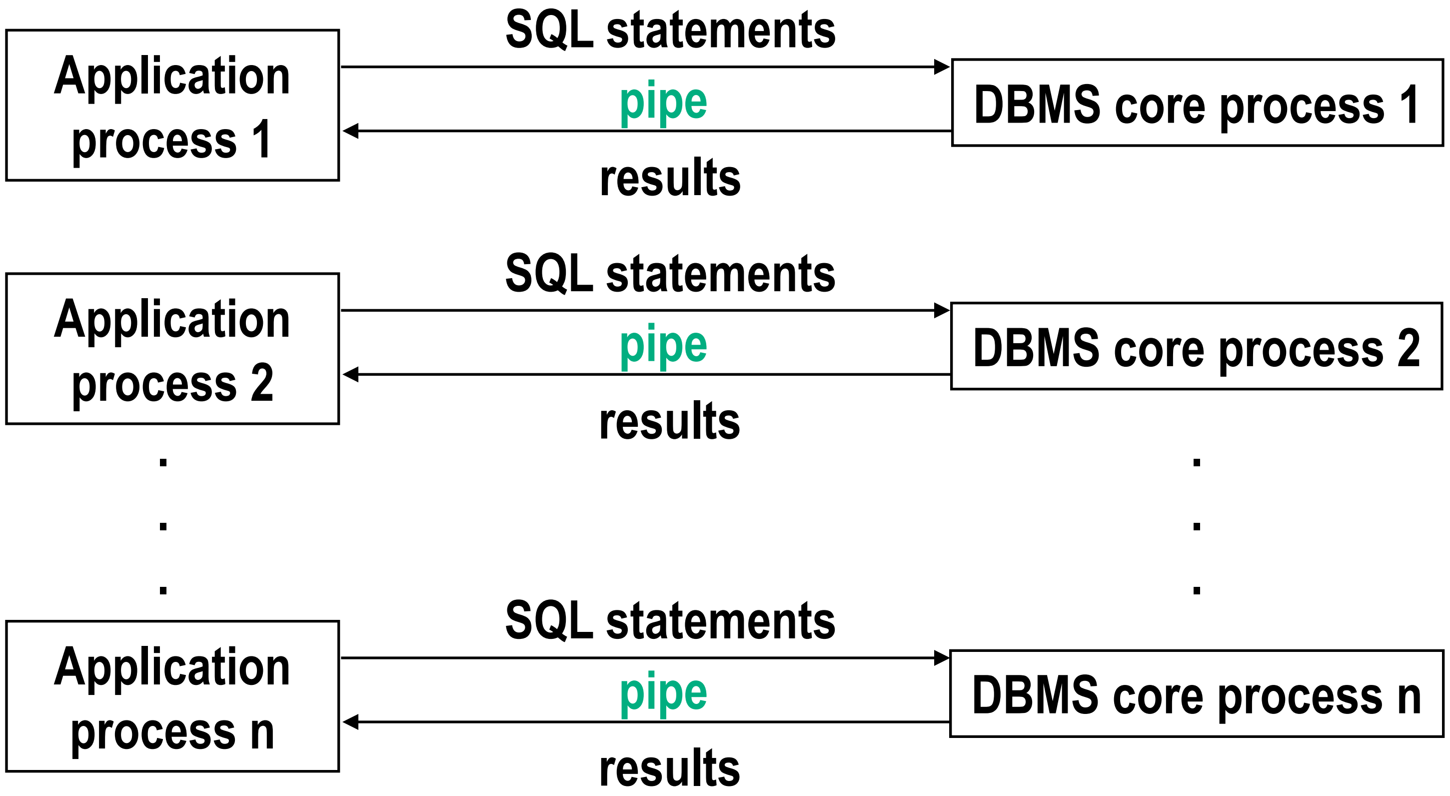
- The application program is compiled with DBMS core as a single .exe file, running as a single process.



.exe file

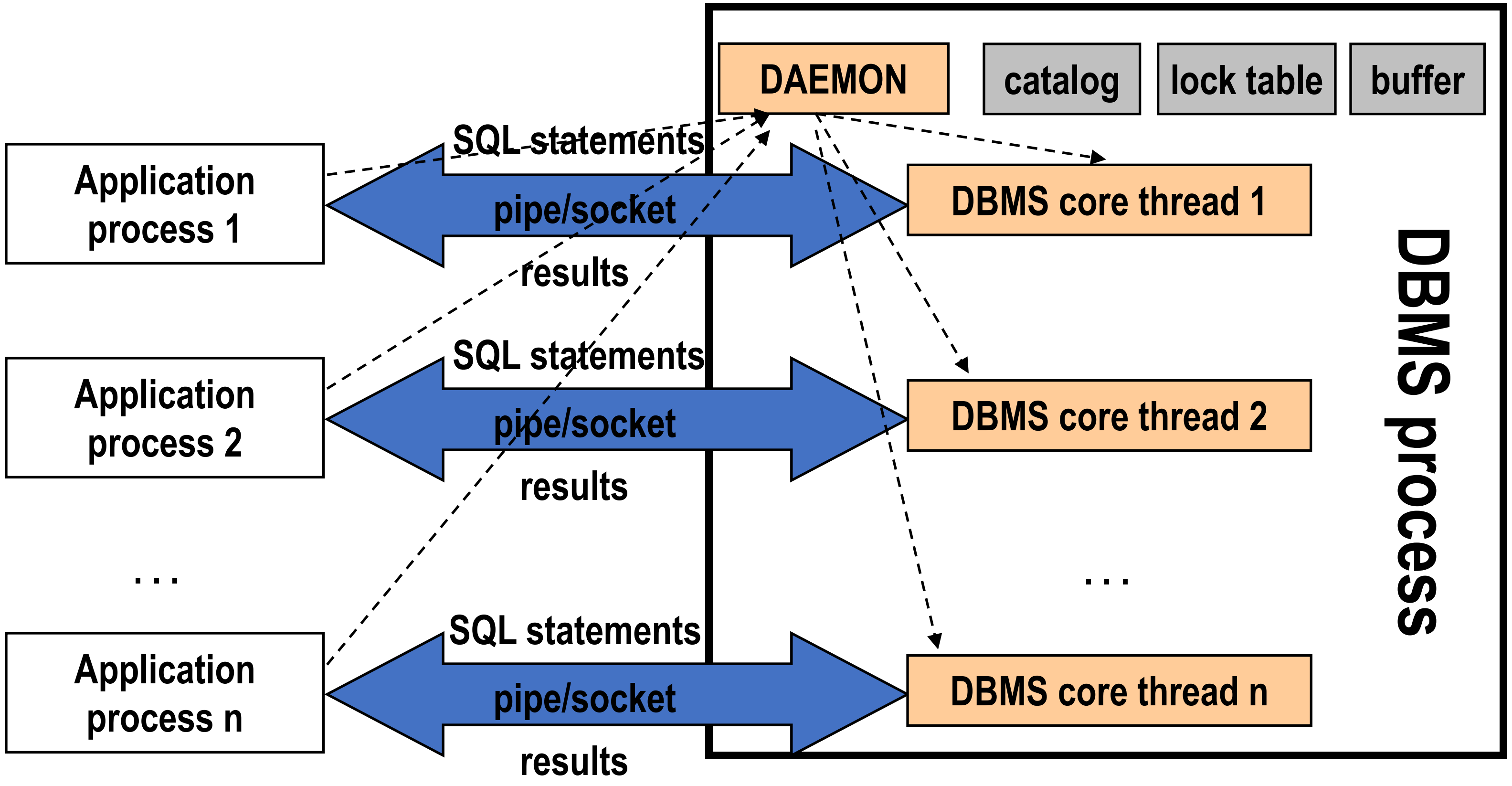
Multi processes structure

- One application process corresponding to one DBMS core process



Multi threads structure

- Only one DBMS process, every application process corresponding to a DBMS core thread.



Big picture:
Architecture of a DBMS

Architecture of a DBMS: SQL Client

- Last few lectures: SQL
- Next:
  - How is a SQL query executed?

DBMS: Parsing & Optimization

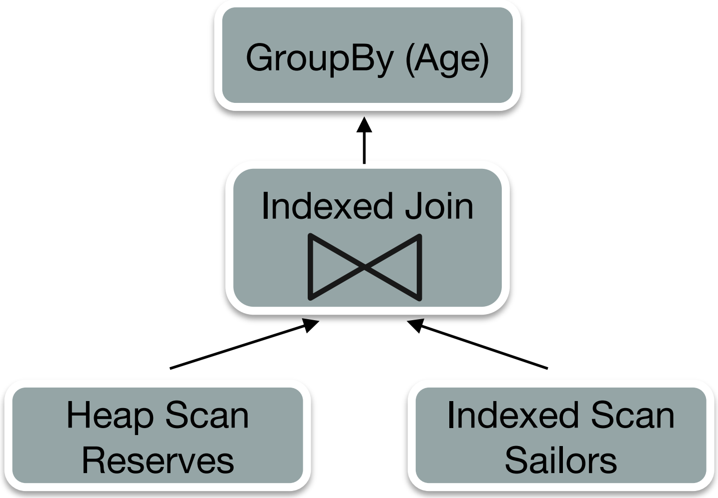
Purpose: Parse, check, and verify the SQL

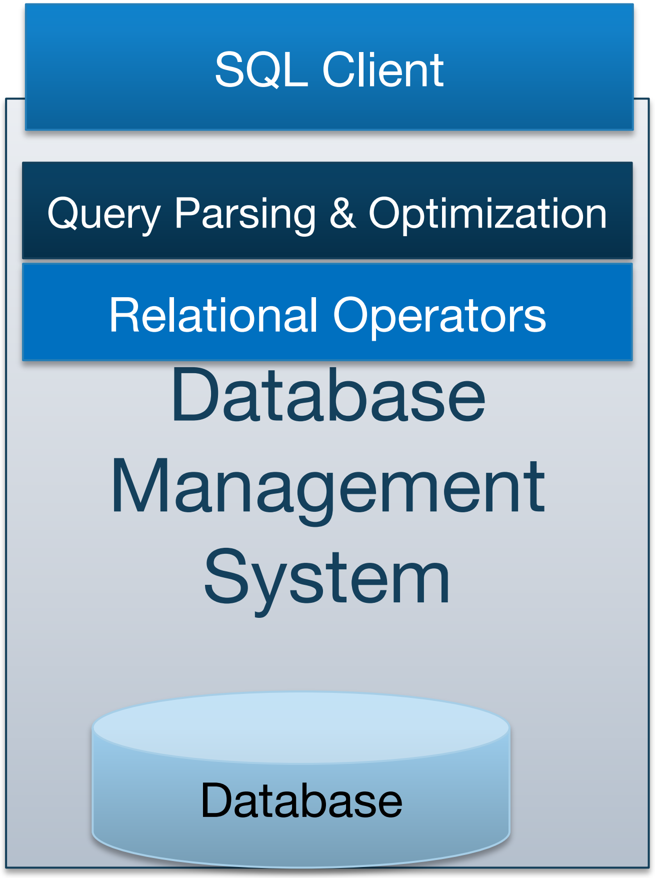
```
SELECT S.sid, S.sname, R.bid
FROM Sailors R, Reserves R
WHERE S.sid = R.sid and S.age > 30
GROUP BY age
```

And translate into an efficient relational query plan.

DBMS: Relational Operators

Purpose: Execute a dataflow by operating on **records** and **files**

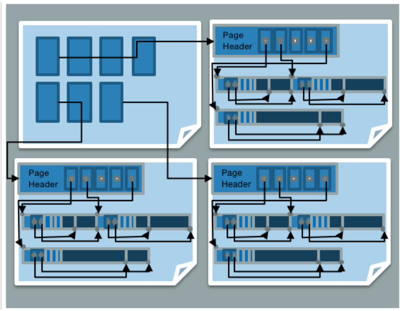


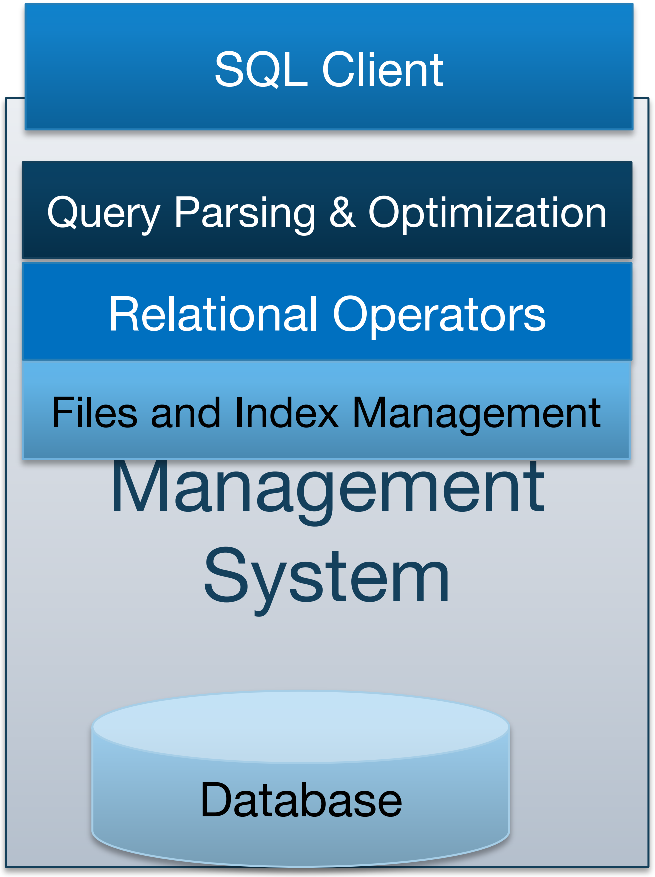


DBMS: Files and Index Management

Purpose: Organize tables and Records as groups of pages in a logical file

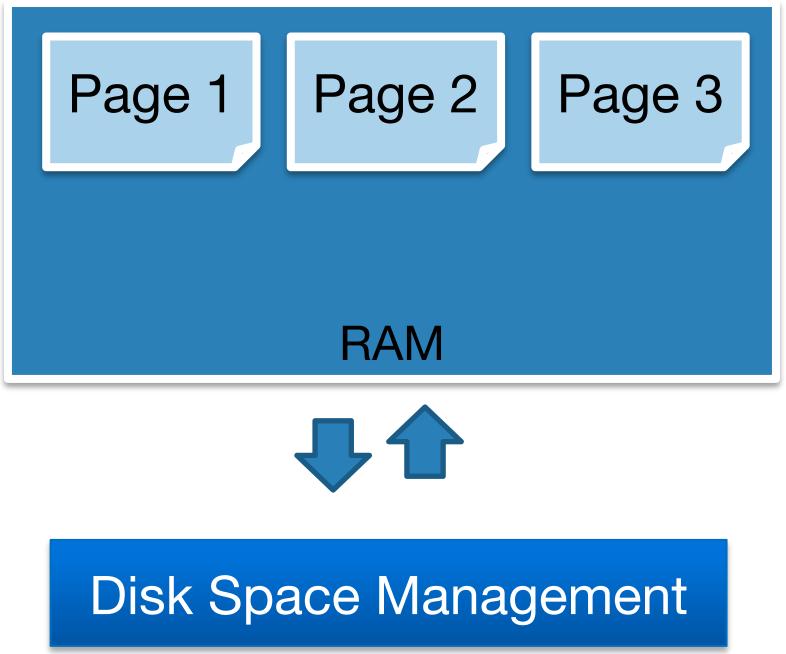
SSN	Last Name	First Name	Age	Salary
123	Adams	Elmo	31	$400
443	Grouch	Oscar	32	$300
244	Oz	Bert	55	$140
134	Sanders	Ernie	55	$400

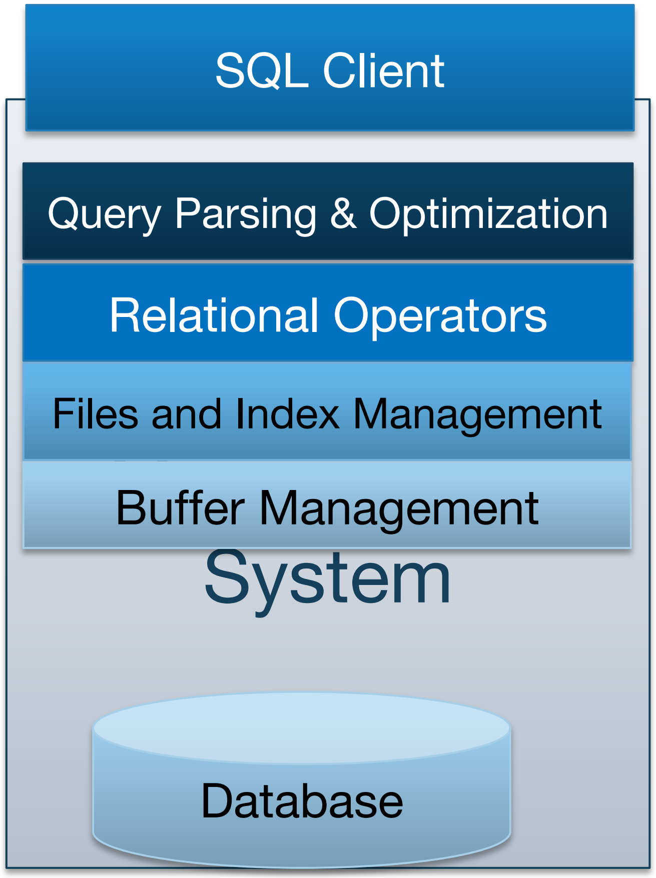




DBMS: Buffer Management

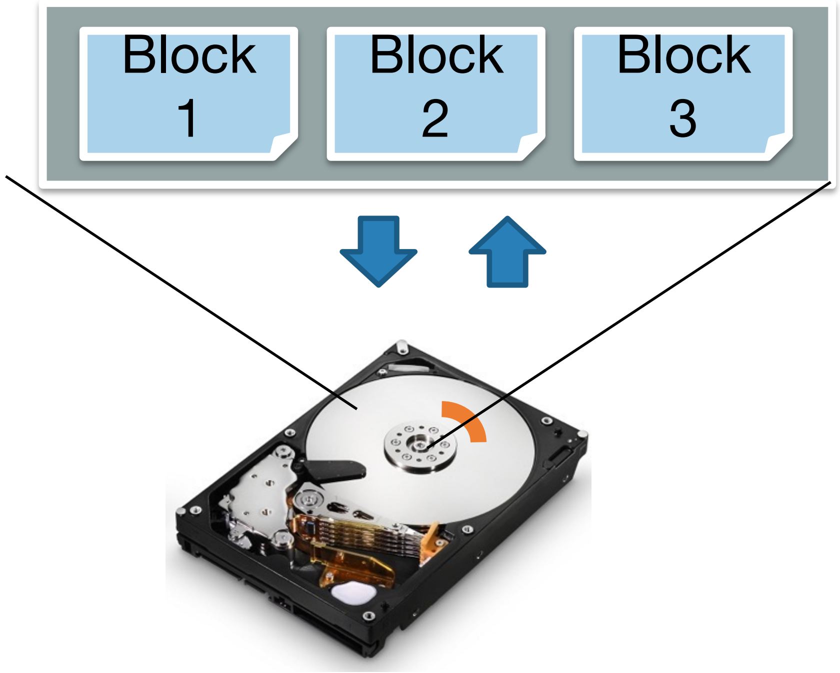
Purpose: Provide the illusion of operating in memory

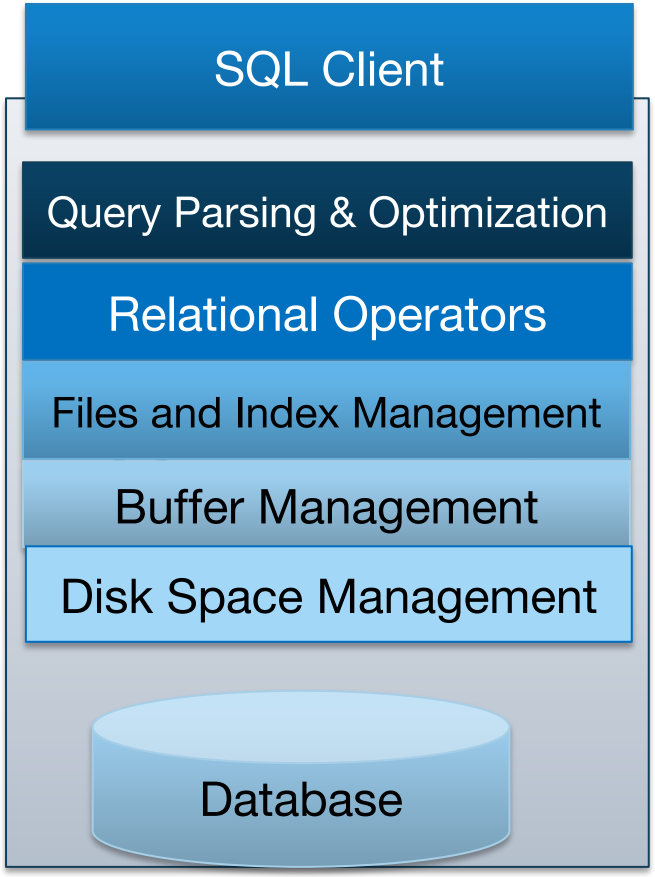




DBMS: Disk Space Management

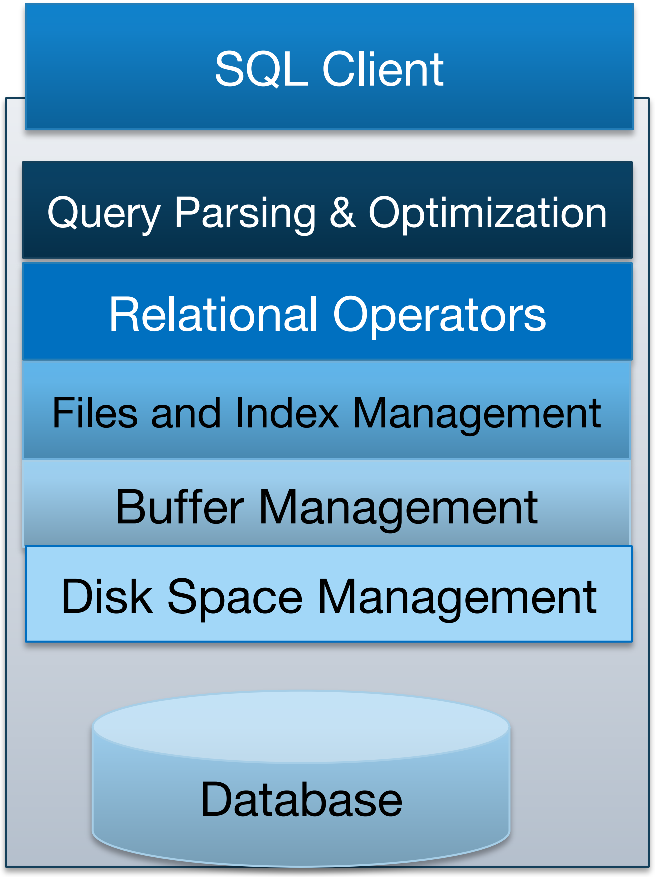
Purpose: Translate page requests into physical bytes on one or more device(s)





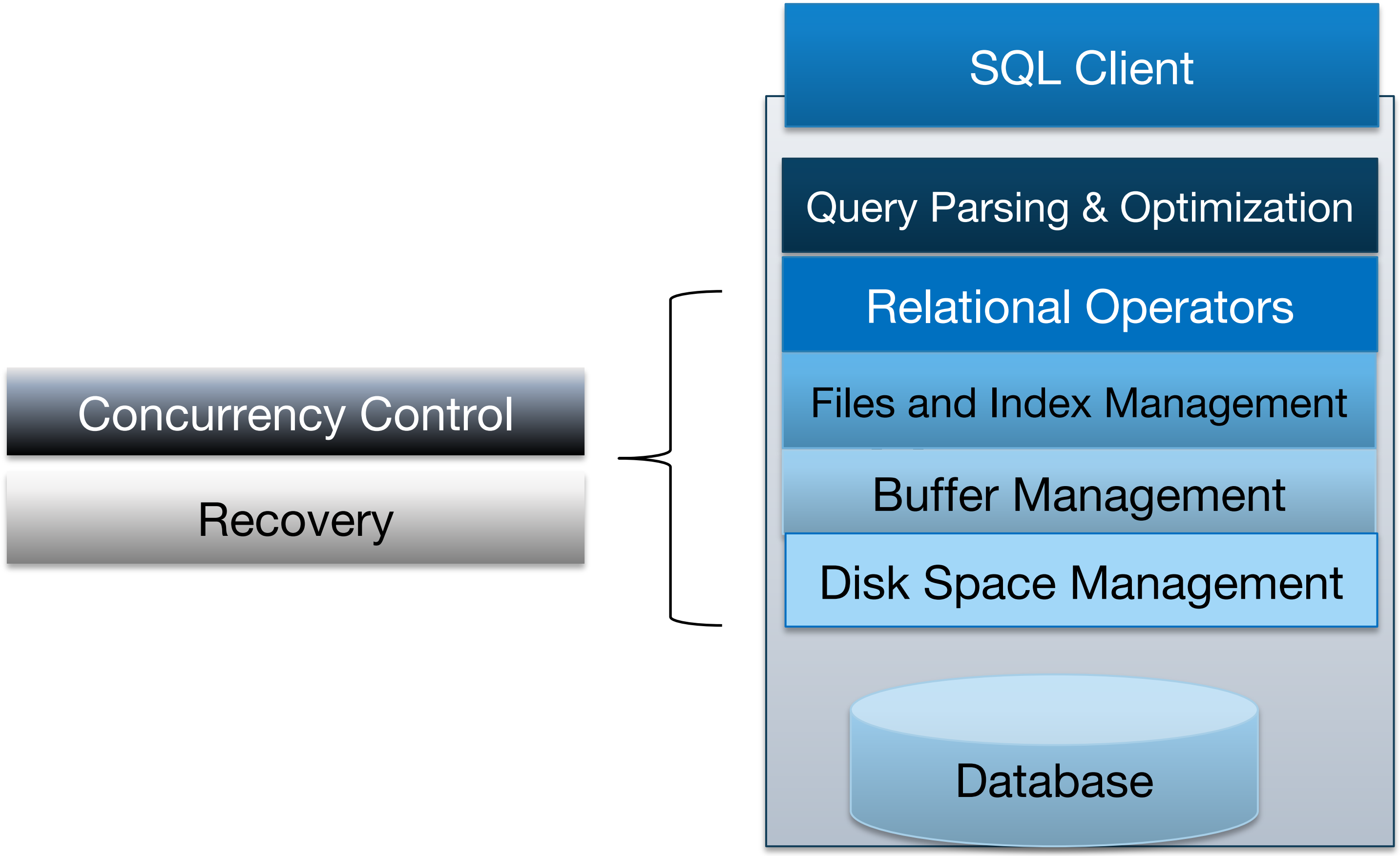
Architecture of a DBMS

- Organized in layers
- Each layer abstracts the layer below
  - Manage complexity
  - Performance assumptions
- Example of good systems design



DBMS: Concurrency & Recovery

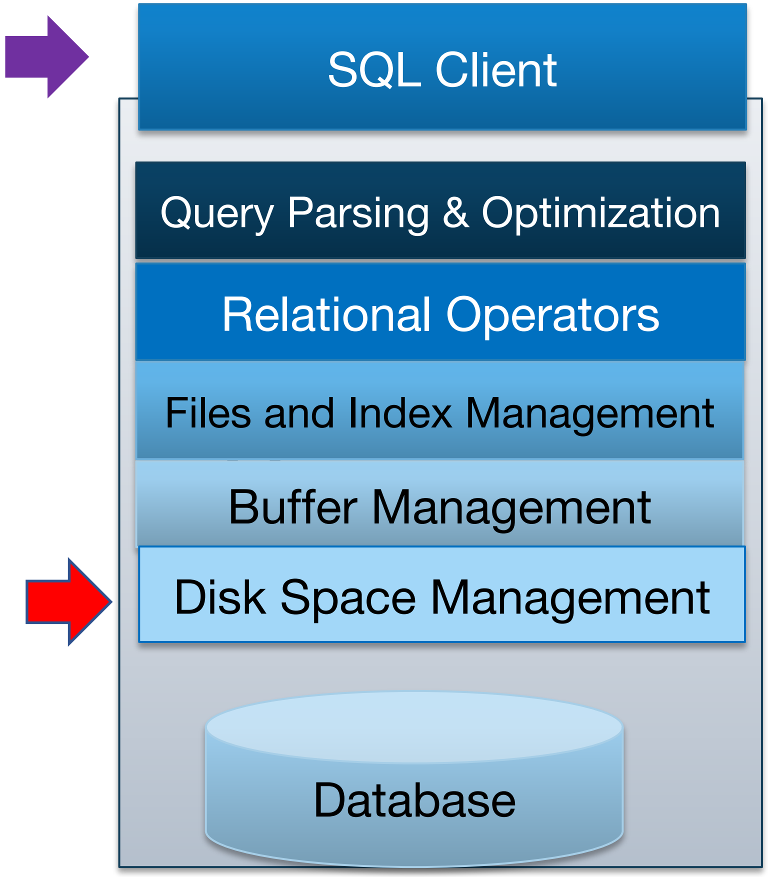
Two cross-cutting issues related to storage and memory management:



Context

Completed

You are Here

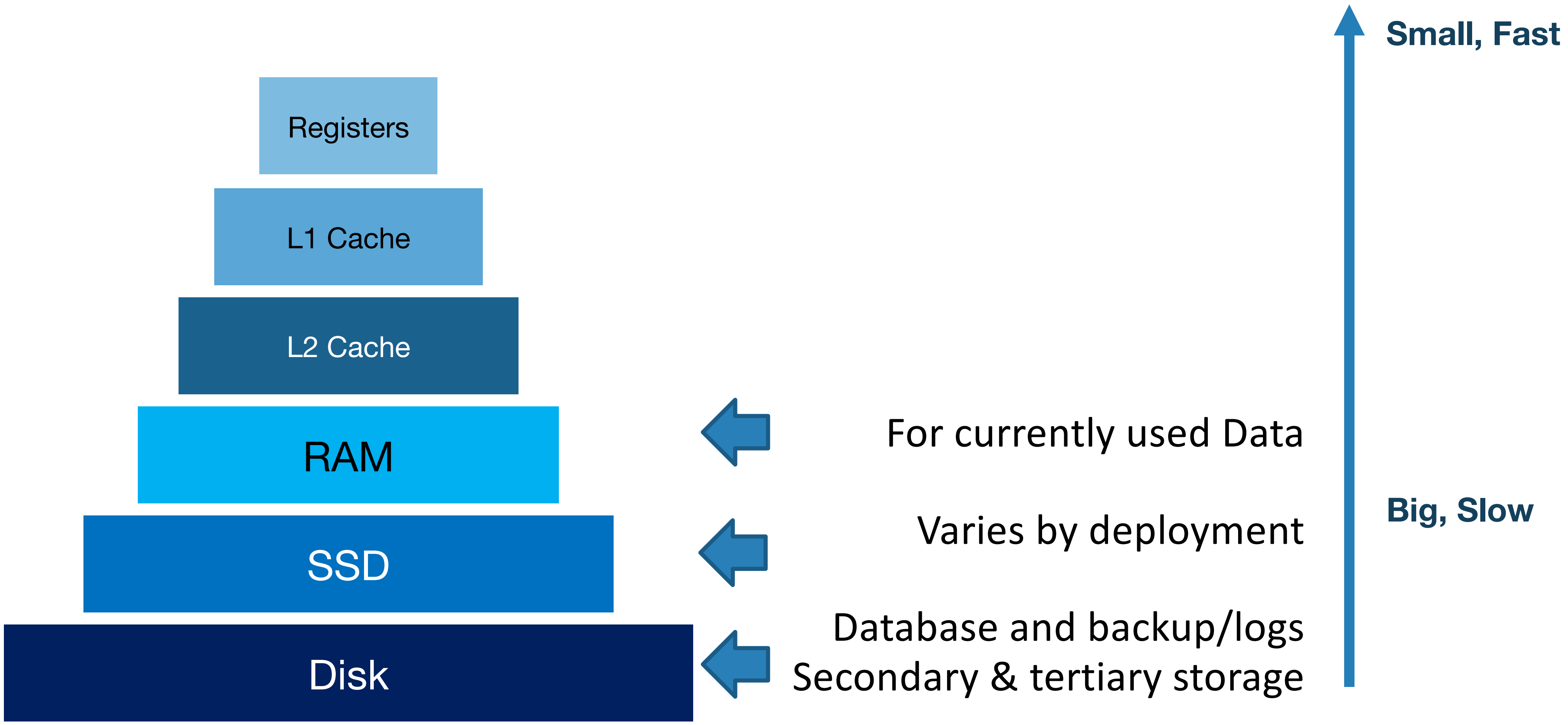


Before We Begin:
Storage Media

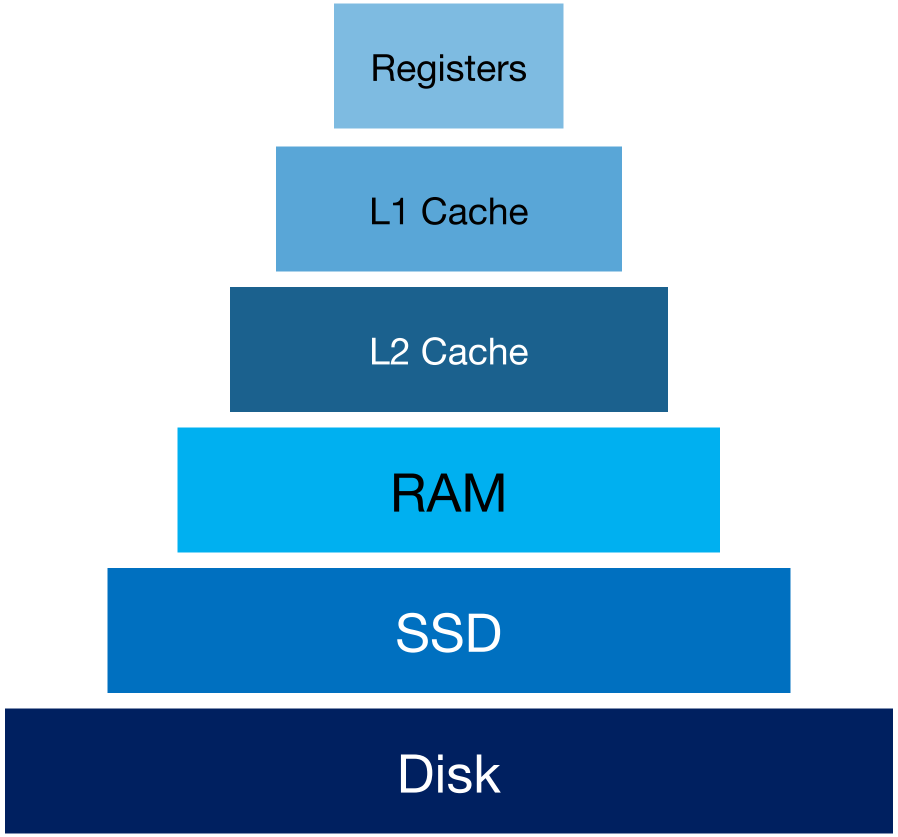
Disks

- Can differentiate storage into:
  - **volatile storage:** loses contents when power is switched off
  - **non-volatile storage**:
    - Contents persist even when power is switched off.
    - Includes secondary and tertiary storage, as well as batter-backed up main-memory.
- Factors affecting choice of storage media include
  - Speed with which data can be accessed
  - Cost per unit of data
  - Reliability

Storage Hierarchy



Hierarchy - Storage Latencies



.5 ns – L1 cache reference

7.0 ns – L2 cache reference

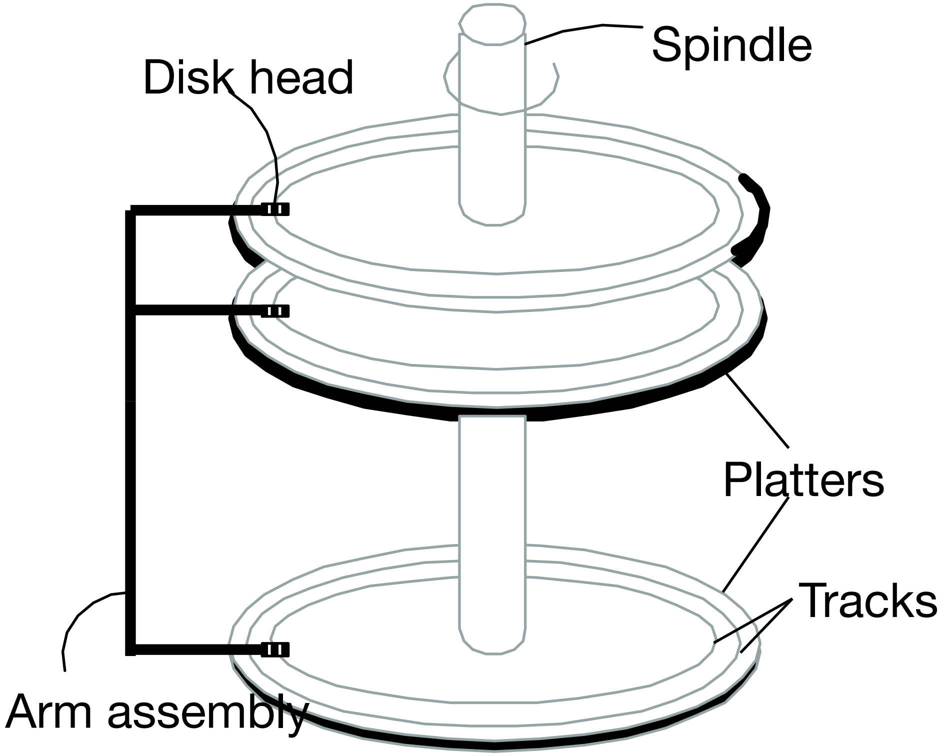
100.0 ns – main memory reference

1,000,000.0 ns – to read 1MB sequentially

20,000,000.0 ns to read 1MB sequentially

Components of a Disk, Pt. 1

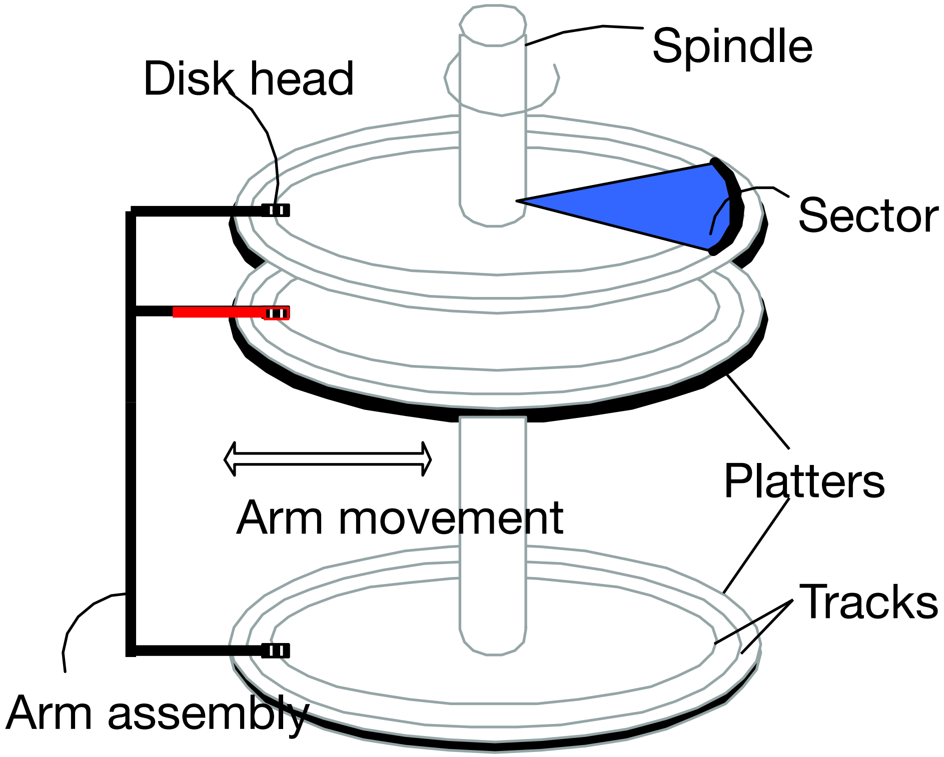
- **Platters** spin (say 15000 rpm)
- **Arm assembly** moved in or out to position a **head** on a desired **track**
  - Tracks under heads make a "cylinder"
- Only one head reads/writes at any one time



[Disk Platters]

Components of a Disk, Pt. 2

- **Platters** spin (say 15000 rpm)
- **Arm assembly** moved in or out to position a **head** on a desired **track**
  - Tracks under heads make a “cylinder”
- Only one head reads/writes at any one time
- Block/page size is a multiple of (fixed) **sector** size

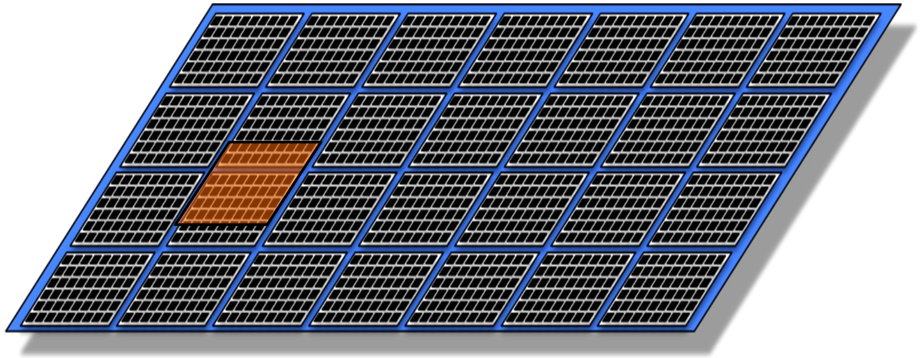


Accessing a Disk page

- Time to access (read/write) a disk block:
  - **seek time** (moving arms to position disk head on track)
    - ~2-3 ms on average
  - **rotational delay** (waiting for block to rotate under head)
    - ~0-4 ms (15000 RPM)
  - **transfer time** (actually moving data to/from disk surface)
    - ~0.25 ms per 64KB page
- Key to lower I/O cost: reduce seek/rotational delays

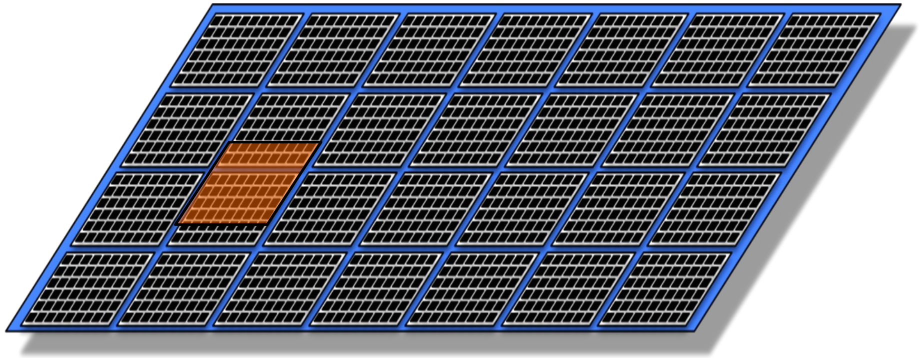
Notes on Flash (SSD)

- Issues in current generation (NAND)
  - Fine-grain reads (4-8K reads), coarse-grain writes (1-2 MB writes)
  - Only 2k-3k erasures before failure,
  - Write amplification: big units, need to reorg for wear & garbage collection



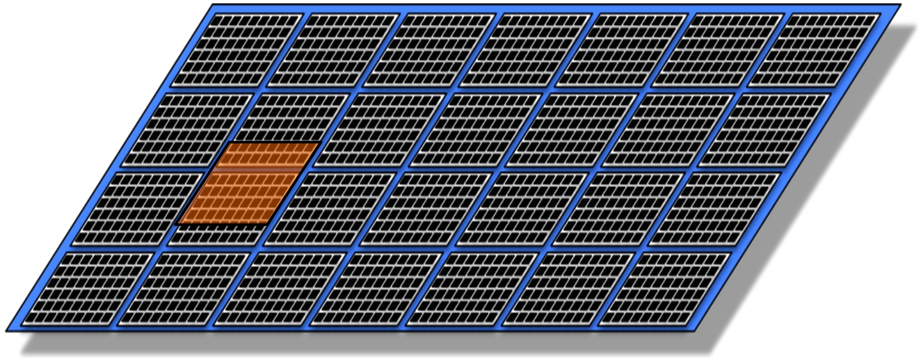
Notes on Flash (SSD), Pt. 2

- So… read is fast and predictable
  - Single read access time: 0.03 ms
  - 4KB random reads: ~500MB/sec
  - Sequential reads: ~525MB/sec
  - 64K: 0.48 ms

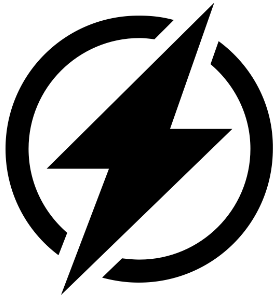


Notes on Flash (SSD), cont

- But... write is not! Slower for random
  - Single write access time: 0.03 ms
  - 4KB random writes: ~120 MB/sec
  - Sequential writes: ~480 MB/sec



Is Flash Faster than Disk?



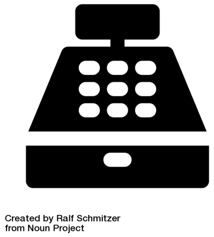
Created by Dima Sho
from Noun Project

- Why of course it is…it’s called “flash”!
  - Can be 1-10x the bandwidth (bytes/sec) of ideal HDD #s
    - Note: Ideal HDD #s hard to achieve.
    - Expect 10-100x bandwidth for non-sequential read.
- Locality” matters for both
  - Reading/writing to “far away” blocks on disk requires slow seek/rotation delay
  - Writing 2 “far away” blocks on SSD can require writing multiple much larger units
  - High-end flash drives are getting much better at this
- And don’t forget:
  - Disk offers about 10x the capacity per $

Storage Trends

- But data sizes grow faster than Moore's Law
  - "Big Data" is real
    - Boeing 787 generates ½ TB of data per flight

    - Walmart handles 1M transactions/hour,
      - maintains 2.5 PetaByte data warehouse



- So…what is the role of disk, flash, RAM

Bottom Line (last few years)

- Very large DBs: relatively traditional
  - Disk still the best cost/MB by a lot
  - SSDs improve performance and performance variance
- Smaller DB story is changing quickly
  - Entry cost for disk is not cheap, so flash wins at the low end
  - Many interesting databases fit in RAM

Bottom Line Pt. 2

- Change brewing on the HW storage tech side
- Mixed answers on the SW/usage side
  - Big Data: Can generate and archive data cheaply and easily
  - Small Data: Many rich data sets have (small) fixed size
- People will continue to worry about magnetic disk for some time yet, typically at large scale

Disk Space Management

Disks and Files

- Recall, most DBMSs stores information on **Disks** and **SSDs**.
  - Disk are a mechanical anachronism (slow!)
  - SSDs faster, **slow relative to memory**, costly writes

Block Level Storage

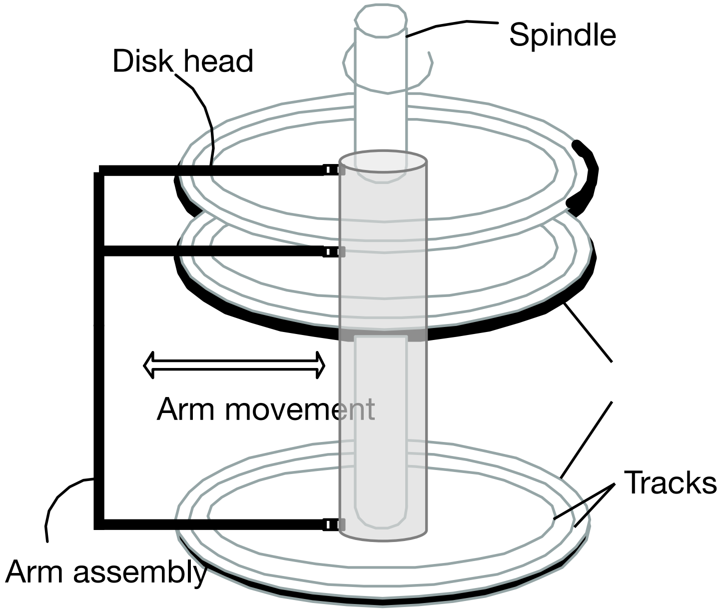
- Read and Write **large chunks of sequential bytes**
- *Sequentially*: “Next” disk block is fastest
- Maximize usage of data per Read/Write
  - “Amortize” seek delays (HDDs) and writes (SSDs):
- Predict future behavior
  - Cache popular blocks
  - Pre-fetch likely-to-be-accessed blocks
  - Buffer writes to sequential blocks
  - More on these as we go

A Note on Terminology

- **Block = Unit of transfer for disk read/write**
  - 64KB – 128KB is a good number today
  - Book says 4KB
- **Page: a common synonym for “block”**
  - In some texts, “page” = a block-sized chunk of RAM

- We’ll treat “block” and “page” as synonyms

Arranging Blocks on Disk



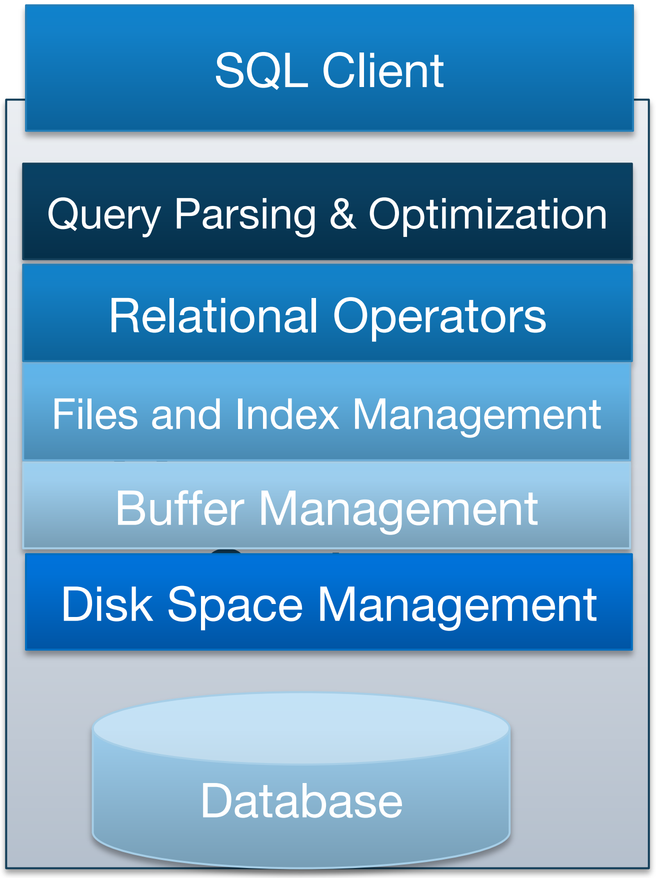
- **‘Next’** block concept:
    - sequential blocks on same track, followed by
    - blocks on same cylinder, followed by
    - blocks on adjacent cylinder

- Arrange file pages sequentially by ‘next’ on disk
    - minimize seek and rotational delay.

- For a **sequential scan**, *pre-fetch*
    - several blocks at a time!
- **Read large consecutive blocks**

Disk Space Management, cont

- **Lowest layer of DBMS, manages space on disk**
- **Purpose:**
  - Map pages to locations on disk
  - Load pages from disk to memory
  - Save pages back to disk & ensuring writes
- Higher levels call upon this layer to:
  - Read/write a page
  - Allocate/de-allocate logical pages



Disk Space Management: Requesting Pages

- Request for a ***sequence*** of pages best satisfied by pages stored sequentially on disk
  - Physical details hidden from higher levels of system
  - Higher levels may “safely” assume **Next Page** is fast, so they will simply expect sequential runs of pages to be quick to scan.

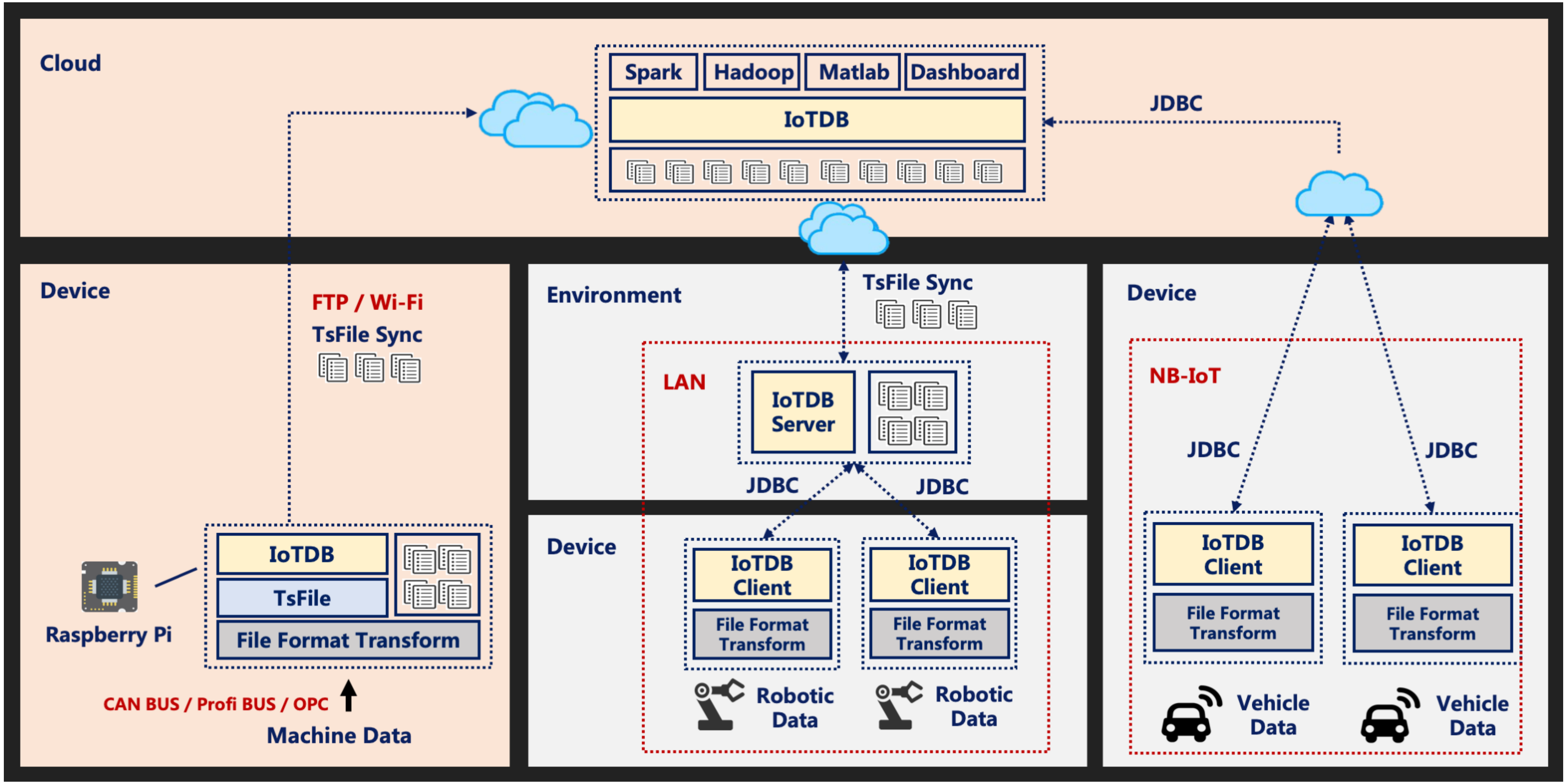
Disk Space Management: Implementation

- **Proposal 1:** Talk to the storage device directly
  - Could be very fast if you knew the device well
  - What happens when devices change?

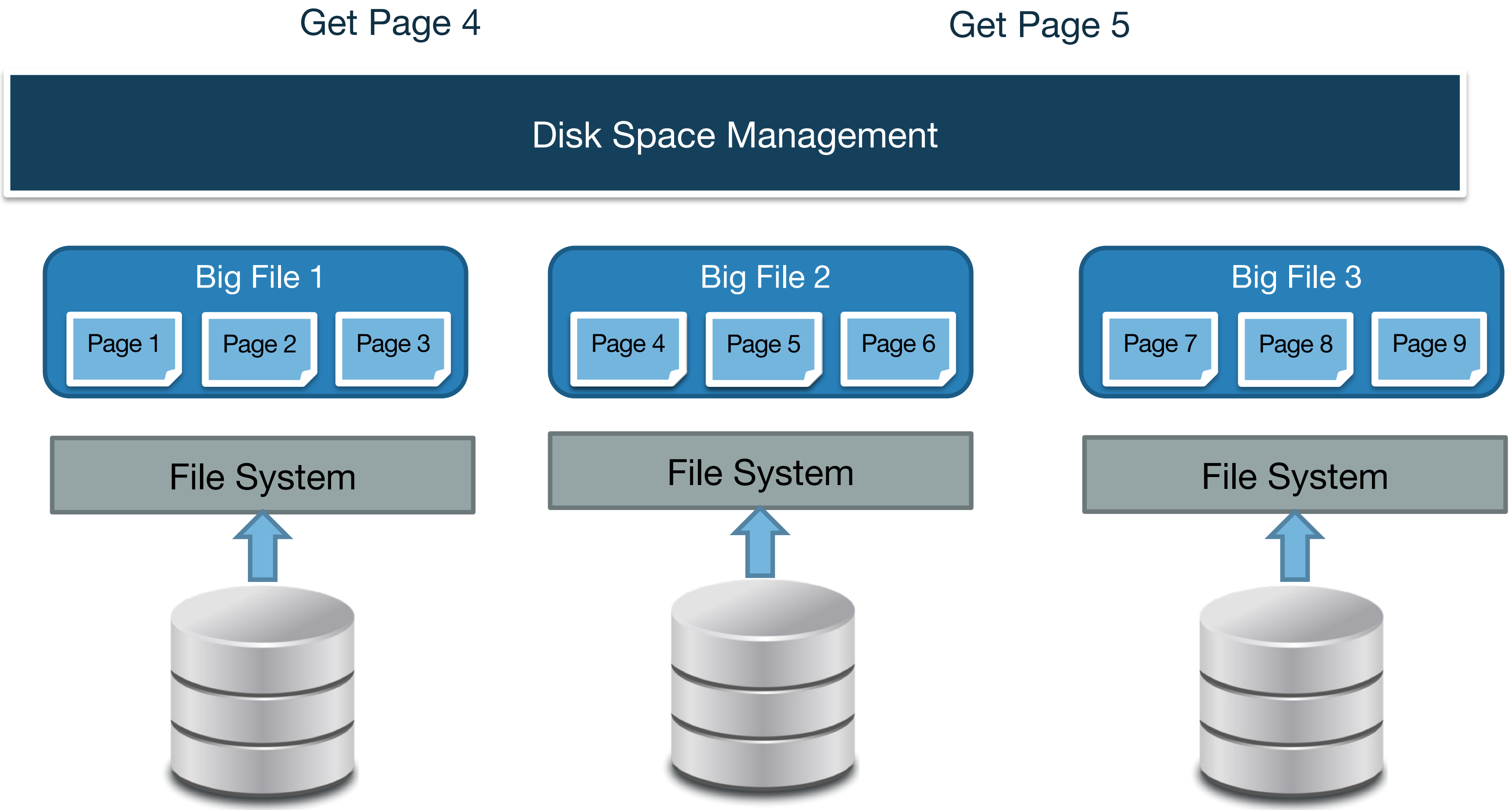
Disk Space Management: Implementation 2

- **Proposal 2**: Run over filesystem (FS)
  - Allocate single large “contiguous” file on a nice empty disk,
    and assume sequential/nearby byte access are fast
  - Most FS optimize disk layout for sequential access
    - Gives us more or less what we want if we start with an empty disk
  - DBMS “file” may span multiple FS files on multiple disks/machines

Example: IoTDB



Using Local Filesystem



Summary: Disk Space Management

- Provide API to read and write pages to device
- Pages: block level organization of bytes on disk
- Provides “next” locality and abstracts FS/device details