

# 1 Fundamentals

## Higher Grade Problem

### Time complexity:

If  $N$  elements in list - compare  $N$  times

If match is found - dequeue element

first case: if match is found for last element

In first case program has to iterate through  $N$  elements ( $O(N)$ ) and call (+1) <sup>räkna ej</sup> dequeue() to iterate through  $N$  elements again to find last element to delete ( $O(N)$ )  
→  $O(2N)$

second case: if match is found for every element

In second case program has to iterate through  $N$  elements ( $O(N)$ ) and then call (+ $N$ ) dequeue() for every  $N$ . dequeue() will always be called for the first element in list, meaning head has to be redirected  $N$  times.

→  $O(3N)$

Worst case:  $O(3N) = O(N)$

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### Memory Complexity:

If  $N$  elements in list -  $N$  memory spaces  
+ constant memory occupants

So...

- $N$  elements
- 1 comparable element
- 1 head node
- 1 counter
- 1 size

which gives:

$N + 4$ , but as  $N \rightarrow \infty$  constants becomes insignificant

$$\rightarrow O(N) = \Omega(N)$$

worst case = every case

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