## Translational Semantics of ASM language in Boogie

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ASM Instruction (S)	Corresponding Boogie Statements ([S])
Stack Handling Instructions	
$push_{\tau} c$	$Stk := [c]::Stk ; (where c is a constant of type \tau \in \{int, bool, string\})$
pop	assert size(Stk)>0; Stk := tl(Stk);
store x	assert size(Stk)>0; $x := hd(Stk)$ ; $Stk := tl(Stk)$ ; (where x is a variable)
load x	Stk := x::Stk; (where x is a variable)
swap	assert size(Stk)>1; Stk := hd(tl(Stk))::hd(Stk)::tl(tl(Stk));
dup	assert size(Stk)>0; Stk := hd(Stk)::Stk;
dup_x1	assert size(Stk)>1; Stk := hd(Stk)::hd(tl(Stk))::hd(Stk)::tl(tl(Stk));
Control Instructions	
if $Stmt_1$ else $Stmt_2$	var cond#: bool;
	$assert size(Stk) > 0 ; cond^{\#} := hd(Stk) ; Stk := tl(Stk) ;$
	if $(\operatorname{cond}^{\#})$ $\llbracket Stmt_1 \rrbracket$ else $\llbracket Stmt_2 \rrbracket$
goto n	goto $l$ ; (where $l$ is a fresh label. It lables the program point which
	corresponds to the ASM instruction offset $n$ )
iter $Stmt_1$ enditer	var col <sup>#</sup> : Collection ;
	assert size(Stk) > 0; col# := hd(Stk); Stk := tl(Stk);
	while (hasNext(col <sup>#</sup> )) $INV$ { Stk := next(col <sup>#</sup> )::Stk ; [ $Stmt_1$ ]]}
pcall sig	$let \ n = arg \ size(sig) \ in$
	$let \ \overline{args} = \overline{tk}(Stk, n), \ ctx = hd(dp(Stk, n)) \ in$
	assert size(Stk) > $n$ ; call invoke(reflect(sig, ctx), $\overline{args}$ );
	Stk := dp(Stk, n+1);
call sig	$let \ n = arg\_size(sig) \ in$
	$  let \overline{args} = tk(Stk, n), ctx = hd(dp(Stk, n)) in$
	var result#: T;
	assert size(Stk) > $n$ ; call result# := invoke(reflect( $sig$ , $ctx$ ), $\overline{args}$ );
	$Stk := result^{\#} : dp(Stk, n+1)$ ;
	(where $T$ is the return type of the reflected method)
Model Handling Instructions	
$\mathbf{new}\ r$	[let mm = hd(Stk), cl = hd(tl(Stk)) in
	$let \ clazz = resolve(mm, cl) \ in :$
	$\operatorname{var} r^{\#} : \operatorname{Ref} ;$
	havoc $r^{\#}$ ; assume $r^{\#} \neq null \land \neg read(heap, r^{\#}, alloc)$ ;
	assert size(Stk) > 1;
	$\textbf{assume typeof}(r^\#) = clazz \; ; \; \text{heap} := \text{update}(\text{heap}, r^\#, alloc, true}); \; \text{Stk} := r^\# :: \text{tl}(\text{tl}(\text{Stk})) \; ;$
get f	$let \ o = hd(Stk) \ in$
	assert size(Stk) > $0 \land o \neq null \land read(heap, o, alloc)$ ;
	Stk := read(heap, o, f) :: tl(Stk);
$\mathbf{set}\ f$	$let \ o = hd(tl(Stk)), \ v = hd(Stk) \ in$
	assert size(Stk) > $1 \land o \neq null \land read(heap, o, alloc)$ ; if (isCollection(f)) { heap := update(heap,read(heap,o,f),read(heap,o,f) $\cup v$ ); }
	else { heap := update(heap, $o,f,v$ ) ; }
	Stk := $tl(tl(Stk))$ ;
	$let \ mm = hd(Stk), \ cl = hd(tl(Stk)) \ in$
findme	assert size(Stk) > 1; Stk := resolve( $mm$ , $cl$ )::tl(tl(Stk));
getasm	Stk := ASM::Stk;
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