

# Happy Number

## Question:

<https://leetcode.com/problems/happy-number/>

Write an algorithm to determine if a number  $n$  is happy.

A happy number is a number defined by the following process:

- Starting with any positive integer, replace the number by the sum of the squares of its digits.
- Repeat the process until the number equals 1 (where it will stay), or it loops endlessly in a cycle which does not include 1.
- Those numbers for which this process ends in 1 are happy.

Return `true` if  $n$  is a happy number, and `false` if not.

## Example 1:

Input:  $n = 19$

Output: `true`

Explanation:

$$1^2 + 9^2 = 82$$

$$8^2 + 2^2 = 68$$

$$6^2 + 8^2 = 100$$

$$1^2 + 0^2 + 0^2 = 1$$

**Example 2:**

Input:  $n = 2$

Output: *false*

## Approach 1:

My approach was to check for a loop. I found this approach by looking at the 2nd Test Case. The second test case moves forward like this: 2 -> 4 -> 16 -> 37 -> 58 -> 89 -> 145 -> 42 -> 20 -> 4 and finally a loop starting from 4.

To check for cycles, I used a hashmap as data fetching takes  $O(1)$  time in a hashmap hence we can check in constant time.

Function for finding the sum of squares of the digits.

```
def check(self, n: int):  
    a = 0  
    while (n > 0):  
        d = n % 10  
        a += d * d  
        n = n // 10  
    return a
```

## Solution 1:

<https://gist.github.com/vermaayush680/574beb6781d419d027e543215c7ca2dd>

```

def isHappy(self, n: int) -> bool:

    d={n}

    while (n!=1):

        n=self.check(n)

        if n in d:

            return False

        d.add(n)

    return True

```

Time Complexity:  $O(n)$

Space Complexity:  $O(n)$

Another approach to finding cycles is Floyd's Cycle-Finding Algorithm.

## Approach 2:

### Floyd's Cycle-Finding Algorithm:

Floyd's cycle-detection algorithm is a two-pointer algorithm used to find the presence of cycles.

Mostly used in linked lists for cycle detection, this algorithm takes constant time.

Explanation:

<https://www.geeksforgeeks.org/detect-loop-in-a-linked-list/>

## Algorithm:

### Use Two Pointers

1. Slow
2. Fast

The fast pointer moves 2 steps at a time while the slow pointer moves 1 step at a time. The idea is that if a cycle is present then at some point in the cycle, the fast pointer will pass the slow pointer and this will prove the existence of a cycle.

If there is no loop then these pointers will never meet and we will reach the end.

Function for finding the sum of squares of the digits. Same as Before.

```
def check(self, n: int) :  
    a=0  
    while (n>0) :  
        d=n%10  
        a+=d**2  
        n=n//10  
    return a
```

## Solution 2:

<https://gist.github.com/vermaayush680/36cf1fecb062e2022b7482149d1e7183>

```
def isHappy(self, n: int) -> bool:  
    slow = n  
    fast = self.check(n)  
    while fast!=1 and fast!=slow:  
        slow=self.check(slow)  
        fast=self.check(self.check(fast))  
    return fast==1
```

Time Complexity:  $O(n)$

Space Complexity:  $O(1)$

### Approach 3:

I found this approach through the discussion forum and think that this is perhaps the best solution among the 3.

After checking various numbers, it was found that there is a single loop that occurs.  
4 -> 16 -> 37 -> 58 -> 89 -> 145 -> 42 -> 20 -> 4

Every number that isn't ending at 1 eventually ends in this loop.

The entry of this loop is 4 so we can just check if the sum of square of digits is 4 if so we return False as it will form a loop.

So similar to approach 1, we check for a loop but instead of checking for repeating numbers, we check for 4 as this is the only sum that forms a loop.

### Solution 3:

<https://gist.github.com/vermaayush680/a710f15143f7143b1b08eee06c63f8c9>

```
def isHappy(self, n: int) -> bool:
    while (n!=1) and n!=4:
        n=self.check(n)
        print(n)
    return n==1
```

Time Complexity:  $O(n)$

Space Complexity:  $O(1)$