

# Inspection of Lisp - Names, Bindings, Scopes and Types

Mahmut Bulut - 14501026  
Computer Engineering Dept., Yıldız Technical University

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## Naming in Lisp

In Lisp symbols are just a name. Naming makes new relations internally. Symbols are datatypes in Lisp. Also symbols are unique identifiers that are identical to other symbols with the same name without restriction with case-sensitivity.

On the other hand Lisp also names functions as variables. This is a distinction with other programming languages.

```
1 => (setq first 'the-bender)
2 -> THE-BENDER
3
4 => (first (list 3 2 1))
5 -> 3
6
7 => first
8 -> THE-BENDER
```

At first and the last one *first* is used as variable. But at second first is used as function. In Lisp a value can have more than one name. This makes aliasing easy and useful.<sup>1</sup>

## Binding in Lisp

Lisp has lexical binding which means that a variable declared with *let* statement will be lexically bound at compile time and cannot be reached outside of lexical scope of *let*.<sup>2</sup> In other words this can be called as static binding.

```
1 (let ((data 1))      ; data is lexically bound.
2   (+ data 3))
3 => 4
4
5 (defun getdata ()
6   data)              ; data is used free in this function.
7
8 (let ((data 1))      ; data is lexically bound.
9   (getdata))
```

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<sup>1</sup><http://psg.com/~dlamkins/sl/chapter03-05.html>

<sup>2</sup>[https://www.gnu.org/software/emacs/manual/html\\_node/elisp/Lexical-Binding.html](https://www.gnu.org/software/emacs/manual/html_node/elisp/Lexical-Binding.html)

```
10 | ;error--> Symbol's value as variable is void: data
```

In addition to lexical binding there is also dynamic binding. In dynamic binding variables lives in one global namespace and can be accessed amongst each other.<sup>3</sup> Dynamic binding makes easier to code greater programs with ease but decrease security with all open namespace, symbol and variable definitions.

## Scope in Lisp

There are several ways to determine the scope:

- Place of the reference in the expression.
- Kind of reference translation takes place.
- Location of reference
- Declaring a variable in local or global scope
- Environment which solves variable bindings within its context.

There are three types of environment in Lisp<sup>4</sup>; global, dynamic and lexical. **Global** environment is like its name global across a program if a declaration made in it, it will be known through the program. **Dynamic** environment is generator like environment. It binds at constructive expressions like *let* block. **Lexical** environment defines a lexical scope within a expression and scope will last as long as this expression lasts.

## Types in Lisp

Data types are generally set of Lisp objects. But lisp has two main types:

- atom
- list

Lists contains atoms or lists of elements. Atoms are seperated in lists from each other with whitespace and they cannot contain anything else except itself<sup>5</sup>. With this seperation expression can be made with combination of lists ad atoms<sup>6</sup>. These types can be included in global scope and it makes easier to infer or coerce between each other. There is already a function named **coerce** in Common Lisp.<sup>7</sup> In addition to it empty data type, which contains no data object are defined by *nil* type. Every type can be created from its relative types and coerce with them. You can see Common Lisp type hierarchy in Fig. 1.

<sup>3</sup><http://emacswiki.org/emacs/DynamicBindingVsLexicalBinding#toc2>

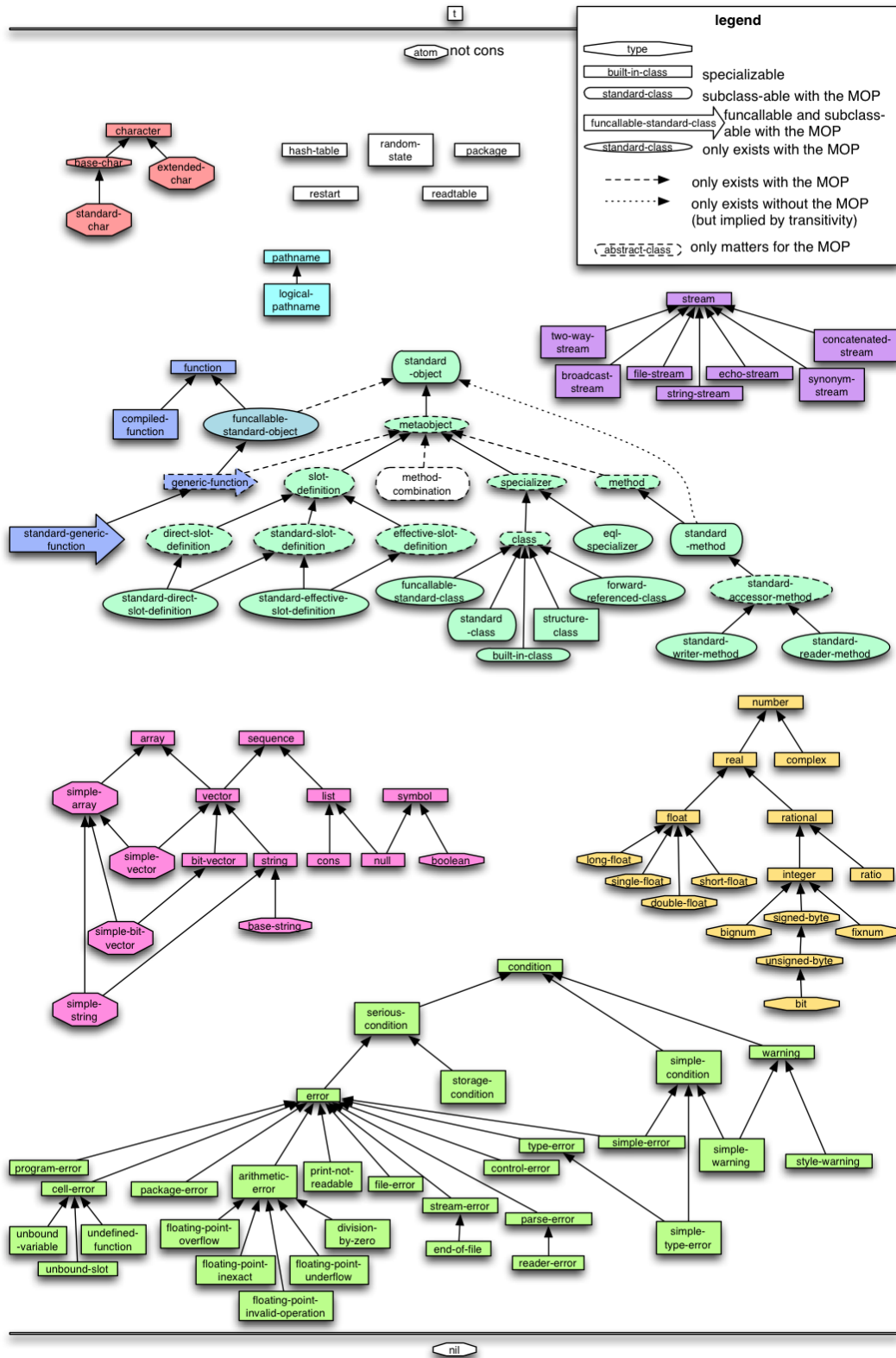
<sup>4</sup>[http://en.wikipedia.org/wiki/Common\\_Lisp#Kinds\\_of\\_environment](http://en.wikipedia.org/wiki/Common_Lisp#Kinds_of_environment)

<sup>5</sup><http://graham.main.nc.us/~bhammel/graham/lisp.html>

<sup>6</sup>[https://www.gnu.org/software/emacs/manual/html\\_node/eintr/Lisp-Atoms.html](https://www.gnu.org/software/emacs/manual/html_node/eintr/Lisp-Atoms.html)

<sup>7</sup><https://www.cs.cmu.edu/Groups/AI/html/cltl/clm/node52.html>

Figure 1: Common Lisp Type Hierarchy



## List of Figures

1	Common Lisp Type Hierarchy . . . . .	3
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