

University of Central Florida

Department of Computer Science

COP 3402: System Software

Spring 2023

Homework #3 (Tiny PL/0 compiler)

Due 3/8/2023 by 11:59 p.m.

This is a solo or team project (Same team as HW2)

REQUIREMENT:

All assignments must compile and run on the Eustis3 server. Please see course website for details concerning use of Eustis3.

Make a copy of lex.c

In the new file lex.c, apply the following changes:

The token list, output HW2, must be kept in the program and or written out to a file (this option will make the parser/codegen slower).

Rename the name of the new copy of lex.c as parsercodegen.c.

Implement the parser/code generator in this file called parsercodegen.c, this means that you will continue inserting code in parsercodegen.c

Objective:

In this assignment, you must implement a Recursive Descent Parser and Intermediate Code Generator for tiny PL/0.

Example of a program written in PL/0:

```
var x, y;  
begin  
    x := y * 2;  
end.
```

Component Descriptions:

The **parser/codegen** must be capable of getting the tokens produced by your Scanner (HW2) and produce, as output, if the program does not follow the grammar, a message indicating the type of error present (**This time: if the scanner step detects an error the compilation process must stop and the error must be indicated, similarly in the parser step, if a syntax error is detected, the compilation process must stop**). A list of the errors that must be considered can be found in Appendix C. In addition, the Parser must populate the Symbol Table, which contains all of the variables and constants names within the PL/0 program. See Appendix E for more information regarding the Symbol Table. If the program is syntactically correct and the Symbol Table is created without error, and code for the virtual machine (HW1) will be generated.

Include always as first entry in the symbol table:

Kind	Name	Value	Level	Address	Mark
3	main	0	0	3	0

For HW3, we will select teams at random to review the compiler. Each team member must know how the compiler and the vm work. If any team member fails in answering a question, a penalty of (-10) will be applied to the whole team in HW3.

Submission Instructions:

1.- Submit via WebCourses:

1. Source code of the tiny- PL/0 compiler (parsercodegen.c).
2. A text file with instructions on how to use your program (readme.txt.).
3. As many Input and output files (cases) to show each one of the errors your compiler can detect, and one correct program. Name them errorin1.text, errorout1.text, errorin2.text, errorout2.text, and so on.
4. All files should be compressed into a single .zip format.
5. Late policy is the same as HW1 and HW2.
6. Only one submission per team: the name of all team members must be written in the source code header file and in the readme document.
7. Include comments in your program
8. Output should print to the screen and should follow the format in Appendix A. A deduction of 5 points will be applied to submissions that do not print to the screen.
9. The input file should be given as a command line argument. A deduction of 5 points will be applied to submissions that do not implement this.

Error Handling

- When your compiler encounters an error, it should print out an error message and stop executing immediately.

Output specifications:

- If you find an error, print it to the screen using the format "Error : <error message>"
- Otherwise, print the assembly code for the virtual machine (HW1) and the symbol table.

See Appendix A

Rubric

15 – Compiles

20 – Produces some instructions before segfaulting or looping infinitely

5 – Follows IO specifications (takes command line argument for input file name and prints output to console)

5 – README.txt containing author names

10 – Correctly create symbol table

10 – Correctly implements expression, term, and factor

10 – Loads and store values correctly

5 – Supports error handling

10 – Correctly implements if statements

10 – Correctly implements while statements

******* If a program does not compile, your grade is zero.**

******* If you do not follow the specifications, your grade is zero. For instance, implementing programming constructs not present in the PL/0 grammar. For example, if you implement procedures, procedure call, it-then-else, your grade will be zero**

Appendix A:

Traces of Execution:

Example 1, if the input is:

```
var x, y;
begin
    x := y * 2;
end.
```

The output should look like:

Assembly Code:

Line	OP	L	M
0	JMP	0	3
1	INC	0	5
2	LOD	0	4
3	LIT	0	2
4	OPR	0	4
5	STO	0	3
6	SYS	0	3

Symbol Table:

Kind	Name	Value	Level	Address	Mark
3	main	0	0	3	1
1	x	0	0	3	1
2	y	0	0	4	1

Example 2, if the input is:

```
var x, y;
begin
    z := y * 2;
end.
```

The output should look like:

Error: undeclared identifier z

Appendix B:

EBNF of tiny PL/0:

program ::= block "." .
 block ::= const-declaration var-declaration statement.
 constdeclaration ::= ["**const**" ident "=" number {"," ident "=" number} ";"].
 var-declaration ::= ["**var**" ident {"," ident} ";"].
 statement ::= [ident ":" expression
 | "**begin**" statement { ";" statement } "**end**"
 | "**if**" condition "**then**" statement
 | "**while**" condition "**do**" statement
 | "**read**" ident
 | "**write**" expression
 | **empty**] .
 condition ::= "**odd**" expression
 | expression rel-op expression.
 rel-op ::= "=" | "<" | "<=" | ">" | ">=" .
 expression ::= ["+" | "-"] term { ("+" | "-") term } .
 term ::= factor { ("*" | "/") factor } .
 factor ::= ident | number | "(" expression ")" .
 number ::= digit { digit } .
 ident ::= letter { letter | digit } .
 digit ::= "0" | "1" | "2" | "3" | "4" | "5" | "6" | "7" | "8" | "9" .
 letter ::= "a" | "b" | ... | "y" | "z" | "A" | "B" | ... | "Y" | "Z" .

Based on Wirth's definition for EBNF we have the following rule:

[] means an optional item.

{ } means repeat 0 or more times.

Terminal symbols are enclosed in quote marks.

A period is used to indicate the end of the definition of a syntactic class.

Appendix C:

Error messages for the tiny PL/0 Parser:

- program must end with period
- const, var, and read keywords must be followed by identifier
- symbol name has already been declared
- constants must be assigned with =
- constants must be assigned an integer value
- constant and variable declarations must be followed by a semicolon
- undeclared identifier
- only variable values may be altered
- assignment statements must use :=
- begin must be followed by end
- if must be followed by then
- while must be followed by do
- condition must contain comparison operator
- right parenthesis must follow left parenthesis
- arithmetic equations must contain operands, parentheses, numbers, or symbols

These are all the error messages you should handle in your parser.

Appendix D: Pseudocode

SYMBOLTABLECHECK (string)

linear search through symbol table looking at name
return index if found, -1 if not

PROGRAM

BLOCK

if token != periodsym

error

emit HALT

BLOCK

CONST-DECLARATION

numVars = VAR-DECLARATION

emit INC ($M = 3 + \text{numVars}$)

STATEMENT

CONST-DECLARATION

if token == const

do

get next token

if token != identsym

error

if SYMBOLTABLECHECK (token) != -1

error

save ident name

get next token

if token != eqlsym

error

get next token

if token != numbersym

error

add to symbol table (kind 1, saved name, number, 0, 0)

get next token

while token == commasym

if token != semicolonsym

error

get next token

VAR-DECLARATION – returns number of variables

numVars = 0

if token == varsym

do

numVars++

get next token

```

        if token != identsym
            error
        if SYMBOLTABLECHECK (token) != -1
            error
        add to symbol table (kind 2, ident, 0, 0, var# + 2)
        get next token
    while token == commasymp
    if token != semicolonsym
        error
    get next token
return numVars

```

STATEMENT

```

    if token == identsym
        symIdx = SYMBOLTABLECHECK (token)
        if symIdx == -1
            error
        if table[symIdx].kind != 2 (not a var)
            error
        get next token
        if token != becomessym
            error
        get next token
        EXPRESSION
        emit STO (M = table[symIdx].addr)
        return
    if token == beginsym
        do
            get next token
            STATEMENT
        while token == semicolonsym
        if token != endsym
            error
        get next token
        return
    if token == ifsym
        get next token
        CONDITION
        jpcIdx = current code index
        emit JPC
        if token != thensym
            error
        get next token
        STATEMENT
        code[jpcIdx].M = current code index
        return
    if token == whilesym

```



```
    get next token
    loopIdx = current code index
    CONDITION
    if token != dosym
        error
    get next token
    jpcIdx = current code index
    emit JPC
    STATEMENT
    emit JMP (M = loopIdx)
    code[jpcIdx].M = current code index
    return
if token == readsym
    get next token
    if token != identsym
        error
    symIdx = SYMBOLTABLECHECK (token)
    if symIdx == -1
        error
    if table[symIdx].kind != 2 (not a var)
        error
    get next token
    emit READ
    emit STO (M = table[symIdx].addr)
    return
if token == writesym
    get next token
    EXPRESSION
    emit WRITE
    return

CONDITION
    if token == oddsym
        get next token
        EXPRESSION
        emit ODD
    else
        EXPRESSION
        if token == eqlsym
            get next token
            EXPRESSION
            emit EQL
        else if token == neqsym
            get next token
            EXPRESSION
            emit NEQ
        else if token == lessym
```

```
        get next token
        EXPRESSION
        emit LSS
    else if token == leqsym
        get next token
        EXPRESSION
        emit LEQ
    else if token == gtrsym
        get next token
        EXPRESSION
        emit GTR
    else if token == geqsym
        get next token
        EXPRESSION
        emit GEQ
    else
        error
```

EXPRESSION

```
    if token == minussym
        get next token
        TERM
        emit NEG
    while token == plussym || token == minussym
        if token == plussym
            get next token
            TERM
            emit ADD
        else
            get next token
            TERM
            emit SUB
```

else

```
    if token == plussym
        get next token
        TERM
    while token == plussym || token == minussym
        if token == plussym
            get next token
            TERM
            emit ADD
        else
            get next token
            TERM
            emit SUB
```

TERM

```
FACTOR
while token == multsym || token == slashsym || token == modsym
    if token == multsym
        get next token
        FACTOR
        emit MUL
    else if token == slashsym
        get next token
        FACTOR
        emit DIV
    else
        get next token
        FACTOR
        emit MOD
```

```
FACTOR
if token == identsym
    symIdx = SYMBOLTABLECHECK (token)
    if symIdx == -1
        error
    if table[symIdx].kind == 1 (const)
        emit LIT (M = table[symIdx].Value)
    else (var)
        emit LOD (M = table[symIdx].addr)
    get next token
else if token == numbersym
    emit LIT
    get next token
else if token == lparensym
    get next token
    EXPRESSION
    if token != rparensym
        error
    get next token
else
    error
```

Appendix E:

Symbol Table

Recommended data structure for the symbol.

```
typedef struct
{
    int kind;           // const = 1, var = 2
    char name[12];      // name up to 11 chars
    int val;            // number
    int level;          // L level
    int addr;           // M address
} symbol;
```

```
symbol_table[MAX_SYMBOL_TABLE_SIZE = 500];
```

For constants, you must store kind, name and value.

For variables, you must store kind, name, L and M.

