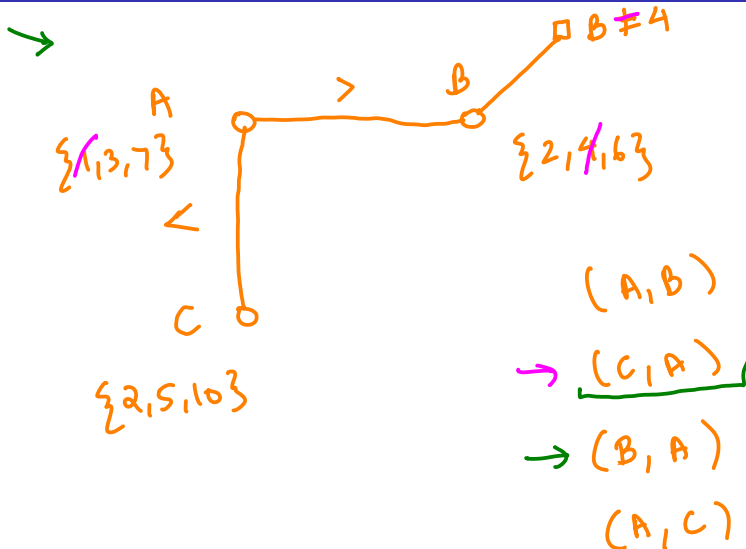


# Finding a solution for a CSP

- ▶ The problem :  $O(d^n)$  possible assignments.
- ▶ Reduce the number of legal values for different variables.
- ▶ Achieve node consistency  
 $WA \neq \text{Green}$
- ▶ Achieve arc consistency

# Binary Constraint and Arc consistency



# AC-3 Algorithm

C 2C

**function** AC-3(*csp*) **returns** false if an inconsistency is found and true otherwise

**inputs:** *csp*, a binary CSP with components ( $X$ ,  $D$ ,  $C$ )

**local variables:** *queue*, a queue of arcs, initially all the arcs in *csp*

**while** *queue* is not empty **do**

→  $(X_i, X_j) \leftarrow \text{REMOVE-FIRST}(\text{queue})$

→ **if** REVISE(*csp*,  $X_i$ ,  $X_j$ ) **then**

**if** size of  $D_i = 0$  **then return** false

**for each**  $X_k$  **in**  $X_i.\text{NEIGHBORS} - \{X_j\}$  **do**

        add  $(X_k, X_i)$  to *queue*

**return** true

---

**function** REVISE(*csp*,  $X_i$ ,  $X_j$ ) **returns** true iff we revise the domain of  $X_i$

*revised*  $\leftarrow$  false

**for each**  $x$  **in**  $D_i$  **do**

**if** no value  $y$  in  $D_j$  allows  $(x, y)$  to satisfy the constraint between  $X_i$  and  $X_j$  **then**

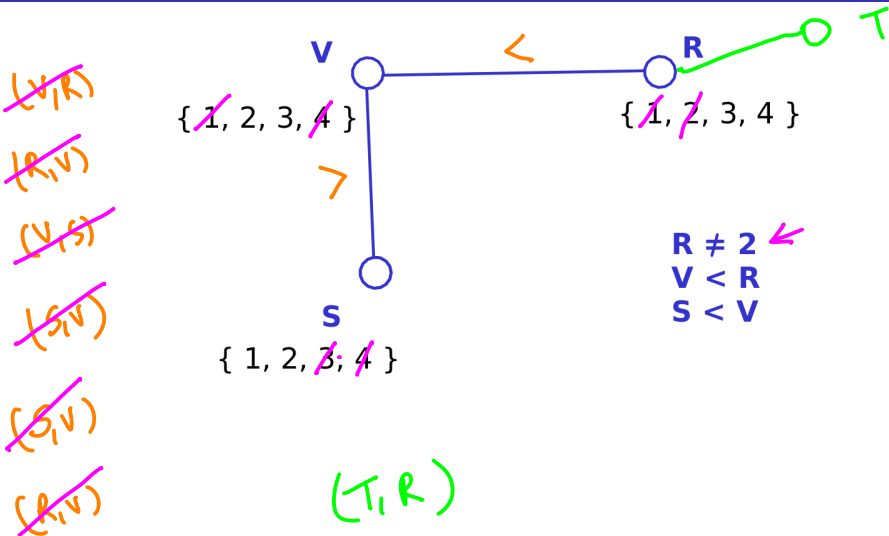
        delete  $x$  from  $D_i$

*revised*  $\leftarrow$  true ←

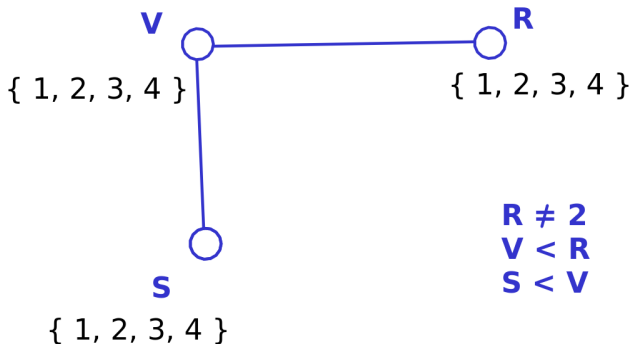
**return** *revised*

# Binary Constraint and Arc consistency

# AC-3 Algorithm Example



# AC-3 Algorithm Example



- The search space is reduced from  $4^3 = 64$  to  $2^3 = 8$ .



# Time complexity of AC-3 algorithm

$$\begin{aligned} & \text{Handwritten notes: } 2c, d^2 \text{ (with arrows pointing to } 2c \times d \text{ and } d^2 \text{)} \\ & \text{Handwritten notes: } (x_k, x_i) \leftarrow d \text{ (with arrows pointing to } x_k, x_i \text{ and } d \text{)} \\ & \text{Handwritten notes: } 2c \times d + d^2 \text{ (with arrows pointing to } 2c \times d \text{ and } d^2 \text{)} \\ & \text{Handwritten notes: } O(cd^3) \end{aligned}$$

- Domain size at most  $d$  and  $c$  binary constraints.

$$\rightarrow (x_k, x_i) \leftarrow d$$

# Time complexity of AC-3 algorithm

- ▶ Domain size at most  $d$  and  $c$  binary constraints.
- ▶ Each arc  $(X_k, X_i)$  can be inserted into the queue at most  $d$  times.



# Time complexity of AC-3 algorithm

- ▶ Domain size at most  $d$  and  $c$  binary constraints.
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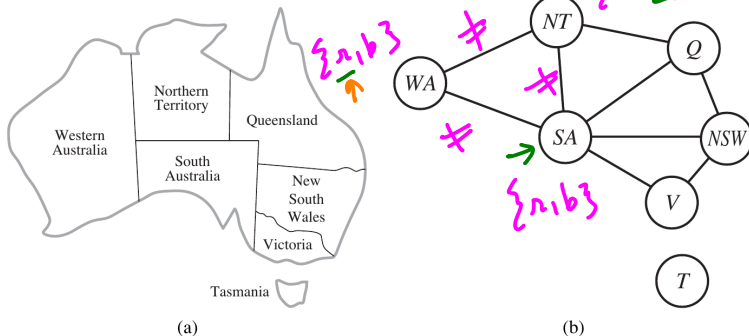
# Time complexity of AC-3 algorithm

- ▶ Domain size at most  $d$  and  $c$  binary constraints.
- ▶ Each arc  $(X_k, X_i)$  can be inserted into the queue at most  $d$  times.
- ▶ Arc consistency can be checked in  $O(d^2)$  time.
- ▶ Worst case time complexity of AC-3 algorithm  $O(cd^3)$ .

$$O(d^n)$$

# A stronger notion of consistency

Can we use node and arc consistency to detect that the map cannot be coloured using two colors:  $\{\text{red}, \text{blue}\}$ ?



**Figure 6.1** (a) The principal states and territories of Australia. Coloring this map can be viewed as a constraint satisfaction problem (CSP). The goal is to assign colors to each region so that no neighboring regions have the same color. (b) The map-coloring problem represented as a constraint graph.

- ▶ A two-variable set  $\{X_i, X_j\}$  is path-consistent with respect to a third variable  $X_m$  if, for every *consistent* assignment  $\{X_i = a, X_j = b\}$ , there is a *consistent* assignment to  $X_m$  such that the constraints on  $\{X_i, X_m\}$  and  $\{X_m, X_j\}$  are satisfied.

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- ▶ We can detect that the australian map cannot be colored using two colors.

# $k$ -consistency

- ▶ A CSP is  $k$ -consistent if, for any consistent assignment to  $k - 1$  variables, there is a consistent assignment for the  $k^{th}$  variable.

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- ▶ A CSP is  $k$ -consistent if, for any consistent assignment to  $k - 1$  variables, there is a consistent assignment for the  $k^{th}$  variable.
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