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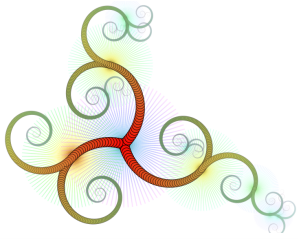
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# **HAMPI**

## **A Solver for String Constraints**

**Vijay Ganesh**  
**MIT**

**(With Adam Kiezun, Philip Guo,  
Pieter Hooimeijer and Mike Ernst)**



# The Problem

Context-free Grammars, Regular expressions,  
string variable



**HAMPI  
String Solver**



String

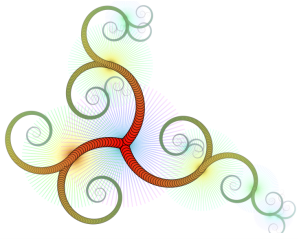
UNSAT

- Emptiness problem for an intersection of Context-free Grammars

$$s \text{ in } (L_1 \cap \dots \cap L_N)$$

where

- $s$  is bounded
- $s$  contains some substring
- Extended Backus-Naur Form
- Different from string matching



# A Problem Instance

Context-free Grammars, Regular expressions,  
string variable



**HAMPI**  
**String Solver**

String

UNSAT

```
var v:4;  
  
cfg E := "()" | E E | "(" E ")";  
  
val q := concat("(" , v , ")");  
  
assert q contains "()()";
```

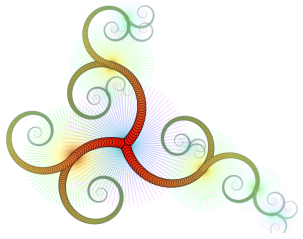
“Find a 4-character string  $v$ , such that:

- $(v)$  has balanced parentheses, and
- $(v)$  contains substring  $()()$ ”

HAMPI finds satisfying assignment

$v = )() ($

Hence,  $q = ()()()$



# HAMPI

## A Novel String Solver

Analyses of string programs

- Formal Methods
- Testing
- Program Analysis

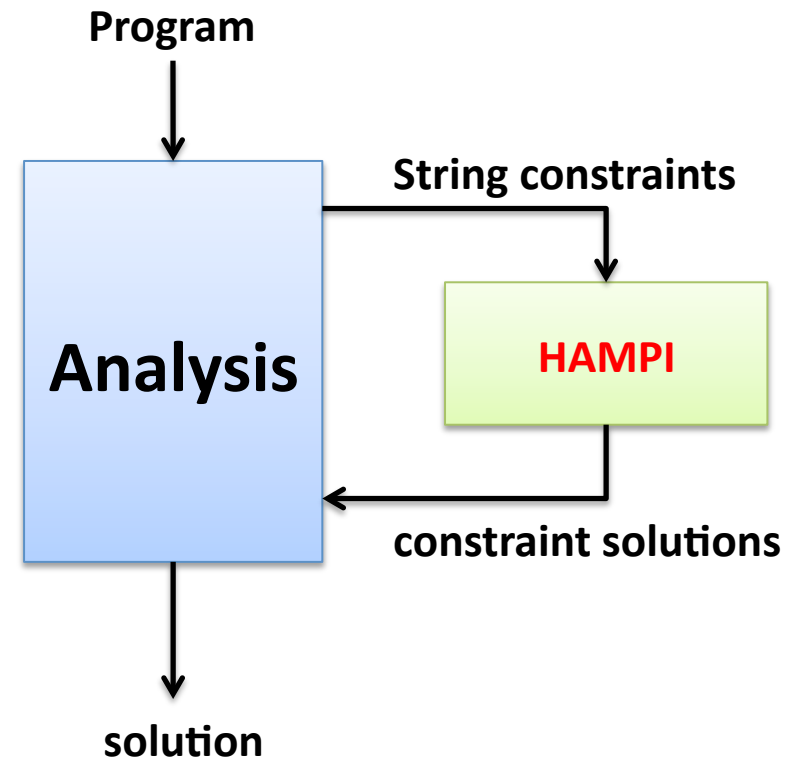
**Efficient**

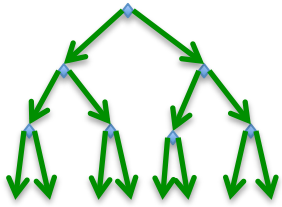
**Expressive**

**Robust** and easy to use

**Tested** on **many diverse apps**

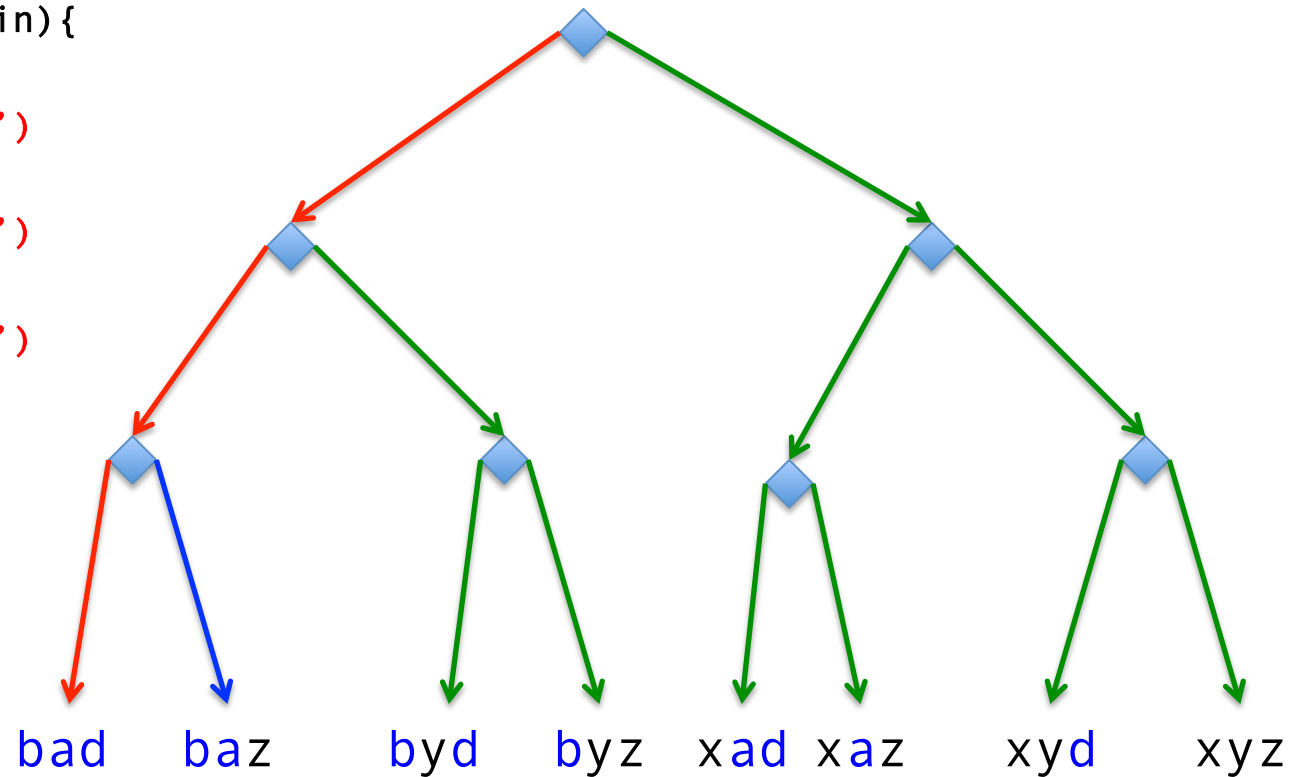
Applied to important and hard problems



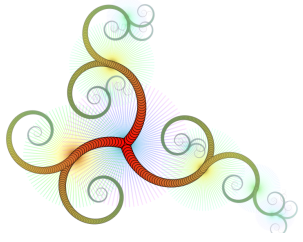


# HAMPI for Dynamic Systematic Testing

```
void main(char[] in){  
    int count=0;  
    if (in[0] == 'b')  
        count++;  
    if (in[1] == 'a')  
        count++;  
    if (in[2] == 'd')  
        count++;  
    if (count == 3)  
        ERROR;  
}
```



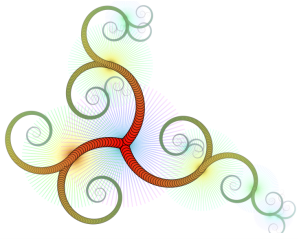
Concolic testing creates inputs for all program paths (assuming finite loops)



# HAMPI for Dynamic Systematic Testing

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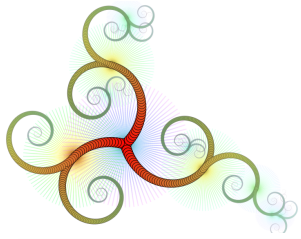
- ✓ **Key to success: Efficient, expressive constraint solver**
- ✓ HAMPI can also be used to produce structured inputs that skip parsing and penetrate deep during concolic testing
- ✓ Engler et al. (EXE, KLEE), Godefroid et al. (DART, SAGE)
- ✓ CUTE, CREST, SmartFuzz etc.



# Typical HAMPI Applications

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- **Constraints generated by symbolic Analyses of string programs**
  - Concolic Testing
  - Formal Methods
  - Program Analysis
  - Checking whether input sanitation code actually works
  - Applied to PHP scripts, JavaScript, C/Java etc.
- **Automatic generation of structured inputs with constraints**
  - String inputs for programs
  - Document files, e.g. PDF, to test PDF reader
  - SQL commands with attack vectors
  - Programming language source files

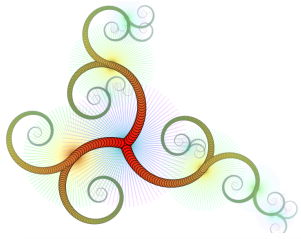


# HAMPI Results Summary

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- ✓ **Expressive:** Supports context-free grammars, regular expression operations
- ✓ **Efficient:** ~7x faster on-average, and  
Often 100s of times faster than CFGAnalyzer
- ✓ **Effectively enabled dynamic systematic testing (concolic testing) of real-world string-manipulating programs**
- ✓ **Already plugged into important analysis infrastructures, such as NASA Java PathFinder**

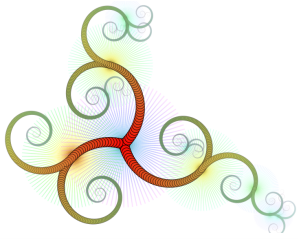




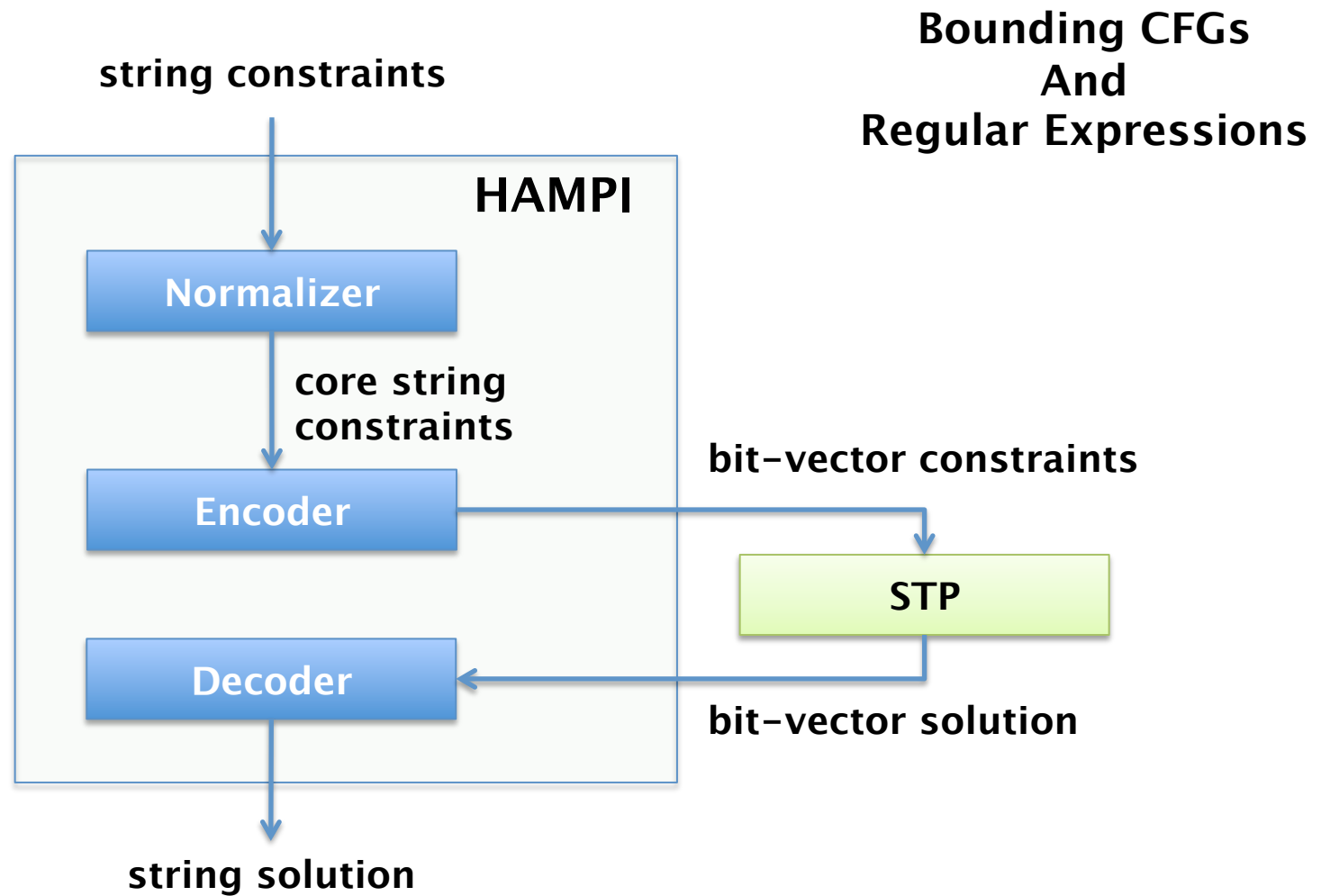
# HAMPI Results Summary

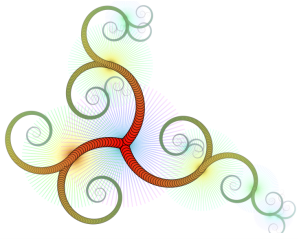
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- ✓ SQL injection vulnerability detection using static analysis
- ✓ SQL injection vulnerability generation using dynamic analysis (**Ardilla tool**, ICSE 2009)
  - ✓ 60 attacks (23 SQL injection, 37 XSS) on 5 PHP applications (300K+ LOC)
- ✓ Automatic generation of structured inputs for **Klee** concolic tester, improved code coverage and bug finding
- ✓ Classic decision problems for context-free grammars

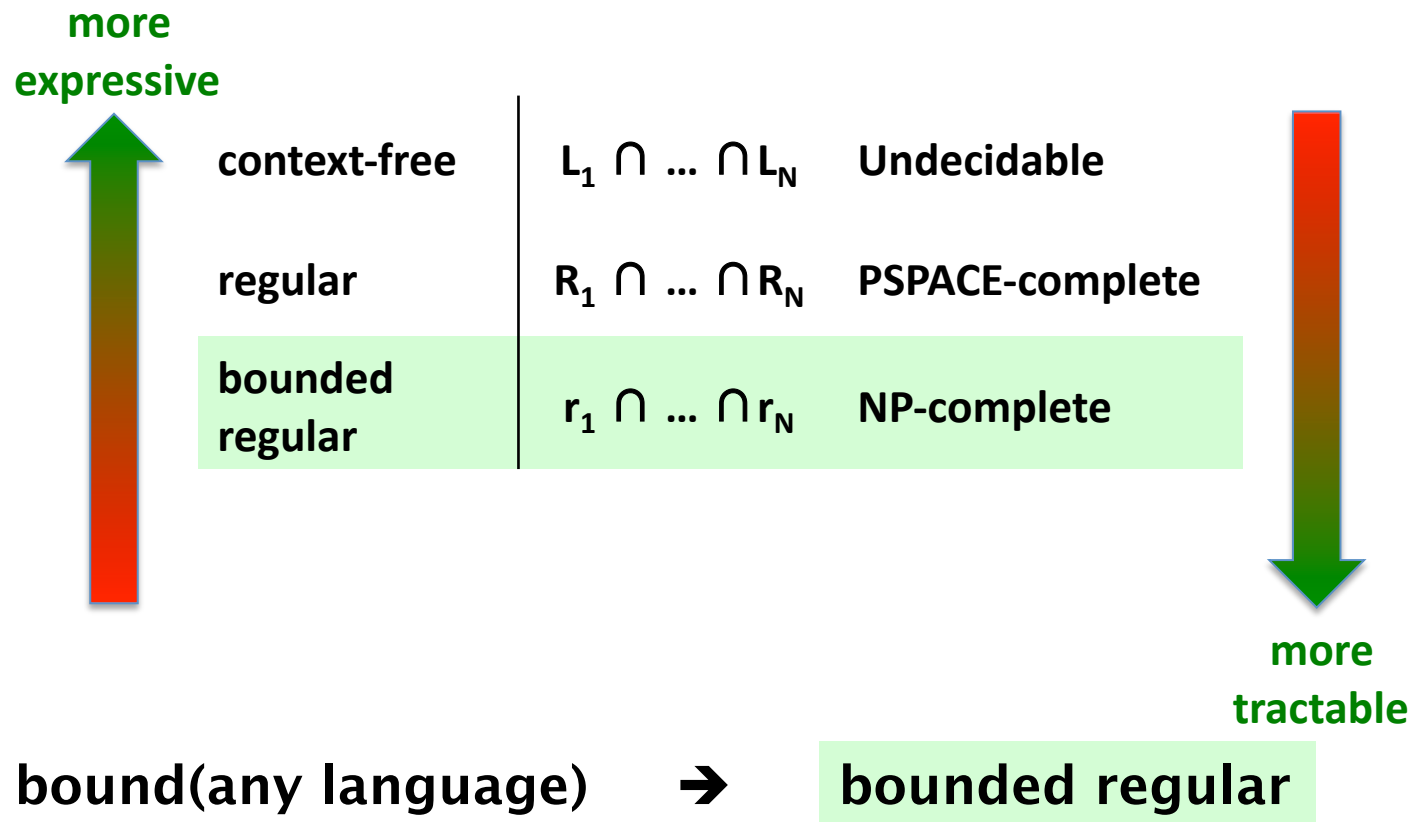


# HAMPI Internals



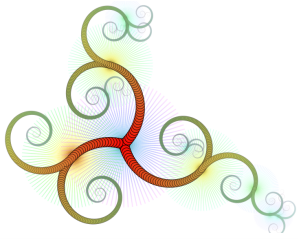


# HAMPI: Bounding is GOOD



## Key HAMPI idea:

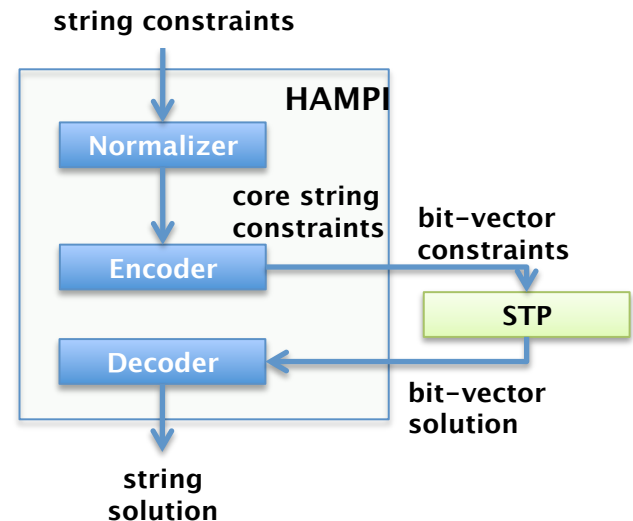
1. Bound length of strings for high expressiveness, efficiency
2. Typical applications require short solutions



# HAMPI Example

```
var v:4;  
  
cfg E := "()" | E E | "(" E " )";  
  
val q := concat("(", v, ")");  
  
assert q in bound(E, 6);  
assert q contains "()()";
```

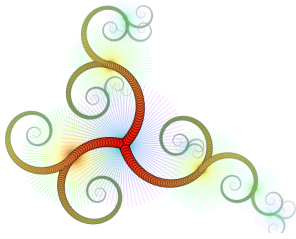
Constraints



“Find a 4-character string  $v$ , such that:

- $(v)$  has balanced parentheses, and
- $(v)$  contains substring  $()()$ ”

HAMPI finds satisfying assignment  $v = )() ($



# HAMPI Normalizer

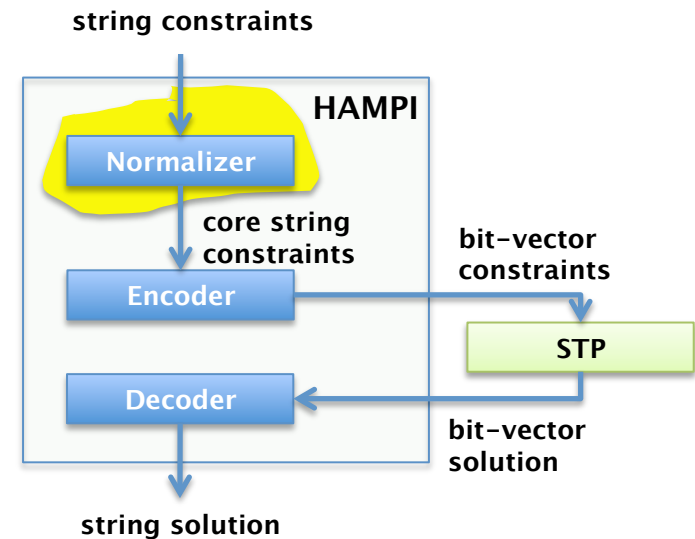
Core string constraint have only regular expressions

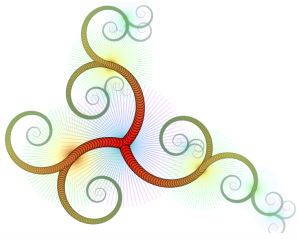
Expand grammars to regexps

- expand nonterminals
- eliminate inconsistencies
- enumerate choices exhaustively
- sub-expression sharing

cfg  $E := "(" E ")" \mid E E \mid "()"$ ;

$\rightarrow \text{bound}(E, 6) \rightarrow ( [ ()() + (() ) ] +$   
 $[ ()() + (() ) ] ( )$



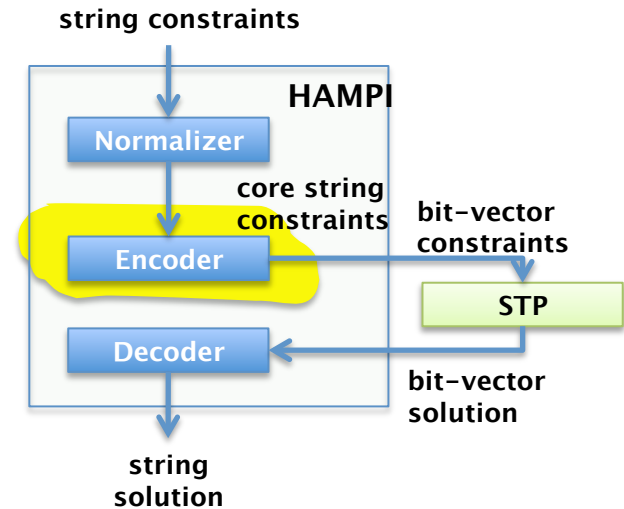


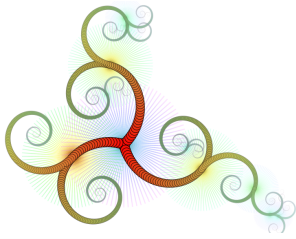
# Bit Vectors Are Ordered, Fixed-Size, Sets Of Bits

Bit vector B (length 6 bits)

0	1	0	1	0	1
---	---	---	---	---	---

$(B[0:1] = B[2:3]) \wedge (B[1:3] = 101)$





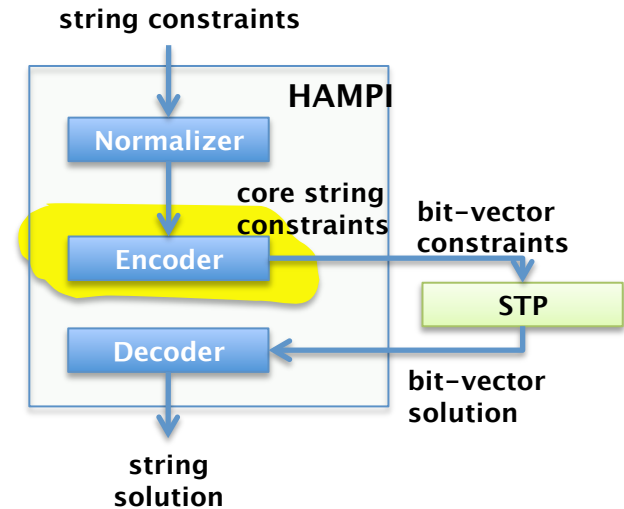
# HAMPI Encodes Input As Bit-Vectors

Map alphabet  $\Sigma$  to bit-vector constants:

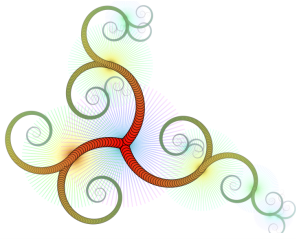
(  $\rightarrow$  0  
 )  $\rightarrow$  1

Compute size of bit-vector B:

$(1+4+1) * 1 \text{ bit} = 6 \text{ bits}$



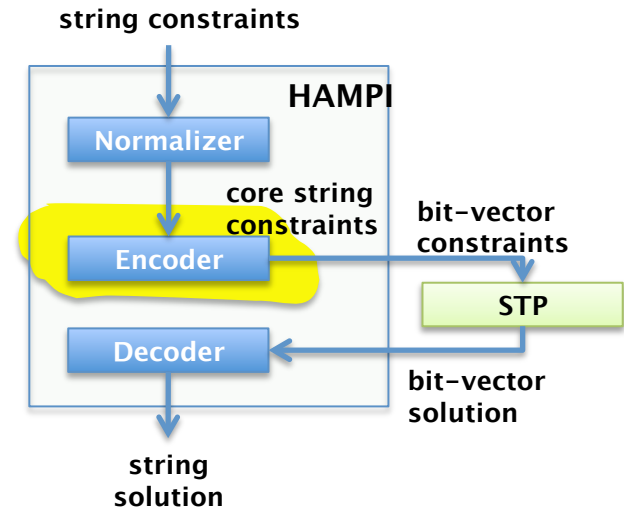
( v )  $\in$  ( ) [ ( ) ( ) + ( ( ) ) ] + [ ( ) ( ) + ( ( ) ) ] ( ) + ( [ ( ) ( ) + ( ( ) ) ] )



# HAMPI Encodes Regular Expressions Recursively

## Encode regular expressions recursively

- union + → disjunction  $\vee$
- concatenation → conjunction  $\wedge$
- Kleene star \* → conjunction  $\wedge$
- constant → bit-vector constant



$$(v) \in () [ () () + (() ) ] + [ () () + (() ) ] () + ( [ () () + (() ) ] )$$

Formula  $\Phi_1$

$\vee$

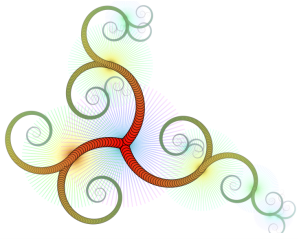
Formula  $\Phi_2$

$\vee$

Formula  $\Phi_3$

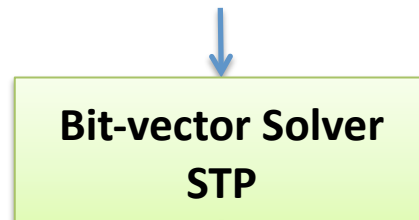
$$B[0]=0 \wedge B[1]=1 \wedge \{ B[2]=0 \wedge B[3]=1 \wedge B[4]=0 \wedge B[5]=1 \vee \dots$$





# HAMPI Uses STP Solver And Decodes Solution

bit-vector constraints

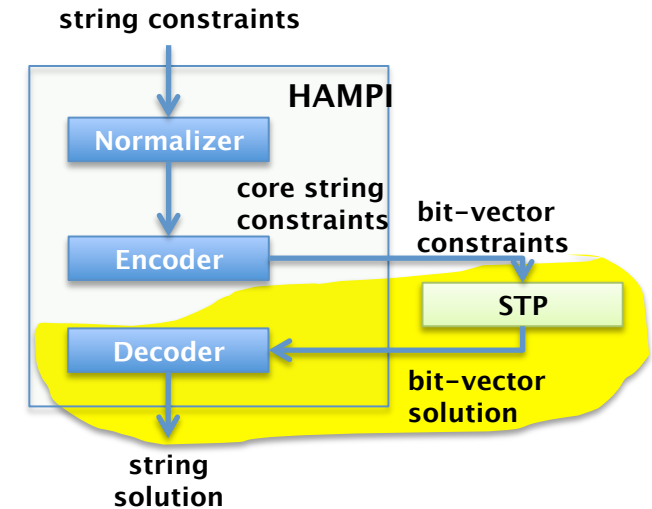


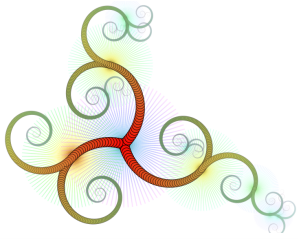
B = 010101



B = ( ) ( ) ( )  
v = ) ( ) (

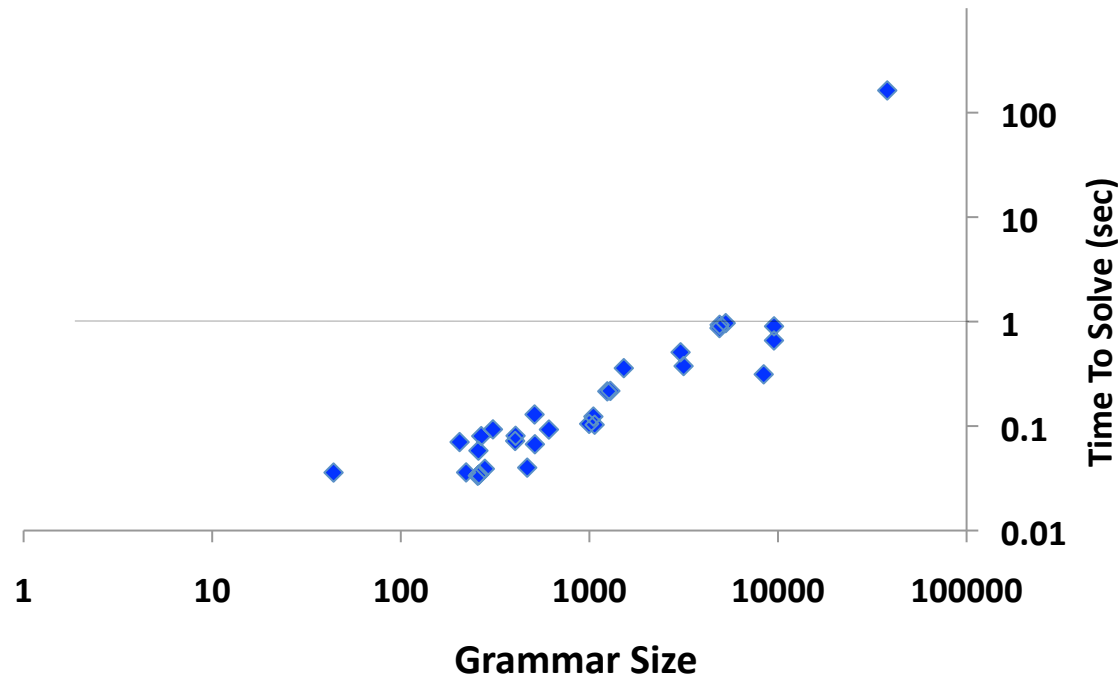
Maps bits back to  
alphabet  $\Sigma$





## Result 1: HAMPI Is Effective In Static SQL Injection Analysis

1367 string constraints from [Wassermann PLDI'07]



HAMPI scales to **large grammars**

HAMPI solved **99.7%** of constraints in **< 1** sec per constraint

All solvable constraints had short solutions  **$N \leq 4$**



## Result 2: HAMPI helps Ardilla Find New Vulnerabilities (Dynamic Analysis)

---

**60** attacks on 5 PHP applications (300K+ LOC)

23 SQL injection

29 XSS

4 cases of data corruption  
19 cases of information leak

216 HAMPI constraints solved

- **46%** of constraints in **< 1 second** per constraint
- **100%** of constraints in **< 10 seconds** per constraint



## **Result 3: HAMPI helps Klee Concolic Tester Find New Bugs**

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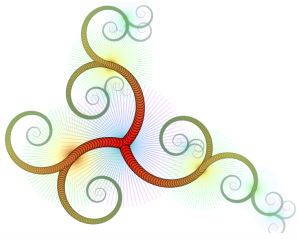
- **Problem: For programs with highly structured inputs, concolic testers can spend too long in the parser**
- **The reason: We may not know which part of input to mark symbolic, and hence mark too much**
- **It is better to generate valid highly structured inputs**
- **Penetrate deep into the program's semantic core**



## Result 3: HAMPI helps Klee Concolic Tester Find New Bugs

Program Name	Marking Input Symbolic <b>Klee style</b> (imperative) legal /total inputs Generated (1 hour)	Marking Input Symbolic <b>HAMPI-2-Klee style</b> (declarative) legal /total inputs generated (1 hour)
Cueconvert (music format converter)	0/14	146/146
Logictree (SAT solver)	70/110	98/98
Bc (calculator)	2/27	198/198

- Improved Code Coverage dramatically (from 30 to 50% with 1 hour work)
- Found 3 new errors in Logictree

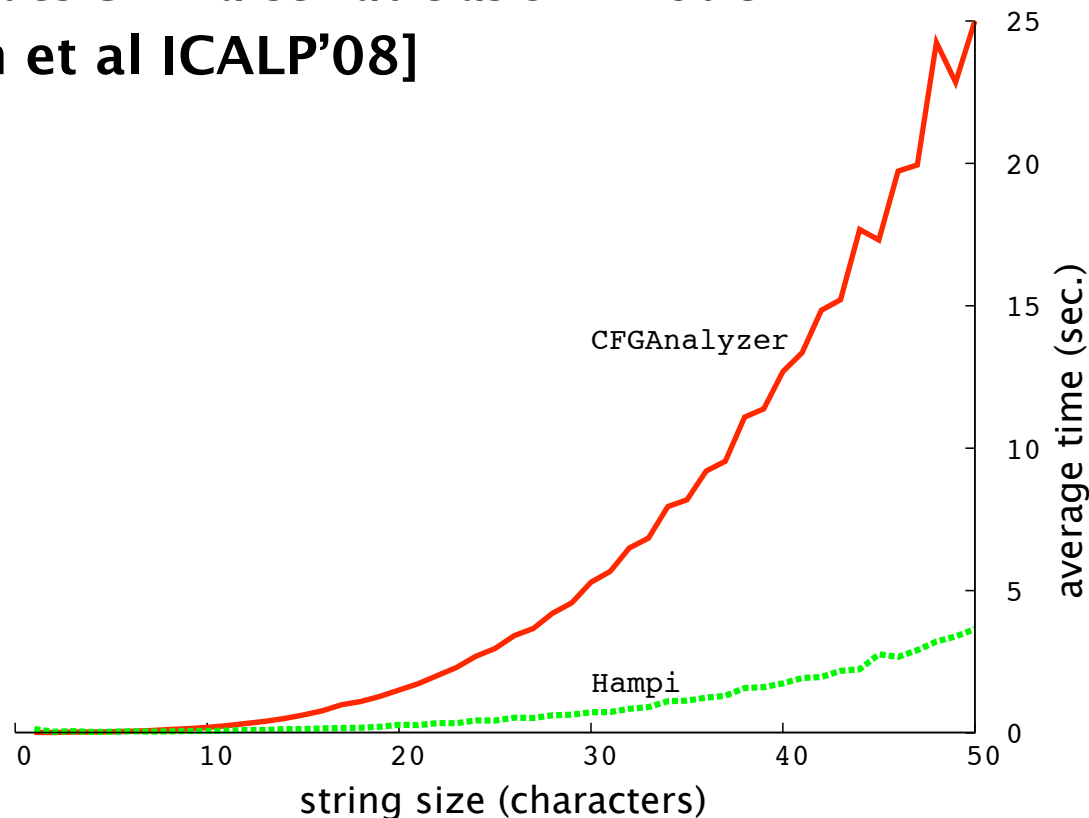


## Result 4: HAMPI Is Faster Than The CFGAnalyzer Solver

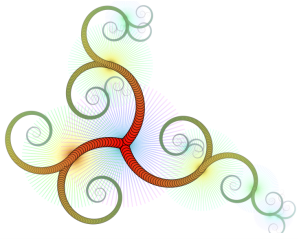
CFGAnalyzer encodes bounded grammar problems in SAT

- Encodes CYK Parse Table as SAT Problem

[Axelsson et al ICALP'08]

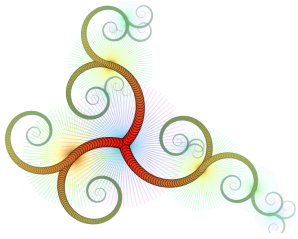


For size 50, HAMPI is **6.8x** faster on average (up to **3000x** faster)



# HAMPI Supports Rich String Constraints

	HAMPI	CFGAnalyzer	Wassermann	Bjorner	Hooijmeier	Emmi	MONA
context-free grammars	●	●	◐				
regular expressions	●	●	◐		●	●	◐
string concatenation	●			●	●		●
stand-alone tool	●	●			●		●
unbounded length			●		●	●	●

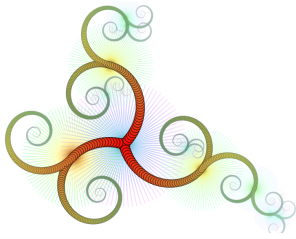


# Conclusions

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- ✓ **HAMPI: A Novel Solver for String Constraints**
- ✓ **Efficient**
- ✓ **Rich Input Language**
- ✓ **Widely applicable: Formal Methods, Testing, Analysis**
- ✓ **Already tested in many real-world apps**
- ✓ **Part of well-known infrastructure: e.g., NASA Java PathFinder**
- ✓ **Download Source + All Experimental Data**
  - ✓ <http://people.csail.mit.edu/akiezun/hampi>
  - ✓ <http://people.csail.mit.edu/vganesh/stp.html>





# Moral of the Story

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- ✓ USE HAMPI and STP
- ✓ **Download Source + All Experimental Data**
  - ✓ <http://people.csail.mit.edu/akiezun/hampi>
  - ✓ <http://people.csail.mit.edu/vganesh/stp.html>



# Analysis for SQL Injection Attacks

---

```
$my_topicid = $_GET['topicid'];  
  
$sqlstmt = "SELECT msg FROM messages WHERE topicid = '$my_topicid';"  
$result = mysql_query($sqlstmt);  
  
WHILE ($row = mysql_fetch_assoc($result)) {  
    echo "Message " . $row['msg'];  
}
```

- PHP Program as part of a website: <http://www.mysite.com/?topicid=1>
- Access Database with command:  
**SELECT msg FROM messages WHERE topicid = 1**
- Attacker can reveal the entire database by URL:  
<http://www.mysite.com/?topicid=' OR '1'='1>



# HAMPI Constraints That Create SQL Injection Attacks

user input  
string {

```
var v : 12;
```

SQL grammar {

```
cfg SqlSmall := "SELECT " [a-z]+ " FROM " [a-z]+ " WHERE " Cond;  
cfg Cond := Val "=" Val | Cond " OR " Cond;  
cfg Val := [a-z]+ | "'" [a-z0-9]* "'" | [0-9]+;
```

bounded  
SQL grammar {

```
reg SqlSmallBounded := bound(SqlSmall, 53);
```

SQL query {

```
val q := concat("SELECT msg FROM messages WHERE topicid=", v, "");
```

SQLI attack  
conditions {

```
assert q in SqlSmallBounded;  
assert q contains "OR '1'='1'";
```

“q is a valid SQL query”

“q contains an attack tautology”

HAMPI finds an attack input:  $v \rightarrow 1' \text{ OR } '1'=1$   
**SELECT msg FROM messages WHERE topicid=1' OR '1'=1'**