

# **Impact of Community Structure on SAT Solver Performance**

**by**

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# Motivation and Problem Statement

## Why are CDCL Solvers efficient for Industrial Instances

- CDCL Solvers are remarkably efficient for large industrial instances
- This is true for industrial instances from a diverse set of apps
- Why is this so?

# Structure in Industrial Instances

## Why are CDCL SAT Solvers efficient for Industrial Instances

- CDCL solvers exploit structure Inherent to SAT instances
- Relevant questions
  - What structure?
  - What evidence connects so-called structure and solver performance?
  - How?
- In this talk we discuss an answer to Question 2

# Community Structure and SAT Solver Performance

## Our Results

- Take-home Message
  - Community structure (whose quality is measured using metric called Q) of SAT instances strongly affect solver performance
- Result #1: Strong correlation between community structure and LBD (Literal Block Distance) in Glucose solver
- Result #2: Hard random instances have low Q ( $0.05 \leq Q \leq 0.13$ )
- Result #3: Number of communities and Q of SAT instances are more predictive of CDCL solver performance than other measures

# Structure in Industrial Instances

## Previous Approaches to Characterizing Structure

- Number of variables ( $N_v$ ) and clauses ( $N_c$ ), and functions over  $N_v$  and  $N_c$  (e.g., clause-variable ratio)
  - Lack explanatory and predictive power for industrial instances
- Large successful predictive model by Xu, Hoos et al.
  - Basis for a machine learning based predictor [XHHL08]
- Wanted
  - Small predictive set of features that forms the basis for a complete explanation (we don't have it yet)

# Community Structure and SAT Solver Performance

## Graphs and their Community Structure

- Our proposal (also by Ansotegui, Levy et al. [AL14])
  - **Community structure, specifically, modularity ( $Q$ ) and number of communities correlate with CDCL solver performance**
- Definition of community in a graph
  - A sub-graph tightly connected internally, but weakly connected externally
- Industrial SAT instances community structure  
(Ansotegui, Levy et al.[AL12, AL13])

# Modularity (Q-factor) and Communities in Graphs

## Community Structure in Graphs

- Modularity of Q lies between 0 and 1
- Q measures quality
  - Higher Q implies “good community structure”, i.e., *highly separable communities*
  - Lower Q implies “bad community structure”, i.e., *one giant hairy ball*
- $Q_{\text{simple}} = \frac{\text{Number of edges inside communities}}{\text{Total number of edges}}$
- Problem
  - The community containing entire graph has  $Q_{\text{simple}} = 1$

# Modularity (Q-factor) and Communities in Graphs

## Community Structure in Graphs

- Problem
  - The community containing the entire graph has  $Q_{\text{simple}} = 1$
- Solution
  - For all community structures over given graph  $G$ 
    - Compute  $Q_{\text{simple}}$
    - Randomize  $G$  and compute  $Q_{\text{randomized}}$
    - Compute  $Q = Q_{\text{simple}} - Q_{\text{randomized}}$
  - $Q$  of graph  $G = \text{Optimal}(Q_{\text{simple}} - Q_{\text{randomized}})$
  - The modularity ( $Q$ ) is the optimal structure furthest from a random-looking version of  $G$

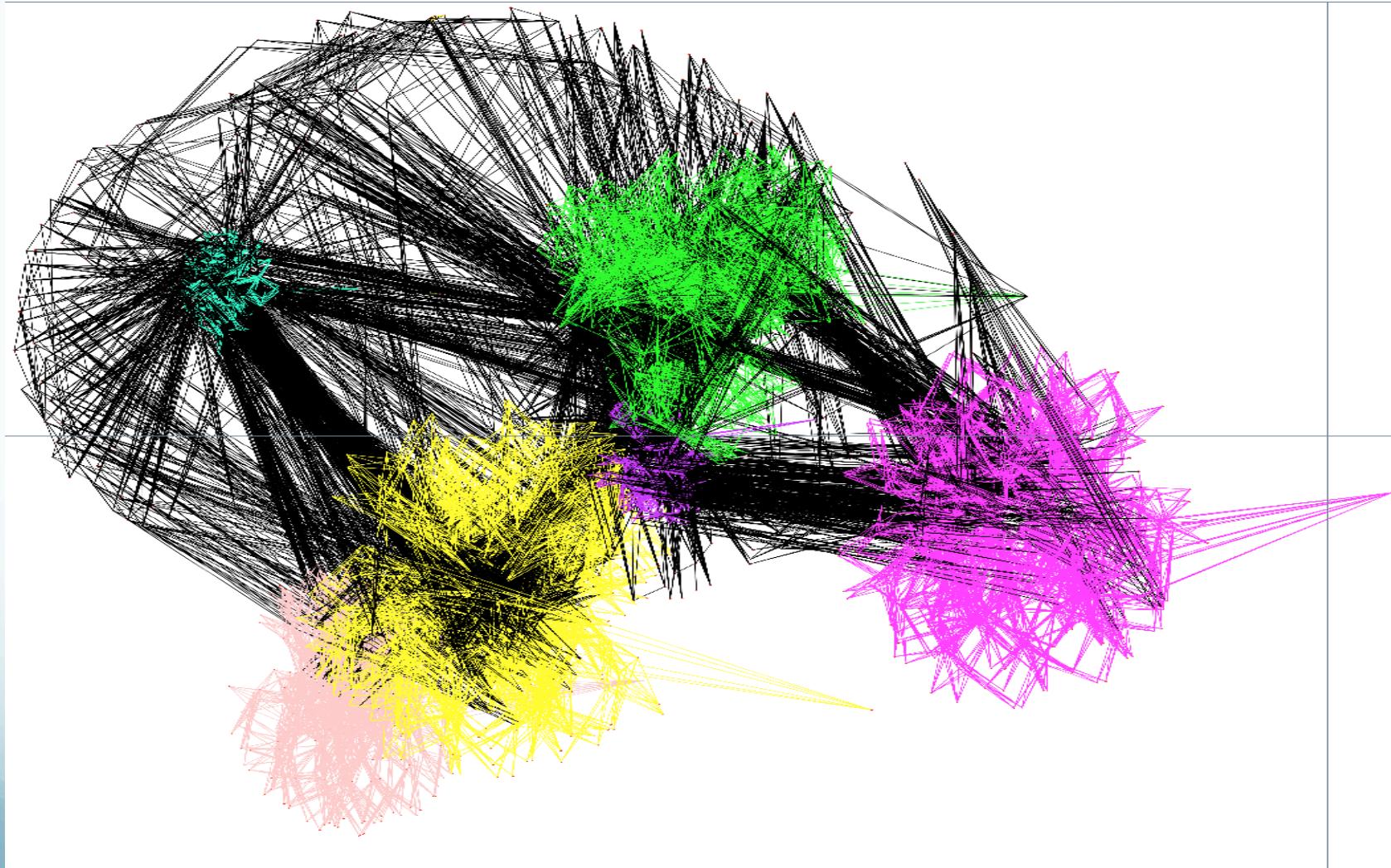
# A Note on Community Structure in Graphs

## Applications

- Community structure [GN03,CNM04,OL13] is used to study all kinds of complex networks, such as,
  - Social networks, e.g., Facebook
  - Internet
  - Protein networks
  - Neural network of the human brain
  - Citation graphs
  - Business networks
  - Populations
  - And more recently the graph of logical formulas

# Community Structure in Graphs

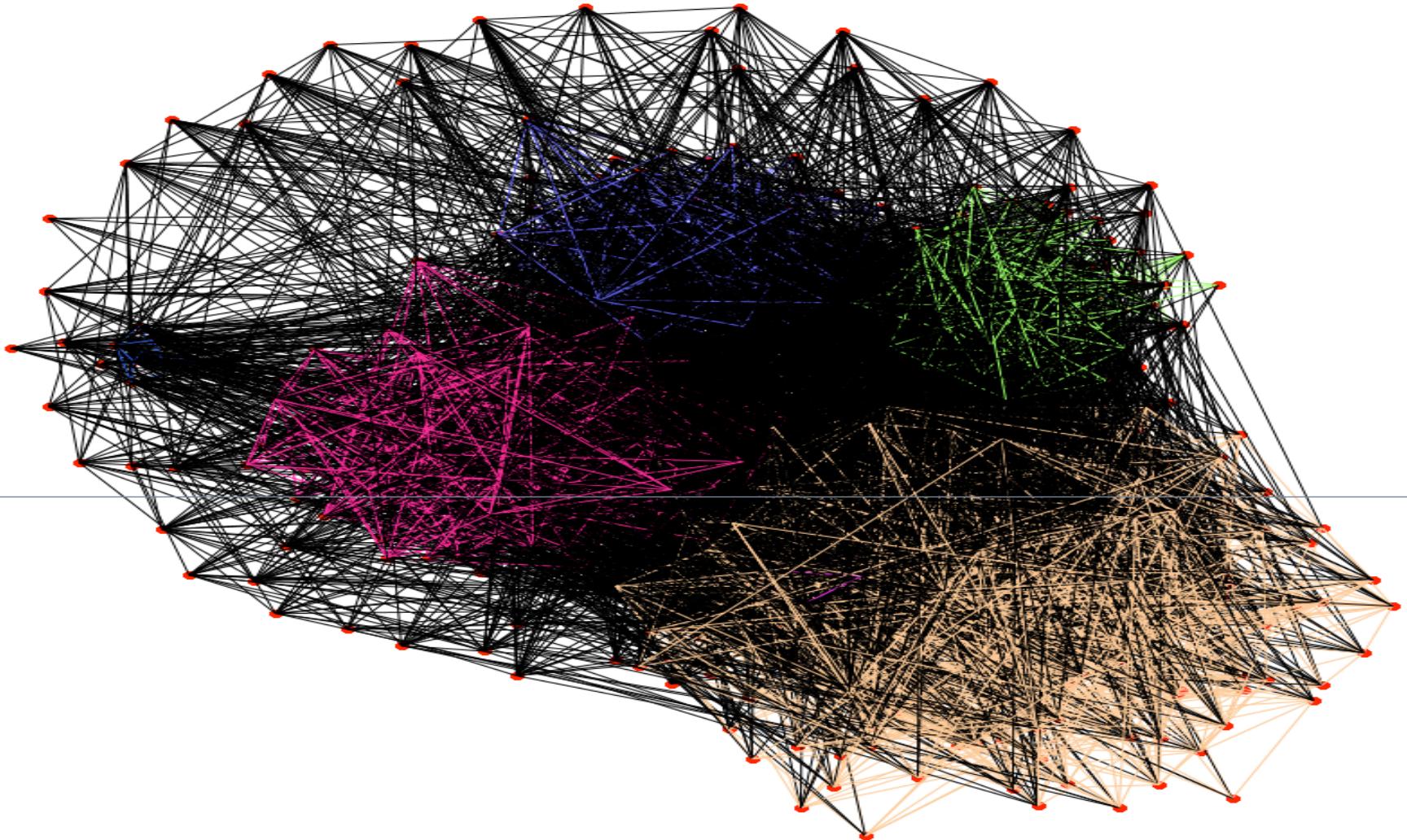
## Variable-incidence Graph of Non-random Formula



SOURCE: mrpp example from SAT 2013 competition viewed using our SATGraf tool

# Community Structure in Graphs

Variable-incidence Graph of Randomly-generated Formula



# Modularity (Q-factor) and Communities in Graphs

## Community Structure in Graphs

- How to compute community structure?
- The decision version of the Q maximization problem is NP-complete  
[Brandes et al., 2006]
- Many efficient *approximate* algorithms proposed, e.g., [CNM04] and [OL13]
- We use the above two algorithms for our experiments
- There are other methods to compute community structure based on graph centrality

# Community Structure and SAT Solver Performance

## Rest of the Talk

- Motivation and Problem Statement
- Results outline
- Preliminaries
  - Definition of communities based on Modularity (aka Q factor)
- **Experiment #1:** Strong correlation between community structure and LBD (Literal Block Distance)
- **Experiment #2:** MiniSAT finds it harder to solve randomly-generated SAT instances whose Q value lies between 0.05 and 0.13, than those that lie outside that range
- **Experiment #3:** Strong correlation between number of communities, Q and the running time of CDCL solvers
  - Hypothesis and experimental setup
  - Reasoning behind the experiments
  - Results
- Speculation on where communities come from in industrial instances
- Conclusion and future directions

# Literal Block Distance (LBD) and Communities

## Experiment #1: Hypothesis and Definitions

### Hypothesis tested

- The number of communities in a conflict clause correlates strongly with its LBD measure

### What is LBD? (Introduced first in Glucose solver [AS09])

- LBD measure  $M$  of a learnt clause  $C$  is a rank based on the number  $N$  of distinct decision levels the vars in  $C$  belong to
- The lower the value of  $N$ , the higher the rank  $M$
- LBD is a powerful measure of the *utility of a conflict clause*

# Literal Block Distance (LBD) and Communities

## Experiment #1: Hypothesis and Definitions

### Clause deletion

- Clause deletion is integral to efficiency of modern solvers
- *Sans* clause deletion, rate of conflict clause production may cause solvers to quickly run out of memory

### Which clauses to delete? LBD to the rescue

- Periodically delete conflict clauses with bad LBD rank
- As we will see, clauses with bad LBD rank are shared by many communities

# Literal Block Distance (LBD) and Communities

## Experiment #1: Intuition

The number of communities in a conflict clause

- The number of communities  $N$  in a conflict clause  $C$  is the number of distinct communities the variable in  $C$  belong to

Intuition behind the hypothesis

- High-quality conflict clauses tend to span very few communities, i.e., the number  $N$  of different communities their variables belong to is small
- High-quality conflict clauses are likely to cause more propagation per decision variable, and hence likely to have low LBD
- LBD picks out high-quality conflict clauses

# Literal Block Distance (LBD) and Communities

## Experiment #1 Setup

- Instances considered
  - 189 SAT 2013 Applications category instances out of 300
  - We were able to compute communities only for these 189
  - The rest caused memory-out
- Step 1 of the experiments
  - For each of the 189 instances in our benchmark compute
    - Community structure
    - The number of communities a learnt clause belongs to
    - LBD of every learnt clause (considered only the first 20,000 learnt clauses due to resource constraints)

# Literal Block Distance (LBD) and Communities

## Experiments Performed (#1)

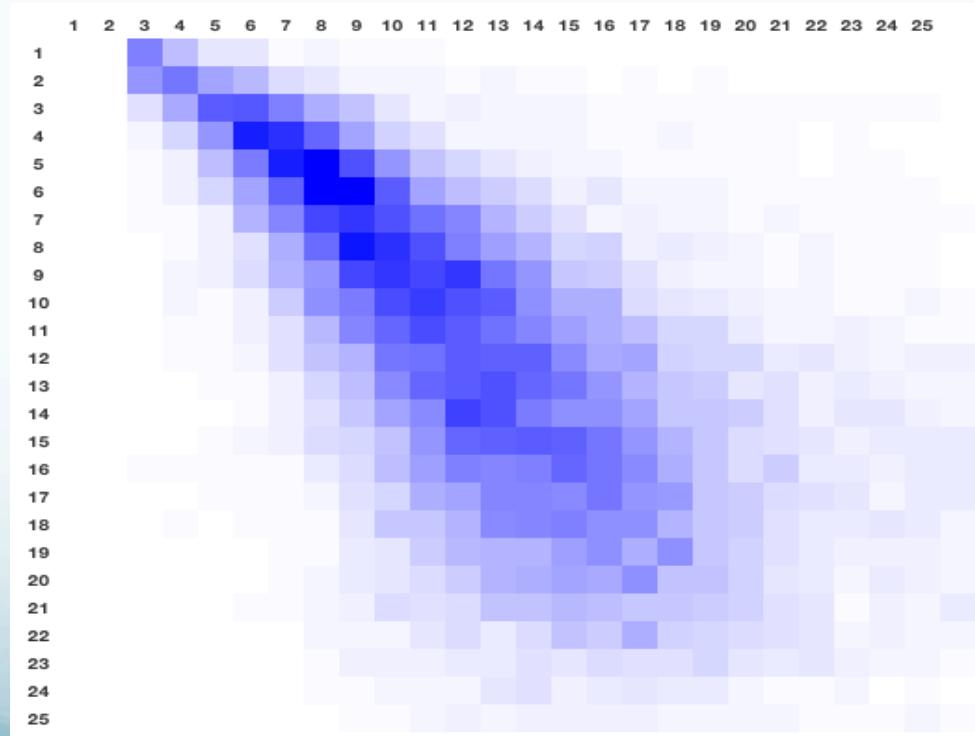
- Step 2 of the experiments
  - LBD of every learnt clause  $L$  considered was correlated with the number of communities  $L$  belongs to
  - Thousands of data points over 189 instances
  - Correlate LBD and num-of-communities using heatmaps
    - Heatmap of LBD and communities of learnt clauses
    - Otherwise difficult to correlate thousands of data points over hundreds of instances
    - One heatmap per SAT instance

# Literal Block Distance (LBD) and Communities

## Experiment #1 Results and Interpretation

### Result #1

Most industrial instances have a very strong (diagonal) relationship between LBD and communities



# Community Structure and Random Instances

## Experiment #2: Hypothesis and Definitions

### Hypothesis tested

- Is there a range of Q-factor values for randomly-generated instances that are hard for CDCL SAT solvers, irrespective of number of variables/clauses
- Are randomly-generated instances outside this range uniformly easy

# Community Structure and Solver Running Time

## Experiment #2 Setup

- Randomly generated 550,000 SAT instances for the experiment
  - Varied  $N_v$  between 500 to 2000 in increments of 100
  - Varied  $N_{cl}$  between 2000 and 10000 in increments of 1000
  - Varied target Q between 0 and 1 in increments of 0.01
  - Varied “Num of communities” between 20 and 400 in increments of 20
- Experiments using MiniSAT
  - Timeout of 900 seconds per run
  - Run solver on inputs in a random order
  - Average the running time over several runs

# Community Structure and Solver Running Time

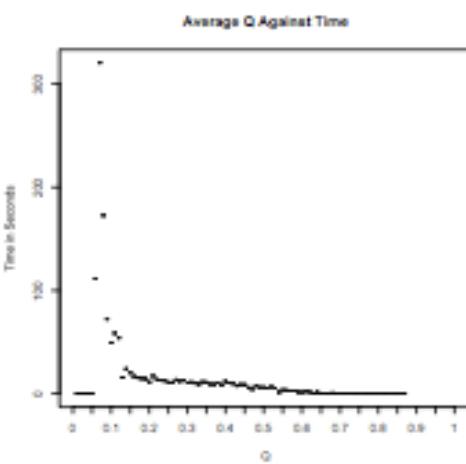
## Experiments Performed (#2)

- Plotted Q against time
- Noticed significant increase in execution time when  $0.05 \leq Q \leq 0.13$
- Also recomputed the results using a stratified sample
  - Used due to high number of instances in target range
  - Randomly sample the data taking 250 results from each 0.1 range of Q between 0 and 0.9
  - Almost same result  $0.05 \leq Q \leq 0.12$

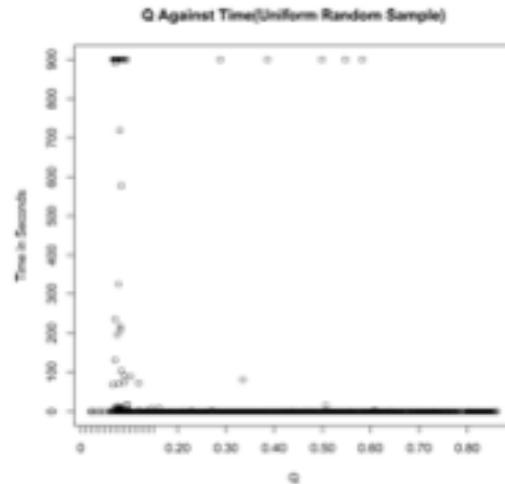
# Community Structure and Solver Running Time

## Results and Interpretation (#2)

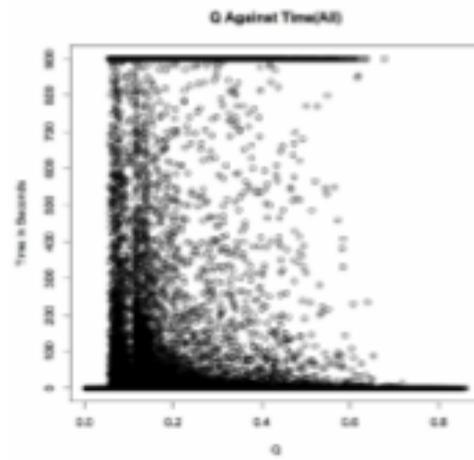
- Huge increase in running time for randomly-generated instances with  $0.05 \leq Q \leq 0.13$



(a) Average Time



(b) Stratified Sample



(c) All instances

# Community Structure and Solver Running Time

## Experiment #3: Hypothesis and Definitions

### Hypothesis tested

- Are the community structure (Q factor) and number of communities better correlated with running time of a CDCL SAT solver than traditional metrics
- Is the correlation is better for industrial instances

# Community Structure and Solver Running Time

## Experiment #3 Setup

### Instances for the experiment

- Approx. 800 instances from SAT 2013 competition. For the remaining we couldn't compute community structure due to resource constraints

Used OL algorithm to compute community structure for the 800 instances

- Much faster and scales better
- Approximates the community structure

All experimental results are for MiniSAT

- Obtained running time of solver from SAT 2013 competition website

Used statistical tool R to perform standard linear regression

# Community Structure and Solver Running Time

## Experiments Performed (#3)

- Performed linear regression on the solver running time data using statistical tool R twice
  - Once with community structure metrics
  - And once without community structure metrics
- Compared the adjusted  $R^2$  (variability) from both experiments
  - Variability measures how “far off” are the model’s prediction from given data
- R tells which of these two models has lower variability, i.e., is a better predictor

# Community Structure and Solver Running Time

## Experiment #3 Results and Interpretation

- R tells that the model with community structure metrics is a better predictor of running time than otherwise
- The model is even better if we only consider industrial instances
- Q was included in all but one significant factors with 99.9% confidence

Factor	Estimate	Std. Error	t value	Pr(>  t )	Sig
CO	-1.237e+00	3.202e-01	-3.864	0.000121	***
CL  $\odot$ Q $\odot$ QCOR	-4.226e+02	1.207e+02	-3.500	0.000492	***
CL  $\odot$ Q	-2.137e+02	6.136e+01	-3.483	0.000523	***
CL  $\odot$ Q $\odot$  CO  $\odot$ QCOR $\odot$ VCLR	-1.177e+03	3.461e+02	-3.402	0.000702	***
CL  $\odot$ Q $\odot$  CO	-6.024e+02	1.774e+02	-3.396	0.000719	***
Q $\odot$ QCOR	3.415e+02	1.023e+02	3.339	0.000881	***
Q	1.726e+02	5.200e+01	3.318	0.000947	***
Q $\odot$  CO  $\odot$ QCOR	9.451e+02	2.927e+02	3.229	0.001292	**

# Impact of Community Structure and Solver Running Time

## Scope for Improvement

- Consider different graph representations for community detection
- Try experiments on more solvers
- Can we construct a highly predictive model
- Compare community structure-based model against graph-width based models
- Community structure of SAT vs. UNSAT cases
- Try different random generation strategy and larger instances
- Hierarchical communities?
- Other solver measures like memory usage, number/rate of conflicts generated,...
- Build visualization tools to confirm/refute hypotheses
- Dynamically track how community structure of learnt clauses evolve

# Community Structure and SAT Solver Performance

## Conclusions and Future Work

- Take-home Message
  - Community structure of SAT instances strongly affect solver performance
- Result #1: Strong correlation between community structure and LBD (Literal Block Distance) in Glucose
- Result #2: Hard random instances have low Q ( $0.05 \leq Q \leq 0.13$ )
- Result #3: Number of communities and Q of SAT instances are more predictive of MiniSAT performance than other measures
- Result #4 (in progress): VSIDS correlates strongly with graph centrality with exponential smoothing average. VSIDS picks out variables central to communities. New VSIDS based on this observation with promising results

# Community Structure in Graphs

## Questions

