Bachelor's Thesis

THESIS TITLE

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Abstract

This thesis proposes ...

Contents

1 Introduction

While the hype on AI is still ongoing there are still people wondering if current AI models are fundamentally able of reasoning. Many other possible models could be better fitted to the task of reasoning. Among others Spiking neural networks present a model closer to the workings of the brain. In fact, they have been called the 3rd generation of AI models, after the 2nd generation that currently drives most successful models.

While the idea behind SNNs is quite old, they have not been as much researched, since it is much more inefficient and harder to train them. Therefore there still remain a lot of open questions about them. In this paper we shall extend on the work done in [Nguyen et al., 2025]. We will add a decaying factor to the input of the neurons and allow recursive connections between neurons in a layer.

We will roughly follow the structure of [Nguyen et al., 2025]. In the second chapter we will formally introduce discrete time leaky-integrate-and-fire SNN, d.t. LIF-SNNs. In section 3 we will give theorems about approximation of continuous functions on compact domains by d.t. LIF-SNNs. The main part will be the following section in which we will see that the number of distinct values a d.t. LIF-SNN can take on only depends on the first hidden layer and grows in particular only quadratically in time. We will further support our findings with experimental data in the last section.

2 Definitions

Our type of SNN should be thought of as a composition of an initial input layer, a number of hidden spiking layers with internal state and an affine-linear layer mapping spikes activations over time, so called spike trains, to the value of the output layer. Like motivated we are adding additional structure to the definitions [Nguyen et al., 2025]. We shall first define the input $i^{[l]}(t)$, the membrane potential before spike $p^{[l]}(t)$, the membrane potential after spike $u^{[l]}(t)$ and the spike activations $s^{[l]}(t)$ of the hidden layers:

Definition 2.1. The **input vector** $i^{[l]}(t) \in \{0,1\}^{n_l}$, the **spike vector** $s^{[l]}(t) \in \{0,1\}^{n_l}$ and the **membrane potential vector** $u^{[l]}(t)$ of a hidden layer $\lambda = (W^{[l]}, b^{[l]}, u^{[l]}(0), i^{[l]}(0), \alpha^{[l]}, \beta^{[l]}, \vartheta^{[l]})$, $l \in [L]$, are recursively defined as

$$i^{[l]}(t) := \alpha^{[l]} i^{[l]}(t-1) + W^{[l]} s^{(l)-1}(t) + V^{[l]} s^{[l]}(t-1) \tag{1}$$

$$p^{[l]}(t) := \beta^{[l]} u^{[l]}(t-1) + i^{[l]}(t) + b^{[l]} \tag{2}$$

$$s^{[l]}(t) := H(p^{[l]}(t) - \vartheta 1_{n_l}) \tag{3}$$

$$u^{[l]}(t) := p^{[l]}(t) - \vartheta s^{[l]}(t) \tag{4}$$

with $s^{[l]}(0) = 0$ and given

- **initial membrane potential**: $u^{[l]}(0) \in \mathbb{R}^{n_l}$,

- **initial input**: $i^{[l]}(0) \in \mathbb{R}^{n_l}$,

- **weight matrices**: $W^{[l]} \in \mathbb{R}^{n_l \times n_{l-1}}$, $V^{[l]} \in \mathbb{R}^{n_l \times n_l}$,

- **bias vectors**: $b^{[l]} \in \mathbb{R}^{n_l}$,

- **leaky terms**: $\alpha^{[l]}, \beta^{[l]} \in [0,1]$,

- **threshold**: $\vartheta^{[l]} \in (0, \infty)$,

where $H := \mathbb{1}_{[0,\infty)}$ is a step function and $T \in \mathbb{N}$ is the number of simulated time steps.

We further define recurrent d.t. LIF-SNN and the function the network realizes:

Definition 2.2. A **recurrent discrete-time LIF-SNN** of **depth** L with **layer-widths** $(n_0, ..., n_{L+1})$ and $T \in \mathbb{N}$ time-steps is given by

$$\Phi := ((W^{[l]}, b^{[l]}, V^{[l]}, u^{[l]}(0), i^{[l]}(0), \alpha^{[l]}, \beta^{[l]}, \vartheta^{[l]})_{l \in [L]}, T, (E, D))$$

where the **input encoder** $E : \mathbb{R}^{n_0} \to \mathbb{R}^{n_0 \times T}$ maps a vector $x \in \mathbb{R}^{n_0}$ to a corresponding first layer spike activation $s^{[0]} = E(x)$ and the **output decoder** $D : \{0,1\}^{n_L \times T} \to \mathbb{R}^{n_{L+1}}$ maps the spike activations of the last hidden layer to real values.

Definition 2.3. A recurrent discrete-time LIF-SNN ϕ **realizes** the function $R(\Phi) : \mathbb{R}^{n_0} \to \mathbb{R}^{n_{L+1}}$:

$$R(\Phi)(x) = D((s^{[L]}(t))_{t \in [T]}) \quad \text{with } s^{[0]} := E(x)$$

Definition 2.4. A recurrent discrete-time LIF-SNN employs **direct encoding** if we have

$$\forall_{t \in [T]} E(x)(t) = x$$

for the input encoder and has **membrane potential outputs** if the output decoder can be written as

$$D((s(t))_{t \in [T]}) = \sum_{t=1}^{T} a_t (W^{[L+1]} s(t) + b^{[L+1]})$$

for some $(a_t)_{t \in [T]} \in \mathbb{R}^T$, $b^{[L+1]} \in \mathbb{R}^{n_{L+1}}$ and $W^{[L+1]} \in \mathbb{R}^{n_{L+1} \times n_L}$.

We will only consider recurrent discrete-time LIF-SNN with direct encoding and membrane potential outputs. In fact, we will use "recurrent discrete-time LIF-SNN" to mean "r. d.t. LIF-SNN with direct encoding and membrane potential".

For clearer construction of our networks we will additionally define neurons:

Definition 2.5. The i-th neuron of a hidden layer $\lambda = (W^{[l]}, b^{[l]}, V^{[l]}, u^{[l]}(0), i^{[l]}(0), \alpha^{[l]}, \beta^{[l]}, \vartheta^{[l]})$ of a r. d.t. LIF-SNN is a tuple (w, b, v, u_0, i_0) with $w \in \mathbb{R}^{n_{l-1}}$, $v \in \mathbb{R}^{n_l}$ and $b, u_0, i_0 \in \mathbb{R}$, such that b, u_0, i_0 are the i-th component of $b^{[l]}$, $u^{[l]}(0)$, $i^{[l]}(0)$ respectively and w, v are the i-th row vector of $W^{[l]}$, $V^{[l]}$ respectively.

The following lemmas will be very helpful for the proofs in the following sections:

Lemma 2.1. *We define the following non-recursive formulas with explicit reference to the previous spikes. Let $1 \leq t \leq T$ and $(s_p^{[l]})_{l \in \{0...L'\}}$ such that $s_p^{[l]} : \{0,1\}^{n_l \times t_l'}$ for $l \in \{0...L'\}$. L' and $\forall_l t_l'$ have to be chosen such that the following terms are well-defined, e.g. $L' = L$ and $\forall_l t_l' = T$*

$$i^{[l]}(t; (s_p^{[l]})_l) := (\alpha^{[l]})^t i^{[l]}(0) + \sum_{k=1}^{t} (\alpha^{[l]})^{t-k} \left(W^{[l]} s_p^{[l-1]}(k) + V^{[l]} s_p^{[l]}(k-1) \right), \tag{5}$$

$$p^{[l]}(t; (s_p^{[l]})_l) := (\beta^{[l]})^t u^{[l]}(0) + \sum_{k=1}^{t} (\beta^{[l]})^{t-k} \left(i^{[l]}(k; (s_p^{[l]})_l) + b^{[l]} \right) - \vartheta \sum_{k=1}^{t-1} (\beta^{[l]})^{t-k} s_p^{[l]}(k), \tag{6}$$

$$s^{[l]}(t; (s_p^{[l]})_l) := H(p^{[l]}(t; (s_p^{[l]})_l) - \vartheta 1_{n_l}) \tag{7}$$

$$u^{[l]}(t; (s_p^{[l]})_l) := (\beta^{[l]})^t u^{[l]}(0) + \sum_{k=1}^{t} (\beta^{[l]})^{t-k} \left(i^{[l]}(k; s_p) + b^{[l]} - \vartheta s_p^{[l]}(k) \right), \tag{8}$$

We use the convention $\forall_{\alpha, \beta \in [0,1]} \alpha^0, \beta^0 = 1$. These formulas are equivalent to the previous recursive definitions, given $s_p^{[l]} = s^{[l]}$ for $l \in \{0...L'\}$.

Proof. Let the time $1 \leq t \leq T$ be non-zero. By induction we immediately get

$$i^{[l]}(t) = (\alpha^{[l]})^t i^{[l]}(0) + \sum_{i=1}^{t} (\alpha^{[l]})^{t-i} \left(W^{[l]} s^{[l-1]}(i) + V^{[l]} s^{[l]}(i-1) \right)$$

By definition of $u^{[l]}$ and $p^{[l]}$ we further obtain

$$u^{[l]}(t) = \beta^{[l]} u^{[l]}(t-1) + i^{[l]}(t) + b^{[l]} - \vartheta s^{[l]}(t)$$

from which

$$u^{[l]}(t) = (\beta^{[l]})^t u^{[l]}(0) + \sum_{i=1}^{t} (\beta^{[l]})^{t-i} \left(i^{[l]}(i) + b^{[l]} - \vartheta s^{[l]}(i) \right)$$

follows by induction. By substituting $u^{[l]}$ in $p^{[l]}$ we further obtain:

$$p^{[l]}(t) = (\beta^{[l]})^t u^{[l]}(0) + \sum_{i=1}^{t} (\beta^{[l]})^{t-i} \left(i^{[l]}(i) + b^{[l]} \right) - \vartheta \sum_{i=1}^{t-1} (\beta^{[l]})^{t-i} s^{[l]}(i)$$

$\square$

Lemma 2.2. *Let $t_0, t_\omega \in [T]$ and $j \in [n_l]$ for an $l \in [L]$ such that $t_0 \leq t_\omega$. If $\forall_{t \in \{t_0...t_\omega\}} i_j^{[l]}(t) + b_j^{[l]} \leq \vartheta$ and $u_j^{[l]}(t_0 - 1) < \vartheta$, then $\forall_{t \in \{t_0...t_\omega\}} u_j^{[l]}(t) < \vartheta$.*

If further $\beta = 1$ and $u_j^{[l]}(t_0 - 1) + \sum_{t=t_0}^{t_\omega} (i_j^{[l]}(t) + b_j^{[l]}) \geq 0$

$$\left(u_j^{[l]}(t_0 - 1) + \sum_{t=t_0}^{t_\omega} (i_j^{[l]}(t) + b_j^{[l]}) \right) - \vartheta \sum_{t=t_0}^{t_\omega} s_j^{[l]}(t) < \vartheta$$

If either $\forall_{t\in\{t_0...t_\omega\}} s_j^{[l]}(t) = 0$ *or* $\forall_{t\in\{t'...t_\omega\}} i_j^{[l]}(t) + b_j^{[l]} \geq 0$, *where* t' *is the latest time* $\leq t_\omega$ *such that* $s_j^{[l]}(t') = 0$, *then we even get*

$$0 \leq \left(u_j^{[l]}(t_0 - 1) + \sum_{t=t_0}^{t_\omega} (i_j^{[l]}(t) + b_j^{[l]}) \right) - \vartheta \sum_{t=t_0}^{t_\omega} s_j^{[l]}(t)$$

Remark 2.1. If $\vartheta = 1$, we get in particular

$$\left\lfloor u_j^{[l]}(t_0 - 1) + \sum_{t=t_0}^{t_\omega} (i_j^{[l]}(t) + b_j^{[l]}) \right\rfloor = \sum_{t=t_0}^{t_\omega} s_j^{[l]}(t)$$

Proof. We proof the first part by induction: Let $t \in \{t_0...t_\omega\}$. By definition of $p^{[l]}$, by the given assumptions and by induction hypothesis we have

$$p_j^{[l]}(t) = \beta^{[l]} u_j^{[l]}(t-1) + i_j^{[l]}(t) + b_j^{[l]} \leq u_j^{[l]}(t-1) + i_j^{[l]}(t) + b_j^{[l]} < 2\vartheta$$

By definition of $u^{[l]}$ and $s^{[l]}$, we further get $u_j^{[l]}(t) < \vartheta$.

Let us further assume $\beta = 1$ and $u_j^{[l]}(t_0 - 1) + \sum_{t=t_0}^{t_\omega} (i_j^{[l]}(t) + b_j^{[l]}) \geq 0$. We now get

$$u^{[l]}(t_\omega) = u^{[l]}(t_0 - 1) + \sum_{k=t_0}^{t_\omega} \left(i^{[l]}(k) + b^{[l]} - \vartheta s^{[l]}(k) \right),$$

due to Lemma 2.1. From which we immediately obtain

$$\left(u_j^{[l]}(t_0 - 1) + \sum_{t=t_0}^{t_\omega} (i_j^{[l]}(t) + b_j^{[l]}) \right) - \vartheta \sum_{t=t_0}^{t_\omega} s_j^{[l]}(t) < \vartheta$$

by the first part of the proof. Let us further proof the left part of the inequality: Let t' be the latest time that $s_j^{[l]}(t') = 1$ has held for $t' \in \{t_0...t_\omega\}$. We may assume that t' exists since we are otherwise immediately finished. We then get $\vartheta \leq p_j^{[l]}(t')$ and therefore

$$0 \leq u_j^{[l]}(t') = u_j^{[l]}(t_0 - 1) + \sum_{k=t_0}^{t'} \left(i_j^{[l]}(k) + b_j^{[l]} - \vartheta s_j^{[l]}(k) \right)$$

by Lemma 2.1. By assumption and choice of t' we can conclude

$$\leq u_j^{[l]}(t_0 - 1) + \sum_{k=t_0}^{t_\omega} \left(i_j^{[l]}(k) + b_j^{[l]} - \vartheta s_j^{[l]}(k) \right)$$

$\square$

Let us now take a look at some simple examples:

Example 2.1. *Let* $T, L \in \mathbb{N}$. *Then there exists a d.t. LIF-SNN with* $\forall_{t\in[T]} s^{[L]}(t) = s^{[0]}(t)$ *for any* $s^{[0]} \in \{0,1\}^{n_0 \times T}$.

We can proof this by using constant width $n_l = n$, *weights* $W^{[l]} = I_n$, *biases* $b^{[l]} = 0$, *initial membrane potential* $u^{[l]}(0) = 0$, *leaky term* $\beta^{[l]} = 0$ *and threshold* $\vartheta^{[l]} = 1$ *for all* $l \in [L]$.

We then get by definition:

$$s^{[l]}(t) = H(s^{[l-1]}(t) - 1_{n_l}) = s^{[l-1]}(t)$$
$$u^{[l]}(t) = s^{[l-1]}(t) - s^{[l]}(t) = 0$$

Example 2.2. *Let $T, L \in \mathbb{N}$. Then there exists a d.t. LIF-SNN with $\forall_{t \in T} s^{[L]}(t) = \max_{t' \in [t-1]}(s^{[0]}(t'))$ for any $s^{[0]} \in \{0,1\}^{n_0 \times T}$: In this network an output neuron switches on when the corresponding input neurons fires and does not switch off later.*

We can proof this by using constant width $n_l = n$, weights $W^{[l]} = T \cdot I_n$, biases $b^{[l]} = 0$, initial membrane potential $u^{[l]}(0) = 0$, leaky term $\beta^{[l]} = 1$ and threshold $\vartheta^{[l]} = 1$ for all $l \in [L]$.

We then get by definition:

$$s^{[l]}(t) = H(\ u^{[l]}(t-1) + T \cdot s^{[l-1]}(t) - 1_{n_l})$$
$$u^{[l]}(t) = \quad u^{[l]}(t-1) + T \cdot s^{[l-1]}(t) - s^{[l]}(t)$$

By adding over all timesteps we obtain

$$s^{[l]}(t) = H(T \sum_{i=1}^{t} s^{[l-1]}(i) - (\sum_{t=i}^{t-1} s^{[l]}(i)) + 1_{n_l}))$$
$$u^{[l]}(t) = \sum_{i=1}^{t} (T \cdot s^{[l-1]}(i) - s^{[l]}(i))$$

Since $\sum_{i=1}^{t-1} s^{[l]}(i) + 1_{n_l} \leq T$, once there is any $t_0 \in [T]$ with $s_i^{[l-1]}(t_0) = 1$ for an i, we get $\forall_{t \geq t_0} s_i^{[l]}(t) = 1$. By induction over the layers we clearly get the required property.

3 Structure of computations in r. d.t. LIF-SNN s

This section concern itself with the approximation of continuous functions by d.t. LIF-SNN.

In [Nguyen et al., 2025] the following theorem was proved:

Theorem 3.1. *Let f be a continuous function on a compact set $\Omega \subset \mathbb{R}^{n_0}$. For all $\varepsilon > 0$, there exists a d.t. LIF-SNN Φ with direct encoding, membrane potential output, $L = 2$ and $T = 1$ such that*

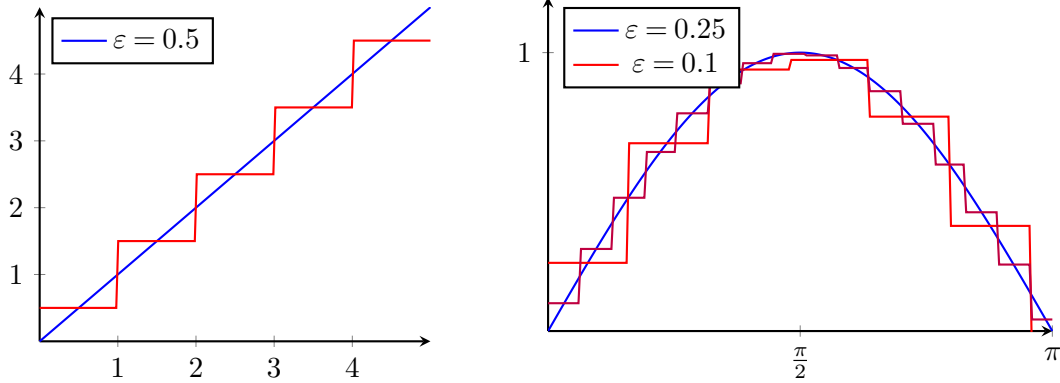$$\|(R(\Phi) - f)|_\Omega\|_\infty \leq \varepsilon$$

Moreover, if f is Γ-Lipschitz, then Φ can be chosen with width parameter $n = (n_1, n_2)$ given by

$$n_1 = \left(\max\left\{ \left\lceil \frac{\mathrm{diam}_\infty(\Omega)}{\varepsilon} \Gamma \right\rceil, 1 \right\} + 1 \right) n_0$$

$$n_2 = \max\left\{ \left\lceil \frac{\mathrm{diam}_\infty(\Omega)}{\varepsilon} \Gamma \right\rceil^{n_0}, 1 \right\}$$

where $\mathrm{diam}_\infty(\Omega) = \sup_{x,y \in \Omega} \|x - y\|_\infty$.

Proof. See [Nguyen et al., 2025]. □

(a) A d.t. LIF-SNN approximating the identity

(b) A d.t. LIF-SNN approximating a sinus wave

The proof of Theorem 3.1 works by first showing that a continuous function can be arbitrarily approximated by step functions, in particular by step functions constant on hypercubes in Ω. Then the d.t. LIF-SNN is constructed by using the first layer to cut the input space by hyperplanes along the cubes and using the second layer to represent the hypercubes.

While quite simple, this construction does not use the unique feature of (r.) d.t. LIF-SNN, the ability of neurons to accumulate state over time. It therefore needs quite a lot more neurons than actually needed for many functions with (almost) linear segments, e.g. a sinus wave, see also Figure 3.1b.

We will now show a more efficient construction for r. d.t. LIF-SNN that uses the fact that r. d.t. LIF-SNN can quite efficiently approximate linear segments. The general intuition behind is to use piece-wise linear functions to approximate continuously differentiable functions and then approximate those piece-wise linear functions by r. d.t. LIF-SNN. In the end, we will see that we can in fact approximate any continuous function like this, by using mollification to extract an continuously differentiable approximation of the function.

We will now provide an alternative version of this theorem which uses an increased latency to reduce the number of neurons:

Theorem 3.2. *Let $f \in \mathcal{C}^1(U, \mathbb{R}^m)$, with $\emptyset \neq U \subset \mathbb{R}^n$ open, be a continuously differentiable function. Let further $\Omega \subset U$ be a compact set. For all $\varepsilon, \xi, \theta > 0$, $\varepsilon = \xi + \theta$, there exists a r. d.t. LIF-SNN Φ with $L = 3$ and*

$$T = (K(\xi) + 1) T_r(\theta)$$
$$n_1 = n + 1$$
$$n_2 = K^n(\xi) \cdot (n + 1)$$

such that

$$\|(R(\Phi) - f)|_\Omega\|_{\infty, 2} \leq \varepsilon.$$

Here $T_r(\theta) := \frac{\mathrm{diam}(\Omega)}{K} \frac{\|f'\|_{\infty, 2}}{2\theta}$. The definition of $K(\xi)$ is given in [Lemma 3.3](#).

Lemma 3.1. *Let $\Omega \subset \mathbb{R}^n$ be compact. Then there exists a half-open cube C with width $\mathrm{diam}_\infty(\Omega)$ such that $\Omega \subset \overline{C}$.*

Proof of [Lemma 3.1](#).. We can first define $x := (\min_{x \in \Omega} x_i)_i$ since Ω is compact We further have $y := x + \mathrm{diam}_\infty(\Omega) \cdot \mathbb{1}_n$. We now have $C := [x, y)$. Suppose now a point $z \in \Omega \setminus \overline{C}$ exists. By definition of x, we have $x \leq z$. By definition of C we further get $z \nleq y$, and therefore $\exists_i y_i < z_i$ by definition of C. But this means $\|x - z\|_\infty > \mathrm{diam}_\infty(\Omega)$. $\qquad\square$

Lemma 3.2. *Let M be a normed vector space and N be a metric space. Let $f : M \to N$ be uniformly continuous with modulus δ. We then have for every $\varepsilon > 0$:*

$$\forall_{x, y \in M}(\|x - y\|_M \leq \delta(\varepsilon) \Rightarrow d_N(x, y) \leq \varepsilon).$$

Proof. Let $\varepsilon > 0$ with modulus $\delta(\varepsilon)$, as well as $x \in M$ be given. Since f is uniformly continuous with modulus $\delta(\varepsilon)$, we have $f(B_\delta(x)) \subset B_\varepsilon(f(x))$. Since f is in particular continuous, we get

$$f(\overline{B}_\delta(x)) = f(\overline{B_\delta(x)}) \subset \overline{f(B_\delta(x))} \subset \overline{B}_\varepsilon(f(x))$$

We need to use the fact that M is a normed vector space in the first equality to get $\overline{B}_\delta(x) = \overline{B_\delta(x)}$. $\qquad\square$

Lemma 3.3. *Let $f \in \mathcal{C}^1(U, \mathbb{R}^m)$, with $\emptyset \neq U \subset \mathbb{R}^n$ open, be a continuously differentiable function. Let further $\Omega \subset U$ be a compact set.*

For very $\varepsilon > 0$ there exists a half-open cube C with $\Omega \subset \overline{C}$ that can be composed in

$$K^n := K(\varepsilon)^n := \min_{\substack{\xi, \theta > 0 \\ \xi\theta = \varepsilon}} \left\{ \left\lceil \frac{\mathrm{diam}_\infty(\Omega)}{\frac{2}{\sqrt{n}} \min(\delta(\xi), \theta)} \right\rceil \right\}^n$$

half-open subcubes $(C_i)_{i=1..K^n}$ such that affine linear functions $g_i : C_i \to \mathbb{R}^m$ exist with $\|f - g\|_{\infty, 2} < \varepsilon$, where g is the continuous extension of $(\sum_{i=1}^m g_i \mathbb{1}_{C_i})|_C$ on $\overline{C}$. Here $\delta(\varepsilon)$ is the modulus of uniform continuity of the total derivative $dF|_\Omega$ (with regard to $\|\cdot\|_2$).

Proof of [Lemma 3.3](#).. By [Lemma 3.1](#) we have a half-open cube C with width $\mathrm{diam}_\infty(\Omega)$ and $\Omega \subset \overline{C}$. Let further $\varepsilon > 0$ be given. Since Ω is compact and $f \in \mathcal{C}^1(\Omega, \mathbb{R}^m)$, dF is uniformly continuous on Ω. Let $\delta(\varepsilon)$ be dFs modulus of uniform continuity. We will now partition C in

$$K^n := \min_{\substack{\xi, \theta > 0 \\ \xi\theta = \varepsilon}} \left\{ \left\lceil \frac{\mathrm{diam}_\infty(\Omega)}{\frac{2}{\sqrt{n}} \min(\delta(\xi), \theta)} \right\rceil \right\}^n$$

smaller half-open cubes. There are ξ, θ such that the minimum in the definition of K^n is obtained, since we take it over the set of natural numbers. The subcubes have width $w := \frac{\text{diam}_\infty(\Omega)}{K} \leq \frac{2}{\sqrt{n}} \min(\delta(\xi), \theta)$.

Let us further define $g_i : C_i \to \mathbb{R}^m$ by $g_i(x) := f(c_i) + df_{c_i}(x - c_i)$ where c_i is the center of C_i, so in particular

$$\|x - c_i\|_2^2 = \sum_{j=1}^n \left\|(x - c_i)_j e_j\right\|_2^2 \leq \frac{w^2}{4} n \leq \min(\delta(\xi), \theta)^2.$$

It suffices now to show $\left\|f|_{C_i} - g_i\right\|_\infty < \varepsilon$. Let $x \in C_i$ and $h(t) := f(t(x - c_i) + c_i)$ with

$$h'(t) = (df_{(x-c_i)t+c_i} \circ d(t \mapsto t(x - c_i) + c_i)_t)(1) = df_{(x-c_i)t+c_i}(x - c_i).$$

We obtain by the Fundamental theorem of calculus:

$$\left\|f(x) - f(c_i) - df_{c_i}(x - c_i)\right\|_2 = \left\|h(1) - h(0) - df_{c_i}(x - c_i)\right\|_2$$
$$= \left\|\int_0^1 df_{(x-c_i)t+c_i}(x - c_i)dt - df_{c_i}(x - c_i)\right\|_2$$

Due to the generalized Minkowski-Inequality we can move the norm inside the integral:

$$\left\|\int_0^1 df_{(x-c_i)t+c_i}(x - c_i)dt - df_{c_i}(x - c_i)\right\|_2 \leq \int_0^1 \left\|df_{(x-c_i)t+c_i}(x - c_i) - df_{c_i}(x - c_i)\right\|_2 dt$$
$$= \int_0^1 \left\|(df_{(x-c_i)t+c_i} - df_{c_i})(x - c_i)\right\|_2 dt$$
$$\leq \int_0^1 \left\|df_{(x-c_i)t+c_i} - df_{c_i}\right\| \|x - c_i\|_2 dt$$
$$\leq \int_0^1 \xi \|x - c_i\|_2 dt$$
$$= \xi \|x - c_i\|_2$$
$$\leq \xi\theta$$
$$= \varepsilon$$

In the fourth step we use $\left\|df_{(x-c_i)t+c_i} - df_{c_i}\right\| \leq \xi$, which holds due to $\forall_{t \in [0,1]} (x - c_i)t + c_i \in C_i$ and Lemma 3.2. $\qquad \square$

Proof of Theorem 3.2. Let a continuously differentiable function $f \in \mathcal{C}^1(\Omega, \mathbb{R}^m)$ on a compact set $\Omega \subset \mathbb{R}^n$ be given. Let there further be $\varepsilon, \xi, \theta > 0$ with $\varepsilon = \xi + \theta$. By Lemma 3.3 we have a half-open cube $C = [x^C, y^C)$ with $\Omega \subset C$ with composition K^n half-open subcubes $(C_i)_{i=1..K^n}$ and linear functions $g_i : C_i \to \mathbb{R}^m$, such that $\|f - g|_\Omega\|_\infty < \xi$ for a $g := \sum_{i=1}^m g_i \mathbb{1}_{C_i}$.

We will now define a r. d.t. LIF-SNN Φ with direct input encoding and membrane-potential outputs such that $\|R(\Phi)|_\Omega - g\|_\infty < \theta$.

Before anything else we shall set the following basic parameters $i^{[l]}(0) = 0$, $\alpha^{[l]} = 0$ and $\beta^{[l]} = \vartheta^{[l]} = 1$ for all layers. We obtain the simplified equations:

$$p^{[l]}(t) = u^{[l]}(t-1) + W^{[l]} s^{[l-1]}(t) + V^{[l]} s^{[l]}(t-1) + b^{[l]}$$
$$s^{[l]}(t) = H(p^{[l]}(t) - 1_{n_l})$$
$$u^{[l]}(t) = p^{[l]}(t) - s^{[l]}(t)$$

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and in particular by Lemma 2.1

$$i^{[l]}(t) := W^{[l]}s^{[l-1]}(t) + V^{[l]}s^{[l]}(t-1)$$

$$p^{[l]}(t) = u^{[l]}(0) + \sum_{k=1}^{t}\left(W^{[l]}s^{[l-1]}(k) + V^{[l]}s^{[l]}(k-1) + b^{[l]}\right) - \sum_{k=1}^{t-1}s^{[l]}(k).$$

The intuitive idea for the construction of the network is the following: In the first layer we have n neurons which capture the position of the input vector regarding C in their respective dimension and one last neuron that acts as an "alarm clock" that shuts down the other neurons of the layer after …

In the second layer we have $n+1$-neurons for each affine linear region C_i. Each of the first n neurons encodes the linear part of a component of the input vector. They are connected to the corresponding neuron of the first layer and the last neuron of their group. That last, additional neuron acts as an activator to the region. It get's enabled by getting a spike of all other neurons of the group and by getting a spike from the clock neuron from the first layer. The other n neurons of the group disable themselves after their first spike.
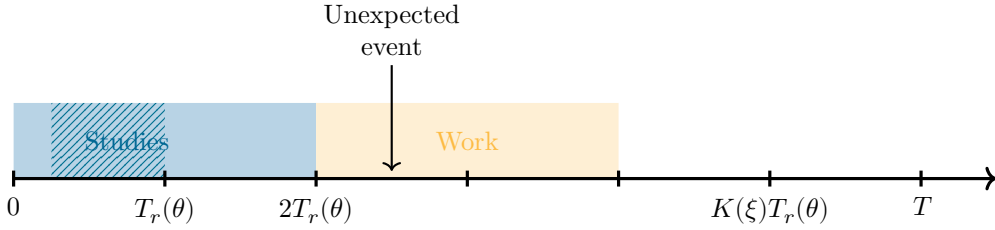

Figure 3.2: Timeline of first-layer neuron

We define the i-th neuron of the n neurons of the first layer by parameters

$$w = \frac{1}{y_i^C - x_i^C}e_i, \quad b = -\frac{x_i^C}{y_i^C - x_i^C}, \quad v = -e_{n+1}, \quad u_0 = 0, \quad i_0 = 0.$$

The last neuron of the first layer, with index $n+1$, is defined by:

$$w = 0, \quad b = \frac{1}{K(\xi)T_r(\theta)}, \quad v = e_{n+1}, \quad u_0 = 0, \quad i_0 = 0.$$

We therefore get:

$$p_{n+1}^{[1]}(t) = \frac{t}{K(\xi)T_r(\theta)} + \sum_{k=1}^{t}s_{n+1}^{[1]}(k-1) - \sum_{k=1}^{t-1}s_{n+1}^{[1]}(k) = \frac{t}{K(\xi)T_r(\theta)}$$

So $s_{n+1}^{[1]}(t) = 1 \Leftrightarrow t \geq K(\xi)T_r(\theta)$. We further get for $i \in [n]$:

$$i_i^{[1]}(t) = \frac{1}{y_i^C - x_i^C}s_i^{[0]}(t) - s_{n+1}^{[1]}(t-1) = \frac{1}{y_i^C - x_i^C}s_i^{[0]}(t) - s_{n+1}^{[1]}(t-1)$$

So we can use Lemma 2.2 for $x_i^C \leq s_i^{[0]}(t) < y_i$ and $t \leq K(\xi)T_r(\theta)$ to obtain

$$\left\lfloor \frac{K(\xi)T_r(\theta)}{y_i^C - x_i^C}(s_i^{[0]}(1) - x_i^C) \right\rfloor = \left\lfloor u_i^{[1]}(0) + \sum_{t=1}^{K(\xi)T_r(\theta)}(i_i^{[1]}(t) + b_i^{[1]}) \right\rfloor = \sum_{t=1}^{K(\xi)T_r(\theta)}s_i^{[1]}(t) \quad (9)$$

The equation clearly also holds for $s_i^{[0]}(t) = y_i$.

Let us know construct the second layer in the following way: For each of the $K^n(\xi)$ subcubes in C we define $n+1$ neurons like so: Let $C_j = [x^{C_j}, y^{C_j})$ be one such subcube with position $q \in \{0, ..., K(\xi)-1\}^{n_1}$ in C. We will write $\iota_j(i) := j(n+1)+i$ to index the first n neurons in the layer and $\omega_j := (j+1)(n+1)$ to index the last neuron of each group.

The i-th neuron of the first n neurons, with index $\iota_j(i)$ in the second layer, has the parameters

$$w = e_i, \quad b = 0,$$

$$v = K(\xi)T_r(\theta)\left(-e_{\iota_j(i)} + e_{\omega_j} - \sum_{\substack{q' \in \{0,...,K(\xi)-1\}^{n_1} \\ q' < q}} e_{\omega_j(q)}\right),$$

$$u_0 = -q_i T_r(\theta) + 1, \quad i_0 = 0.$$

where $j(q)$ is the index of the subcube at position q. We get

$$p^{[2]}_{\iota_j(i)}(t) = -q_i T_r(\theta) + 1 + \sum_{k=1}^{t} s^{[1]}_i(k) - (K(\xi)T_r(\theta)+1)\sum_{k=1}^{t-1} s^{[2]}_{\iota_j(i)}(k).$$

for $t \le K(\xi)T_r(\theta)$, since as we will later show, $s^{[1]}_{\omega_j}(t) = 0$ for all $j \in [K^n]$ and $t \le K(\xi)T_r(\theta)$. We therefore have $s^{[2]}_{\iota_j(i)}(t) = 1$ for a $t \le K(\xi)T_r(\theta)$ exactly if $\sum_{k=1}^{K(\xi)T_r(\theta)} s^{[1]}_i(k) \ge q_i T_r(\theta)$ which is equivalent to by (9)

We further have

The final neuron of the group, with index ω_j in its layer, has the parameters

$$w = \frac{1}{n+1}e_{n+1}, \quad b = 0, \quad v = e_{\omega_j} + \frac{1}{n+1}\sum_{i=1}^{n} e_{\iota_j(i)},$$

$$u_0 = 0, \quad i_0 = 0.$$

We get

$$p^{[2]}_{\omega_j}(t) = u^{[2]}_{\omega_j}(t-1) + s^{[2]}_{\omega_j}(t-1) + \frac{1}{n+1}\left(s^{[1]}_{n+1}(t) + \sum_{i=1}^{n} s^{[2]}_{\iota_j(i)}(t-1)\right)$$

We further define the parameters of the output decoder by $a_t = 0$, for $t \le K(\xi)T_r(\theta)$ and otherwise $a_t = 1$. We further set $b^{[L+1]} = 0$ and $(W^{[L+1]})_{l,\iota_j(i)} = g(x^{C_j} + \frac{1}{T_\theta(\theta)}e_i) - g(x^{C_j})$ for $l \in [m]$, $j \in K^n(\xi)$ and $i \in [n]$, as well as $(W^{[L+1]})_{l,\omega_j} = \frac{1}{T_\theta(\theta)}g(x^{C_j})$ for $l \in [m]$ and $j \in K^n(\xi)$.

Let now $s^{[0]}(t) = x \in C$. We will proof $\|R(\Phi)(x) - g(x)\|_{\infty,2} \le \theta$.

Let now $s^{[0]}(t) = x \in \overline{C} \setminus C$.

$\square$

Sadly the size of the network in this construction is not always smaller than the one from Theorem 3.1. A concrete counter example is a sinus wave with high frequency and small amplitude, like $f(x) := \frac{\sin(nx)}{n}$ with $n \in \mathbb{N}$ on $\Omega = [0, 2\pi]$. Since $f'(x) = \cos(nx)$ and $\|f'|_\Omega\|_\infty = 1$, we can choose $L = 1$ as a Lipschitz constant for f. At the same time, since $f''(x) = n\sin(nx)$ and $\|f''|_\Omega\|_\infty = n$, the biggest possible modulus of uniform continuity on Ω we can give for f' is $\delta(\varepsilon) := \frac{\varepsilon}{n}$. So we get

$$\max_{\substack{\xi,\theta>0 \\ \xi\theta=\varepsilon}} \min(\delta(\xi),\theta) = \max_{\xi>0} \min(\frac{\xi}{n}, \frac{\varepsilon}{\xi}) = \sqrt{\frac{\varepsilon}{n}}$$

So we get that $K(\varepsilon)$ has a value of $\lfloor \frac{\pi\sqrt{n}}{\sqrt{\varepsilon}} \rfloor$ by definition. We therefore get for the layer sizes:

Theorem 3.1	*Theorem* 3.2
$n_1 = \lceil \frac{2\pi}{\varepsilon} \rceil + 1$	$n_1 = 2$
$n_2 = \lceil \frac{2\pi}{\varepsilon} \rceil$	$n_2 = 2 \lfloor \frac{\pi\sqrt{n}}{\sqrt{\varepsilon}} \rfloor$

While the first and second layer are clearly far smaller than the first layer of the other construction, especially for small ε, the third layer of our construction is arbitrarily bad for $n \to \infty$ compared to the last layer of the other construction.

But there is one big difference. Since the size of last layer grows only proportionally to $\frac{1}{\sqrt{\varepsilon}}$ and not proportionally to $\frac{1}{\varepsilon}$, it is arbitrarily smaller than the other constructions last layer for $\varepsilon \to 0$ and any $n \in \mathbb{N}$.

We will generalize this observation with the following theorem:

Theorem 3.3. *Let $f \in \mathcal{C}^1(U, \mathbb{R})$, with $\emptyset \neq U \subset \mathbb{R}^n$ open, be a continuously differentiable function. Let further $\Omega \subset U$ be a compact set.*

Let $(\xi_\varepsilon)_{\varepsilon > 0}$ be given with $\forall_{\varepsilon > 0} \xi_\varepsilon < \varepsilon$, such that

$$\lim_{\varepsilon \to 0} \left(\frac{K(\xi_\varepsilon)}{\left\lceil \frac{\operatorname{diam}_\infty(\Omega)}{\varepsilon} \Gamma \right\rceil} \right) = 0$$

where Γ is the Lipschitz-constant of f.

Proof. It suffices to show there exist $(\xi_\xi)_{\xi > 0}$ and $(\theta_\xi)_{\xi > 0}$ with $\forall_{\xi > 0} \xi_\xi \theta_\xi = 1$ and

$$\lim_{\varepsilon \to 0} \frac{\varepsilon}{\min(\delta(\xi_\xi), \theta_\xi)} = 0$$

$\square$

4 Complexity of input partitions

In the following we will analyze what shape the graph of a r. d.t. LIF-SNN has. Since the output of the following layers only depends on the spike trains of the first hidden layer, those layers are only able to merge different output regions of the first hidden layer. We will therefore only study the output landscape of the first hidden layer.

We can therefore simplify the notation in this section by writing $W, b, V, \beta, \vartheta, u, s$ for $W^{[1]}, b^{[1]}, V^{[1]}, \beta^{[1]}, \vartheta^{[1]}, u^{[1]}, s^{[1]}$ respectively. Since we are using direct encoding, we will write x for $s^{[0]} \in \mathbb{R}^{n_0}$:

$$i(t) = \alpha i(t-1) + Wx + Vs(t-1) \tag{10}$$
$$p(t) = \beta u(t-1) + i(t) + b \tag{11}$$
$$s(t) = H(p(t) - \vartheta 1_n) \tag{12}$$
$$u(t) = p(t) - \vartheta s(t) \tag{13}$$

By using the simplified notation, we obtain the following explicitely formulated definitions from Lemma 2.1 for $x \in \mathbb{R}^{n_0}$ and $s : \{0,1\}^{n_1 \times t'}$ for t' chosen such that the following equations are well-defined:

$$i(t; x; s_p) = \alpha^t i(0) + \sum_{k=1}^{t} \alpha^{t-k} \left(Wx + Vs_p(k-1) \right), \tag{14}$$

$$u(t; x; s_p) = \beta^t u(0) + \sum_{k=1}^{t} \beta^{t-k} \left(i(k; x; s_p) + b - \vartheta s_p(k) \right), \tag{15}$$

$$p(t; x; s_p) = \beta^t u(0) + \sum_{k=1}^{t} \beta^{t-k} \left(i(k; x; s_p) + b \right) - \vartheta \sum_{k=1}^{t-1} \beta^{t-k} s_p(k), \tag{16}$$

$$s(t; x; s_p) = H(p(t; x; s_p) - \vartheta 1_{n_1}) \tag{17}$$

Remark 4.1. We will sometimes just write $i(t; x)$ instead of $i(t; x; s_p)$, with which we imply to use s (at x) for s_p. The same holds for u, p, s.

It is furthermore quite obvious that i grows linear in x, given fixed s_p. The same holds therefore for u and p. Further, the interdependence of the different components in $i/u/p/s$ is gone using a constant s_p. We can therefore split e.g. p into functions $p_{i;s_p}$ such that $p(t, x) = p_{1;s_p}(t; x_1) \times ... \times p_{n_1;s_p}(t; x_{n_1})$.

The functions $i_{i;s_p}, u_{i;s_p}, p_{i;s_p}$ are not only linear in x, but also grow monotonically in x since $\alpha, \beta \geq 0$. If $t \geq 1$, then $i_{i;s_p}, u_{i;s_p}, p_{i;s_p}$ grow even strictly monotonically in x.

So we in particular have that p to is growing monotonically (regarding by-component ordering) on fixed s_p. Since H is a monotonically growing function, s is also growing monotonically with fixed s_p.

We will examine the graph by investigating the shape and location of the constant regions of a r. d.t. LIF-SNN :

Definition 4.1. The set of constant regions of a r. d.t. LIF-SNN Φ is defined as the partition

$$C_\Phi := \{ R(\Phi)^{-1}(\{y\}) \mid y \in \text{im}(R(\Phi)) \}$$

of $\mathbb{R}^{n_0}$. A constant region with spike train $s' \in \{0,1\}^{n_1 \times T}$, of the first layer of a r. d.t. LIF-SNN Φ is defined as

$$C_{s'} := \{ x \in \mathbb{R}^{n_0} \mid \forall_{t \in [T]} s(t; x) = s'(t) \}$$

We further notate the set of these regions by $C_{\Phi,1} := \{ C_{s'} \mid s' \in \{0,1\}^{n_1 \times T}, C_{s'} \neq \emptyset \}$.

Proposition 4.1. *The constant regions of the first layer of a r. d.t. LIF-SNN Φ with $W = I_{n_1}$ are half-open cuboids.*

Lemma 4.1. $\bigcap_{j \in J} \prod_{i \in [n]} M_{i,j} = \prod_{i \in [n]} \bigcap_{j \in J} M_{i,j}$ *for an index set J and sets $(M_{i,j})_{i \in [n], j \in J}$.*

Proof of Lemma 4.1. For every $x = (x_i)_{i \in [n]} \in M := \prod_{i \in [n]} \bigcup_{j \in J} M_{i,j}$ holds:

$$x \in \bigcap_{j \in J} \prod_{i \in [n]} M_{i,j} \Leftrightarrow \forall_{j \in J} x \in \prod_{i \in [n]} M_{i,j} \Leftrightarrow \forall_{j \in J} \forall_{i \in [n]} x_i \in M_{i,j}$$

On the other hand, we have

$$x \in \prod_{i \in [n]} \bigcap_{j \in I} M_{i,j} \Leftrightarrow \forall_{i \in [n]} x_i \bigcap_{j \in I} M_{i,j} \Leftrightarrow \forall_{i \in [n]} \forall_{j \in J} x_i \in M_{i,j}$$

In total we get

$$\bigcap_{j \in J} \prod_{i \in [n]} M_{i,j} = M \cap \bigcap_{j \in J} \prod_{i \in [n]} M_{i,j} = M \cap \prod_{i \in [n]} \bigcap_{j \in I} M_{i,j} = \prod_{i \in [n]} \bigcap_{j \in I} M_{i,j}$$

$\square$

Lemma 4.2. *The intersection $C_1 \cap C_2$ of half-open cuboids $C_1, C_2 \subset \mathbb{R}^n$ is a half-open cuboid.*

Proof of Lemma 4.2. By definition of a half-open cuboid, we have half-open intervals $[c_{i,j}, d_{i,j})$ with $j \in [n]$ such that $C_i := \prod_{j \in [n]} [c_{i,j}, d_{i,j})$ for $i \in [2]$. We therefore have

$$C_1 \cap C_2 = \prod_{j \in [n]} ([c_{1,j}, d_{1,j}) \cap [c_{2,j}, d_{2,j})) = \prod_{j \in [n]} ([\max(c_{1,j}, c_{2,j}), \min(d_{1,j}, d_{2,j})))$$

by Lemma 4.1

$\square$

Lemma 4.3. *Let $[c, d)$ be a half-open interval and $\varphi : \mathbb{R} \to \mathbb{R}$ be an affine linear function. Then $\varphi^{-1}([c, d))$ is a half-open interval.*

Proof. If φ is constant, then $\varphi^{-1}([c, d))$ is either $\mathbb{R} = [-\infty, \infty)$ or $\emptyset = [0, 0)$. If φ is not constant, then $\varphi = ax + b$ with $a > 0$ and $\varphi^{-1}([c, d)) = [\frac{c-b}{a}, \frac{d-b}{a})$, if $c, d \in \mathbb{R}$; if $c = -\infty$ or $d = \infty$ the lower limit or upper limit of $\varphi^{-1}([c, d))$ is $-\infty$ or ∞ respectively. $\square$

Proof of Proposition 4.1. Let $C \in C_{\Phi,1}$ be such a region and s_C the corresponding spike train. We then get

$$C = \bigcap_{i \in [n_1], t \in [T]} \{x \mid (s_C)_i(t) = s_i(t; x)\}$$

$$= \bigcap_{i \in [n_1], t \in [T]} \{x \mid (s_C)_i(t) = s_i(t; x; s_C)\}$$

$$= \left(\bigcap_{\substack{i \in [n_1], t \in [T] \\ (s_C)_i(t) = 0}} \{x \mid p_i(t; x; s_C) < \vartheta\} \right) \cap \left(\bigcap_{\substack{i \in [n_1], t \in [T] \\ (s_C)_i(t) = 1}} \{x \mid p_i(t; x; s_C) \geq \vartheta\} \right)$$

So C is an intersection of pre-images of half-open intervals regarding functions $p_i(t; \cdot; s_C)$. Since $p_i(t; x; s_C) = p_{i;s_C}(t; \pi_i(x); s_C)$ for all $x \in \mathbb{R}^{n_0}$, and since $p_{i;s_C}$ is linear and growing monotonically, we can apply Lemma 4.3. A pre-image of a half-open interval regarding $p_i(t; \cdot; s_C)$ is therefore a pre-image of a half-open interval regarding π_i, so a half-open cuboid.

We can now apply Lemma 4.2 and obtain that C is indeed a half-open cuboid.

$\square$

We will add an ordering to spike trains based on lexicographical ordering. We define $s' \leq_l s''$ for $s', s'' \in \{0,1\}^{n_1 \times T}$ to mean that either $s' = s''$ or $s'(t) < s''(t)$ holds for the minimal time t such that the spike trains have different values.

Lemma 4.4. *Let us regard a r. d.t. LIF-SNN Φ with $W = I_{n_1}$. Let further $x, y \in \mathbb{R}^{n_0}$. We then have $x \leq y \Rightarrow s(\cdot; x) \leq s(\cdot; y)$.*

Proof. Let $x \leq y$ and $s(\cdot; x) \neq s(\cdot; y)$. We then have a minimal $t \in [T]$ such $s(t; x) \neq s(t; y)$. Now due to $\forall_{t' < t} s(t'; x) = s(t'; y)$ and Remark 4.1 we have $s(t; x) \leq s(t; y)$ and therefore $s(t; x) < s(t; y)$. $\qquad\square$

Lemma 4.5. *Let $s' \in \{0,1\}^{n_1 \times T}$ such that $\forall_{i \in [n_1]} \forall_{t, t' \in [T]} s'_i(t) = s'_i(t')$. Let further $C_{s'} = [x^{s'}, y^{s'})$ be the corresponding region. Then $C_{s'}$ is non-empty, i.e. $\forall_{i \in [n_1]} x_i^{s'} < y_i^{s'}$, and $x_i^{s'} = -\infty$ exactly if $\forall_{t \in [T]} s'_i(t) = 0$ and $y_i^{s'} = \infty$ exactly if $\forall_{t \in [T]} s'_i(t) = 1$ for all $i \in [n_1]$.*
Furthermore, $C_{s'}$ has exactly one finite vertex.

Proof. By Remark 4.1 $p_{i; s_p}(t; x)$ is a non-constant linear function for $t \geq 1$ and fixed s_p. We have therefore $\exists_z s_{i; s_p}(t; (-\infty, z_i)) = 0$ and $\exists_z s_{i; s_p}(t; [z_i, \infty)) = 1$ for any fixed $t \in [T]$ and s_p. Since there are only finitely many time-steps $t \in [T]$ and spike trains s_p and, we get $s_{i; s_p}(t; (-\infty, z_i)) = 0$ or $s_{i; s_p}(t; [z_i, \infty)) = 1$ respectively for any s_p and $t \in [T]$, as well as any z with z_i small or large enough.

By choosing the component z_i correctly small or large enough we therefore get a $z \in C_{s'}$. We further get for any $i \in [n_1]$ that if $\forall_t s'_i(t) = 0$, then by construction $s_i(t; (-\infty, z_i)) = 1$, so $\forall_{\delta > 0} z - \delta e_i \in C_{s'}$ and therefore $x_i = -\infty$. If on the other hand $\forall_t s'_i(t) = 1$, then $s_i(t; [z_i, \infty)) = 0$, so $\forall_{\delta \geq 0} z + \delta e_i \in C_{s'}$ and therefore $y_i = \infty$.

Since $T \geq 1$ by definition, we have either $x_i^{s'} = -\infty$ or $y_i^{s'} = \infty$ by definition of s' for each $i \in [n_1]$. By choosing the only finite one from $\{x_i^{s'}, y_i^{s'}\}$ for each $i \in [n_1]$ we get the only finite vertex of $C_{s'}$. $\qquad\square$

Lemma 4.6. *If $q \in \mathbb{R}^{n_1}$ is a (finite) vertex of a region $C_{s'} = [x^{s'}, y^{s'}) \in C_{\Phi, 1}$, $s' \in \{0,1\}^{n_1 \times T}$, then $\forall_{i \in [n_1]} \forall_{\delta > 0} s_i(\cdot; q) \neq s_i(\cdot; q - \delta(\partial_q)_i e_i)$. Here $\partial_q := 2H(x^{s'} + y^{s'} - 2q) - 1$ points is a vector pointing from q roughly in the direction of the opposite vertex of the cube $C_{s'}$ (We allow $x^{s'}, y^{s'}$ to have $-\infty$ or ∞ as components here).*

Proof. Proof by contradiction: If the statement is wrong, than there exists an $i \in [n_1]$, such that for every $\delta > 0$ with $s_i(\cdot; q) = s_i(\cdot; q - \delta(\partial_q)_i e_i)$. Since the other components of $p(\cdot; x)$ as well as $s(\cdot; x)$ are only affected by the value of x_i through $s_i(\cdot; x)$ we further get $s(\cdot; q) = s(\cdot; q - \delta(\partial_q)_i e_i)$. So $q - \delta(\partial_q)_i e_i \in C_{s'}$.

On the other hand $x_i^{s'} \leq q_i - \delta(\partial_q)_i < y_i^{s'}$ cannot be true: Assume that it is. Since q is a (finite) vertex of $[x^{s'}, y^{s'})$, we have $q_i \in \{x_i^{s'}, y_i^{s'}\}$. Case distinction: Suppose $q_i = x_i^{s'}$. Then $(\partial_q)_i = 1$ and $x_i^{s'} \leq q_i - \delta(\partial_q)_i$ is wrong. Suppose on the other hand $q_i = y_i^{s'}$. Then $(\partial_q)_i = -1$ and $q_i - \delta(\partial_q)_i < y_i^{s'}$ is wrong. $\qquad\square$

Lemma 4.7. *Let $M \subset \mathbb{R}^n$ be convex. Then the closure $\overline{M}$ is convex as well.*

Proof. Let $x, y \in \mathrm{M}$ and $t \in [0, 1]$. We will proof $z := x + (y - x)t \in \mathrm{M}$. By definition of the closure have sequences $(x_n)_{n \in \mathbb{N}}, (y_n)_{n \in \mathbb{N}} \subset M$ with $x_n \to x$, $y_n \to y$ for $n \to \infty$. We get a corresponding sequence $z_n := x_n + (y_n - x_n)t$ with $z_n \to z$, since

$$
\begin{aligned}
|z - z_n| &= |(x - x_n)(1 - t) + (y - y_n)t| \\
&\leq (1 - t)|x - x_n| + t|y - y_n| \\
&\to 0
\end{aligned}
$$

for $n \to \infty$. $\qquad\square$

Lemma 4.8. *Let $M \subset \mathbb{R}^n$ and $x \in \mathbb{R}^m \setminus \overline{\mathrm{conv}(M)}$. Then there exists an affine hyperplane $H \subset \mathbb{R}^m$ with $x \in H$ and $H \cap \mathrm{conv}(M) = \emptyset$.*

Proof. Let now $\delta = \mathrm{dist}(x, \mathrm{conv}(M)) := \inf_{y \in \mathrm{conv}(M)}(\|x - y\|_2)$. We clearly have $\delta > 0$, since otherwise $x \in \mathrm{conv(M)}$. Let $(y_n)_{n \in \mathbb{N}}$ be a sequence with $\|x, y_n\|_2 \to \delta$ for $n \to \infty$. Since $\mathbb{R}^n$ is complete we can assume w.l.o.g. that $(y_n)_n$ converges to $y \in \mathbb{R}^n$. We now choose $H := x + \{z \in \mathbb{R}^n \mid \langle z, x - y \rangle = 0\}$. This is indeed a hyperplane, since $\|x - y\|_2 = \delta > 0$ and therefore $x \neq y$.

Suppose there is a $z \in H \cap \mathrm{conv}(M)$. We then have $y' := y + (z - y)\frac{\langle x - y, z - y\rangle}{\|z - y\|_2^2}$. By using $\|\cdot\|_2^2 = \langle \cdot, \cdot \rangle$ we get

$$\|x - y'\|_2^2 = \left\| x - y - (z - y)\frac{\langle x - y, z - y\rangle}{\|z - y\|_2^2} \right\|_2^2 = \|x - y\|_2^2 - \left\| (z - y)\frac{\langle x - y, z - y\rangle}{\|z - y\|_2^2} \right\|_2^2$$

since $x - y' \perp y' - y$:

$$\langle x - y', (z - y)\frac{\langle x - y, z - y\rangle}{\|z - y\|_2^2} \rangle$$
$$= \frac{\langle x - y, z - y\rangle}{\|z - y\|_2^2}\langle x - y - (z - y)\frac{\langle x - y, z - y\rangle}{\|z - y\|_2^2}, z - y\rangle$$
$$= \frac{\langle x - y, z - y\rangle}{\|z - y\|_2^2}\langle x - y, z - y\rangle - \frac{\langle x - y, z - y\rangle^2}{\|z - y\|_2^4}\langle z - y, z - y\rangle$$
$$= 0$$

So $\|x - y'\|_2 < \|x - y\|_2 = \delta$, since $x \neq y$ and $z \neq y$ ($y \notin H$). We further have $\mathrm{dist}(y', \mathrm{conv}(M)) = 0$:

First notice that $0 \leq \frac{\langle x - y, z - y\rangle}{\|z - y\|_2^2} < 1$. Since $z \in H$, we have $\langle z - x, x - y\rangle = 0$ and therefore $\|x - y\|_2^2 + \|z - x\|_2^2 = \|z - y\|_2^2$, so we get

$$\langle x - y, z - y\rangle \leq \|z - y\|\|x - y\| < \|z - y\|^2$$

by the Cauchy-Schwarz-Inequality and $z \neq x$. Notice that we can proof $\frac{\langle z - x, z - y\rangle}{\|z - y\|_2^2} < 1$ in the same way. Since we have $\frac{\langle x - y, z - y\rangle}{\|z - y\|_2^2} + \frac{\langle z - x, z - y\rangle}{\|z - y\|_2^2} = 1$ we get $0 \leq \frac{\langle x - y, z - y\rangle}{\|z - y\|_2^2}$.

So $y' \in [y, z] \subset \overline{\mathrm{conv}(M)}$ by Lemma 4.7.

So there are $y'' \in \mathrm{conv}(M)$ with $\|y'' - y'\|_2 < \delta - \|x - y'\|_2$ and we get $\|x - y''\| < \delta$. This contradicts the definition of δ. $\square$

Lemma 4.9. *If $H \subset \mathbb{R}^n$ is an affine hyperplane, then there exists a dimension $i \in [n]$ such that $\pi_{[n]\setminus\{i\}}(H) = \mathbb{R}^{n-1}$.*

Proof. Suppose this is not the case. Let $x \in H$. Then $0 \neq \ker(\pi_{[n]\setminus\{i\}}|_{H-x}) \subset \ker(\pi_{[n]\setminus\{i\}}|_{\mathbb{R}^n}) \cong \mathbb{R}$, and therefore $e_i \in H - x$. But this contradicts $H - x$ having dimension $n - 1$. $\square$

Proposition 4.2. *Let P be the set of all finite vertices of regions $C_{s'}$, $s' \in \{0, 1\}^{n_1 \times T}$, with $\forall_{i \in [n_1]} \forall_{t, t' \in [T]} s_i'(t) = s_i'(t')$. Then all (finite) vertices of the constant regions $C_{s'} \in C_{\Phi, 1}$ with $W = I_{n_1}$ are contained in the convex hull of P.*

We further have $\mathrm{conv}(P) \subset [a, b]$ where

$$a :=$$
$$b := \max(-\beta u(0) - \alpha i(0),) + (-b) + \vartheta \cdot \mathbb{1}_{n_1}$$

Proof. We will first proof that if $x \in \mathbb{R}^{n_1}$ is a point with $\exists_{i \in [n_1]} \exists_{t,t' \in [T]} s_i'(t) \neq s_i'(t')$, then $x \in \text{conv}(P)$. Suppose $x \notin \text{conv}(P)$. Then there is an affine hyperplane $H \subset \mathbb{R}^{n_1} \setminus \text{conv}(P)$ with $x \in H$ by Lemma 4.8. We further have by Lemma 4.9, that there is a dimension $i' \in [n]$ such that $\pi_{[n] \setminus \{i'\}}(H) = \mathbb{R}^{n-1}$. So $x + e_{i'} \notin H$ and we have a vector $y \in H$ such that $\forall_{i \in [n] \setminus \{i'\}} y_i \neq 0$.

Let us now regard the region C_{s_y}, where we define s_y such that $\forall_{i \in [n] \setminus \{i\}} \forall_{t \in [T]} (s_y)_i(t) = 1 \Leftrightarrow y_i \geq 0 \Leftrightarrow y_i > 0$, and $\forall_t (s_y)_{i'}(t) = 1$ exactly when $x + e_{i'}$ and $\text{conv}(P) \subset \mathbb{R} \setminus H$ are on different connected components of $\mathbb{R}^n \setminus H$, i.e. there does not exist a path $c : [0,1] \to \mathbb{R} \setminus H$ with $c(0) = x + e_{i'}$ and ending with $c(1) \in \text{conv}(P)$. Since convex sets are trivially path connected, if such a path exists for any point in $\text{conv}(P)$, it exists for every point of $\text{conv}(P)$.

By definition of s_y we can apply Lemma 4.5 and obtain that C_{s_y} has one finite vertex $z \in P$, that $C_{s_y} = [x^{s_y}, y^{s_y})$ with $\forall_{i \in [n_1]} (x^{s_y})_i < (y^{s_y})_i$ and $x_i^{s_y} = -\infty$ if $\forall_{t \in [T]} (s_y)_i(t) = 0$ and $y_i^{s_y} = \infty$ if $\forall_{t \in [T]} (s_y)_i(t) = 1$ for all $i \in [n_1]$. So by definition of s_y there exists a $\delta > 0$ such that $\forall_{i \in [n] \setminus i'} \delta y_i \in [x_i^{s_y}, y_i^{s_y})$. The same holds for $i = i'$: Suppose $(s_y)_{i'}(t) = 0$. Then there exists a path between $x + e_{i'}$ and z

$\square$

While in theory we would expect the number of constant regions to grow exponentially with time, it grows only quadratically.

Theorem 4.1. *A r. d.t. LIF-SNN with $W = I_{n_1}$, ... has at a maximum $(\frac{T^2+T}{2} + 1)^n$ different constant regions.*

Proof. Since the number of constant regions of a r. d.t. LIF-SNN are just unions of the constant regions of the corresponding first layer, it suffices to compute the maximum number of constant regions of that layer.

Let us consider by how much the regions can increase going from $t-1$ to $t \in [T]$. We can categorize the regions at $t-1$ by the number of spikes they have in each component. We shall write $C_{\sigma,t-1}$ for the region with sums $(\sigma_1, ..., \sigma_{n_1}) = \sigma \in [t-1]_0^{n_1}$ at time-step $t-1$. Let further $C_{\Sigma,t-1} := \{C_{\sigma,t-1} \mid \sigma \in [t-1]_0^{n_1}\}$. By definition we have $|C_{\Sigma,t-1}| \leq t^{n_1}$.

We will now show in each region $C_{\sigma,t-1}$ only

$\square$

Corollary 4.1.

Proof.

$\square$

5   Experimental results

To try to better understand the landscape of the input of a r. d.t. LIF-SNN to proof Theorem 4.1, we have created a few programs.

5.1   Computing the number of regions

We assume $W = I_n$ yet again, since we the algorithms become much simpler and more efficient. Also once we understand the situation for $W = I_n$, we are hopeful to proof it for arbitrary W.

For the first algorithm we define a function $g_i(t; s_p)$ such that we have $\forall_{i \in [n_1]} s_i(t; x; s_p) = 1 \Leftrightarrow x_i \geq g_i(t; s_p)$ (using notation from section 4) We compute for a $i \in [n_1]$:

$$H(p_i(t; x; s_p) - \vartheta)$$

$$= H\left(\sum_{k=1}^{t} \beta^{t-k} \left(i_i(k; x; s_p) + b_i \right) + \underbrace{\beta^t u_i(0) - \vartheta \left(1 + \sum_{k=1}^{t-1} \beta^{t-k} (s_p)_i(k) \right)}_{(*)} \right)$$

$$= H\left(\sum_{k=1}^{t} \beta^{t-k} \left(\alpha^k i_i(0) + \sum_{l=1}^{k} \alpha^{k-l} \left(x_i + (V s_p(l-1))_i \right) + b_i \right) + (*) \right)$$

$$= H\left(\sum_{k=1}^{t} \beta^{t-k} \sum_{l=1}^{k} \alpha^{k-l} x_i + \sum_{k=1}^{t} \beta^{t-k} \left(\alpha^k i_i(0) + b_i + \sum_{l=1}^{k} \alpha^{k-l} (V s_p(l-1))_i \right) + (*) \right)$$

$$= H\left(x_i + \frac{\sum_{k=1}^{t} \beta^{t-k} \left(\alpha^k i_i(0) + b_i + \sum_{l=1}^{k} \alpha^{k-l} (V s_p(l-1))_i \right) + (*)}{\sum_{k=1}^{t} \beta^{t-k} \sum_{l=1}^{k} \alpha^{k-l}} \right)$$

So we get

$$g(t; s_p) = -\frac{\sum_{k=1}^{t} \beta^{t-k} \left(\alpha^k i(0) + b + \sum_{l=1}^{k} \alpha^{k-l} V s_p(l-1) \right) + (*)}{\sum_{k=1}^{t} \beta^{t-k} \sum_{l=1}^{k} \alpha^{k-l}}$$

Algorithm 1 Recursive Region Computation

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    function COMPUTEREGIONSSTARTINGWITHST(t, st, x^C, y^C)
        if $t > T$ **then**
            return $\{(st, x^C, y^C)\}$
        end if
        $x \leftarrow g(t; st)$
        $subRegions \leftarrow \{[x', y') \mid x' < y', \ x'_i \in \{x^C_i, x_i\}, \ y'_i \in \{y^C_i, x_i\}\}$
        for $C \in subRegions$ **do**
            $newST \leftarrow \text{append}(st, [(\mathbf{1}_{\{x_i \leq x'_i\}})_{i \in [n_1]}])$
            COMPUTEREGIONSSTARTINGWITHST$(t+1, newST, x', y')$
        end for
    end function
    function COMPUTEREGIONS()
        COMPUTEREGIONSSTARTINGWITHST$(1, [(0, ..., 0)], (-\infty, ..., -\infty), (\infty, ..., \infty))$
    end function

Bibliography

References

[Nguyen et al., 2025] Nguyen, D. A., Araya, E., Fono, A., and Kutyniok, G. (2025). Time to spike? understanding the representational power of spiking neural networks in discrete time.