Branch-and-Bound

- Typically used to look for an optimal solution by searching an implicit graph
 - The graph is implicit and typically acyclic or even a tree
- During the search, we calculate at each node a bound on the possible value of any solution that might lie further in the graph
 - If the bound guarantees us the further solution is worse than the best solution so far, we need not to explore this part of graph anymore
- Implementation
 - Priority list

Some facts about the problem

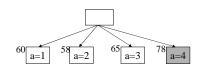
- Some initial bounds
 - One upper bound
 - total cost along diagonals: min(73, 87) = 73
 - One lower bound
 - Sum of minimum cost on each row (best job): 58
 - Sum of minimum cost on each row (best person): 49
 - We have an initial bound [58, 73]

	1	2	3	4
a	11	12	18	40
b	14	15	13	22
С	11	17	19	23
d	17	14	20	28

The assignment problem

- *n* agents are to be assigned *n* tasks. Each agent has exactly one task to perform
 - If agent *i*, $1 \le i \le n$, is assigned task *j*, $1 \le j \le n$, the cost to perform the task is c_{ii} .
 - Given the complete matrix of costs, assign the agents to tasks so as to minimize the total cost of executing the n tasks
- Can you come out a brutal-force algorithm? How much does it cost?
 - Some similarity to the 8 queens problem, but we need consider the cost

Example



Initial bound [58..73]

- At a time, we expand a node and estimate the minimum possible cost, using the heuristic that each job has to be assigned to an agent.
- We disable a node if its cost is higher than the current upper bound

	1	2	3	4
a	11	12	18	40
b	14	15	13	22
С	11	17	19	23
d	17	14	20	28

