

Average Case Analysis

- We need to know instance distribution
 - Still have to model for some distribution
 - Sometimes we make ideal assumption that all instances of any given sizes are equally distributed
 - Randomize algorithms
- Probabilistic Analysis for average-case or expect time

The hiring problem

- You are to hire a new office assistant
- A recruiter send you one out of n candidates each day
- You interview the candidate being sent
 - The interview cost is c_i , which is small
- You fire current assistant and hire the candidate if he/she is the best so far
 - The firing/hiring cost is c_h , which is high
- You like to know the expect “cost” (not execution time here) of the hiring

The algorithm for the hiring problem

```
HIRE-ASSISTANT( $n$ )
{
   $best = 0$ ;
  for ( $i=1$ ;  $i \leq n$ ;  $i++$ ) {
    interview candidate  $i$ ;           ←  $cost\ c_i$ 
    if (candidate  $i$  is better than candidate  $best$ ) {
      fire candidate  $best$ ;
       $best = i$ ;
      hire candidate  $i$ ;             ←  $cost\ c_h$ 
    }
  }
}
```

What is the worst case cost? Best case?
What is the expected cost?

Probabilistic Analysis and Randomized Algorithm

- Probabilistic analysis
 - Use probability in the analysis
 - Need to know to make assumption the distribution of inputs/instances
 - Computing expected running time
 - Averaging over all possible inputs
- Randomized algorithm
 - Randomize the input so the behavior of algorithm depends on the input as well as the random number generator

The randomized algorithm for the hiring problem

```

HIRE-ASSISTANT(n)
{
    best = 0;
    for (i=1; i<=n; i++) {
        randomly to choose a candidate c to interview;
        if (candidate c is better than candidate best) {
            fire candidate best;
            best = c;
            hire candidate c;
        }
    }
}

```

Indicator random variables

- Given a sample space S and an event A . The indicator variable $I\{A\}$ associated with event A is

$$I\{A\} = \begin{cases} 1 & \text{if } A \text{ occurs} \\ 0 & \text{if } A \text{ does not occur} \end{cases}$$

- Example: flip a fair coin. $S = \{H, T\}$.
 – X_H ($I\{H\}$) is the indicator variable associate with the event H
 $X_H = I\{H\} = \begin{cases} 1 & \text{if } H \text{ occurs} \\ 0 & \text{if } T \text{ occurs} \end{cases}$

Expected value of indicator variable

- Theorem:
 - Given a sample space S and an event A , let $X_A = I\{A\}$. Then $E[X_A] = \Pr\{A\}$.
- Example:
 - X_H is the indicator variable associate with the event H (head) $E[X_H] = \Pr(H) = 1/2$
 - What is the expected number of heads in n coin flips
 - Let X be the random variable denoting the total number of head in the n coin flips, $X_i = I\{\text{the } i_{th} \text{ flip results in the event } H\}$

$$X = \sum_{i=1}^n X_i$$

$$E[X] = E\left(\sum_{i=1}^n X_i\right) = \sum_{i=1}^n E[X_i] = \sum_{i=1}^n \frac{1}{2} = \frac{n}{2}$$

Analysis of the hiring problem

- Let X be the random variable denoting the total number of times a new assistant is hired
- $X_i = I\{\text{the } i_{th} \text{ candidate is hired}\}$

$$E[X_i] = \Pr\{\text{the } i_{th} \text{ candidate is hired}\} = \frac{1}{i}$$

$$X = \sum_{i=1}^n X_i$$

$$E[X] = E\left(\sum_{i=1}^n X_i\right) = \sum_{i=1}^n E[X_i] = \sum_{i=1}^n \frac{1}{i} = \ln n + O(1)$$

- Assuming that the candidates are presented in a random order, algorithm HIRE_ASSISTANT has a total hiring cost $O(c_h \ln n)$

Insertion sort

- Assume that
 - all $n!$ permutations of the array elements are equally likely
 - Or we randomize the array in the beginning
 - The n elements of the array are distinct

```
void insertionSort(int A[0..n-1])
{
    int i, j, tmp;

    for (i=1; i<n; i++) {
        tmp=A[i];
        j = i-1;
        while (j>=0 && tmp<A[j]) {
            A[j+1] = A[j];
            j--;
        }
        A[j+1] = tmp;
    }
}
```

Average case analysis of insertion sort

- Choose the while loop condition check as the barometer
- Define the *partial rank* of $A[i]$ be the position of $A[i]$ if the sub array $A[0..i]$ is sorted
 - The partial rank of $A[i]$ is equally likely to take any value between 0 and i
- Note that $A[0..i-1]$ is sorted. Let the partial rank of $A[i]$ be k . The barometer executes $i-k+1$ times. The average number of times of the test is

$$c_i = \frac{1}{i+1} \sum_{k=0}^i (i-k+1) = \frac{i}{2} + 1$$

- The total average is $\sum_{i=1}^{n-1} c_i = \sum_{i=1}^{n-1} \left(\frac{i}{2} + 1\right) = \frac{(n-1)(n+4)}{4}$

Analysis using indicator variable

- Let X_{ij} be the indicator variable that associated with event $A[i] < A[j]$
- Let X_i be the number of times of the test at iteration i (from 1 to $n-1$), we have $X_i = \sum_{j=0}^{i-1} X_{ij} + 1$

$$E[X_{ij}] = \frac{1}{2}$$

$$E[X_i] = E\left[\sum_{j=0}^{i-1} X_{ij}\right] + 1 = \frac{i}{2} + 1$$