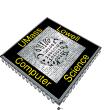
# COMP 3080 – OPL Exception-Based Syntax Error Recovery (for Assignment 2)

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```
program \longrightarrow stmt\_list $$
stmt\_list \longrightarrow stmt stmt\_list \mid \epsilon
stmt \longrightarrow id := expr \mid read id \mid write expr
expr → term term_tail
term\_tail \longrightarrow add\_op term term\_tail \mid \epsilon
term → factor factor_tail
factor\_tail \longrightarrow mult\_op factor factor\_tail \mid \epsilon
factor \longrightarrow (expr) \mid id \mid number
add\_op \longrightarrow + | -
mult\_op \longrightarrow * | /
```

```
void factor () {
    switch (input token) {
        case t id:
            printf ("predict factor --> id\n");
            match (t id);
            break;
        case t literal:
            printf ("predict factor --> literal\n");
            match (t literal);
            break;
        case t lparen:
            printf ("predict factor --> lparen expr rparen\n");
            match (t lparen);
            expr ();
            match (t rparen);
            break;
        default: error ();
```

```
void error () {
    printf ("syntax error\n");
    exit (1);
void match (token expected) {
    if (input token == expected) {
        printf ("matched %s", names[input token]);
        if (input token == t id || input token == t literal)
            printf (": %s", token image);
        printf ("\n");
        input token = scan ();
    else error ();
```



## [From the description of Assignment 2:]

Implement exception-based syntax error recovery, as described in Section 2.3.5 on the textbook's companion site. At the least, you should attach handlers to statements, relations, and expressions. [...]

When match sees a token other than the one it expects, it could simply throw a Syntax Error exception. The resulting algorithm would recover by deletion only: The exception handler will delete tokens until it finds something in either the FIRST set or the FOLLOW set of the nonterminal corresponding to the current recursive descent routine. An attractive alternative is to mirror Wirth's recovery algorithm, and have match insert what it expects and continue (presumably after printing an error message). You may implement either strategy.



# Exception-Based Recovery in Recursive Descent

# Pseudo-code for statement-level error recovery:

```
procedure statement()
    try
                              — code to parse a statement
    except when syntax_error
        loop
             if next\_token \in FIRST(statement)
                 statement() -- try again
                 return
             elsif next_token \in FOLLOW(statement)
                 return
             else get_next_token()
```



#### **FIRST**

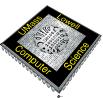
```
program {id, read, write, $$}
stmt_list {id, read, write}
stmt {id, read, write}
expr {(, id, number}
term_tail {+, -}
term {(, id, number}
factor_tail {*, /}
factor {(, id, number}
add_op {+, -}
mult_op {*, /}
```

### **FOLLOW**

```
program Ø
stmt_list {$$}
stmt {id, read, write, $$}
expr {), id, read, write, $$}
term_tail {), id, read, write, $$}
term {+, -, ), id, read, write, $$}
factor_tail {+, -, ), id, read, write, $$}
factor {+, -, *, /, ), id, read, write, $$}
add_op {(, id, number}
mult_op {(, id, number)}
```

#### PREDICT

- 1.  $program \longrightarrow stmt\_list \$\$ \{id, read, write, \$\$\}$
- 2. stmt\_list → stmt stmt\_list {id, read, write}
- 3.  $stmt\_list \longrightarrow \epsilon \{\$\$\}$
- 4.  $stmt \longrightarrow id := expr\{id\}$
- 5.  $stmt \longrightarrow read id \{read\}$
- 6.  $stmt \longrightarrow write expr\{write\}$
- 7.  $expr \longrightarrow term \ term\_tail \{ (, id, number \} \}$
- 8.  $term\_tail \longrightarrow add\_op \ term \ term\_tail \{+, -\}$
- 9.  $term\_tail \longrightarrow \epsilon$  {), id, read, write, \$\$}
- 10.  $term \longrightarrow factor \ factor\_tail \{(, id, number)\}$
- 11.  $factor\_tail \longrightarrow mult\_op \ factor \ factor\_tail \ \{*, /\}$
- 12.  $factor\_tail \longrightarrow \epsilon \{+, -, \}, id, read, write, \$\}$
- 13.  $factor \longrightarrow (expr) \{(\}$
- 14.  $factor \longrightarrow id \{id\}$
- 15.  $factor \longrightarrow number \{number\}$
- 16.  $add\_op \longrightarrow + \{+\}$
- 17.  $add\_op \longrightarrow -\{-\}$
- 18.  $mult\_op \longrightarrow * \{*\}$
- 19.  $mult\_op \longrightarrow / \{/\}$



```
void error () {
    printf ("syntax error\n");
    return;
void match (token expected) throw (Syntax Error) {
    if (input token == expected) {
        printf ("matched %s", names[input_token]);
        if (input token == t id || input token == t literal)
            printf (": %s", token image);
        printf ("\n");
        input token = scan ();
    else {
        error ();
        input token = expected; // Implement Wirth's method
        throw Syntax Error (/* parameters? */);
```

```
void factor () {
 try {
    switch (input token) {
        case t id:
            printf ("predict factor --> id\n");
            match (t id);
            break;
        case t literal:
            printf ("predict factor --> literal\n");
            match (t literal);
            break;
        case t lparen:
            printf ("predict factor --> lparen expr rparen\n");
            match (t lparen);
            expr ();
            match (t rparen);
            break;
        default: error (); throw Syntax Error( /* parameters? */ );
  catch (Syntax Error /* parameters? */) {
    /* Implement recovery loop from pseudo-code - see next slide */
```

```
catch (Syntax Error /* parameters? */) {
  while (true) {
    switch (input token)
    /* FIRST (factor) */
    case t lparen: case t id: case t_literal:
      factor (); // try again
      return;
    /* FOLLOW (factor) */
    case t add: case t sub: case t mul: case t div:
    case t rparen: case t id: case t read: case t write:
    case t eof:
      return;
    default:
      input token = scan(); // get next token
```