

Q1

The DFA takes in an input symbol, a state and returns the next state. The NFA takes in an empty string, a state and input ~~symbols~~ ^{string}, then returns a set of next states. The DFA has a unique output and input. NFA can accept an empty string.

A terminal ^{symbol} is a set ~~of~~ of characters in an alphabet that appear in strings and are generated by the grammar.

A nonterminal variable is a set of ~~for~~ patterns of terminals generated by the nonterminal symbol.

- Used to check to see if languages are ^{not regular} ~~Must be recognized by a finite automata~~.
1. ~~Must be recognized by a finite automata~~
 2. An alphabet is a finite set of symbols denoted by Σ
 3. A string contains symbols of the Alphabet that combine to form
 4. A language is something that can be detected by a finite automata
 5. A sequence is a pattern of symbols following ^{some rule} ~~some~~ ~~condition~~
 6. K-tuple determines how many rules a machine will have determined by K.
So for a GFG, is a 4 tuple (V, Σ, R, S)
 7. Ordered pair contains two symbols that are different $(0,1)$ $(1,0)$
 8. Unordered pair can be the same two symbols $(1,1)$ $(2,2)$
 9. The Domain is the input of a function $f(n)$
 10. The Range is the output of a function $f(n) = 2$

GFG is a 4 tuple also known as $\{V, \Sigma, R, S\}$

V is ~~the~~ a finite set called variable
 Σ is a finite set of variable V , called termin
 R is ~~the~~ a finite set of rules
 S is the state variable

PDA is a 6 tuple also known as a Push down automata, $\{Q, \Sigma, P, q_0, \Delta, F\}$

Q is the set of states
 Σ is the set of input symbols alphabet
 P is the stack symbols alphabet
 q_0 is the start state
 Δ is the transition function
 F is the ^{set of} accept states

Regular Expression:

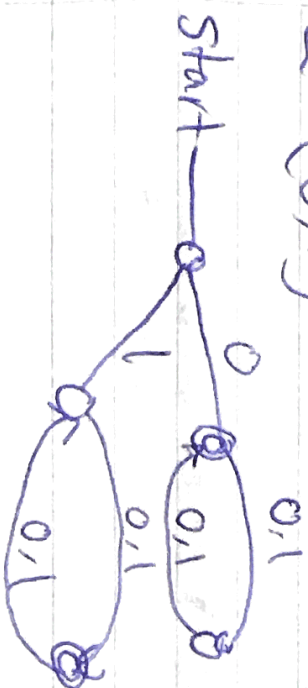
$\{a\}$, $\{e\}$, \emptyset

Assuming... \rightarrow

$\{R_1 \cup R_2\}$ R_1 and R_2 are regular expressions
 $\{R_1 \circ R_2\}$ R_1 and R_2 are regular expressions
 $\{R_1^*\}$ R_1 is a regular expression

DFA

3. $\Sigma = \{0, 1\}$



4. $\Sigma = \{0, 1\}$ Alphabet
 $Q = \{q_0, q_1, q_2, q_3, q_4\}$ States

	0	1	
q_0	q_1	q_2	$q_0 = q_1$ (start state)
q_1	q_1	q_3	$F = \{q_2, q_3, q_4\}$ (final state)
q_2	q_1	q_3	
q_3	q_1	q_4	
q_4	q_1	q_4	$L = \{w \mid w \text{ starts with two } 1\}$

5. The PDA is a 6 tuple and contains stack ~~calculated~~ P and the DFA doesn't.

The PDA is similar to the NFA. The ~~PDA~~ DFA can only take in ~~regular~~ regular languages and a PDA can take in ~~regular~~ non-regular languages. The DFA can only have unique one input state and output state. The PDA can accept empty strings. One is called push down automata and DFA is called Non-deterministic automata

6h. Expression = Expression + terms / terms

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Selection

C.a. ~~Expression~~-stmt

Expression-stmt \Rightarrow var-stmt

Select ion-stmt

Expression-stmt \Rightarrow var-stmt

C.b.

~~Expression~~