L2910-5335 Final Exam VerB

Phong Vo TOTAL POINTS 96 / 100 **QUESTION 10 QUESTION 1** 10 Priority inversion 2/2 1 Microkernel 2/2 √ - 0 pts Correct √ - 0 pts Correct **QUESTION 11** QUESTION 2 11 Named pipes 0/2 2 Policy 2 / 2 √ - 2 pts Incorrect (correct answer is "named pipe") √ - 0 pts Correct **QUESTION 12 QUESTION 3** 12 Pipe = file 2 / 2 3 Stack section 2/2 √ - 0 pts Correct √ - 0 pts Correct **QUESTION 13 QUESTION 4** 13 Critical section requirements 2/2 4 Child process 2/2 √ - 0 pts Correct √ - 0 pts Correct **QUESTION 14 QUESTION 5** 14 Signal() before wait() 2 / 2 5 Modules 2/2 √ - 0 pts Correct √ - 0 pts Correct **QUESTION 15** QUESTION 6 15 Logical address -> page number 2 / 2 6 Message passing 2/2 √ - 0 pts Correct √ - 0 pts Correct **QUESTION 16 QUESTION 7** 16 Inverted or hashed page table 2/2 7 Spinlocks 2/2 √ - 0 pts Correct (either "inverted page table" or √ - 0 pts Correct "Hashed page table" is acceptable) **QUESTION 8 QUESTION 17** 8 Thread signal handling 2/2 17 Safety algorithm 2/2 √ - 0 pts Correct √ - 0 pts Correct **QUESTION 9 QUESTION 18** 9 Thread-local storage 2/2 18 SCAN 2 / 2 √ - 0 pts Correct √ - 0 pts Correct

QUESTION 19 19 Running state 2/2 √ - 0 pts Correct **QUESTION 20** 20 Mach 2/2 √ - 0 pts Correct **QUESTION 21** 21 Amdahl's law 2/2 √ - 0 pts Correct **QUESTION 22** 22 Kernel 2 / 2 √ - 0 pts Correct **QUESTION 23** 23 API 2/2 √ - 0 pts Correct **QUESTION 24** 24 mutex=binary semaphore 2/2 √ - 0 pts Correct **QUESTION 25** 25 non-preemptive kernel 2/2 √ - 0 pts Correct **QUESTION 26** 26 hashed page tables 0/2 √ - 2 pts Incorrect (correct answer is False) **QUESTION 27** 27 TLB miss 2 / 2 √ - 0 pts Correct **QUESTION 28** 28 deadlock-free / starvation 2 / 2 √ - 0 pts Correct

QUESTION 30 30 LRU 2/2 √ - 0 pts Correct **QUESTION 31** 31 OS modes 10 / 10 √ - 0 pts Correct **QUESTION 32** 32 Working set 10 / 10 √ - 0 pts Correct (or close enough for full credit!) QUESTION 33 33 RAID levels 10 / 10 √ - 0 pts Correct (or close enough for full credit!) **QUESTION 34** 34 Intel five levels 10 / 10 √ - 0 pts Correct

QUESTION 29

√ - 0 pts Correct

29 thread components 2/2

Name:		PHONG VO	Student ID: _	01790283					
*									
		Fina	al Exam						
	COMF	2.3080-201-202 — Operating Sy	ystems; Decemb	per 19, 2019 – Dr. Wilkes					
Note	: This e		ccept for up to th tes (no photocop	ree 8.5x11" sheets of paper with ies).					
Multip	ole Cho	ice Questions – 2 points each:	Mark the corre	ct answer(s).					
1.	(MAR	K A SINGLE ANSWER) A microke	ernel is a kernel _	*					
				d to reduce resident memory size					
	O that is compressed before loading in order to reduce its resident memory size								
	0	that is compiled to produce the	e smallest size po	ssible when stored to disk.					
	•	that is stripped of all nonessen	tial components						
2.	(MARI	(A SINGLE ANSWER) Policy	*						
		determines how to do somethi	ng						
	•	determines what will be done							
		is not likely to change across p	The second secon	•					
		is not likely to change over tim							
3.	•	(A SINGLE ANSWER) The	-	-					
	2	on parameters, return addresses	s, and local varial	oles.					
		data section							
		program counter							
	4	stack section							

4. (MARK ALL THAT APPLY) When a child process is created, which of the following is a

possibility in terms of the execution or address space of the child process?

5. (MARK A SINGLE ANSWER) _____ allow operating system services to be loaded

6. (MARK A SINGLE ANSWER) Microkernels use _____ for communication.

The child process is a duplicate of the parent process.
The child process has a new program loaded into it.
The child process runs concurrently with the parent.

O text section

☐ None of the above

O Virtual machines

message passingo shared memoryo system callso virtualization

O Graphical user interfaces

ModulesFile systems

dynamically.

Name:		Student ID:	(b)
7.	(MARI	K A SINGLE ANSWER) Which of the following statements is true?	
		A binary semaphore can never be used as a counting semaphore.	
		A counting semaphore can never be used as a binary semaphore.	
	-	Counting semaphores can be used to control access to a resource with a finite number of instances.	
	0		
	O	Spinlocks can be used to prevent busy waiting in the implementation of semaphore.	
8.	(MARI	(A SINGLE ANSWER) Which of the following would be an acceptable signal	
	handli	ng scheme for a multithreaded program?	
	0	Deliver the signal to the thread to which the signal applies.	
	0	Deliver the signal to every thread in the process.	
	0	Deliver the signal to only certain threads in the process.	
	•	All of the above	
9.	(MARI	K A SINGLE ANSWER) Thread-local storage is data that	
	0	is generated by the thread independent of the thread's process	
	0	is not associated with any process	
	. 🔴	is unique to each thread	
	0	is not associated with any process	
10.	(MARI	(A SINGLE ANSWER) occurs when a higher-priority process need	ls
	to acce	ess a data structure that is currently being accessed by a lower-priority process.	
	0	A critical section	
	0	A race condition	
	0	Deadlock	
	•	Priority inversion	
11.	(MARI	(A SINGLE ANSWER) A(n) allows several unrelated processes t	to
	use th	e pipe for communication.	
	•	anonymous pipe	
	0	LIFO pipe	
	0	named pipe	
	0	ordinary pipe	
12.	(MARI	(A SINGLE ANSWER) Child processes inherit UNIX ordinary pipes from their	
	parent	process because:	
	0	All IPC facilities are shared between the parent and child processes.	
	•	A pipe is treated as a file descriptor and child processes inherit open file	
		descriptors from their parents.	
	0	The pipe is part of the code and children inherit code from their parents.	
	0	The STARTUPINFO structure establishes this sharing.	
13.	(MARI	(A SINGLE ANSWER) A solution to the critical section problem does not have to	
	satisfy	which of the following requirements?	
		Atomicity	
^\	, O	Bounded waiting	
	0	Mutual exclusion	
	0	Progress	

Process ID	Maximum need	Currently allocated
P0	10	4
P1	3	1
P2	6	4

Which of the following correctly characterizes the state of this system?

- O It is safe.
- It is not safe.
- O It is an impossible state.
- O The state cannot be determined.

Name:	Student ID:	(b)
		_ (1)

- 18. (MARK A SINGLE ANSWER) Consider a disk queue holding requests to the following cylinders in the listed order: 116, 22, 3, 11, 75, 185, 100, 87. Using the SCAN scheduling algorithm, what is the order that the requests are serviced, assuming the disk head initially is at cylinder 88 and moving upward (toward higher cylinder numbers)?
 - O 87 75 100 116 185 22 11 3
 - O 100 116 185 3 11 22 75 87
 - **9** 100 116 185 87 75 22 11 3
 - O 116 22 3 11 75 185 100 87

Name:	PHONGVO	Student ID: _	01790283	(b)
True/f	False Questions – 2 points each: Ma	rk the correct an	swer,	
19.	For a single-processor/single-core sys Running state. O True False	stem, there can be	multiple processes in the	
20.	The Mach operating system treats syn True O False	stem calls with me	essage passing.	
21.	Amdahl's Law describes performance parallel component. True O False	gains for applicat	ions with both a serial and	
22.	The operating system kernel consists computer. O True False	of all system and	application programs in a	
23.	Application programmers typically usTrueFalse	e an API rather th	an directly invoking system c	alls.
	Mutex locks and binary semaphores a True False 	are essentially the	same thing.	
25.	A nonpreemptive kernel is safe from True O False	race conditions or	n kernel data structures.	
26.	. Hashed page tables are not useful wh True O False	nen handling addre	esses larger than 32 bits.	*
27.	A page fault must be preceded by a T True True	LB miss.		
28.	A deadlock-free solution eliminates tTrueFalse	he possibility of st	arvation.	
29.	Each thread has its own register set aTrueFalse	and stack.		

Name:	:	 Student ID:	(b)

- 30. When using an LRU approximation page-replacement algorithm, if the page-fault rate is too high, the process may have too many frames.
 - O True
 - False

Name: Student ID:	ame:	16 VO	Student ID: 017-90 283	(b)
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Short Answer Questions – 10 points each: Write your answer in the space provided.

- and user mode, where the OS mus, is also known as privileged on supervisor mode. The software in leavel mode can directly access the handware and can control the switching between CPU modes. Rest of the software muse in user mode. Handware direct access is restricted in user mode. The program in user mode has enough memory to do its job. I has given an address space which is only visible to them.
 - 32. Many operating systems (including Microsoft Windows 10) use the "working set" model to determine which logical (virtual) pages of a process should be held in physical memory frames at a given time in order to minimize the number of page faults. Briefly explain the working set model, including what happens to the page fault rate when a process transitions between regions of locality in the program code.
- Working set is used to allocate the location of logical (virtual) pages will be held in physical memory frames. Acress to a page marked invalid causes a page fault. The working set will notice that the invalid bit is set, cashing a trap to the os. This trap is the result of the os's failure to bring the desired page into memory.

Name:	Student ID:	(h)
TVOICE.	student is.	(n)

33. Describe RAID levels 0, 1, and 5. Include sample diagrams for each of these three RAID levels showing how data blocks, as well as parity blocks (if applicable), are distributed among the disks. (In your diagrams, use the minimum number of disks necessary for each level.)

RAID				国	国	1	Ÿ		must be identical. 5: P P P P P Tokutica 10: P P P P P (HOOS) 10: P P P P P P P P P P P P P P P P P P P
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RAID	Ö :	Йo	n - F	Redun	t dan	t Sta	ipping		中国一国一国(HOOS
PAID		M	no	ring	end	shadi	owing		图图图图图
RAID	2 :			Inter	denec	d Dist	ributed	1 pa	aity

34. Intel's original "IA-32e mode" memory addressing scheme had a hierarchical page table with four levels. Explain why Intel recently added a fifth level of hierarchy to the x86-64 (IA-32e) page tables.

With 5 level paging, the vintual address size increases from a 25678 maximum to 128 PB while the physical address size threshold goes from 6478 to 4PB. This big set of patches increases the vintual on physical address space capacity of the Linux kernel for the future Intel X8-64 hard wome which can connently be found via this kernel mailing list thread.