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**Vérification d'une méthodologie pour la conception de systèmes
numériques critiques**

Présenté par Vincent IAMPIETRO

Le Date de la soutenance

**Sous la direction de David Delahaye
et David Andreu**

Devant le jury composé de

[Nom Prénom], [Titre], [Labo]	[Statut jury]
[Nom Prénom], [Titre], [Labo]	[Statut jury]
[Nom Prénom], [Titre], [Labo]	[Statut jury]



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The acknowledgments and the people to thank go here, don't forget to include your project advisor...

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List of Abbreviations

SITPN	S ynchronously executed I nterpreted T ime P etri N et with priorities
VHDL	V ery high speed integrated circuit H ardware D escription L anguage
PCI	P lace C omponent I nstance
TCI	T ransition C omponent I nstance
GPL	G eneric P rogramming L anguage
HDL	H ardware D escription L anguage
LRM	L anguage R eference M anual

For/Dedicated to/To my...

Chapter 1

\mathcal{H} -VHDL: a target hardware description language

- Make a remark on the notation of design states, e.g. \mathcal{E}'' refers to the set of events of design state σ'' when there is no ambiguity.

1.1 Presentation of the VHDL language

The intent here is to give an overview of the VHDL language, its purpose, its main syntactical constructs, and an informal description of its semantics as presented in the Language Reference Manual (LRM) [12]. As explained in Chapter ??, the hilecop transformation generates a VHDL design implementing the input SITPN model. To do so, the transformation generates and connects the component instances of two previously defined VHDL designs: the place design that implements the concept of a SITPN place, and the transition design that implements a SITPN transition. These designs were defined by the INRIA team at the creation of the HILECOP methodology. Here, we use excerpts of the definition of the place and transition designs to illustrate the content of VHDL programs. The reader will find the source code of the place and transition designs in concrete and abstract syntax in Appendices B and B.

1.1.1 Main concepts

The VHDL acronym stands for Very high speed integrated circuit Hardware Description Language. As its name indicates, the main purpose of the VHDL language is to describe hardware circuits. In VHDL, the concept of program is called *design*. A design corresponds to the description of a hardware circuit. A VHDL design is composed of two descriptive parts. The first part is called the entity and describes the interfaces of circuit, namely: the input and output ports, and the generic constants. Listing is an excerpt of the transition design's entity that defines the generic constants, the input and output port interfaces of the design. Figure is a visual representation of the interfaces of the transition design.

```

1  entity transition is
2      generic(
3          transition_type : transition_t := NOT_TEMPORAL;

```

```

4   input_arcs_number : natural := 1;
5   conditions_number : natural := 1;
6   maximal_time_counter : natural := 1
7   );
8   port(
9     clock : in std_logic;
10    reset_n : in std_logic;
11    input_conditions : in std_logic_vector(conditions_number-1 downto 0);
12    time_A_value : in natural range 0 to maximal_time_counter;
13    time_B_value : in natural range 0 to maximal_time_counter;
14    input_arcs_valid : in std_logic_vector(input_arcs_number-1 downto 0);
15    reinit_time : in std_logic_vector(input_arcs_number-1 downto 0);
16    priority_authorizations : in std_logic_vector(input_arcs_number-1 downto 0);
17    fired : out std_logic
18  );
19  end transition;

```

LISTING 1.1: The entity part of the transition design in concrete VHDL syntax.

The entity generic clause holds the declaration of the generic constants. The purpose of generic constants is either to represent some dimensions of the design (e.g. the size of ports, internal signals...) or to represent constant values used throughout the design. In Listing ??, one can see that the `conditions_number` generic constant gives a dimension to the type of the `input_conditions` input port, which is an array of Boolean values with indexes ranging from 0 to `conditions_number-1` (that is the meaning of `std_logic_vector (conditions_number-1 downto 0)`). The port clause holds the declaration of input and output ports of the design. The `in` keyword indicates the declaration of an input port and the `out` indicates the declaration of an output port.

FIGURE 1.1: A representation of the transition design entity.

The second part is called the architecture and corresponds to the description of the internal behavior of the design.

1.2 Choosing a formal semantics for VHDL

When considering the choice of a formal semantics for the language VHDL, it is important to remember the first aim of our work. We want to prove that the HILECOP transformation that generates VHDL code from SITPNs preserves the behavior of the initial model (i.e, the SITPN model) into the generated VHDL program. A formal semantics for the VHDL language is therefore a necessary element to be able to reason about the generated VHDL programs, and moreover to be able to compare their behaviors with the behaviors of the source SITPN models. Keeping that in mind, which formal semantics should we consider for VHDL?

The same holds for any research task: there is a tradeoff between finding a tool designed by others that will fit our needs, and creating our own tool that will mitigate the gaps between

our needs and what is available in the literature. In the present case, the tool is a formal semantics for VHDL. Adopting a fully-set semantics found in the literature as a baseground for the implementation of a formal semantics for VHDL has multiple perks. First, it reduces the formalization effort, which is not a lesser point considering that the proof ahead might be long and must still be completed within the time span of the thesis. Still, the semantics would need to be implemented in Coq, if no implementation exists (or not written in Coq). Second, the formal semantics of programming languages found in the literature are often general in their approach, this to provide a generic framework to reason about programs. However, we must not lose sight of our goal which is to prove behavior preservation; a generic formal semantics could turn out to be too complex, or necessitate too much tweaking and thus hinder the fulfillment of our task. On the other side, creating our own formal semantics for VHDL, based on the work of others, is the best way to fit our needs in compliance with our final aim. However, the pitfalls are that the resulting semantics might prove to be very specific, therefore preventing others from using it. Also, a work of formalization would be necessary which, as we already stated, would be time consuming. In order to determine whether we ought to use an existing semantics or design a new one, we must first clearly specify our needs pertaining to the VHDL language.

1.2.1 Specifying our needs: HILECOP and VHDL

Two elements are of major influence to the specification of our needs for a formal semantics: first, the context of HILECOP and the specificities of the VHDL programs that are generated; second, the context of theorem proving. These two aspects entail the following considerations.

The need for coverage

The HILECOP methodology generates particular VHDL programs. Even if some transformations can be operated on the generated programs to simplify them, the looked-for formal semantics must be able to deal with a certain subset of the VHDL language. Especially, this subset must include:

- 0-delay (or δ -delay) signal assignments (equivalent to transport-delay signal assignment with a “0 ns” after clause).
- component instantiation statements with generic constant and port mapping.
- entity’s generic constant clauses (declaration of generic constants in a design entity).

HILECOP’s VHDL programs only deal with 0-delay signal assignments because they are the only kind of signal assignments that can be synthesized. As a matter of fact, the industrial compiler/synthesizer used in the HILECOP methodology only accepts VHDL programs with no timing constructs (i.e, no wait statements or delayed signal assignments) as input.

Concerning component instantiation statement, the VHDL LRM describes a way to transform these statements into equivalent process statements and block constructs [12, p. 141] which are a part of the elaboration of the design. However, we want to preserve the hierarchical structure provided by the component instance statements arguing that it will be easier

to compare the state of a given SITPN model with a VHDL design state with an explicit hierarchical structure. Indeed, there exists a mapping between places and transitions of an SITPN and their mirror (generated by the transformation) place and transition component instances. This one-to-one correspondence might turn out to be handy to perform the proof of behavior preservation. Obviously, the semantics must cover the evaluation of process statements which are the core concurrent statements of VHDL programs.

The types of signals and variables used in HILECOP VHDL designs must have finite ranges of values. For instance, a VHDL signal that ranges over \mathbb{N} cannot be synthesized on a physical circuit. Indeed, \mathbb{N} has an infinite number of values, and would therefore require an infinite number of latches to be physically implemented. Moreover, as the number of latches used to implement a digital circuit greatly impacts the power consumption of the circuit, the types of signals and variables must be as constrained as possible to optimize the dimensioning of the circuit. The generic constant clauses of the HILECOP place and transition designs (see Section Presentation of the VHDL language) help limiting the type ranges. The generic constants define the bound of the array and natural range types for the different signals and variables declared in the place and transition designs' architecture. When a place or a transition component is instantiated, that is during the transformation of the SITPN model into VHDL code, its generic constants receive values via a generic map; we call it the dimensioning of the component instance. Therefore, generic constant clauses must belong to the subset of the VHDL language covered by the semantics (even though it seems a trivial formalization task compared to the core of the VHDL semantics).

The need for synchronous execution

The second property of HILECOP's generated VHDL programs is their synchronous execution. Indeed, the digital circuits designed with the HILECOP methodology are all synchronously executed on physical target. The generated VHDL designs declare a clock signal as an input port of their entity port interface. Thus, the behavioral part of the designs contains two kinds of processes: *synchronous* processes, i.e processes that are sensitive to the clock signal, and *combinational* processes, i.e processes that are not sensitive to the clock signal, and that are permanently running until the stabilization of the signal values. Synchronous processes react to the events of the clock signal, i.e the rising and the falling edge, and possess blocks of sequential statements that are only executed at the precise moment of the clock event¹. Therefore, we are in a strong need of a semantics that deal with synchronism, and that explicitly integrates the synchronization with a clock signal into the expression of the simulation cycle. Thus, the two dimensions of time, respectively delta time and "real" time [18], that are part of the VHDL language, must be handled by our semantics. Delta time pertains to delta-cycles, which are triggered by the execution of delta-delay signal assignments, that is, the only kind of signal assignments found in HILECOP VHDL programs. Real time pertains to time-steps. When all signal values are stable (when there are no more delta-cycles), the simulation advances the current time value to the next time point where a relevant event happens. A relevant event could be the execution of the next pending signal assignment that was associated with a delay, or the end of a `wait for x` statement (where `x` specifies a time delay). In our case, neither delayed

¹These blocks are guarded by the expressions `rising_edge(clk)` and `falling_edge(clk)`.

signal assignments nor wait statements are part the covered VHDL language subset. The only two relevant events to which the simulation advances during a time-step are the rising edge and the falling edge of the clock signal.

Other considerations

Considering the kind of proof that needs to be established, we would rather consider an operational semantics for VHDL than a denotational one. The reason is that, in the CompCert project [14], which is one of our major inspiration source, the whole C compiler toolchain is verified by reasoning over the operational semantics of the source and target languages. A last consideration pertains to whether or not the VHDL semantics must explicitly handle errors. As the SITPN semantics does not include the production of error values, the handling of errors by the VHDL semantics is not a mandatory aspect.

Qualifying criterions

We here give the list of the qualifying criterions that will help to analyze the different VHDL semantics encountered in the literature and presented in the next section. The three most relevant criterions are:

- *Synchronism*. We distinguish three levels for this criterion:
 - Synchronism is not expressible in the considered VHDL semantics.
 - Synchronism is expressible in the considered VHDL semantics. Synchronism is expressible if time-steps are handle in the semantics, at least to be able to represent clock events.
 - Synchronism is explicit, i.e. the simulation loop is built around the occurrences of clock events.

We will promote the semantics that explicitly formalize a synchronized execution of a VHDL design.

- *Component instantiation*. Either the semantics handle the component instantiation statement in its simulation rules, or component instantiation statements must be transformed in order to be executed. We will promote the semantics that handle component instantiation statements without transformation.
- *Elaboration*. This criterion expresses the ability of the semantics to handle constrained types, i.e. arrays and natural ranges, and generic constant clauses that are both dealt with during *elaboration* phase. Either the semantics handle these constructs or it does not. Of course, we will promote the first kind of semantics.

1.2.2 Looking for an existing formal semantics

Here, we give a summary of the work found in the literature pertaining to the formalization of the VHDL language semantics. Articles are gathered and presented depending on the type of semantics that is used in the formalization (operational, denotational, axiomatic...). Each semantics is analyzed regarding the needs that were previously expressed.

Denotational semantics

Some authors have been interested in giving a formal denotational semantics to VHDL. In a general manner, these authors want to reason about VHDL programs: prove properties over a VHDL program, prove that two programs are equivalent... Such tasks that are more fit for a denotational semantics.

In [9], the authors give a denotational semantics to the VHDL language, this leveraging the Focus [6] method for the development of distributed systems. Signal values and their evolution through time are represented as streams of values. Statements are denoted as stream-processing functions. Processes are stream-processing functions that takes input signal streams (signals of the sensitivity list) and yields transaction traces (i.e, waveforms) over output signals (i.e, signal that are written by the process). Transaction traces are merged together as the result of the concurrent execution of processes. Resolution functions are used in case of multiply-driven signals (i.e, signals that receive a value from multiple processes). The authors only consider 0-delay signal assignment in their semantics, stating that it is sufficient to “consider time at a logical level to model both synchronous and asynchronous designs”. However, it necessitates some transformations on a design that has a synchronous execution to express it only with 0-delay signal assignments. Therefore, this semantics does not express synchronism of execution in an explicit manner. Moreover, the component instantiation statement is not treated by the semantics.

In [3], the authors give a denotational, yet relational, semantics for VHDL. A state of a VHDL design is represented by a function binding signals to values; a worldline is a time-ordered list of states. Statements (including processes) are denoted in the semantics by a relation that binds an input couple, composed of a time point and a worldline, to an output couple of the same type. Multiple input and output couples possibly satisfy the relation denoting a particular statement; thus, the semantics is undeterministic. The semantics tries to abstract from the formalization of the simulation cycle as it is done in the LRM. The authors want to establish a semantics that is abstract enough to be able to compare all other works of formalization with the authors semantics. The authors also give an axiomatic semantics (i.e, in the Hoare logic style) which is proved to be sound and complete with the first denotational semantics. A Prolog [5] implementation of the axiomatic semantics is given. Regarding our needs, the semantics only deals with unit-delay signal assignments (no 0-delay), and therefore does not cover the VHDL subset that we are interested in. The hierarchical structure of designs is also not preserved, and, although expressible, the semantics does explicitly express a synchronous simulation cycle.

The denotational semantics expressed in [17] uses interval temporal logic as an underlying model. Leveraging this underlying model, the authors are interested in proving some properties over VHDL designs to help compilers to optimize the code, for instance, by using rewrite rules proved to be valid against the model. Some of the proofs laid out by the authors are embedded in PVS [16]. The expression of the dynamic model uses many concepts described in the LRM, like drivers, port association, driving and effective values for signals. The semantics deals with both unit-delay and δ -delay. The semantics works on fully-elaborated designs, therefore, it does not deal with component instance statements. Moreover, interval temporal logic is useful to reason on the VHDL designs in the presence of delays, however, it loses its

interest for designs presenting only 0-delay assignments.

Joining the common opinion, in [1], the author states that “denotational semantics is more adequate for mathematical reasoning”. The author formalizes the VHDL semantics to prove the equivalence between VHDL programs (for instance, a specification and an implementation). What is of major interest regarding our needs is that the author is interested in expressing a simulation cycle for synchronous designs. Therefore, a distinction is made between combinational and synchronous processes in the abstract syntax. Moreover, this work formalizes the elaboration part of a VHDL design former to the simulation; also, the elaboration keeps the hierarchical setting of the VHDL design, that is component instantiation statements are not replaced by processes. Due to the time abstraction, the semantics only deals with 0-delay. It is explained by the fact that the reference time-unit is the clock period (i.e, the only known time-step), and the advancing of time, happening during the simulation cycle as described in the LRM, is captured within the setting of the simulation cycle. Also, the semantics takes primary inputs into account (i.e, input ports of the top-level design); to preserve a synchronous behavior for the simulated design, the hypothesis is made that the values of the primary inputs are stable between two clock events. The only critic that can be made to this semantics regarding our needs is that it is expressed in denotational style.

Operational semantics

Multiple works formalize an operational semantics for VHDL. Naturally, these works are interested in the formal description of the VHDL simulator, more or less closely to the description of the LRM.

In [2], a formal description of a *functional* semantics for VHDL is laid out based on stream-processing functions. The semantics is expressed with the functional programming language Gofer [13], thus enabling the computation of execution traces, that is, the computation of the streams representing the values taken by signals over time. As in the former work of the same author [3], only unit-delay signal assignments are dealt with, however, this time the author describes a deterministic operational semantics. Regarding our needs, this work is neither interested in preserving the hierarchical structure of VHDL designs, and no mention is made regarding how a design is elaborated, nor in expressing an explicit synchronous simulation cycle.

In [4], the authors formalize the simulation loop of the LRM using Evolving Algebra machines (EA-machines). All important constructs of the VHDL language are represented as records; processes are represented as concurrent agents running pseudo-codes, and the simulation control flow is passed to and fro between the kernel process (i.e, the simulation orchestrator) and the rest of the processes that execute the design behavior. This semantics implements closely the simulation loop as described in the LRM. Therefore, it is very rich and deals with most of the VHDL constructs, including the two time paradigms of the language. Moreover, the semantics works on fully-elaborated designs, therefore, component instantiation statements are omitted. However, a synchronous execution is fully expressible even if not explicitly embedded in the expression of the simulation loop.

In [20], the author presents a natural semantics for VHDL. The simulation loop is expressed by inference rules, and the execution of processes is based on the events over signals of their

corresponding sensitivity lists. The execution of statements computes transaction traces, that is, the projected waveforms for signals over the future of the simulation. The semantics deals both with unit and delta delays. Regarding our needs, this semantics covers the subset of the VHDL language that we are interested in, even if, it also covers some constructs pertaining to unit delays that are irrelevant to us (like wait and unit-delay signal assignment statements). A synchronous execution is expressible within the semantics, although it would be hidden in the inference rule formalizing the generic simulation loop. Also, the semantics does not provide its simulation loop with a simulation horizon (a maximum number of simulation cycle to be computed). The simulation ends when signal values evolve no more. The question of the influence of the environment, measured through the values of the primary inputs of a design, is not discussed.

In [10], the author presents an operational semantics for VHDL in the small-step style. The semantics follows closely the simulation cycle described in the LRM; however it is very concise and clear. The covered VHDL subset comprises arbitrary wait statements, and both unit and delta-delay signal assignments. There is an interested discussion about the non-determinism of VHDL, since it is a concurrent programming language: it entails that non-determinism is only existent at the processes level, that is, internal sequential statement of processes can be executed in an undeterministic manner (referred to by the author as *A* actions, that is, *internal* actions). But at every delta or time step (referred to as δ and *T* actions) of the execution, the design state can be computed in a deterministic manner, since all processes have reached a wait statement that stalled the execution of their inner body. The author is interested in the comparison of the behavior, and therefore, the equivalence between VHDL programs. He describes two strategies to compare VHDL programs. The first one is bisimulation; it is based on the comparison of the sequence of actions (either *A*, δ or *T* actions) performed by the two programs. The second one is observational equivalence; it is based on the observation of the value of the output signals of two VHDL programs (the observees), that receive values in their input signals from another VHDL program (the observer). The observer stimulates the entries of the observees and reaches a success state based on its observations of the value of the outputs. Regarding our needs, this semantics permits the description of our synchronous simulation cycle. However, like most of the semantics presented here, the component instantiation statement is not supported as it stands, but it is rather transformed into the equivalent processes statements. Small-step semantics is not needed in our case because we are only interested in the values of signals at the delta and time steps (for us, time steps correspond to clock events). We are not interested to capture the design states in the middle of the execution of a process body. We are more interested in "weak bisimulation", therefore forsaking the internal actions (i.e., *A* actions, execution of a process body that does not end in a wait statement) performed by a VHDL program. Indeed, a natural operational semantics in the style of Van Tassel's [20] is sufficient in our case. In [19], the authors extend the work of [10], especially by handling shared variables, in the presence of which a VHDL program can have a concrete underterministic behavior. The authors are also interested in the equivalence between two VHDL programs, and they are interested in determining a unique meaning property for VHDL programs. The unique meaning property states that the execution of a VHDL design in the presence of shared variables is unique. This work is interesting as it points out the fact that VHDL is not only subject to benign undeterminism. However, we are not interested in dealing

with constructs so advanced as shared variables or postponed processes, therefore, this work is not really relevant to us.

Translational semantics

Another kind of semantics, called “translational”, is interested in establishing a formal semantics for VHDL by translating a VHDL design into another formal model. Thus, the semantics of VHDL is modeled by the translation and the formal semantics of the target model. The target model has the ability to model concurrency, which is one of the specificity of VHDL. Moreover, target models are chosen because of the tools that they provide for analysis, and thus, a translational semantics for VHDL is often related to model checking considerations.

In [18], the author expresses the formal semantics of VHDL by translating a VHDL design into a corresponding flowgraph. All VHDL constructs, ranging from sequential statements to concurrent processes, are expressed with individual flowgraphs that are then composed together through their interfaces. The simulation cycle of VHDL is also encoded by means of connected flow graphs: one for the “execution part” of the semantics, that is, all processes run until blocked in a wait configuration, and one for the update part (i.e, the kernel process in the semantics of [4]). Flowgraphs come with a large amount of tools for analysis, and this translational semantics is involved in the setting of a framework to reason about VHDL programs using multiple technics (automatic theorem proving, model checking...). All these technics lean on the flowgraph formalism.

In [8], the author introduces a translational semantics for VHDL based on deterministic finite-state automaton. Again, the reason for using such automaton lies in the existence of many analysis tools. Moreover, forcing the generation of deterministic automaton improves the time execution of model-checking technics. The translation is performed on an elaborated VHDL design; a data space stores the values of signals and variables, and automaton represent the control-flow of VHDL statements. Each VHDL statement is associated to a specific automaton; sequence of statements are achieved by automaton composition. The simulation kernel is also represented by a specific automaton. Processes are composed together with respect to synchronization states, i.e. states that permit to pass the control from one process to another (for instance, after a wait statement), therefore achieving determinism in the control flow of the overall automaton.

In [15], the author presents a translation from VHDL to Coloured Petri Nets (CPNs) thus giving a formal semantics to VHDL constructs. The author approach to VHDL semantics is a strict translation of the “event-based” VHDL simulator by means of Petri nets. The author translates VHDL execution models (sea of processes) into CPNs, and also translates the kernel process into a CPN. The kernel process has previously been expressed as a VHDL process so that the translation into CPN is similar to the translation of other processes. Signals are not represented in the subnets, instead, three shared variables depict the signal states: one variable for the driving, one for the effective and one for the current value of a given signal. Colour domains of places in the subnets represent the different types of VHDL domains. Variables are represented by tokens. Values in drivers are represented by sequences of transactions (equivalent to waveforms); the author defines a set of functions that are convenient to handle sequences of transactions. Sequential statements are partitioned into two kinds: control flow

(if, loop, case...) and notation (operations on signals and variables) nets. Processes subnets are made by the fusing of each sequential statements in the process body. There is a special *Resume* place that can be set by the kernel process to resume the activity of a process. Concurrency is not discussed here, as the Petri net models are inherently concurrent models. The kernel process is a broad CPN having some of its places interfaced with the process subnets. The decoloration of the Petri net enables the analysis of the model and the detection of dead-locks.

In [7], the author gives a formal semantics to VHDL by transforming a VHDL design into an abstract machine, i.e defined by a set of inputs, outputs, states and transition function over states and outputs. The author is interested in the verification of properties over VHDL designs (temporal properties) or to prove equivalence between designs (bisimulation). To operate this transformation, only a subset of VHDL is considered, otherwise a finite-state representation is not reachable. The covered VHDL subset consists of objects with finite types, and no quantitative timing constructs (no after clause in signal assignments or *for* clause in wait statements). The author claims that a VHDL design is implemented by an abstract machine if they have the same observational behavior, i.e, for the same value in their inputs they yield the same values in their outputs. Each process statement part is transformed into a decision diagram (control flow graph); then, the decision diagram encodes the transition functions over states and outputs in the abstract machine implementing the corresponding process. Process statements are composed in relation to some composition operator. Moreover, the article lays out a method to transform a block statement into an abstract machine. The initiative is to be noticed as there are few papers of the VHDL semantics that are interested in such hierarchical constructs as block or component instantiation statements. The article concludes with an expression of the space of complexity entailed by the transformation of a VHDL design into an abstract machine.

Although the translational semantics described above meet most of the qualifying criterions in relation to our needs, we are not especially interested in implementing one of these. The main reason being that it would necessitate the implementation of the transformation from the abstract VHDL syntax to the target model in addition to the implementation of the semantics of the target model.

Table 1.1 summarizes the analysis of the VHDL semantics encountered during our literature review. Table 1.1 compares the different VHDL semantics in relation to the qualifying criterions (see Section 1.2.1).

		Fuchs and Mandler [9]	Breuer et al. [3]	Pandey et al. [17]	Borrione and Salem [1]	Breuer et al. [2]	Böger et al. [4]	Van Tassel [20]	Goossens [10]	Reetz and Kropf [18]	Döhnen and Hermann [8]	Olcoz [15]	Déharbe and Borrione [7]
Semantics	Kind	D	D, A	D	D	O	O	O	O	T	T	T	T
Description	Purpose	AR, ATP	AR	AR	AR	SS	SS	SS, ITP	SS, MC	ATP, MC, ITP	MC, ITP	MC	MC
Qualifying	Component Instantiation	T	T	T	N	T	T	T	T	T	T	T	N
Criteria	Synchronism	NE	NE	NE	Ex	E	E	E	E	E	E	E	NE
	Elaboration	×	×	×	✓	×	×	✓	×	×	×	×	✓
Extra. Informations.	Impl. Technology	Focus [6]	Prolog [5]	PVS [16]	?	Gofer [13]	?	HOL [11]	?	HOL [11]	?	?	?
	Particular Model or Data Types	Stream Processing	No	Interval Temporal Logic	No	Stream Processing	Evolving Algebra Machines	Natural Semantics (big-step)	Structural Semantics (small-step)	Flow Graphs	Finite-State Automata	Colored Petri Nets	Abstract Machines and Decision Diagrams

TABLE 1.1: A comparative summary on VHDL formal semantics.

- Kind : D (Denotational) - A (Axiomatic) - O (Operational) - T (Translational).
- Purpose : AR (Abstract Reasoning) - ATP (Automatic Theorem Proving) - SS (Simulator Specification) - ITP (Interactive Theorem Proving) - MC (Model Checking).
- Component Instantiation : T (statement is *Transformed* into equivalent processes) - N (statement is *Natively* taken into account in the semantics).
- Synchronism : E (Expressible within the semantics) - NE (Not Expressible within the semantics) - Ex (Explicitly built in the semantics).

To summarize, we are interesting in a semantics with an operational setting, built for the purpose of interactive theorem proving (ideally, with an existing implementation in the Coq proof assistant). Most important, the formal semantics must be able to deal with the expression of synchronous designs, that is, designs synchronized with a clock signal. Therefore, a synchronous simulation cycle must be at least expressible within the semantics. Moreover, the semantics must handle component instantiation statements as they are, that is, without transforming them into equivalent processes. As a bonus, the semantics should formalize the elaboration part of VHDL semantics.

In Table 1.1, cells are colored in green when the cell's content foster the adoption of the semantics, in yellow when the content does not go towards the adoption of the semantics but is not disqualifying, and red when the content is a disqualifying criterion. Cell are labelled in light grey when their content is neutral in relation to the semantics adoption. Now comparing the entries of Table 1.1 with the expression of our needs, we can discard the semantics with a cell labelled in red, that is most of the denotational semantics; moreover, all translational semantics are let aside for the reasons cited before. The candidate semantics are the operational semantics, plus the denotational semantics by Borrione and Salem [1], the only semantics that

formalizes an explicitly synchronous simulation cycle. The semantics that is the most likely to be adopted is the Borriore and Salem's semantics. However, we prefer an operational setting for our semantics because it is more fit to our task. To lower down the complexity of proofs, we really need a semantics that builds the synchronism into its simulation cycle, therefore putting aside all the intricacies of the full-blown VHDL simulation cycle. Moreover, the big-step style for an operational semantics is more relevant to us; as stated before, we are not interested in the intermediary states of computation that a small-step style semantics considers. Based on the observations, we have decided to formalize our own VHDL semantics inspired from the semantics of Borriore and Salem's [1] and Van Tassel's [20]. The following sections are dedicated to the presentation of the syntax and semantics of a subset of VHDL called \mathcal{H} -VHDL. \mathcal{H} -VHDL embeds the subset of VHDL that we are interested in when considering the VHDL designs generated by the HILECOP transformation.

1.3 Abstract syntax for \mathcal{H} -VHDL

Terminals of the language will be written in typewriter font, or will be enclosed in simple quotes for symbols with no typewriter representation. The a^* denotes a possibly empty repetition of the element a ; the a^+ denotes a non-empty repetition of a .

1.3.1 Design declaration

As in [20], we define the *design* construct in the \mathcal{H} -VHDL's abstract syntax which has no equivalent in the concrete syntax of VHDL.

```

design ::= design ide ida gens ports sigs cs
gens   ::= gdecl*
ports  ::= pdecl*
sigs   ::= sdecl*
```

1.3.2 Generic constant, port and internal signal declaration.

```

gdecl ::= (id,  $\tau$ , e)
pdecl ::= ((in|out), id,  $\tau$ )
sdecl ::= (id,  $\tau$ )
```

1.3.3 Concurrent statements.

```
cs ::= psstmt | cistmt | cs || cs | null
```

Process statement.

```

psstmt ::= process (idp, sl, vars, ss)
sl      ::= id*
vars    ::= vdecl*
vdecl   ::= (id,  $\tau$ )
```

Component instantiation statement.

$$\begin{aligned}
\text{cistmt} &::= \text{comp}(\text{id}_c, \text{id}_e, \text{gmap}, \text{ipmap}, \text{opmap}) \\
\text{gmap} &::= \text{assoc}_g^* \\
\text{ipmap} &::= \text{assoc}_{ip}^* \\
\text{opmap} &::= \text{assoc}_{op}^* \\
\text{assoc}_g &::= (\text{id}, e) \\
\text{assoc}_{ip} &::= (\text{name}, e) \\
\text{assoc}_{op} &::= (\text{id}, (\text{name} \mid \text{open})) \mid (\text{id}(e), \text{name})
\end{aligned}$$
1.3.4 Sequential statement.

$$\begin{aligned}
\text{ss} &::= \text{name} \Leftarrow e \mid \text{name} := e \mid \text{if}(e) \text{ ss} [\text{ss}] \mid \text{for}(\text{id}, e, e) \text{ ss} \\
&\quad \mid \text{falling ss} \mid \text{rising ss} \mid \text{rst ss ss}' \mid \text{ss}; \text{ss} \mid \text{null}
\end{aligned}$$
1.3.5 Expressions, names and types.

$$\begin{aligned}
e &::= e \text{ and } e \mid e \text{ or } e \mid \text{not } e \mid e = e \mid e \neq e \\
&\quad \mid e < e \mid e \leq e \mid e > e \mid e \geq e \mid e + e \mid e - e \\
&\quad \mid \text{name} \mid \text{natural} \mid \text{boolean} \mid (e^+) \\
\text{name} &::= \text{id} \mid \text{id}(e) \\
\text{boolean} &::= \text{true} \mid \text{false} \\
\tau &::= \text{boolean} \mid \text{natural}(e, e) \mid \text{array}(\tau, e, e)
\end{aligned}$$

Under the expression entry, the natural non-terminal represents the set of natural numbers (\mathbb{N}). The id non-terminal represents the set of identifiers.

1.4 Preliminaries to \mathcal{H} -VHDL semantics**1.4.1 Semantics Domains**

Let *id* denote the set of identifiers in the semantical domain. We write *prefix-id* to denote arbitrary subsets of the *id* set. The *type* and *value* semantical types are defined as follows:

$$\begin{aligned}
\text{type} &::= \text{bool} \mid \text{nat}(\text{nat}, \text{nat}) \mid \text{array}(\text{type}, \text{nat}, \text{nat}) \\
\text{value} &::= \text{bool} \mid \text{nat} \mid \text{array} \\
\text{bool} &::= \text{'}\top\text{' } \mid \text{'}\bot\text{' } \\
\text{nat} &::= 0 \mid 1 \mid \dots \mid \text{NATMAX} \\
\text{array} &::= (\text{value}^+)
\end{aligned}$$
TABLE 1.2: The *type* and *value* semantical types.

In Table 1.2, the *type* type is in any way similar to the τ entry of the abstract syntax, however, all constraint bounds in the *nat* and *array* types have been evaluated to natural numbers. NATMAX denotes the maximum value for a natural number. The NATMAX value depends on the implementation of the VHDL language; NATMAX must at least be equal to $2^{31} - 1$. Note that the *array* value contains at least one value as an array's index range contains at least one index (that is index 0).

Notation 1 (Partial functions). *Here, we present our notations pertaining to partial functions:*

- The \rightarrow arrow denotes a partial function.
- The \rightarrow denotes an application (i.e, a total function).
- For all $f \in A \rightarrow B$, $x \in f$ states that x is in the domain of function f .
- For all $f \in A \rightarrow B$ and $g \in A \rightarrow C$, $f \subseteq g$ states that the domain of f is a subset of the domain of g .
- For all $X \subset A$ and $f \in A \rightarrow B$, $X \subseteq f$ states that X is a subset of the domain of f .

Now, let us define the structure of an elaborated design which is a structure bound to a given \mathcal{H} -VHDL design and to a design store, i.e a global environment mapping identifiers to \mathcal{H} -VHDL designs. Only the designs referenced into the global design store can be instantiated as subcomponents of a given design. The elaborated design structure is used in the expression of the elaboration relation presented in Section **Elaboration rules**, as well as in the expression of the simulation rules. Let $ElDesign(d, \mathcal{D})$ be the set of the elaborated designs for a given \mathcal{H} -VHDL design d and a design store \mathcal{D} . An elaborated design is a composite environment built out of multiple sub-environments. Each sub-environment is a table, represented as a partial function, mapping identifiers of a certain category of constructs (e.g, input port identifiers) to their declaration information (e.g, type indication for input ports). We represent an elaborated design as a record where the fields are the sub-environments. An elaborated design is defined as follows:

Definition 1 (Elaborated Design). *For a given \mathcal{H} -VHDL design $d \in design$ s.t. $d = design\ id_{ent}\ id_{arch}\ gens\ ports\ sigs\ behavior$ and a given design store $\mathcal{D} \in entity-id \rightarrow design$, an elaborated design $\Delta \in ElDesign(d, \mathcal{D})$ is a record $\langle Gens, Ins, Outs, Sigs, Ps, Comps \rangle$ where:*

- $Gens \in generic-id \rightarrow (type \times value)$ where $generic-id = \{id \mid (id, \tau, e) \in gens\}$, is the partial function yielding the type and the value of generic constants.
- $Ins \in input-id \rightarrow type$ where $input-id = \{id \mid (in, id, \tau) \in ports\}$, is the partial function yielding the type of input ports.
- $Outs \in output-id \rightarrow type$ where $output-id = \{id \mid (out, id, \tau) \in ports\}$, the partial function yielding the type of output ports.
- $Sigs \in declared-signal-id \rightarrow type$ where $declared-signal-id = \{id \mid (id, \tau) \in sigs\}$, the partial function yielding the type of declared signals.

- $P_s \in \text{process-id} \rightarrow (\text{variable-id}(id_p) \rightarrow (\text{type} \times \text{value}))$ where $\text{process-id} = \{id_p \mid \text{process}(id_p, sl, vars, ss) \in \text{behavior}\}$, the partial function associating processes to their local environment. Local environments are functions mapping local variable identifiers to their corresponding type and value. Therefore, each set of local variable identifiers $\text{variable-id}(id_p)$ depends on the process identifier (represented by id_p) passed as the first argument of the P_s function.
- $\text{Comps} \in \text{component-id} \rightarrow \text{ElDesign}(d_e, \mathcal{D})$, where $\text{component-id} = \{id_c \mid \text{comp}(id_c, id_e, gm, ipm, opm) \in \text{behavior}\}$, the partial function mapping component instance ids to their elaborated design version. The set $\text{ElDesign}(d_e, \mathcal{D})$ depends on the design d_e from which the component identifier id_c , passed as the first argument of the Comps function, is an instance. Design d_e is retrieved from the design store \mathcal{D} s.t. $d_e = \mathcal{D}(id_e)$.

We assume that there are no overlapping between the identifiers of the sub-environments (i.e, an identifier belongs to at most one sub-environment). We note $\Delta(x)$ to denote the value returned for identifier x , where x is looked up in the appropriate field of Δ . We note $x \in \Delta$ to state that identifier x is defined in one of Δ 's fields. We note $\Delta(x) \leftarrow v$ the overriding of the value associated to identifier x with value v in the appropriate field of Δ , $\Delta \cup (x, v)$ to note the addition the mapping from identifier x to value v in the appropriate field of Δ , that assuming $x \notin \Delta$. We note $x \in \mathcal{F}(\Delta)$, where \mathcal{F} is a field of Δ , when more precision is needed regarding the lookup of identifier x in the record Δ .

Let $\Sigma(\Delta)$ be the set of design states for a given elaborated design Δ . A design state of Δ is defined as follows:

Definition 2 (Design state). A design state $\sigma \in \Sigma(\Delta)$, for a given design $d \in \text{design}$, a given design store \mathcal{D} and an elaborated design $\Delta \in \text{ElDesign}(d, \mathcal{D})$, is a record $\langle \mathcal{S}, \mathcal{C}, \mathcal{E} \rangle$ where:

- $\mathcal{S} \in \text{signal-id} \rightarrow \text{value}$, the partial function yielding the current values of the design's signals (ports and declared signals).
- $\mathcal{C} \in \text{component-id} \rightarrow \Sigma(\Delta_c)$, the partial function yielding the current state of design's component instances, where $\Delta_c = \Delta(id_c)$ and $id_c \in \text{component-id}$ is the component identifier passed to function \mathcal{C} .
- $\mathcal{E} \subseteq \text{signal-id} \cup \text{component-id}$, the set of signal and component instance ids that generated an event at the current simulation cycle.

The *signal-id* subset is the disjoint union of *input-id*, *output-id* and *declared-signal-id*. We use $\sigma(id)$ to denote the value associated to an identifier in the signal store \mathcal{S} or in the component store \mathcal{C} fields. $id \in_{\mathcal{E}} \sigma$ states that an identifier belongs to the event set \mathcal{E} , whereas $id \in \sigma$ states that an identifier is defined in either the signal store \mathcal{S} or the component store \mathcal{C} fields. $\sigma \cup_{\mathcal{E}} \{id\}$ denotes a new design state, in all points similar to σ but enriched with the identifier id in its events set.

Notation 2 (No events design state). For a given \mathcal{H} -VHDL design d , a design store \mathcal{D} , and an elaborated design $\Delta \in \text{ElDesign}(d, \mathcal{D})$, the function $\text{NoEv} \in \Sigma(\Delta) \rightarrow \Sigma(\Delta)$ returns a design state similar to the one passed in parameter only with an empty set of events. I.e, for all design state $\sigma \in \Sigma(\Delta)$ s.t. $\sigma = \langle \mathcal{S}, \mathcal{C}, \mathcal{E} \rangle$, $\text{NoEv}(\sigma) = \langle \mathcal{S}, \mathcal{C}, \emptyset \rangle$.

1.5 Elaboration rules.

The goal of the elaboration phase is to build an elaborated design Δ along with a *default* state σ_e , out of a \mathcal{H} -VHDL design d . The elaboration relation also covers type-checking operations for the declarative and behavioral parts of design d .

1.5.1 Design elaboration.

Definition 3 (Design store). $\mathcal{D} \in \text{entity-id} \rightarrow \text{design}$ is the partial function mapping entity ids to their corresponding representation in abstract syntax. As a prerequisite to the elaboration of HILECOP-generated designs (i.e, resulting from the transformation of a SITPN into an \mathcal{H} -VHDL design), a particular design store $\mathcal{D}_{\mathcal{H}}$ is defined. Design store $\mathcal{D}_{\mathcal{H}}$ holds the description of the *place* and *transition* designs which full definition in abstract syntax are given in appendices ?? and ??.

At the beginning of the elaboration phase, a function $\mathcal{M}_g \in \text{generic-id} \rightarrow \text{value}$ mapping the top-level design's generic constants to values is passed as an element of the environment. The \mathcal{M}_g function is referred to as the *dimensioning* function.

DESIGNELAB

$$\begin{array}{c}
 \Delta_{\emptyset}, \mathcal{M}_g \vdash \text{gens} \xrightarrow{egens} \Delta \\
 \Delta, \sigma_{\emptyset} \vdash \text{ports} \xrightarrow{eports} \Delta', \sigma \\
 \Delta', \sigma \vdash \text{sigs} \xrightarrow{esigs} \Delta'', \sigma' \\
 \mathcal{D}, \Delta'', \sigma' \vdash \text{cs} \xrightarrow{ebenh} \Delta''', \sigma'' \\
 \hline
 \mathcal{D}, \mathcal{M}_g \vdash \text{design id}_e \text{id}_a \text{gens ports sigs cs} \xrightarrow{elab} \Delta''', \sigma''
 \end{array}$$

where Δ_{\emptyset} denotes an empty elaborated design, that is an elaborated design initialized with empty fields (empty tables). In the same manner, σ_{\emptyset} denotes an empty design state. The effect of the *egens*, *eports*, *esigs* and *ebenh* that respectively deal with the elaboration of the generic constants, the ports, the architecture declarative part and the behavioral part of the design, are explicited in the following sections.

1.5.2 Generic clause elaboration.

Premises

- $e\text{type}_g$ transforms a subtype indication, specifically attached to a generic constant declaration, into a *type* instance and checks its well-formedness (see 1.5.5).
- The e relation evaluates expression e to value v (see 1.7.9).
- SE_l states that an expression is *locally* static (see 1.6.2).

- $v \in_c T$ and $\mathcal{M}(\text{id}_g) \in_c T$ checks that the default value and the value of yielded by the dimensioning function belongs to the type of the declared generic constant (see 1.6.1).

GENELABDIMEN

$$\frac{\vdash \tau \xrightarrow{\text{etype}_g} T \quad \Delta_\emptyset, \sigma_\emptyset, \Lambda_\emptyset \vdash e \xrightarrow{e} v \quad SE_l(e) \quad \mathcal{M}(\text{id}_g) \in_c T \quad \text{id}_g \notin \Delta}{\Delta, \mathcal{M} \vdash (\text{id}_g, \tau, e) \xrightarrow{\text{egens}} \Delta \cup (\text{id}_g, (T, \mathcal{M}(\text{id}_g)))} \quad \text{id}_g \in \mathcal{M}$$

GENELABDEFAULT

$$\frac{\vdash \tau \xrightarrow{\text{etype}_g} T \quad \Delta_\emptyset, \sigma_\emptyset, \Lambda_\emptyset \vdash e \xrightarrow{e} v \quad SE_l(e) \quad v \in_c T \quad \text{id}_g \notin \Delta}{\Delta, \mathcal{M} \vdash (\text{id}_g, \tau, e) \xrightarrow{\text{egens}} \Delta \cup (\text{id}_g, (T, v))} \quad \text{id}_g \notin \mathcal{M}$$

GENELABCOMP

$$\frac{\Delta, \mathcal{M} \vdash \text{gdecl} \xrightarrow{\text{egens}} \Delta' \quad \Delta', \mathcal{M} \vdash \text{gens} \xrightarrow{\text{egens}} \Delta''}{\Delta, \mathcal{M} \vdash \text{gdecl}; \text{gens} \xrightarrow{\text{egens}} \Delta''}$$

1.5.3 Port clause elaboration.

Premises

The *defaultv* relation yields the implicit *default* value for a given type.

INPORTELAB

$$\frac{\Delta \vdash \tau \xrightarrow{\text{etype}} T \quad \Delta \vdash T \xrightarrow{\text{defaultv}} v \quad \text{id} \notin \Delta \quad \text{id} \notin \sigma}{\Delta, \sigma \vdash (\text{in}, \text{id}, \tau) \xrightarrow{\text{eports}} \Delta \cup (\text{id}, T), \sigma \cup (\text{id}, v)}$$

OUTPORTELAB

$$\frac{\Delta \vdash \tau \xrightarrow{\text{etype}} T \quad \Delta \vdash T \xrightarrow{\text{defaultv}} v \quad \text{id} \notin \Delta \quad \text{id} \notin \sigma}{\Delta, \sigma \vdash (\text{out}, \text{id}, \tau) \xrightarrow{\text{eports}} \Delta \cup (\text{id}, T), \sigma \cup (\text{id}, v)}$$

PORTELABCOMP

$$\frac{\Delta, \sigma \vdash \text{pdecl} \xrightarrow{\text{eports}} \Delta', \sigma' \quad \Delta', \sigma' \vdash \text{ports} \xrightarrow{\text{eports}} \Delta'', \sigma''}{\Delta, \sigma \vdash \text{pdecl}; \text{ports} \xrightarrow{\text{eports}} \Delta'', \sigma''}$$

$$\frac{\text{SIGELAB} \quad \Delta \vdash \tau \xrightarrow{etype} T \quad \Delta \vdash T \xrightarrow{defaulttv} v \quad \begin{array}{l} \text{id} \notin \Delta \\ \text{id} \notin \sigma \end{array}}{\Delta, \sigma \vdash (\text{id}, \tau) \xrightarrow{esigs} \Delta \cup (\text{id}, T), \sigma \cup (\text{id}, v)}$$

1.5.5 Type indication elaboration.

$$\frac{\text{ETYPEBOOL}}{\Delta \vdash \text{boolean} \xrightarrow{etype} \text{bool}} \quad \frac{\text{ETYPENAT} \quad \Delta \vdash (\mathbf{e}, \mathbf{e}') \xrightarrow{econstr} (v, v')}{\Delta \vdash \text{natural}(\mathbf{e}, \mathbf{e}') \xrightarrow{etype} \text{nat}(v, v')}$$

Remark 1 (Type of constraints). *As the VHDL language reference stays unclear about the type of range and index constraints [VHDL00], we add the restriction that range and index constraints must have bounds of the `nat` type (i.e, value of type `nat`).*

Premises

SE_g states that an expression is *globally* static (see 1.6.2).

$$\begin{array}{c}
 \text{ECONSTR} \\
 \frac{\Delta \vdash SE_g(e) \quad \Delta, \sigma_\emptyset, \Lambda_\emptyset \vdash e \xrightarrow{e} v \quad v \in_c \text{nat}(0, \text{NATMAX}) \quad \Delta \vdash SE_g(e') \quad \Delta, \sigma_\emptyset, \Lambda_\emptyset \vdash e' \xrightarrow{e} v' \quad v' \in_c \text{nat}(0, \text{NATMAX})}{\Delta \vdash (e, e') \xrightarrow{econst} (v, v')} \quad v \leq v'
 \end{array}$$

When considering a type indication in a generic constant declaration, the definition of well-formedness differs slightly from the general definition. A type indication τ associated to a generic constant declaration is well-formed if τ denotes the boolean keyword, or the nat keyword with a *well-formed* constraint.

The $etype_g$ relation is specially defined to check the well-formedness of a subtype indication associated with a generic constant declaration.

$$\begin{array}{c}
 \text{ETYPEGBL} \qquad \text{ETYPEGNAT} \\
 \frac{}{\vdash \text{boolean} \xrightarrow{etype} \text{bool}} \quad \frac{\Delta \vdash (e, e') \xrightarrow{econst_g} (v, v')}{\vdash \text{natural}(e, e') \xrightarrow{etype} \text{nat}(v, v')}
 \end{array}$$

The $econst_g$ relation checks that a *generic* constraint (i.e, a constraint appearing in a type indication associated with a generic constant declaration) is well-formed and evaluates the constraint bounds.

A *generic* constraint c is well-formed if:

- its bounds are locally static expressions [VHDL00] of the nat type.
- its lower bound value is inferior or equal to its upper bound value.

$$\begin{array}{c}
 \text{ECONSTRG} \\
 \frac{SE_l(e) \quad \Delta_\emptyset, \sigma_\emptyset, \Lambda_\emptyset \vdash e \xrightarrow{e} v \quad v \in_c \text{nat}(0, \text{NATMAX}) \quad SE_l(e') \quad \Delta_\emptyset, \sigma_\emptyset, \Lambda_\emptyset \vdash e' \xrightarrow{e} v' \quad v' \in_c \text{nat}(0, \text{NATMAX})}{\vdash (e, e') \xrightarrow{econst_g} (v, v')} \quad v \leq v'
 \end{array}$$

1.5.6 Behavior elaboration.

Elaboration of concurrent statements.

As it is, the elaboration of the composition of concurrent statements is performed in a sequential manner.

CS_{PARELAB}

$$\frac{\mathcal{D}, \Delta, \sigma \vdash \text{cs} \xrightarrow{ebch} \Delta', \sigma' \quad \mathcal{D}, \Delta', \sigma' \vdash \text{cs}' \xrightarrow{ebch} \Delta'', \sigma''}{\mathcal{D}, \Delta, \sigma \vdash \text{cs} \parallel \text{cs}' \xrightarrow{ebch} \Delta'', \sigma''}$$

CS_{NULLLAB}

$$\frac{}{\mathcal{D}, \Delta, \sigma \vdash \text{null} \xrightarrow{ebch} \Delta, \sigma}$$

Process elaboration.**Premises**

valid_{ss} states that a sequential statement is well-typed.

Side conditions

$sl \subseteq \text{Ins}(\Delta) \cup \text{Sigs}(\Delta)$ indicates that the sensitivity list sl must only contain signal identifiers that are readable, that is, *input* ports and declared signals.

PS_{ELAB}

$$\frac{\Delta, \Lambda_{\emptyset} \vdash \text{vars} \xrightarrow{evars} \Lambda \quad \Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(ss) \quad \text{id}_p \notin \Delta}{\mathcal{D}, \Delta, \sigma \vdash \text{process}(\text{id}_p, sl, \text{vars}, ss) \xrightarrow{ebch} \Delta \cup (\text{id}_p, \Lambda), \sigma \quad sl \subseteq \text{Ins}(\Delta) \cup \text{Sigs}(\Delta)}$$

Process declarative part elaboration.

The *evars* relation builds a local environment out of a process declarative part.

VAR_{ELAB}

$$\frac{\Delta \vdash \tau \xrightarrow{etype} T \quad \vdash T \xrightarrow{defaultv} v \quad \text{id} \notin \Lambda, \text{id} \notin \Delta}{\Delta, \Lambda \vdash (\text{id}, \tau) \xrightarrow{evars} \Lambda \cup (\text{id}, (T, v))}$$

VAR_{ELABCOMP}

$$\frac{\Delta, \Lambda \vdash \text{vdecl} \xrightarrow{evars} \Lambda' \quad \Delta, \Lambda' \vdash \text{vars} \xrightarrow{evars} \Lambda''}{\Delta, \Lambda \vdash \text{vdecl}, \text{vars} \xrightarrow{evars} \Lambda''}$$

Component instantiation elaboration.**Premises**

- The *emapg* relation binds a generic map to a function $\mathcal{M} \in id \nrightarrow value$ (see definition below).
- valid_{ipm} (resp. valid_{opm}) states that a in port map (resp. out port map) is valid, i.e. well-formed and well-typed (see 1.6.3).

Side conditions

$\mathcal{M} \subseteq \text{Gens}(\Delta_c)$ checks that the generic map gmap contains references to known generic constant identifiers only.

COMPELAB

$$\frac{\mathcal{M}_\emptyset \vdash \text{gmap} \xrightarrow{\text{emapg}} \mathcal{M} \quad \mathcal{D}, \mathcal{M} \vdash \mathcal{D}(\text{id}_e) \xrightarrow{\text{elab}} \Delta_c, \sigma_c \quad \begin{array}{l} \Delta, \Delta_c, \sigma \vdash \text{valid}_{ipm}(\text{ipmap}) \\ \Delta, \Delta_c \vdash \text{valid}_{opm}(\text{opmap}) \end{array} \quad \begin{array}{l} \text{id}_c \notin \Delta, \text{id}_c \notin \sigma \\ \text{id}_e \in \mathcal{D} \end{array}}{\mathcal{D}, \Delta, \sigma \vdash \text{comp}(\text{id}_c, \text{id}_e, \text{gmap}, \text{ipmap}, \text{opmap}) \xrightarrow{\text{ebch}} \Delta \cup (\text{id}_c, \Delta_c), \sigma \cup (\text{id}_c, \sigma_c) \quad \mathcal{M} \subseteq \text{Gens}(\Delta_c)}$$

A port map is a mapping between signals coming from an embedding design (Δ) and ports of an internal component instance (Δ_c). The formal part of a port map entry (i.e, left of the arrow) belongs to the internal component, whereas the actual part (i.e, right of the arrow) refers to the embedding design. Therefore, we need both Δ and Δ_c to verify if a port map is well-typed leveraging the valid_{pm} predicate.

Remark 2 (Valid generic map). *Note that we are not checking the validity of the generic map. In case of an ill-formed generic map, a inconsistent mapping \mathcal{M} is generated by the emapg that will make the elab relation, taking \mathcal{M} as a parameter, never derivable. Therefore, the elab relation does an implicit validity check on the generic map.*

ASSOCGELAB

GMAPELABCOMP

$$\frac{SE_l(e) \quad \Delta_\emptyset, \sigma_\emptyset, \Lambda_\emptyset \vdash e \xrightarrow{e} v \quad \text{id}_g \notin \mathcal{M}}{\mathcal{M} \vdash \text{id}_g \Rightarrow e \xrightarrow{\text{emapg}} \mathcal{M} \cup (\text{id}_g, v)} \quad \frac{\mathcal{M} \vdash \text{assoc}_g \xrightarrow{\text{emapg}} \mathcal{M}' \quad \mathcal{M}' \vdash \text{gmap} \xrightarrow{\text{emapg}} \mathcal{M}''}{\mathcal{M} \vdash \text{assoc}_g, \text{gmap} \xrightarrow{\text{emapg}} \mathcal{M}''}$$

An assoc_g entry doesn't allow indexed id in its formal part, due to the restriction of generic constants to scalar types. Note that this restriction is not imposed by the VHDL language reference; it is our stance towards a simplification of the VHDL semantics.

1.5.7 Implicit default value.

According to the VHDL reference, when declaring a port, a signal or a variable, these items must receive an implicit default value depending on their types [VHDL00]. The defaultv relation determines the default value for a given type.

$$\frac{\text{DEFAULTVBOOL}}{\text{bool} \xrightarrow{\text{defaultv}} \perp} \quad \frac{\text{DEFAULTVCNAT}}{\text{nat}(n, m) \xrightarrow{\text{defaultv}} n} \quad n \leq m$$

$$\frac{\text{DEFAULTVCARR} \quad T \xrightarrow{\text{default}v} v}{\text{array}(T, n, m) \xrightarrow{\text{default}v} \text{create_array}(\text{size}, T, v)} \quad \begin{array}{l} n \leq m \\ \text{size} = (m - n) + 1 \end{array}$$

$\text{create_array}(\text{size}, T, v)$ creates an array of size size , containing element of type T , where each element is initialized with the value v .

1.6 Static type-checking rules.

1.6.1 Typing relation.

$$\frac{\text{ISBOOL}}{b \in_c \text{bool}} \quad b \in \mathbb{B} \quad \frac{\text{ISCNAT}}{n \in_c \text{nat}(l, u)} \quad n \in [l, u] \quad \frac{\text{ARRAY} \quad v_i \in_c T \quad i = 1, \dots, n}{\Delta \vdash (v_1, \dots, v_n) \in_c \text{array}(T, l, u)} \quad n = (u - l) + 1$$

1.6.2 Static expressions.

Static expressions are either locally static or globally static.

Locally static expressions.

An expression is *locally* static if:

- It is composed of operators and operands of a *scalar* type (i.e, natural or boolean).
- It is a *literal* of a scalar type.

$$\frac{\text{LSENOT}}{SE_l(n)} \quad n \in \mathbb{N} \quad \frac{\text{LSEBOOL}}{SE_l(b)} \quad b \in \mathbb{B} \quad \frac{\text{LSENOT}}{SE_l(\text{not } e)} \quad SE_l(e) \quad \frac{\text{LSEBINOP}}{SE_l(e \text{ op } e')} \quad SE_l(e) \quad SE_l(e') \quad \text{op} \in \{ +, -, =, \neq, <, \leq, >, \geq, \text{and}, \text{or} \}$$

Globally static expressions.

An expression is *globally* static in the context Δ if:

- It is a generic constant.
- It is an array aggregate composed of globally static expressions.
- It is a locally static expression.

$$\frac{\text{GSELOCAL}}{\Delta \vdash SE_g(e)} \quad SE_l(e) \quad \frac{\text{GSEGEN}}{\Delta \vdash SE_g(\text{id}_g)} \quad \text{id}_g \in \text{Gens}(\Delta) \quad \frac{\text{GSEAGGREGATE}}{\Delta \vdash SE_g((e_1, \dots, e_n))} \quad \Delta \vdash SE_g(e_i) \quad i = 1, \dots, n$$

1.6.3 Valid port map.

The valid_{ipm} predicate states that an *in* port map is valid in the context Δ, Δ_c , where Δ is the embedding design structure and Δ_c denotes the component instance owner of the port map, if:

- All ports defined in Δ_c are exactly mapped once in the portmap.
- For each port map entry, the formal and actual part are of the same type.

Premises

- list_{ipm} builds a set $\mathcal{L} \subset id \sqcup (id \times \mathbb{N})$ out of the port map association list.
- check_{pm} checks the validity of a port map based on the corresponding port list and the set built by the list_{ipm} relation.

$$\frac{\text{VALIDIPM} \quad \Delta, \Delta_c, \sigma, \mathcal{L}_\emptyset \vdash \text{ipmap} \xrightarrow{\text{list}_{ipm}} \mathcal{L} \quad \text{check}_{pm}(\text{Ins}(\Delta_c), \mathcal{L})}{\Delta, \Delta_c, \sigma \vdash \text{valid}_{ipm}(\text{ipmap})}$$

Side conditions

$\text{id}_p \notin \mathcal{L}$ checks that the port identifier id_p is not already mapped.

LISTIPMSIMPLE

$$\frac{\Delta, \sigma \vdash e \xrightarrow{e} v \quad v \in_c T}{\Delta, \Delta_c, \sigma, \mathcal{L} \vdash \text{id}_p \Rightarrow e \xrightarrow{\text{list}_{ipm}} \mathcal{L} \cup \{\text{id}_p\}} \quad \begin{array}{l} \text{id}_p \notin \mathcal{L}, \text{id}_p \in \text{Ins}(\Delta_c) \\ \Delta_c(\text{id}_p) = T \end{array}$$

Premises

$v_i \in_c \text{nat}(n, m)$ checks that the index value stays in the array bounds.

Side conditions

$\text{id}_p \notin \mathcal{L}$ and $(\text{id}_p, v_i) \notin \mathcal{L}$ checks that neither the port identifier id_p nor the couple port identifier id_p and index v_i are already mapped.

LISTIPMPARTIAL

$$\begin{array}{c}
\frac{SE_I(e_i) \quad \Delta, \sigma \vdash e \xrightarrow{e} v \quad \begin{array}{l} \vdash e_i \xrightarrow{e} v_i \quad v_i \in_c \mathbf{nat}(n, m) \\ v \in_c T \end{array} \quad \begin{array}{l} \text{id}_p \notin \mathcal{L}, (\text{id}_p, v_i) \notin \mathcal{L} \\ \text{id}_p \in \text{Ins}(\Delta_c) \end{array}}{\Delta, \Delta_c, \sigma, \mathcal{L} \vdash \text{id}_p(e_i) \Rightarrow e \xrightarrow{\text{list}_{ipm}} \mathcal{L} \cup \{ (\text{id}_p, v_i) \} \quad \Delta_c(\text{id}_p) = \text{array}(T, n, m)}
\end{array}$$

LISTIPMCONS

$$\frac{\Delta, \Delta_c, \sigma, \mathcal{L} \vdash \text{assoc}_{ip} \xrightarrow{\text{list}_{ipm}} \mathcal{L}' \quad \Delta, \Delta_c, \sigma, \mathcal{L}' \vdash \text{ipmap} \xrightarrow{\text{list}_{ipm}} \mathcal{L}''}{\Delta, \Delta_c, \sigma, \mathcal{L} \vdash \text{assoc}_{ip}, \text{ipmap} \xrightarrow{\text{list}_{ipm}} \mathcal{L}''}$$

The $\text{check}_{pm}(\text{Ports}, \mathcal{L})$ predicate states that all port identifiers referenced in the domain of $\text{Ports} \in id \multimap \text{type}$ appear in \mathcal{L} as a simple identifier, or if the port identifier is of type array, then all couples (id, i) must belong to \mathcal{L} , where i denotes all indexes of the array range and id , the port id.

$$\text{check}_{pm}(\text{Ports}, \mathcal{L}) \equiv \forall \text{id}_p \in \text{dom}(\text{Ports}), \text{id}_p \in \mathcal{L} \vee (\text{Ports}(\text{id}_p) = \text{array}(T, n, m) \wedge \forall i \in [n, m], (\text{id}_p, i) \in \mathcal{L})$$

The valid_{opm} predicate states that an *output* port map is valid in the context Δ, Δ_c , where Δ is the embedding design structure and Δ_c denotes the component instance owner of the port map, if:

- An output port appears at most once in the output port map.
- Two different output ports cannot be connected to the same signal.
- For each port map entry, the formal and the actual part are of the exact same type (i.e, in the sense of the Leibniz equality).

We allow partially connected output port map; i.e, an output port map where all output ports might not be present in the mapping. Such output ports are open by default.

Premises

list_{opm} builds two sets $\mathcal{L}, \mathcal{L}_{ids} \subseteq id \sqcup (id \times \mathbb{N})$ out of the port map opmap . \mathcal{L}_{ids} is built incrementally to check that there are no multiply-driven signals resulting of the port map connection.

VALIDOPM

$$\frac{\Delta, \Delta_c, \mathcal{L}_\emptyset, \mathcal{L}_{ids\emptyset} \vdash \text{opmap} \xrightarrow{\text{list}_{opm}} \mathcal{L}, \mathcal{L}_{ids}}{\Delta, \Delta_c \vdash \text{valid}_{opm}(\text{opmap})}$$

Side conditions

- $\text{id}_p \notin \mathcal{L}$ checks that the port identifier id_p is not already mapped (i.e, is not already in the formal part of the port map).
- $\text{id}_s \notin \mathcal{L}_{ids}$ checks that the signal identifier id_s is not already mapped (i.e, is not already in the actual part of the port map).
- $\Delta_c(\text{id}_p) = \Delta(\text{id}_s) = T$ checks that id_p and id_s are exactly of the same type.

LISTOPMSIMPLETOSIMPLE

$$\frac{\Delta, \Delta_c, \mathcal{L}, \mathcal{L}_{ids} \vdash \text{id}_p \Rightarrow \text{id}_s \xrightarrow{\text{list}_{opm}} \mathcal{L} \cup \{\text{id}_p\}, \mathcal{L}_{ids} \cup \{\text{id}_s\}}{\begin{array}{l} \text{id}_p \notin \mathcal{L}, \text{id}_s \notin \mathcal{L}_{ids} \\ \text{id}_p \in \text{Outs}(\Delta_c) \\ \text{id}_s \in \text{Sigs}(\Delta) \cup \text{Outs}(\Delta) \\ \Delta_c(\text{id}_p) = \Delta(\text{id}_s) = T \end{array}}$$

Side conditions

$\text{Outs}_c(\text{id}_p) = T$ and $\text{Sigs}(\text{id}_s) = \text{array}(T, n, m)$ checks that the type of id_p and the type of the elements of id_s are the same. Note that id_s must denote an array as id_p is mapped to one of id_s 's partial.

LISTOPMSIMPLETOPARTIAL

$$\frac{\begin{array}{l} SE_l(e_i) \quad \vdash e_i \xrightarrow{e} v_i \quad v_i \in_c \text{nat}(n, m) \end{array}}{\Delta, \Delta_c, \mathcal{L}, \mathcal{L}_{ids} \vdash \text{id}_p \Rightarrow \text{id}_s(e_i) \xrightarrow{\text{list}_{opm}} \mathcal{L} \cup \{\text{id}_p\}, \mathcal{L}_{ids} \cup \{(\text{id}_s, v_i)\}} \quad \begin{array}{l} \text{id}_p \notin \mathcal{L}, \text{id}_s, (\text{id}_s, v_i) \notin \mathcal{L}_{ids} \\ \text{id}_p \in \text{Outs}(\Delta_c) \\ \text{id}_s \in \text{Sigs}(\Delta) \cup \text{Outs}(\Delta) \\ \Delta_c(\text{id}_p) = T \\ \Delta(\text{id}_s) = \text{array}(T, n, m) \end{array}}$$

LISTOPMSIMPLETOOPEN

$$\frac{\Delta, \Delta_c, \mathcal{L}, \mathcal{L}_{ids} \vdash \text{id}_p \Rightarrow \text{open} \xrightarrow{\text{list}_{opm}} \mathcal{L} \cup \{\text{id}_p\}, \mathcal{L}_{ids}}{\begin{array}{l} \text{id}_p \notin \mathcal{L} \\ \text{id}_p \in \text{Outs}(\Delta_c) \end{array}}$$

Remark 3 (Unconnected output port.). We forbid the case where an indexed formal part corresponding to the subelement of a composite output port is unconnected, i.e $\text{id}_p(e_i) \Rightarrow \text{open}$, as it could lead to the case where some subelements of a composite output port are connected while others are not (error case in [VHDL00]).

LISTOPMPARTIALTOSIMPLE

$$\frac{\begin{array}{l} SE_l(e_i) \quad \vdash e_i \xrightarrow{e} v_i \quad v_i \in_c \text{nat}(n, m) \end{array}}{\Delta, \Delta_c, \mathcal{L} \vdash \text{id}_p(e_i) \Rightarrow \text{id}_s \xrightarrow{\text{list}_{opm}} \mathcal{L} \cup \{(\text{id}_p, v_i)\}, \mathcal{L}_{ids} \cup \{\text{id}_s\}} \quad \begin{array}{l} \text{id}_p, (\text{id}_p, v_i) \notin \mathcal{L}, \text{id}_s \notin \mathcal{L}_{ids} \\ \text{id}_p \in \text{Outs}(\Delta_c) \\ \text{id}_s \in \text{Sigs}(\Delta) \cup \text{Outs}(\Delta) \\ \Delta_c(\text{id}_p) = \text{array}(T, n, m) \\ \Delta(\text{id}_s) = T \end{array}$$

LISTOPMPARTIALTOPARTIAL

$$\begin{array}{c}
\begin{array}{c}
SE_l(e'_i) \quad \vdash e'_i \xrightarrow{e} v'_i \quad v'_i \in_c \mathbf{nat}(n', m') \\
SE_l(e_i) \quad \vdash e_i \xrightarrow{e} v_i \quad v_i \in_c \mathbf{nat}(n, m)
\end{array} \\
\hline
\Delta, \Delta_c, \mathcal{L} \vdash \text{id}_p(e_i) \Rightarrow \text{id}_s(e'_i) \xrightarrow{\text{list}_{opm}} \mathcal{L} \cup \{(id_p, v_i)\}, \mathcal{L}_{ids} \cup \{(id_s, v'_i)\}
\end{array}
\quad
\begin{array}{l}
\text{id}_p, (id_p, v_i) \notin \mathcal{L}, \text{id}_s, (id_s, v'_i) \notin \mathcal{L}_{ids} \\
\text{id}_p \in \text{Outs}(\Delta_c) \\
\text{id}_s \in \text{Sigs}(\Delta) \cup \text{Outs}(\Delta) \\
\Delta_c(\text{id}_p) = \mathbf{array}(T, n, m) \\
\Delta(\text{id}_s) = \mathbf{array}(T, n', m')
\end{array}$$

LISTOPMCONS

$$\begin{array}{c}
\Delta, \Delta_c, \mathcal{L}, \mathcal{L}_{ids} \vdash \text{assoc}_{po} \xrightarrow{\text{list}_{opm}} \mathcal{L}', \mathcal{L}'_{ids} \quad \Delta, \Delta_c, \mathcal{L}', \mathcal{L}'_{ids} \vdash \text{opmap} \xrightarrow{\text{list}_{opm}} \mathcal{L}'', \mathcal{L}''_{ids} \\
\hline
\Delta, \Delta_c, \mathcal{L}, \mathcal{L}_{ids} \vdash \text{assoc}_{po}, \text{opmap} \xrightarrow{\text{list}_{opm}} \mathcal{L}'', \mathcal{L}''_{ids}
\end{array}$$

1.6.4 Valid sequential statements.

The valid_{ss} predicate states that a sequential statement is well-typed.

Well-typed signal assignment.

WELLTYPEDSIGASSIGN

$$\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad v \in_c T \quad \text{id}_s \in \text{Sigs}(\Delta) \cup \text{Outs}(\Delta)}{\Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(\text{id}_s \leftarrow e) \quad \Delta(\text{id}_s) = T}$$

WELLTYPEDIDXSIGASSIGN

$$\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad v \in_c T \quad \Delta, \sigma, \Lambda \vdash e_i \xrightarrow{e} v_i \quad v_i \in_c \mathbf{nat}(n, m) \quad \text{id}_s \in \text{Sigs}(\Delta) \cup \text{Outs}(\Delta)}{\Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(\text{id}_s(e_i) \leftarrow e) \quad \Delta(\text{id}_s) = \mathbf{array}(T, n, m)}$$

Well-typed variable assignment.

WELLTYPEDVARASSIGN

$$\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad v \in_c T \quad \text{id}_v \in \Lambda}{\Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(\text{id}_v := e) \quad \Lambda(\text{id}_v) = (T, \text{val})}$$

WELLTYPEDIDXVARASSIGN

$$\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad v \in_c T \quad \Delta, \sigma, \Lambda \vdash e_i \xrightarrow{e} v_i \quad v_i \in_c \mathbf{nat}(n, m) \quad \text{id}_v \in \Lambda}{\Delta, \Lambda \vdash \text{valid}_{ss}(\text{id}_v(e_i) := e) \quad \Lambda(\text{id}_v) = (\mathbf{array}(T, n, m), \text{val})}$$

Well-typed if statements.

WELLTYPEDIF

$$\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad v \in_c \mathbf{bool} \quad \Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(\text{ss})}{\Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(\text{if } (e) \text{ ss})}$$

WELLTYPEDIFELSE

$$\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad v \in_c \mathbf{bool} \quad \Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(\text{ss}) \quad \Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(\text{ss}')}{\Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(\text{if } (e) \text{ ss ss'})}$$

Well-typed loop statement.

WELLTYPEDLOOP

$$\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad v \in_c \text{nat}(0, \text{NATMAX}) \quad \Delta, \sigma, \Lambda \vdash e' \xrightarrow{e} v' \quad v' \in_c \text{nat}(0, \text{NATMAX}) \quad \Delta, \sigma, \Lambda' \vdash \text{valid}_{ss}(ss)}{\Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(\text{for}(\text{id}_v, e, e') ss)} \quad \Lambda' = \Lambda \cup (\text{id}_v, (\text{nat}(v, v'), v))$$

Well-typed rising and falling edge blocks

WELLTYPEDRISING

$$\frac{\Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(ss)}{\Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(\text{rising } ss)}$$

WELLTYPEDFALLING

$$\frac{\Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(ss)}{\Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(\text{falling } ss)}$$

Well-typed rst blocks

WELLTYPEDRST

$$\frac{\Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(ss) \quad \Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(ss')}{\Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(\text{rst } ss \ ss')}$$

Well-typed null statement

WELLTYPEDNULL

$$\frac{}{\Delta, \sigma, \Lambda \vdash \text{valid}_{ss}(\text{null})}$$

1.7 Simulation rules.**1.7.1 Full Simulation**

The full simulation process is decomposed in two steps. The first step is the elaboration phase that builds an elaborated version of a \mathcal{H} -VHDL design along with its default state, and type-checks the design. Previous to the elaboration phase, the top-level design receives a value for each of its generic constant; we refer to it as the *dimensioning* of the top-level design. The second step is the simulation phase that executes the behavioral part of the top-level design starting from an initial state. The simulation is decomposed into simulation cycles. Each simulation cycle is divided in four parts entailed by the *synchronous* execution of \mathcal{H} -VHDL top-level designs, i.e designs whose behavior depend on a clock signal. The four parts are, first, the execution of concurrent statements responding to the rising edge of the clock signal, then, a phase of signal stabilization followed by the execution of concurrent statements responding to the falling edge of the clock signal, and finally another phase of signal stabilization. At each clock event, the

value of the primary inputs of the design being currently simulated are captured and injected in the simulation; primary inputs receive values from the design environment. Here, the environment is represented by a function mapping input port identifiers to values depending on the current count of simulation cycles and the considered clock event. This leads to the following hypothesis:

Hypothesis 1 (Stable primary inputs). *The values of primary inputs (i.e, input ports of the top-level design) are captured at each clock event, and therefore are stable (i.e, their values do not change) between two contiguous clock events.*

Hypothesis **Stable primary inputs** arises from the fact that the clock signal sample rate respects the Nyquist-Shannon sampling theorem. Therefore, the sample rate of the design's clock is sufficient to capture all events possibly arising in the environment. We only need to settle the values of the primary inputs at the clock edges.

Also, after each clock event phase follows a signal stabilization phase in the proceedings of a simulation cycle. One more hypothesis is needed here:

Hypothesis 2 (Stabilization). *All signals have enough time to stabilize during the signal stabilization phase that happens between two clock events.*

As a \mathcal{H} -VHDL design represents a physical circuit, one can assume that the represented circuit is analyzed former to the simulation. Therefore, one knows exactly how much time is needed to propagate signal values through the longest physical path; as a consequence, a proper clock frequency is set ensuring signal stabilization between two clock events. Thus, Hypothesis **Stabilization** arises from the previous facts.

The *full* simulation relation takes in parameter a top-level design d , a design store $\mathcal{D} \in id \rightarrow design$, an elaborated design $\Delta \in ElDesign(d)$, a dimensioning function $\mathcal{M}_g \in Gens(\Delta) \rightarrow value$, a primary input environment $E_p \in (\mathbb{N} \times Clk) \rightarrow (Ins(\Delta) \rightarrow value)$, a simulation cycle count $\tau \in \mathbb{N}$, and a simulation trace $\theta \in list(\Sigma(\Delta))$, corresponding to the list of states yielded by design d after τ simulation cycles. Note that we use the pointed notation to access the behavioral part of design d , written $d.cs$. It is this part of the design that is executed during the simulation, and therefore is passed as a parameter of the initialization and simulation relations.

FULLSIM

$$\frac{\mathcal{D}, \mathcal{M}_g \vdash d \xrightarrow{elab} \Delta, \sigma \quad \mathcal{D}, \Delta, \sigma \vdash d.cs \xrightarrow{init} \sigma_0 \quad \mathcal{D}, E_p, \Delta, \tau, \sigma_0 \vdash d.cs \rightarrow \theta}{\mathcal{D}, \Delta, \mathcal{M}_g, E_p, \tau \vdash d \xrightarrow{full} (\sigma_0 :: \theta)}$$

where:

- $\mathcal{M}_g \in Gens(\Delta) \rightarrow value$, the function yielding the values of generic constants for a given top-level design, referred to as the *dimensioning* function.
- $E_p \in (\mathbb{N} \times Clk) \rightarrow (ident \rightarrow value)$, the function yielding a mapping from primary inputs (i.e, input ports of the top-level design) to values at a given simulation cycle count (i.e, the \mathbb{N} argument), and a given clock event (i.e, the Clk argument, where $Clk = \{\uparrow, \downarrow\}$).
- τ , the number of simulation cycles to execute. The value of τ is decremented at each clock cycle until it reaches zero (see Section 1.7.2).

1.7.2 Simulation loop.

$$\begin{array}{c}
 \text{SIMEND} \\
 \hline
 \mathcal{D}, E_p, \Delta, 0, \sigma \vdash cs \rightarrow []
 \end{array}
 \quad
 \begin{array}{c}
 \text{SIMLOOP} \\
 \hline
 \mathcal{D}, E_p, \Delta, \tau, \sigma \vdash cs \xrightarrow{\uparrow, \downarrow} \sigma', \sigma'' \quad \mathcal{D}, E_p, \Delta, \tau - 1, \sigma'' \vdash cs \rightarrow \theta
 \end{array}
 \quad
 \frac{}{\mathcal{D}, E_p, \Delta, \tau, \sigma \vdash cs \rightarrow (\sigma' :: \sigma'' :: \theta)} \tau > 0$$

1.7.3 Simulation cycle.

To ease the reading of forward simulation rules, we need to introduce two notations.

Notation 3 (Overriding union). For all partial function $f, f' \in X \rightharpoonup Y$, $f \overset{\leftarrow}{\cup} f'$ denotes the overriding

$$\text{union of } f \text{ and } f' \text{ such that } f \overset{\leftarrow}{\cup} f'(x) = \begin{cases} f'(x) & \text{if } x \in \text{dom}(f') \\ f(x) & \text{otherwise} \end{cases}$$

Notation 4 (Differentiated intersection domain). For all partial function $f, f' \in X \rightharpoonup Y$, $f \overset{\neq}{\cap} f'$ denotes the intersection of the domain of f and f' for which f and f' yields different values. That is, $f \overset{\neq}{\cap} f' = \{ x \in \text{dom}(f) \cap \text{dom}(f') \mid f(x) \neq f'(x) \}$.

Definition 4 (Input port values update). Given an \mathcal{H} -VHDL design $d \in \text{design}$, a design store $\mathcal{D} \in \text{id} \rightarrow \text{design}$, an elaborated design $\Delta \in \text{ElDesign}(d, \mathcal{D})$, a simulation environment $E_p \in (\mathbb{N} \times \{\uparrow, \downarrow\}) \rightarrow (\text{Ins}(\Delta) \rightarrow \text{value})$, let us define the relation expressing the update of the values of the input ports of Δ at a given design state $\sigma \in \Sigma(\Delta)$, clock cycle count $\tau \in \mathbb{N}$, and clock event $\text{clk} \in \{\uparrow, \downarrow\}$, and thus resulting in a new state $\sigma_i \in \Sigma(\Delta)$. The relation is written $\text{Inject}_{\text{clk}}(\sigma, E_p, \tau, \sigma_i)$ and verifies that: $\sigma = \langle \mathcal{S}, \mathcal{C}, \mathcal{E} \rangle$ and $\sigma_i = \langle \mathcal{S} \overset{\leftarrow}{\cup} E_p(\tau, \text{clk}), \mathcal{C}, \mathcal{E} \rangle$.

The cycle relation states that the design state σ'' is the result of the execution of concurrent statement cs over one simulation cycle, this starting from state σ . As told in Hypothesis 1, we update the input port values at each clock event. New input port values are coming from the environment E_p . The updates are made in the definitions of σ_i and σ'_i . These definitions are expressed as side conditions.

$$\begin{array}{c}
 \text{SIMCYC} \\
 \hline
 \mathcal{D}, \Delta, \sigma_i \vdash cs \xrightarrow{\uparrow} \sigma_{\uparrow} \quad \mathcal{D}, \Delta, \sigma_{\uparrow} \vdash cs \xrightarrow{\rightsquigarrow} \sigma' \\
 \mathcal{D}, \Delta, \sigma'_i \vdash cs \xrightarrow{\downarrow} \sigma_{\downarrow} \quad \mathcal{D}, \Delta, \sigma_{\downarrow} \vdash cs \xrightarrow{\rightsquigarrow} \sigma'' \quad \text{Inject}_{\uparrow}(\sigma, E_p, \tau, \sigma_i) \\
 \hline
 \mathcal{D}, E_p, \Delta, \tau, \sigma \vdash cs \xrightarrow{\uparrow, \downarrow} \sigma', \sigma'' \quad \text{Inject}_{\downarrow}(\sigma', E_p, \tau, \sigma'_i)
 \end{array}$$

Remark 4 (Input ports and signal store). For a given $\Delta \in \text{Design}$, $\sigma \in \Sigma(\Delta)$, $E_p \in \mathbb{N} \rightarrow \text{Clk} \rightarrow (\text{Ins}(\Delta) \rightarrow \text{value})$, $\tau \in \mathbb{N}$, $\text{clk} \in \text{Clk}$, we have $\text{dom}(E_p(\tau, \text{clk})) \subseteq \text{dom}(\mathcal{S})$, where $\sigma = \langle \mathcal{S}, \mathcal{C}, \mathcal{E} \rangle$. Indeed, the input ports of Δ that constitutes the domain of $E_p(\tau, \text{clk})$ are a subset of the set of signals. The set of signals constitutes the domain of the signal store of σ (i.e. \mathcal{S}); thus we have $\text{dom}(E_p(\tau, \text{clk})) \subseteq \text{dom}(\mathcal{S})$.

1.7.4 Initialization rules

$$\frac{\text{INIT} \quad \mathcal{D}, \Delta, \sigma \vdash \text{cs} \xrightarrow{\text{runinit}} \sigma' \quad \mathcal{D}, \Delta, \sigma' \vdash \text{cs} \xrightarrow{\sim} \sigma''}{\mathcal{D}, \Delta, \sigma \vdash \text{cs} \xrightarrow{\text{init}} \sigma''}$$

At the initialization phase, the block of sequential instructions of all processes is executed exactly once (*runinit*), then a stabilization phase follows (*stabilize*). It is during the initialization phase that the first part of *rst* blocks is executed. A block (*rst ss ss'*) is equivalent to (*if* (*rst* = *false*) *then* *ss* *else* *ss'*) where *rst* is a reserved signal identifier. Therefore, when considering a *rst* block, the *runinit* relation executes the *ss* block; at every other moment of the simulation, the *ss'* block is executed. This mimicks the conventional proceeding of a simulation where the *rst* signal (for *reset* signal) is set to *false* during the initialization (only during the *runinit* phase, not during the stabilization phase), and then is set to *true* for the rest of the simulation.

Evaluation of a process statement

Premises

- The *i* flag of the *ss_i* relation indicates that all sequential instructions responding to the initialization phase (i.e, *rst* blocks) will be executed.
- The set of events of state σ is emptied (*NoEv*(σ)) before the evaluation of the process statement body, and the resulting state is the starting state that will be written through the execution of the process statement body.

$$\frac{\text{PSRUNINIT} \quad \Delta, \sigma, \text{NoEv}(\sigma), \Lambda \vdash \text{ss} \xrightarrow{\text{ss}_i} \sigma', \Lambda'}{\mathcal{D}, \Delta, \sigma \vdash \text{process}(\text{id}_p, \text{sl}, \text{vars}, \text{ss}) \xrightarrow{\text{runinit}} \sigma'} \quad \Delta(\text{id}_p) = \Lambda$$

Evaluation of a component instantiation statement

Side conditions

If σ_c'' has a non-empty set of events, then the state of component id_c must be overridden in the target state of the *runinit* relation. Moreover, id_c must be added to the set of events of the embedding design. This is expressed by the side condition: $\sigma'' = (\sigma'(\text{id}_c) \leftarrow \sigma_c'') \cup_{\mathcal{E}} \{\text{id}_c\}$

COMPRUNINIT

$$\begin{array}{c}
\Delta, \Delta_c, \sigma, \sigma_c \vdash i \xrightarrow{\text{mapip}} \sigma'_c \\
\mathcal{D}, \Delta_c, \sigma'_c \vdash \mathcal{D}(\text{id}_e).cs \xrightarrow{\text{runinit}} \sigma''_c \\
\Delta, \Delta_c, \text{NoEv}(\sigma), \sigma''_c \vdash o \xrightarrow{\text{mapop}} \sigma' \quad \text{id}_e \in \mathcal{D} \\
\hline
\mathcal{D}, \Delta, \sigma \vdash \text{comp}(\text{id}_c, \text{id}_e, g, i, o) \xrightarrow{\text{runinit}} \sigma'' \quad \begin{array}{l} \Delta(\text{id}_c) = \Delta_c, \sigma(\text{id}_c) = \sigma_c \\ \sigma'' = \langle S', C'(\text{id}_c) \leftarrow \sigma''_c, \mathcal{E}' \cup (\mathcal{C} \dot{\cap} C') \rangle \end{array}
\end{array}$$

Evaluation of the composition of concurrent statements

PARRUNINIT

$$\frac{\mathcal{D}, \Delta, \sigma \vdash cs \xrightarrow{\text{runinit}} \sigma' \quad \mathcal{D}, \Delta, \sigma \vdash cs' \xrightarrow{\text{runinit}} \sigma''}{\mathcal{D}, \Delta, \sigma \vdash cs \parallel cs' \xrightarrow{\text{runinit}} \text{merge}(\sigma, \sigma', \sigma'')} \quad \mathcal{E}' \cap \mathcal{E}'' = \emptyset$$

NULLRUNINIT

$$\frac{}{\Delta, \sigma \vdash \text{null} \xrightarrow{\text{runinit}} \text{NoEv}(\sigma)}$$

The merge function computes a new state based on the original state o , and the states s and s' yielded by the computation of two concurrent statements. In the resulting state, the signal value store S_m is a function merging together the signal store functions at state o , s and s' . S_m yields values from the signal store \mathcal{S} (resp. \mathcal{S}') for all signal that belongs to the set of events at state s (resp. s'), and yields values from the original signal store \mathcal{S}_o for all unchanged signals. The same goes for the resulting component instance state store C_m . The new set of events \mathcal{E}_m is the union between set of events at state s and s' . The merge correctly merges the state o , s and s' only if the set of events of s and s' are disjoint. Fortunately, the PARRUNINIT rule that calls to the merge function defines the condition of disjoint set of events as a side condition.

```

1 Definition merge(o,s,s') :=
2   let o = (So, Co, Eo) in
3   let s = (S, C, E) in
4   let s' = (S', C', E') in
5   let Sm = λid. if id ∈ E then S(id) else if id ∈ E' then S'(id) else So(id)
6   let Cm = λid. if id ∈ E then C(id) else if id ∈ E' then C'(id) else Co(id)
7   let Em = E ∪ E' in (Sm, Cm, Em).

```

Remark 5 (No multiply-driven signals). For all states $\sigma = (S, C, E)$ and $\sigma' = (S', C', E')$ resulting from the execution of two concurrent statements cs and cs' , $E \cap E' = \emptyset$. Otherwise, there exists some multiply-driven signals, which are forbidden in our semantics.

1.7.5 Clock phases rules

The following rules express the evaluation of concurrent statements at clock phases, i.e, the \uparrow and \downarrow phases. The clock signal, triggering the evaluation of synchronous process statements, is represented by the reserved signal identifier `clk`. Thus, synchronous processes are processes containing the `clk` in their sensitivity list.

Evaluation of a process statement

$$\frac{\text{PsRENoCLK}}{\mathcal{D}, \Delta, \sigma \vdash \text{process}(\text{id}_p, \text{sl}, \text{vars}, \text{ss}) \xrightarrow{\uparrow} \sigma} \quad \text{clk} \notin \text{sl}$$

Premises

The \uparrow flag in the ss_{\uparrow} relation indicates that rising blocks will be executed.

$$\frac{\text{PsRECLK} \quad \Delta, \sigma, \text{NoEv}(\sigma), \Lambda \vdash \text{ss} \xrightarrow{\text{ss}_{\uparrow}} \sigma', \Lambda'}{\mathcal{D}, \Delta, \sigma \vdash \text{process}(\text{id}_p, \text{sl}, \text{vars}, \text{ss}) \xrightarrow{\uparrow} \sigma'} \quad \begin{array}{l} \text{clk} \in \text{sl} \\ \Delta(\text{id}_p) = \Lambda \end{array}$$

$$\frac{\text{PsFENoCLK}}{\mathcal{D}, \Delta, \sigma \vdash \text{process}(\text{id}_p, \text{sl}, \text{vars}, \text{ss}) \xrightarrow{\downarrow} \sigma} \quad \text{clk} \notin \text{sl}$$

Premises

The \downarrow flag in the ss_{\downarrow} relation indicates that falling blocks will be executed.

$$\frac{\text{PsFECLK} \quad \Delta, \sigma, \text{NoEv}(\sigma), \Lambda \vdash \text{ss} \xrightarrow{\text{ss}_{\downarrow}} \sigma', \Lambda'}{\mathcal{D}, \Delta, \sigma \vdash \text{process}(\text{id}_p, \text{sl}, \text{vars}, \text{ss}) \xrightarrow{\downarrow} \sigma'} \quad \begin{array}{l} \text{clk} \in \text{sl} \\ \Delta(\text{id}_p) = \Lambda \end{array}$$

Evaluation of a component instantiation statement

$$\frac{\text{COMPRe} \quad \begin{array}{l} \Delta, \Delta_c, \sigma, \sigma_c \vdash \text{i} \xrightarrow{\text{mapip}} \sigma'_c \\ \mathcal{D}, \Delta_c, \sigma'_c \vdash \mathcal{D}(\text{id}_e).cs \xrightarrow{\uparrow} \sigma''_c \\ \Delta, \Delta_c, \sigma, \sigma''_c \vdash \text{o} \xrightarrow{\text{mapop}} \sigma' \end{array}}{\mathcal{D}, \Delta, \sigma \vdash \text{comp}(\text{id}_c, \text{id}_e, \text{g}, \text{i}, \text{o}) \xrightarrow{\uparrow} \sigma''} \quad \begin{array}{l} \text{id}_e \in \mathcal{D} \\ \Delta(\text{id}_c) = \Delta_c, \sigma(\text{id}_c) = \sigma_c \\ \sigma'' = \langle \mathcal{S}', \mathcal{C}'(\text{id}_c) \leftarrow \sigma''_c, \mathcal{E}' \cup (\mathcal{C} \setminus \mathcal{C}') \rangle \end{array}$$

COMPFE

$$\begin{array}{c}
\Delta, \Delta_c, \sigma, \sigma_c \vdash i \xrightarrow{mapip} \sigma'_c \\
\mathcal{D}, \Delta_c, \sigma'_c \vdash \mathcal{D}(id_e).cs \xrightarrow{\downarrow} \sigma''_c \\
\Delta, \Delta_c, \sigma, \sigma''_c \vdash o \xrightarrow{mapop} \sigma' \quad id_e \in \mathcal{D} \\
\hline
\mathcal{D}, \Delta, \sigma \vdash \text{comp}(id_c, id_e, g, i, o) \xrightarrow{\downarrow} \sigma'' \quad \Delta(id_c) = \Delta_c, \sigma(id_c) = \sigma_c \\
\sigma'' = \langle S', \mathcal{C}'(id_c) \leftarrow \sigma''_c, \mathcal{E}' \cup (\mathcal{C} \setminus \mathcal{C}') \rangle
\end{array}$$

Evaluation of the composition of concurrent statements

PARFE

$$\frac{\mathcal{D}, \Delta, \sigma \vdash cs \xrightarrow{\downarrow} \sigma' \quad \mathcal{D}, \Delta, \sigma \vdash cs' \xrightarrow{\downarrow} \sigma''}{\mathcal{D}, \Delta, \sigma \vdash cs \parallel cs' \xrightarrow{\downarrow} \text{merge}(\sigma, \sigma', \sigma'')} \quad \mathcal{E}' \cap \mathcal{E}'' = \emptyset$$

NULLFE

$$\frac{}{\Delta, \sigma \vdash \text{null} \xrightarrow{\downarrow} \sigma}$$

PARRE

$$\frac{\mathcal{D}, \Delta, \sigma \vdash cs \xrightarrow{\uparrow} \sigma' \quad \mathcal{D}, \Delta, \sigma \vdash cs' \xrightarrow{\uparrow} \sigma''}{\mathcal{D}, \Delta, \sigma \vdash cs \parallel cs' \xrightarrow{\uparrow} \text{merge}(\sigma, \sigma', \sigma'')} \quad \mathcal{E}' \cap \mathcal{E}'' = \emptyset$$

NULLRE

$$\frac{}{\Delta, \sigma \vdash \text{null} \xrightarrow{\uparrow} \sigma}$$

1.7.6 Stabilization rules

STABILIZEEND

$$\frac{\mathcal{D}, \Delta, \sigma \vdash cs \xrightarrow{comb} \sigma}{\mathcal{D}, \Delta, \sigma \vdash cs \xrightarrow{\rightsquigarrow} \sigma} \quad \mathcal{E} = \emptyset$$

STABILIZELOOP

$$\frac{\mathcal{D}, \Delta, \sigma \vdash cs \xrightarrow{comb} \sigma' \quad \mathcal{D}, \Delta, \sigma' \vdash cs \xrightarrow{\rightsquigarrow} \sigma''}{\mathcal{D}, \Delta, \sigma \vdash cs \xrightarrow{\rightsquigarrow} \sigma''} \quad \begin{array}{l} \mathcal{E} \neq \emptyset \\ \mathcal{E}'' = \emptyset \end{array}$$

Evaluation of a process statement**Premises**

The c flag (for *combinational*) on the ss_c relation indicates that instructions responding to clock events (falling and rising blocks) and instructions executed during the initialization phase only (rst blocks) will not be considered.

PSCOMB

$$\frac{\Delta, \sigma, NoEv(\sigma), \Lambda \vdash ss \xrightarrow{ss_c} \sigma', \Lambda'}{\mathcal{D}, \Delta, \sigma \vdash \text{process}(id_p, sl, vars, ss) \xrightarrow{comb} \sigma'} \quad \Delta(id_p) = \Lambda$$

Evaluation of a component instantiation statement

COMPCOMB

$$\begin{array}{c}
\Delta, \Delta_c, \sigma, \sigma_c \vdash i \xrightarrow{\text{mapip}} \sigma'_c \\
\mathcal{D}, \Delta_c, \sigma'_c \vdash \mathcal{D}(id_e).cs \xrightarrow{\text{comb}} \sigma''_c \\
\Delta, \Delta_c, NoEv(\sigma), \sigma''_c \vdash o \xrightarrow{\text{mapop}} \sigma' \quad id_e \in \mathcal{D} \\
\hline
\Delta(id_c) = \Delta_c, \sigma(id_c) = \sigma_c \\
\mathcal{D}, \Delta, \sigma \vdash \text{comp}(id_c, id_e, g, i, o) \xrightarrow{\text{comb}} \sigma'' \quad \sigma'' = \langle S', \mathcal{C}'(id_c) \leftarrow \sigma''_c, \mathcal{E}' \cup (\mathcal{C} \setminus \mathcal{C}') \rangle
\end{array}$$

Evaluation of the composition of concurrent statements

PARCOMB

$$\frac{\mathcal{D}, \Delta, \sigma \vdash cs \xrightarrow{\text{comb}} \sigma' \quad \mathcal{D}, \Delta, \sigma \vdash cs' \xrightarrow{\text{comb}} \sigma''}{\mathcal{D}, \Delta, \sigma \vdash cs \parallel cs' \xrightarrow{\text{comb}} \text{merge}(\sigma, \sigma', \sigma'')} \quad \mathcal{E}' \cap \mathcal{E}'' = \emptyset$$

NULLCOMB

$$\frac{}{\Delta, \sigma \vdash \text{null} \xrightarrow{\downarrow} NoEv(\sigma)}$$

1.7.7 Evaluation of input and output port maps

MAIPSIMPLE

$$\frac{\Delta, \sigma \vdash e \xrightarrow{e} v \quad v \in_c T \quad \Delta_c(id_s) = T}{\Delta, \Delta_c, \sigma, \sigma_c \vdash id_s \Rightarrow e \xrightarrow{\text{mapip}} \langle S', \mathcal{C}, \mathcal{E} \rangle} \quad \sigma_c = \langle S, \mathcal{C}, \mathcal{E} \rangle \\
S' = S(id_s) \leftarrow v$$

MAIIPARTIAL

$$\frac{\Delta, \sigma \vdash e \xrightarrow{e} v \quad v \in_c T \quad \vdash e_i \xrightarrow{e} v_i \quad v_i \in_c \text{nat}(n, m)}{\Delta, \Delta_c, \sigma, \sigma_c \vdash id_s(e_i) \Rightarrow e \xrightarrow{\text{mapip}} \langle S', \mathcal{C}, \mathcal{E} \rangle} \quad \Delta_c(id_s) = \text{array}(T, n, m) \\
\sigma_c = \langle S, \mathcal{C}, \mathcal{E} \rangle \quad S' = S(id_s) \leftarrow \text{set_at}(v, v_i, S(id_s))$$

MAIPCOMP

$$\frac{\Delta, \Delta_c, \sigma, \sigma_c \vdash \text{assoc}_{ip} \xrightarrow{\text{mapip}} \sigma'_c \quad \Delta, \Delta_c, \sigma, \sigma'_c \vdash \text{ipmap} \xrightarrow{\text{mapip}} \sigma''_c}{\Delta, \Delta_c, \sigma, \sigma_c \vdash \langle \text{assoc}_{ip}, \text{ipmap} \rangle \xrightarrow{\text{mapip}} \sigma''_c}$$

Remark 6 (Out ports and e). We can not use the e relation to interpret the values of output ports, because output ports are write-only constructs. We append the flag o to the e relation (i.e, e_o) to permit the evaluation of output port identifiers as regular signal identifier expressions.

MAPOPOPEN

$$\Delta, \Delta_c, \sigma, \sigma_c \vdash \text{id}_f \Rightarrow \text{open} \xrightarrow{\text{mapop}} \sigma_c$$

MAPOPSIMPLETOSIMPLE

$$\frac{\Delta_c, \sigma_c \vdash \text{id}_f \xrightarrow{e_o} v \quad v \in_c T}{\Delta, \Delta_c, \sigma, \sigma_c \vdash \text{id}_f \Rightarrow \text{id}_a \xrightarrow{\text{mapop}} \langle S', \mathcal{C}, \mathcal{E}' \rangle} \quad \begin{array}{l} \text{id}_a \in \text{Sigs}(\Delta) \cup \text{Outs}(\Delta) \\ \Delta(\text{id}_a) = T \\ \sigma = \langle S, \mathcal{C}, \mathcal{E} \rangle \\ S' = S(\text{id}_a) \leftarrow v, \mathcal{E}' = \mathcal{E} \cup (S \overset{\neq}{\cap} S') \end{array}$$

MAPOPSIMPLETOPARTIAL

$$\frac{\begin{array}{l} \vdash e_i \xrightarrow{e} v_i \quad v \in_c T \\ \Delta_c, \sigma_c \vdash \text{id}_f \xrightarrow{e_o} v \quad v_i \in_c \text{nat}(n, m) \end{array}}{\Delta, \Delta_c, \sigma, \sigma_c \vdash \text{id}_f \Rightarrow \text{id}_a(e_i) \xrightarrow{\text{mapop}} \langle S', \mathcal{C}, \mathcal{E}' \rangle} \quad \begin{array}{l} \text{id}_a \in \text{Sigs}(\Delta) \cup \text{Outs}(\Delta) \\ \Delta(\text{id}_a) = \text{array}(T, n, m) \\ \sigma = \langle S, \mathcal{C}, \mathcal{E} \rangle \\ S' = S(\text{id}_a) \leftarrow \text{set_at}(v, v_i, S(\text{id}_a)) \\ \mathcal{E}' = \mathcal{E} \cup (S \overset{\neq}{\cap} S') \end{array}$$

MAPOPPARTIALTOSIMPLE

$$\frac{\Delta_c, \sigma_c \vdash \text{id}_f(e'_i) \xrightarrow{e_o} v \quad v \in_c T}{\Delta, \Delta_c, \sigma, \sigma_c \vdash \text{id}_f(e'_i) \Rightarrow \text{id}_a \xrightarrow{\text{mapop}} \langle S', \mathcal{C}, \mathcal{E}' \rangle} \quad \begin{array}{l} \text{id}_a \in \text{Sigs}(\Delta) \cup \text{Outs}(\Delta) \\ \Delta(\text{id}_a) = T \\ \sigma = \langle S, \mathcal{C}, \mathcal{E} \rangle \\ S' = S(\text{id}_a) \leftarrow v, \mathcal{E}' = \mathcal{E} \cup (S \overset{\neq}{\cap} S') \end{array}$$

MAPOPPARTIALTOPARTIAL

$$\frac{\begin{array}{l} \vdash e_i \xrightarrow{e} v_i \quad v \in_c T \\ \Delta_c, \sigma_c \vdash \text{id}_f(e'_i) \xrightarrow{e_o} v \quad v_i \in_c \text{nat}(n, m) \end{array}}{\Delta, \Delta_c, \sigma, \sigma_c \vdash \text{id}_f(e'_i) \Rightarrow \text{id}_a(e_i) \xrightarrow{\text{mapop}} \langle S', \mathcal{C}, \mathcal{E}' \rangle} \quad \begin{array}{l} \text{id}_a \in \text{Sigs}(\Delta) \cup \text{Outs}(\Delta) \\ \Delta(\text{id}_a) = \text{array}(T, n, m) \\ \sigma = \langle S, \mathcal{C}, \mathcal{E} \rangle \\ S' = S(\text{id}_a) \leftarrow \text{set_at}(v, v_i, S(\text{id}_a)) \\ \mathcal{E}' = \mathcal{E} \cup (S \overset{\neq}{\cap} S') \end{array}$$

MAPOPCOMP

$$\frac{\Delta, \Delta_c, \sigma, \sigma_c \vdash \text{assoc}_{po} \xrightarrow{\text{mapop}} \sigma' \quad \Delta, \Delta_c, \sigma', \sigma_c \vdash \text{opmap} \xrightarrow{\text{mapop}} \sigma''}{\Delta, \Delta_c, \sigma, \sigma_c \vdash \langle \text{assoc}_{po}, \text{opmap} \rangle \xrightarrow{\text{mapop}} \sigma''}$$

The e_o relation is only defined to retrieve the value of out ports from a store signal S under a design state $\sigma = \langle S, \mathcal{C}, \mathcal{E} \rangle$.

$$\frac{\text{OUTO}}{\Delta, \sigma \vdash \text{id}_s \xrightarrow{e_o} \sigma(\text{id}_s)} \quad \begin{array}{l} \text{id}_s \in \text{Outs}(\Delta) \\ \text{id}_s \in \sigma \end{array}$$

IDXOUTO

$$\frac{\begin{array}{l} \vdash e_i \xrightarrow{e} v_i \quad v_i \in_c \text{nat}(n, m) \end{array}}{\Delta, \sigma \vdash \text{id}_s(e_i) \xrightarrow{e_o} \text{get_at}(i, \sigma(\text{id}_s))} \quad \begin{array}{l} \text{id}_s \in \text{Outs}(\Delta) \\ \text{id}_s \in \sigma \\ \Delta(\text{id}_s) = \text{array}(T, n, m) \\ i = v_i \bmod n \end{array}$$

1.7.8 Evaluation of sequential statements

The ss symbol indicates that the evaluation of the considered sequential statement does not depend on a specific flag (i.e, the c , i , \uparrow or \downarrow flag). In the rules of the ss relation, a ss flag is tranferred from the conclusion to the premises when an sequential statement is composed of inner sequential blocks.

Signal assignment statement

A signal assignment generates a new design state with a modified signal store and a new set of events. Note that there are two states on the left side of the thesis symbol. σ represents the state holding the current values of signals, and σ_w holds the new values of signals (i.e. the *written* state).

Premises

The premise $\mathcal{S}(id_s) \in_c T$ checks that the value associated to signal id_s in the signal store of σ complies with type T , where T is the type associated with signal id_s in Δ .

$$\begin{array}{c}
 \text{SIGASSIGN} \\
 \frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad v \in_c T}{\Delta, \sigma, \sigma_w, \Lambda \vdash id_s \leftarrow e \xrightarrow{ss} \langle \mathcal{S}'_w, \mathcal{C}_w, \mathcal{E}'_w \rangle, \Lambda} \quad \begin{array}{l} id_s \in Sigs(\Delta) \cup Outs(\Delta) \\ \Delta(id_s) = T \\ \mathcal{S}'_w = \mathcal{S}_w(id_s) \leftarrow v \\ \mathcal{E}'_w = \mathcal{E}_w \cup (\mathcal{S}_w \overset{\neq}{\cap} \mathcal{S}'_w) \end{array} \\
 \\
 \text{IDXSIGASSIGN} \\
 \frac{\begin{array}{l} \Delta, \sigma, \Lambda \vdash e_i \xrightarrow{e} v_i \quad v_i \in_c T \\ \Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad v_i \in_c \text{nat}(n, m) \end{array}}{\Delta, \sigma, \sigma_w, \Lambda \vdash id_s(e_i) \leftarrow e \xrightarrow{ss} \langle \mathcal{S}'_w, \mathcal{C}_w, \mathcal{E}'_w \rangle, \Lambda} \quad \begin{array}{l} id_s \in Sigs(\Delta) \cup Outs(\Delta) \\ \Delta(id_s) = \text{array}(T, n, m) \\ \mathcal{S}'_w = \mathcal{S}_w(id_s) \leftarrow \text{set_at}(v, v_i, \mathcal{S}_w(id_s)) \\ \mathcal{E}'_w = \mathcal{E}_w \cup (\mathcal{S}_w \overset{\neq}{\cap} \mathcal{S}'_w) \end{array}
 \end{array}$$

Variable assignment statement

A variable assignment statement modifies the variable values in the local environment.

$$\begin{array}{c}
 \text{VARASSIGN} \\
 \frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad v \in_c T}{\Delta, \sigma, \sigma_w, \Lambda \vdash id_v := e \xrightarrow{ss} \sigma_w, \Lambda(id_v) \leftarrow (T, v)} \quad \begin{array}{l} id_v \in \Lambda \\ \Lambda(id_v) = (T, val) \end{array} \\
 \\
 \text{IDXVARASSIGN} \\
 \frac{\begin{array}{l} \Delta, \sigma, \Lambda \vdash e_i \xrightarrow{e} v_i \quad v_i \in_c \text{nat}(n, m) \\ \Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad v \in_c T \end{array}}{\Delta, \sigma, \sigma_w, \Lambda \vdash id_v(e_i) := e \xrightarrow{ss} \sigma_w, \Lambda(id_v) \leftarrow (T, \text{set_at}(v, v_i, val))} \quad \begin{array}{l} id_v \in \Lambda \\ \Lambda(id_v) = (\text{array}(T, n, m), val) \end{array}
 \end{array}$$

If statement

$$\begin{array}{c}
\text{IF}\top \\
\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} \top \quad \Delta, \sigma, \sigma_w, \Lambda \vdash ss \xrightarrow{ss} \sigma'_w, \Lambda'}{\Delta, \sigma, \sigma_w, \Lambda \vdash \text{if } (e) \text{ ss } \xrightarrow{ss} \sigma'_w, \Lambda'} \\
\\
\text{IF}\perp \\
\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} \perp}{\Delta, \sigma, \sigma_w, \Lambda \vdash \text{if } (e) \text{ ss } \xrightarrow{ss} \sigma_w, \Lambda} \\
\\
\text{IFELSE}\top \\
\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{expr} \top \quad \Delta, \sigma, \sigma_w, \Lambda \vdash ss \xrightarrow{ss} \sigma'_w, \Lambda'}{\Delta, \sigma, \sigma_w, \Lambda \vdash \text{if } (e) \text{ ss ss'} \xrightarrow{ss} \sigma'_w, \Lambda'} \\
\\
\text{IFELSE}\perp \\
\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} \perp \quad \Delta, \sigma, \sigma_w, \Lambda \vdash ss' \xrightarrow{ss} \sigma'_w, \Lambda'}{\Delta, \sigma, \sigma_w, \Lambda \vdash \text{if } (e) \text{ ss ss'} \xrightarrow{ss} \sigma'_w, \Lambda'}
\end{array}$$

Loop statement

$$\begin{array}{c}
\text{LOOP}\perp \\
\frac{\Delta, \sigma, \sigma_w, \Lambda_i \vdash ss \xrightarrow{ss} \sigma'_w, \Lambda' \quad \Delta, \sigma, \Lambda_i \vdash \text{id}_v = e' \xrightarrow{e} \perp \quad \Delta, \sigma, \sigma'_w, \Lambda' \vdash \text{for } (\text{id}_v, e, e') \text{ ss } \xrightarrow{ss} \sigma''_w, \Lambda'' \quad \text{id}_v \in \Lambda}{\Delta, \sigma, \sigma_w, \Lambda \vdash \text{for } (\text{id}_v, e, e') \text{ ss } \xrightarrow{ss} \sigma''_w, \Lambda''} \quad \begin{array}{l} \Lambda(\text{id}_v) = (T, val) \\ \Lambda_i = \Lambda(\text{id}_v) \leftarrow (T, val + 1) \end{array} \\
\\
\text{LOOP}\top \\
\frac{\Delta, \sigma, \Lambda_i \vdash \text{id}_v = e' \xrightarrow{e} \top}{\Delta, \sigma, \sigma_w, \Lambda \vdash \text{for } (\text{id}_v, e, e') \text{ ss } \xrightarrow{ss} \sigma_w, \Lambda \setminus (\text{id}_v, \Lambda(\text{id}_v))} \quad \begin{array}{l} \text{id}_v \in \Lambda \\ \Lambda(\text{id}_v) = (T, val) \\ \Lambda_i = \Lambda(\text{id}_v) \leftarrow (T, val + 1) \end{array} \\
\\
\text{LOOPINIT} \\
\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad \Delta, \sigma, \Lambda \vdash e' \xrightarrow{e} v' \quad \Delta, \sigma, \sigma_w, \Lambda_i \vdash \text{for } (\text{id}_v, e, e') \text{ ss } \xrightarrow{ss} \sigma'_w, \Lambda' \quad \text{id}_v \notin \Lambda}{\Delta, \sigma, \sigma_w, \Lambda \vdash \text{for } (\text{id}_v, e, e') \text{ ss } \xrightarrow{ss} \sigma'_w, \Lambda'} \quad \Lambda_i = \Lambda \cup (\text{id}_v, (\text{nat}(v, v'), v))
\end{array}$$

Rising and falling edge block statements

$$\begin{array}{c}
\text{RISINGEDGEDEFAULT} \\
\frac{}{\Delta, \sigma, \sigma_w, \Lambda \vdash \text{rising ss } \xrightarrow{ss_f} \sigma_w, \Lambda} \quad f \neq \uparrow \\
\\
\text{FALLINGEDGEDEFAULT} \\
\frac{}{\Delta, \sigma, \sigma_w, \Lambda \vdash \text{falling ss } \xrightarrow{ss_f} \sigma_w, \Lambda} \quad f \neq \downarrow \\
\\
\text{RISINGEDGEEXEC} \\
\frac{\Delta, \sigma, \sigma_w, \Lambda \vdash ss \xrightarrow{ss_\uparrow} \sigma'_w, \Lambda'}{\Delta, \sigma, \sigma_w, \Lambda \vdash \text{rising ss } \xrightarrow{ss_\uparrow} \sigma'_w, \Lambda'} \\
\\
\text{FALLINGEDGEEXEC} \\
\frac{\Delta, \sigma, \sigma_w, \Lambda \vdash ss \xrightarrow{ss_\downarrow} \sigma'_w, \Lambda'}{\Delta, \sigma, \sigma_w, \Lambda \vdash \text{falling ss } \xrightarrow{ss_\downarrow} \sigma'_w, \Lambda'}
\end{array}$$

Rst block statement

$$\begin{array}{c}
\text{RSTDEFAULT} \\
\frac{\Delta, \sigma, \sigma_w, \Lambda \vdash ss' \xrightarrow{ss_f} \sigma'_w, \Lambda'}{\Delta, \sigma, \sigma_w, \Lambda \vdash \text{rst } ss \text{ } ss' \xrightarrow{ss_f} \sigma'_w, \Lambda'} \quad f \neq i \\
\text{RSTEXEC} \\
\frac{\Delta, \sigma, \sigma_w, \Lambda \vdash ss \xrightarrow{ss_i} \sigma'_w, \Lambda'}{\Delta, \sigma, \sigma_w, \Lambda \vdash \text{rst } ss \text{ } ss' \xrightarrow{ss_i} \sigma'_w, \Lambda'}
\end{array}$$

Composition of sequential statements and null statement

$$\begin{array}{c}
\text{SEQSTMT} \\
\frac{\Delta, \sigma, \sigma_w, \Lambda \vdash ss \xrightarrow{ss} \sigma'_w, \Lambda' \quad \Delta, \sigma, \sigma'_w, \Lambda' \vdash ss' \xrightarrow{ss} \sigma''_w, \Lambda''}{\Delta, \sigma, \sigma_w, \Lambda \vdash ss; ss' \xrightarrow{ss} \sigma''_w, \Lambda''} \quad \text{NULLSTMT} \\
\frac{}{\Delta, \sigma, \sigma_w, \Lambda \vdash \text{null} \xrightarrow{ss} \sigma_w, \Lambda}
\end{array}$$

1.7.9 Evaluation of expressions

$$\begin{array}{c}
\text{NAT} \\
\frac{}{\Delta, \sigma, \Lambda \vdash n \xrightarrow{e} n} \quad n \in \mathbb{N} \quad n \leq \text{NATMAX} \\
\text{FALSE} \\
\frac{}{\Delta, \sigma, \Lambda \vdash \text{false} \xrightarrow{e} \perp} \\
\text{TRUE} \\
\frac{}{\Delta, \sigma, \Lambda \vdash \text{true} \xrightarrow{e} \top}
\end{array}$$

$$\begin{array}{c}
\text{AGGREG} \\
\frac{\Delta, \sigma, \Lambda \vdash e_i \xrightarrow{e} v_i}{\Delta, \sigma, \Lambda \vdash (e_1, \dots, e_n) \xrightarrow{e} (v_1, \dots, v_n)} \quad i = 1, \dots, n
\end{array}$$

$$\begin{array}{c}
\text{SIG} \\
\frac{}{\Delta, \sigma, \Lambda \vdash \text{id}_s \xrightarrow{e} \sigma(\text{id}_s)} \quad \text{id}_s \in \text{Sigs}(\Delta) \cup \text{Ins}(\Delta) \\
\text{VAR} \\
\frac{}{\Delta, \sigma, \Lambda \vdash \text{id}_v \xrightarrow{e} v} \quad \text{id}_v \in \Lambda \quad \Lambda(\text{id}_v) = (T, v)
\end{array}$$

$$\begin{array}{c}
\text{GEN} \\
\frac{}{\Delta, \sigma, \Lambda \vdash \text{id}_g \xrightarrow{e} v} \quad \text{id}_g \in \text{Gens}(\Delta) \quad \Lambda(\text{id}_g) = (T, v)
\end{array}$$

$$\begin{array}{c}
\text{IDXSIG} \\
\frac{\Delta, \sigma, \Lambda \vdash e_i \xrightarrow{e} v_i \quad v_i \in_c \text{nat}(n, m) \quad \text{id}_s \in \text{Sigs}(\Delta) \cup \text{Ins}(\Delta) \quad \Lambda(\text{id}_s) = \text{array}(T, n, m)}{\Delta, \sigma, \Lambda \vdash \text{id}_s(e_i) \xrightarrow{e} \text{get_at}(i, \sigma(\text{id}_s))} \quad i = v_i \bmod n
\end{array}$$

$$\begin{array}{c}
\text{IDXVAR} \\
\frac{\Delta, \sigma, \Lambda \vdash e_i \xrightarrow{e} v_i \quad v_i \in_c \text{nat}(n, m) \quad \text{id}_v \in \Lambda \quad \Lambda(\text{id}_v) = (\text{array}(T, n, m), v)}{\Delta, \sigma, \Lambda \vdash \text{id}_v(e_i) \xrightarrow{e} \text{get_at}(i, v)} \quad i = v_i \bmod n
\end{array}$$

where $\text{get_at}(i, a)$ is a function returning the i -th element of array a .

NATADD

$$\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad \Delta, \sigma, \Lambda \vdash e' \xrightarrow{e} v'}{\Delta, \sigma, \Lambda \vdash e + e' \xrightarrow{e} v +_{\mathbb{N}} v'} \quad v +_{\mathbb{N}} v' \leq \text{NATMAX}$$

NATSUB

$$\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad \Delta, \sigma, \Lambda \vdash e' \xrightarrow{e} v'}{\Delta, \sigma, \Lambda \vdash e - e' \xrightarrow{e} v -_{\mathbb{N}} v'} \quad v \geq v'$$

ORDOP

$$\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad \Delta, \sigma, \Lambda \vdash e' \xrightarrow{e} v'}{\Delta, \sigma, \Lambda \vdash e \text{ op}_{\text{ordn}} e' \xrightarrow{e} v \text{ op}_{\text{ordn}} v'} \quad \text{op}_{\text{ordn}} \in \{<, \leq, >, \geq\}$$

BOOLBINOP

$$\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad \Delta, \sigma, \Lambda \vdash e' \xrightarrow{e} v'}{\Delta, \sigma, \Lambda \vdash e \text{ op}_{\text{bool}} e' \xrightarrow{e} v \text{ op}_{\text{bool}} v'} \quad \text{op}_{\text{bool}} \in \{\text{and}, \text{or}\}$$

NOTOP

$$\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v}{\Delta, \sigma, \Lambda \vdash \text{not } e \xrightarrow{e} \neg v}$$

EQOP

$$\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v \quad \Delta, \sigma, \Lambda \vdash e' \xrightarrow{e} v'}{\Delta, \sigma, \Lambda \vdash e = e' \xrightarrow{e} \text{eq}(v, v')}$$

DIFFOP

$$\frac{\Delta, \sigma, \Lambda \vdash e = e' \xrightarrow{e} v}{\Delta, \sigma, \Lambda \vdash e \neq e' \xrightarrow{e} \neg v}$$

where eq is the equality relation established for all types defined in the semantics.

PARENTH

$$\frac{\Delta, \sigma, \Lambda \vdash e \xrightarrow{e} v}{\Delta, \sigma, \Lambda \vdash (e) \xrightarrow{e} v}$$

Appendix A

The place design in concrete and abstract VHDL syntax

```

1  entity place is
2    generic(
3      input_arcs_number : natural := 1;
4      output_arcs_number : natural := 1;
5      maximal_marking : natural := 1
6    );
7    port(
8      clock : in std_logic;
9      reset_n : in std_logic;
10     initial_marking : in natural range 0 to maximal_marking;
11     input_arcs_weights : in weight_vector_t(input_arcs_number1 downto 0);
12     output_arcs_types : in arc_vector_t(output_arcs_number1 downto 0);
13     output_arcs_weights : in weight_vector_t(output_arcs_number1 downto 0);
14     input_transitions_fired : in std_logic_vector(input_arcs_number1 downto 0);
15     output_transitions_fired : in std_logic_vector(output_arcs_number1 downto 0);
16     output_arcs_valid : out std_logic_vector(output_arcs_number1 downto 0);
17     priority_authorizations : out std_logic_vector(output_arcs_number1 downto 0);
18     reinit_transitions_time : out std_logic_vector(output_arcs_number1 downto 0);
19     marked : out std_logic
20   );
21 end place;
22
23 architecture place_architecture of place is
24
25   subtype local_weight_t is natural range 0 to maximal_marking;
26
27   signal s_input_token_sum : local_weight_t;
28   signal s_marking : local_weight_t;
29   signal s_output_token_sum : local_weight_t;
30
31 begin
32

```

```

33
34 input_tokens_sum : process(input_arcs_weights, input_transitions_fired)
35     variable v_internal_input_token_sum : local_weight_t;
36 begin
37     v_internal_input_token_sum := 0;
38
39     for i in 0 to input_arcs_number - 1 loop
40         if (input_transitions_fired(i) = '1') then
41
42             v_internal_input_token_sum := v_internal_input_token_sum + input_arcs_weights(i)
43             ;
44
45         end if;
46     end loop;
47
48     s_input_token_sum <= v_internal_input_token_sum;
49 end process input_tokens_sum;
50
51 output_tokens_sum : process(output_arcs_types, output_arcs_weights,
52     output_transitions_fired)
53     variable v_internal_output_token_sum : local_weight_t;
54 begin
55     v_internal_output_token_sum := 0;
56
57     for i in 0 to output_arcs_number - 1 loop
58         if (output_transitions_fired(i) = '1' and output_arcs_types(i) = arc_t(BASIC)) then
59             v_internal_output_token_sum := v_internal_output_token_sum +
60             output_arcs_weights(i);
61         end if;
62     end loop;
63
64     s_output_token_sum <= v_internal_output_token_sum;
65 end process output_tokens_sum;
66
67 marking : process(clock, reset_n, initial_marking)
68 begin
69     if (reset_n = '0') then
70         s_marking <= initial_marking;
71     elsif rising_edge(clock) then
72         s_marking <= s_marking + (s_input_token_sum - s_output_token_sum);
73     end if;
74 end process marking;
75
76 determine_marked : process(s_marking)
77 begin
78     if (s_marking = 0) then
79         marked <= '0';

```



```

77     else
78         marked <= '1';
79     end if;
80 end process determine_marked;
81
82 marking_validation_evaluation : process(output_arcs_types, output_arcs_weights,
83     s_marking)
84 begin
85     for i in 0 to output_arcs_number - 1 loop
86         if ( ((output_arcs_types(i) = arc_t(BASIC)) or (output_arcs_types(i) = arc_t(TEST)))
87             and (s_marking >= output_arcs_weights(i)) )
88             or ( (output_arcs_types(i) = arc_t(INHIBITOR)) and (s_marking <
89                 output_arcs_weights(i)) ) )
90         then
91             output_arcs_valid(i) <= '1';
92         else
93             output_arcs_valid(i) <= '0';
94         end if;
95     end loop;
96 end process marking_validation_evaluation;
97
98 priority_evaluation : process(output_arcs_types, output_arcs_weights,
99     output_transitions_fired, s_marking)
100 variable v_saved_output_token_sum : local_weight_t;
101 begin
102     v_saved_output_token_sum := 0;
103
104     for i in 0 to output_arcs_number - 1 loop
105         if (s_marking >= v_saved_output_token_sum + output_arcs_weights(i)) then
106             priority_authorizations(i) <= '1';
107         else
108             priority_authorizations(i) <= '0';
109         end if;
110
111         if ((output_transitions_fired(i) = '1') and (output_arcs_types(i) = arc_t(BASIC)))
112         then
113             v_saved_output_token_sum := v_saved_output_token_sum + output_arcs_weights(i);
114         end if;
115     end loop;
116 end process priority_evaluation;
117
118 reinit_transitions_time_evaluation : process(clock, reset_n)
119 begin
120     if (reset_n = '0') then
121         reinit_transitions_time <= (others => '0');
122     elsif rising_edge(clock) then

```

```
119   for i in 0 to output_arcs_number - 1 loop
120       if ( (((output_arcs_types(i) = arc_t(BASIC)) or (output_arcs_types(i) = arc_t(TEST)))
121           and (s_marking s_output_token_sum < output_arcs_weights(i))
122           and (s_output_token_sum > 0))
123         or output_transitions_fired(i) = '1' )
124       then
125         reinit_transitions_time(i) <= '1';
126       else
127         reinit_transitions_time(i) <= '0';
128       end if;
129     end loop;
130   end if;
131   end process reinit_transitions_time_evaluation;
132
133 end place_architecture;
```

LISTING A.1: The place design in concrete VHDL syntax.

Appendix B

The transition design in concrete and abstract VHDL syntax

```

1  entity transition is
2    generic(
3      transition_type : transition_t := NOT_TEMPORAL;
4      input_arcs_number : natural := 1;
5      conditions_number : natural := 1;
6      maximal_time_counter : natural := 1
7    );
8    port(
9      clock : in std_logic;
10     reset_n : in std_logic;
11     input_conditions : in std_logic_vector(conditions_number-1 downto 0);
12     time_A_value : in natural range 0 to maximal_time_counter;
13     time_B_value : in natural range 0 to maximal_time_counter;
14     input_arcs_valid : in std_logic_vector(input_arcs_number-1 downto 0);
15     reinit_time : in std_logic_vector(input_arcs_number-1 downto 0);
16     priority_authorizations : in std_logic_vector(input_arcs_number-1 downto 0);
17     fired : out std_logic
18   );
19 end transition;
20
21 architecture transition_architecture of transition is
22
23   signal s_condition_combination : std_logic;
24   signal s_enabled : std_logic;
25   signal s_firable : std_logic;
26   signal s_firing_condition : std_logic;
27   signal s_priority_combination : std_logic;
28   signal s_reinit_time_counter : std_logic;
29   signal s_time_counter : natural range 0 to maximal_time_counter;
30
31 begin
32

```

```

33 condition_evaluation : process(input_conditions)
34     variable v_internal_condition : std_logic;
35 begin
36     v_internal_condition := '1';
37
38     for i in 0 to conditions_number - 1 loop
39         v_internal_condition := v_internal_condition and input_conditions(i);
40     end loop;
41
42     s_condition_combination <= v_internal_condition;
43 end process condition_evaluation;
44
45 enable_evaluation : process(input_arcs_valid)
46     variable v_internal_enabled : std_logic;
47 begin
48     v_internal_enabled := '1';
49
50     for i in 0 to input_arcs_number - 1 loop
51         v_internal_enabled := v_internal_enabled and input_arcs_valid(i);
52     end loop;
53
54     s_enabled <= v_internal_enabled;
55 end process enable_evaluation;
56
57 reinit_time_counter_evaluation : process(reinit_time, s_enabled)
58     variable v_internal_reinit_time_counter : std_logic;
59 begin
60     v_internal_reinit_time_counter := '0';
61
62     for i in 0 to input_arcs_number - 1 loop
63         v_internal_reinit_time_counter := v_internal_reinit_time_counter or reinit_time(i);
64     end loop;
65
66     s_reinit_time_counter <= v_internal_reinit_time_counter;
67 end process reinit_time_counter_evaluation;
68
69 time_counter : process(reset_n, clock)
70 begin
71     if (reset_n = '0') then
72         s_time_counter <= 0;
73     elsif falling_edge(clock) then
74         if ((s_enabled = '1') and (transition_type /= transition_t(NOT_TEMPORAL))) then
75             if (s_reinit_time_counter = '0') then
76                 if (s_time_counter < maximal_time_counter) then
77                     s_time_counter <= s_time_counter + 1;
78                 end if;

```

```

79         else
80             s_time_counter <= 1;
81         end if;
82     else
83         s_time_counter <= 0;
84     end if;
85 end if;
86 end process time_counter;
87
88 firing_condition_evaluation : process (s_enabled, s_condition_combination,
89     s_reinit_time_counter, s_time_counter)
90 begin
91     if ((s_condition_combination = '1')
92         and (s_enabled = '1')
93         and ((transition_type = transition_t(NOT_TEMPORAL))
94
95             or ((transition_type = transition_t(TEMPORAL_A_B))
96                 and (s_reinit_time_counter = '0')
97                 and (s_time_counter >= (time_A_value1))
98                 and (s_time_counter < time_B_value)
99                 and (time_A_value /= 0)
100                 and (time_B_value /= 0) )
101
102             or ((s_reinit_time_counter = '0')
103                 and (time_A_value /= 0)
104                 and (((transition_type = transition_t(TEMPORAL_A_A))
105                     and (s_time_counter = (time_A_value1)))
106                     or ((transition_type = transition_t(TEMPORAL_A_INFINITE))
107                         and (s_time_counter >= (time_A_value1)) )
108                 )
109
110             or ((transition_type /= transition_t(NOT_TEMPORAL))
111                 and (s_reinit_time_counter = '1')
112                 and (time_A_value = 1) )
113         ) then
114         s_firing_condition <= '1';
115     else
116         s_firing_condition <= '0';
117     end if;
118 end process firing_condition_evaluation;
119
120 priority_authorization_evaluation : process(priority_authorizations)
121     variable v_priority_combination : std_logic;
122 begin
123     v_priority_combination := '1';
124

```

```

125   for i in 0 to input_arcs_number - 1 loop
126       v_priority_combination := v_priority_combination and priority_authorizations(i);
127   end loop;
128
129   s_priority_combination <= v_priority_combination;
130 end process priority_authorization_evaluation;
131
132 firable : process(reset_n, clock)
133 begin
134     if (reset_n = '0') then
135         s_firable <= '0';
136     elsif falling_edge(clock) then
137         s_firable <= s_firing_condition;
138     end if;
139 end process firable;
140
141 fired_evaluation : process (s_firable, s_priority_combination)
142 begin
143     fired <= s_firable and s_priority_combination;
144 end process fired_evaluation;
145
146 end transition_architecture;

```

LISTING B.1: The transition design in concrete VHDL syntax.

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