

This is the Theory of Computation Course.

CS2410 CS2420 for IITH.

This is one of the most fundamental courses that you will see as part of the CSE curriculum.

We will try to find out

\* What is computation?

\* What are the questions for which computers can provide an answer?

What are the fundamental capabilities and limitations of computers?

Though computers have evolved significantly over the years, the theory of computation has remained applicable throughout.

Modern digital computers were developed

Modern digital computers were developed around and after the second world war. The theoretical basis was provided by Alan Turing in 1936.

"On Computable Numbers, with an Application to the Entscheidungsproblem."

Algorithms } How fast / efficiently can we  
Complexity Theory } solve a problem?

Computability Theory } Is something "computable" at all?

Automata Theory } Some computation models with some very limited resources.

\* Very little math as you know it.

But mostly discrete maths. You

... with Turing

But many more ...  
should get comfortable with theorems  
and proofs.

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Today is lecture 1. We will move to  
pre-recorded lectures and "Q&A meetings".  
All recordings and notes will be provided.  
We may not meet for all slots each week.

Evaluation — Weekly Quizzes  
2 or 3 exams (CS 2410 &  
CS 2420  
together)

Exam distribution to be announced next  
class.

Text! Michael Sipser: Introduction to  
Theory of Computation.

Communication through Google Classroom. All  
announcements to be posted there.

Exercise before next lecture: Read Chapter 0

Exercise before next lecture. read chapters  
of the book. Sets, Functions, Relations,  
Graphs, Logic, Theorem, Proofs etc. Try out  
some problems.

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Every Computational Problem is represented  
as a string input to a computer.

String will be over alphabet, say  $\Sigma$ .

Examples of Alphabet:  $\Sigma_1 = \{0, 1\}$  Binary

$\Sigma_2 = \{0, 1, 2, 3, \dots, 9\}$  Decimal

$\Sigma_3 = \{a, b, c, \dots, z\}$  English.

$\Sigma_4 = \{a, b, 0, 1, 2\}$ .

A string over an alphabet  $\Sigma$ , is a finite  
sequence of symbols written one after  
another.

Example: 011011 is a string in binary  
abc is a string in English

string is a string in English  
839 is a string in decimal

We say  $w$  is a string over  $\Sigma$ .

$|w|$  denotes the length of  $w$ .

$$w = 110110 \Rightarrow |w| = 6$$

$$x = 825 \Rightarrow |x| = 3$$

$$y = \text{hello} \Rightarrow |y| = 5$$

$$\epsilon \rightarrow \text{empty string} \Rightarrow |\epsilon| = 0$$

Q: Is "hello there!" a string over  
 $\{a, b, c, d, \dots, x, y, z\}$ ?  $\cup \{ \_ , ! \}$

$$w^R = \text{reverse of } w.$$

Concatenation:  $x$  is a string of length  $m$   
 $y$  is a string of length  $n$ .

$$x = \text{"top"} \quad y = \text{"hat"}$$

$$xy = \text{"tophat"}$$

$xy = \text{"tophat"}$

$x^R y = \text{"pot hat"}$

$x^k = \underbrace{xx \dots x}_{k \text{ times}}$

Kleene Star  $\left\{ x^* = \text{Set of all } x^k = \{x^k \mid k \geq 0\} \right\}.$

$\Sigma^* = \text{Set of all strings over } \Sigma.$

Substring:  $v$  is a substring of  $w$  if there exists strings  $x$  and  $y$  such that  $w = xvy$ .

A language over  $\Sigma$  is a set of strings over  $\Sigma$ . That is, a language over  $\Sigma$  is a subset  $A \subseteq \Sigma^*$ .

Example: Set of all binary strings with odd number of 1's is a language over  $\{0,1\}$ .

languages can be specified in several ways.

1. Brute Force listing.  $\{a, ab, abb, abbb, \dots\}$
2. language operations.  $ab^*$ .
3. Other set theoretic descriptions.

## Finite Automata

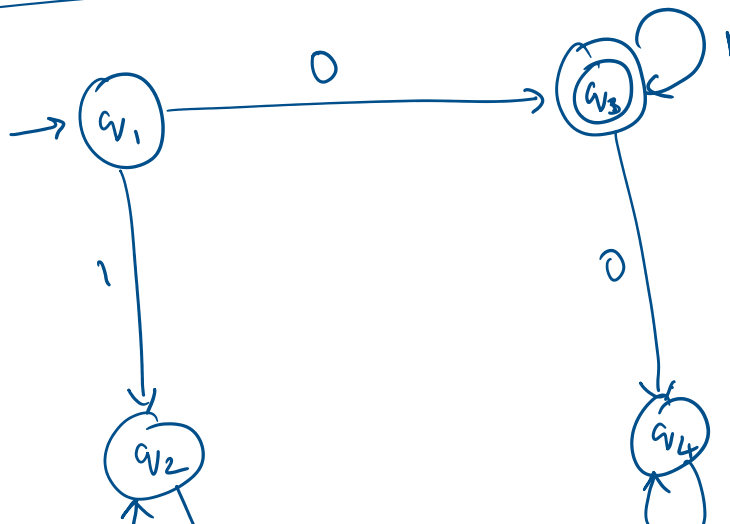
Computers with limited amount of memory.

State based devices.

Examples: Timers, door open/close controller,  
Thermostat.

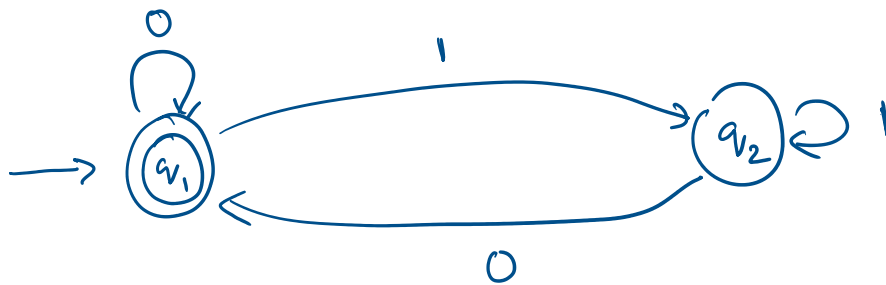
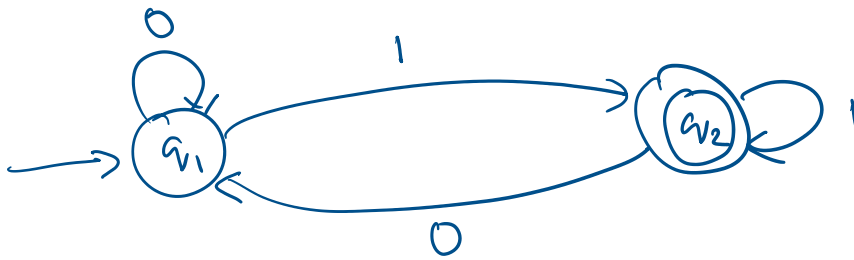
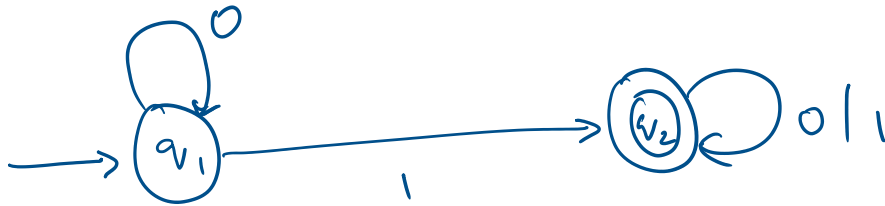
Why? There are abstract models that help us gain understanding.

## Deterministic Finite Automata





$$A = \{0, 01, 011, 0111, 01111, \dots\}$$



## DFA.

Def 1.5 : A deterministic finite automaton (DFA) is a 5-tuple  $(Q, \Sigma, \delta, q_0, F)$  where .

$Q$  is a finite set of states



when .

1.  $Q$  is a finite set of states
2.  $\Sigma$  is a finite alphabet .
3.  $\delta: Q \times \Sigma \rightarrow Q$  is the transition function
4.  $q_0 \in Q$  is the start state
5.  $F \subseteq Q$  is the set of accepting states .

Read example 1.11, 1.13, 1.17, 1.21.

Read definition 1.16.