Design and Analysis of Algorithms

CS202

Grading Policy

Weekly Quizzes: 20 marks

Programming + Written Assignments: 20 marks

• Exams: 3*20 = 60 marks

1. Introduction

2. Inductive Design

3. Divide and Conquer Paradigm

Algorithmic Problems

- 1. Given a collection of webpages, and a keyword, find the webpages most relevant to the keyword.
- 2. Given my current location, find all petrol pumps near me.
- 3. Schedule time-table for CSE courses at IITH.
- 4. Given a positive integer n, check if n is a prime.
- 5. Given a sequence of elements, arrange them in increasing order.

Goals:

- 1. Design correct and efficient algorithms
- 2. Understand and apply standard algorithms
- 3. Learn broad design techniques

E.g. 1: Closest pair problem

Input:
$$(x_1, y_1), (x_2, y_2), \dots, (x_n, y_n)$$

Output: A pair (x_i, y_i) , (x_j, y_j) such that $sqrt[(x_i-y_i)^2+(x_j-y_j)^2]$ is minimized.

Instance: (0,0), (0.3,0.4), (-0.5,3.6), (3,4)

Input Size: n

Algorithm 1: Closest pair

For each pair of points, compute the distance between them.

Find the minimum of all these distances, and return the corresponding points.

Algorithm 1: Closest pair

```
Min_so_far = -\infty

For each i in 1 to n-1

For each j>i

Find d= sqrt[(x_i-y_i)^2+(x_j-y_j)^2]

Min_so_far=min(d,Min_so_far)
```

Time Complexity of algorithm A

Input size: n

The maximum number of instructions made by A over all inputs of size n-> T(n)

[Worst-case complexity]

Time Complexity of Linear Search

<u>Input:</u> A[1,...,n], x

Output: Some position of x in A, if present

Lin_Search: For i=1 to n, if x=A[i] return i.

T(n)=n

Analysis of correctness and time complexity

Number of instructions<3n^2
Processor which can execute 3*10^9 instructions per second.

n: 10^6 Time: 1000s (3*10^12/3*10^9) ~ 16 min

n: 10^7 Time: 1600 min ~1.16 days

n: 10^8 Time: 116 days

Closest pair admits O(n log n) time algorithm.

n: 10^6 Time: 20 ms

n: 10⁷ Time: 0.23 s

n: 10^8 Time: 2.7 s



RAM model

Input Size:

Input: A positive integer n

Output: Find √n

Accurate size: Number of bits

log, n

Input Size:

```
Input: Two positive integers m,n
Output: Find mn
Log 2m + log 2n
Input size in terms of space (#bits) to store m,n.
1421424
8342414 12 k bits: 00000 (k times) to 11111 (k
times): 0 to 2^k-1 = 2^k-1 = 
    56
```

#hits to represent an integer $m \cdot \log 2 (m+1)$

RAM Model:

- 1. Random access: Unit cost for read/write of A[i]
- 2. Each word of data is limited in size (#bits)
- 3. Each instruction is unit cost: arithmetic operations (+,-,*), comparison.

Find the input size

Input: x, y in {1,2,...,n}. Output: xy
 log 2 n

2. Input: A,B: n by n matrix. Output: AB (n^2)

Input size: 2n^2 (number of elements)* size of each element (measured in bits)

Find the time complexity of the natural algo

1. Input: x, y in {1,2,...,n}. Output: xy

2. Input: A,B: n by n matrix. Output: AB

Paradigm 0: Induction

Examples: Linear search, Adding x to all the elements, calculating Fibonacci numbers, finding the minimum element of an unsorted array.

Inductive Design

```
Input: A[1,2,...,n]
Output: f(A[1,2,...,n])
```

```
Algorithm:
For i=1 to n,
Update value of f on seeing A[i]
Compute f(A[1]),f(A[1,2]),f(A[1,2,3]),...
```

E.g. 2: Finding the Minimum in an array

Input: A[1,2,...,n]

Output: Min(A[1,2,...,n])

Example Input: 14,3,-7,5,0,-2,8,5,10,4

Example Output: -7

Algorithm to find minimum

```
Min_so_far=A[1]
For i=2 to n
Set Min_so_far=min(Min_so_far,A[i])
```

Divide & Conquer Strategy

E.g. 1: Given a positive integer n, find \sqrt{n} .

```
i=1
While (i*i<n)
   Increment I
Return (i-1)</pre>
```

E.x. 1:

Finding a local minimum

3,7,0,-2,12,8,10,1,4,2

A local minimum is an element A[i] such that it is:

A[i] <= A[i-1] and A[i] <= A[i+1]