Database Management Systems (DBMS)

Lec 5: Relational model of data (Cont.)

Ramesh K. Jallu

IIIT Raichur

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Recap

- Some warmup exercise
- The join operation
 - $R \bowtie_{< \text{join condition}>} S = \sigma_{< \text{join condition}>} (R \times S)$

Today's plan

- Answers to the take-home queries
- The other variations of the join operation
 - Equijoin and Natural join
- The division operation (÷)

Exercise: Bank database

- Branch(Name, Assets, City)
- Customer(**ID**, Customer_Name, Street, City)
- Loan (**Loan_Number**, Branch_name, Amount)
- Borrower(BC_ID, BL_Number)
- Account (Account_Number, AB_name, Balance)
- Depositor (**DC_ID** , **A_number**)

Take-home queries

- 1. Find the customer name for the last three queries with the same requirement
 - Find all customers IDs who have loan 10,000/- or account balance > 5,000/-
 - Find the IDs of all customers who have a loan at the Raichur branch
 - Find the IDs of all customers who have a loan at the Raichur branch but do not have an account at any branch of the bank
- 2. Find the largest account balance
 - Hint: Use Rename, Cartesian, and Set Minus operations

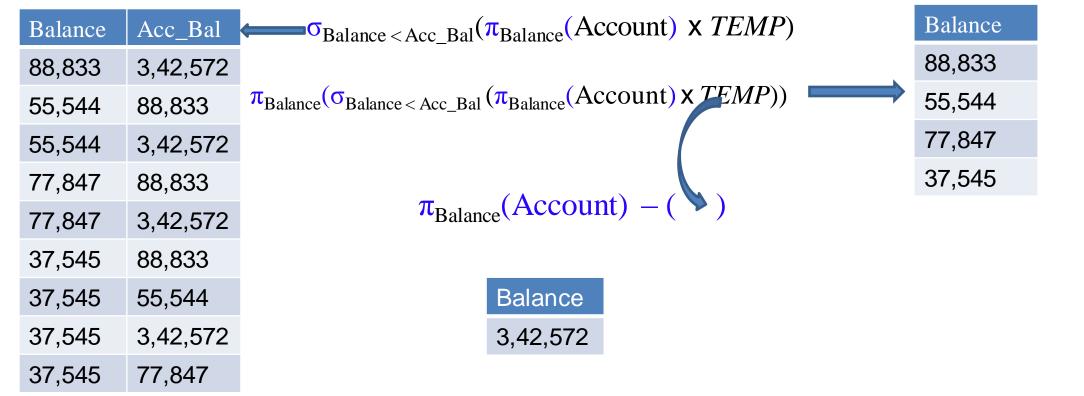
Answers

- 1. Use the join operation with the previous solutions with join condition
 - π_{BC_ID} ($\sigma_{Amount = 10000}$ (Borrower)) $\cup \pi_{DC_ID}$ ($\sigma_{Account_Number = A_Number}$ (Account x Depoister))
 - π_{BC_ID} ($\sigma_{Branch_Name = 'Raichur'}$ AND $BL_Number = Loan_Number$ (Borrower x Loan))
 - TEMP $\leftarrow \pi_{BC_ID}$ ($\sigma_{Branch_Name = 'Raichur' AND BL_Number = Loan_Number}$ (Borrower x Loan)) TEMP $- \pi_{DC \ ID}$ (Depoister)
- 2. Strategy:
 - 1. Find those balances that are *not* the largest
 - Use set difference to find those account balances that were not found in the Step 1

Account (Account_Number, AB_name, Balance)

Account TEMP

Account_Number	AB_name	Balance	Acc_Bal	
63165065258	Raichur	88,833	88,833	
54674498364	Hyderabad	55,544	55,544	ρ_{TEMP} (Acc_Bal)($\pi_{Balance}$ (Account))
59369845487	Delhi	3,42,572	3,42,572	
36465594547	Pune	77,847	77,847	
76994489457	Chennai	37,545	37,545	



Variations of join: Equi join and Natural join

- A general join condition: <cond> AND<cond>AND....AND<cond>
 - Theta join: $A_i \theta B_j$ such that $dom(A_i) = dom(B_j)$, and θ is one of the comparison operators $\{=, <, \le, >, \ge, \ne\}$
 - Equi join: Join operation with join condition with equality comparisons, I.e., θ is "="

Example from last class

EMPLOYEE

Fname	Minit	Lname	Ssn	Bdate	Address		Salary	Super_ssn	Dno
John	В	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	М	30000	333445555	5
Franklin	Т	Wong	333445555	1955-12-08	638 Voss, Houston, TX	М	40000	888665555	5
Alicia	J	Zelaya	999887777	1968-01-19	3321 Castle, Spring, TX	F	25000	987654321	4
Jennifer	S	Wallace	987654321	1941-06-20	291 Berry, Bellaire, TX	F	43000	888665555	4
Ramesh	K	Narayan	666884444	1962-09-15	975 Fire Oak, Humble, TX	М	38000	333445555	5
Joyce	Α	English	453453453	1972-07-31	5631 Rice, Houston, TX	F	25000	333445555	5
Ahmad	V	Jabbar	987987987	1969-03-29	980 Dallas, Houston, TX	М	25000	987654321	4
James	Е	Borg	888665555	1937-11-10	450 Stone, Houston, TX	М	55000	NULL	1

DEPARTMENT

Dname	Dnumber	Mgr_ssn	Mgr_start_date
Research	5	333445555	1988-05-22
Administration	4	987654321	1995-01-01
Headquarters	1	888665555	1981-06-19

Query: Retrieve the name of the manager of each department

DEPT_MGR \leftarrow DEPARTMENT $\bowtie_{Mgr_ssn=Ssn}$ EMPLOYEE RESULT $\leftarrow \pi_{Dname, Lname, Fname}(DEPT_MGR)$

DEPT MGR

Dname	Dnumber	Mgr_ssn	 Fname	Minit	Lname	Ssn	
Research	5	333445555	 Franklin	Т	Wong	333445555	
Administration	4	987654321	 Jennifer	S	Wallace	987654321	
Headquarters	1	888665555	 James	E	Borg	888665555	

RESULT

Dname	Lname	Fname
Research	Wong	Franklin
Administration	Wallace	Jennifer
Headquarters	Borg	James

Natural join (*)

- An equi join operation followed by the removal of the superfluous attributes
- Notation: **R*****S**
- The natural join of two relations $R(A_1,A_2,...,A_n)$ and $S(B_1,B_2,...,B_m)$, where Ai=Bj, is a new relation $Q(A_1,A_2,...,A_i,...,A_n,B_1,B_2,...,B_{j-1},B_{j+1},B_m)$
 - Q contains those tuples satisfying equi join condition
 - Q removes B_i

Observations

- Join attribute: Natural Join requires that the two join attributes have the same name in both relations
- If this is not the case, a renaming operation is applied first
- There can be a list of join attributes from each relation, and each corresponding pair must have the same name

EMPLOYEE

Fname	Minit	Lname	Ssn	Bdate	Address		Salary	Super_san	Dno
John	В	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	М	30000	333445555	5
Franklin	Т	Wong	333445555	1955-12-08	638 Voss, Houston, TX M		40000	888665555	5
Alicia	J	Zelaya	999887777	1968-01-19	3321 Castle, Spring, TX	F	25000	987654321	4
Jen nifer	S	Wallace	987654321	1941-06-20	291 Berry, Bellaire, TX	F	43000	888665555	4
Ramesh	K	Narayan	666884444	1962-09-15	975 Fire Oak, Humble, TX	М	38000	333445555	5
Joyce	Α	English	453453453	1972-07-31	5631 Rice, Houston, TX	F	25000	333445555	5
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James	Е	Borg	888665555	1937-11-10	450 Stone, Houston, TX	М	55000	NULL	1

DEPARTMENT

Dname	Dnumber	Mgr_san	Mgr_start_date
Research	5	333445555	1988-05-22
Administration	4	987654321	1995-01-01
Headquarters	1	888665555	1981-06-19

DEPT_LOCATIONS

Diocation
Houston
Stafford
Bellaire
Sugarland
Houston

WORKS_ON

Essn	Pno	Hours
123456789	1	32.5
123456789	2	7.5
666884444	3	40.0
453453453	1	20.0
453453453	2	20.0
333445555	2	10.0
333445555	3	10.0
333445555	10	10.0
333445555	20	10.0
999887777	30	30.0
999887777	10	10.0
987987987	10	35.0
987987987	30	5.0
987654321	30	20.0
987654321	20	15.0
888665555	20	NULL

PROJECT

Pname	Pnumber	Plocation	Dnum
ProductX	1	Bellaire	5
ProductY	2	Sugarland	5
ProductZ	3	Houston	5
Computerization	10	Stafford	4
Reorganization	20	Houston	1
Newbenefits	30	Stafford	4

DEPENDENT

Essn	Dependent_name	Sex	Bdate	Relationship
333445555	Alice	F	1986-04-05	Daughter
333445555	Theo dore	М	1983-10-25	Son
333445555	Joy	F	1958-05-03	Spouse
987654321	Abner	М	1942-02-28	Spouse
123456789	Michael	М	1988-01-04	Son
123456789	Alice	F	1988-12-30	Daughter
123456789	Elizabe th	F	1967-05-05	Spouse

PROJ_DEPT ← PROJECT * (Dname, Dnum, Mgr_ssn, Mgr_start_date) (DEPARTMENT

PROJ DEPT

Pname	<u>Pnumber</u>	Plocation	Dnum	Dname	Mgr_ssn	Mgr_start_date
ProductX	1	Bellaire	5	Research	333445555	1988-05-22
ProductY	2	Sugarland	5	Research	333445555	1988-05-22
ProductZ	3	Houston	5	Research	333445555	1988-05-22
Computerization	10	Stafford	4	Administration	987654321	1995-01-01
Reorganization	20	Houston	1	Headquarters	888665555	1981-06-19
Newbenefits	30	Stafford	4	Administration	987654321	1995-01-01

DEPT_LOCS ← DEPARTMENT * DEPT_LOCATIONS

DEPT_LOCS

Dname	Dnumber	Mgr_ssn	Mgr_start_date Location	
Headquarters	1	888665555	1981-06-19 Hous	
Administration	4	987654321	1995-01-01 Staf	
Research	5	333445555	1988-05-22	Bellaire
Research	5	333445555	1988-05-22	Sugarland
Research	5	333445555	1988-05-22	Houston

The division operation

- Notation: R ÷ S
- Used to deal queries which contain the keyword All/every
 - Retrieve students IDs who enrolled in every course
 - Retrieve details of a customer who has account in all the banks in a city
 - Retrieve the names of employees who work on all the projects that 'John Smith' works on

How it works?

- 1. Division is applied to two relations R(X) and S(Y) such that $Y \subseteq X$
- 2. Let Z = X Y l.e., $X = Y \cup Z$
- 3. $T(Z) = R(X) \div S(Y)$
 - For every tuple t in T(Z), the values in t must appear in R in combination with every tuple in S

Explanation

- 1. R(X) such that $X = \{A_1, A_2, ..., A_n, B_1, B_2, ..., B_m\}$ and S(Y) such that $Y = \{B_1, B_2, ..., B_m\}$
- 2. Let $Z = X Y = \{A_1, A_2, ..., A_n\}$
- 3. $R(X) \div S(Y)$ Contains attributes $A_1, A_2, ..., A_n$ with tuple values $\langle a_1, a_2, ..., a_n \rangle$ such that for every tuple $\langle b_1, b_2, ..., b_m \rangle$ in S, $\langle a_1, a_2, ..., a_n, b_1, b_2, ..., b_m \rangle$ is in R

Example

STUDENT

Student_ID	Course_ID
01CS19	MA101
16CS19	CS201
01CS19	MA203
01CS19	CS101
16CS19	MA101
18CS20	MA101
01CS19	CS201

COURSE

Course_ID
MA101
CS201
MA203
CS101

RESULT

Student_ID 01CS19 R

Α	В		
a1	b1		
a2	b1		
аЗ	b1		
a4	b1		
a1	b2		
аЗ	b2		
a2	b3		
a3	b3		
a4	b3		
a1	b4		
a2	b4		

a3

b4

S

Α
a1
a2
аЗ

Т

В
b1
b4

RESULT = STUDENT + COURSE

Exercise

Query: Retrieve the names of employees who work on **all** the projects that 'John Smith' works on

- **Strategy:** 1. Retreive the projects that 'John Smith' works on: **R**₁
 - 2. Retreive the employees who work on which projects: R₂
 - 3. $R_2 \div R_1$

EMPLOYEE

Fname	Minit	Lname	San	Bdate	Address	Sex	Salary	Super_san	Dno
John	В	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	М	30000	333445555	5
Franklin	Т	Wong	333445555	1955-12-08	638 Voss, Houston, TX	М	40000	888665555	5
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James	Е	Borg	888665555	1937-11-10	450 Stone, Houston, TX M 55000 NULL		1		

DEPARTMENT

Dname	Dnumber	Mgr_san	Mgr_start_date
Research	5	333445555	1988-05-22
Administration	4	987654321	1995-01-01
Headquarters	1	888665555	1981-06-19

Pname	Pnumber	Plocation	Dnum
ProductX	1	Bellaire	
ProductY	2	Sugarland	5
ProductZ	3	Houston	5
Computerization	10	Stafford	4
Reorganization	20	Houston	1
Newbenefits	30	Stafford	4

DEPT_LOCATIONS

Diocation

Houston Stafford Bellaire Sugarland

Houston

Dnumber

DEPENDENT

Esan	De pendent_name	Sex	Bdate	Relationship	
333445555	Alice	F	1986-04-05	Daughter	
333445555	Theo dore	М	1983-10-25	Son	
333445555	Joy	F	1958-05-03	Spouse	
987654321	Abner	М	1942-02-28	Spouse	
123456789	Michael	М	1988-01-04	Son	
123456789	Alice	F	1988-12-30	Daughter	
123456789	Elizabe th	F	1967-05-05	Spouse	

SMITH ← or Fname = 'John' AND Lname = 'Smith' (EMPLOYEE)

 $SMITH_{PNOS} \leftarrow \pi_{Pno}(WORKS_{ON} \bowtie_{Essn = Ssn} SMITH)$

 $SSN_PNOS \leftarrow \pi_{Essn, Pno} (WORKS_ON)$

SSNS (Ssn) ← SSN_PNOS ÷ SMITH_PNOS

RESULT $\leftarrow \pi_{\text{Fname, Lname}}$ (SSNS * EMPLOYEE)

SSN PNOS

Essn	Pno
123456789	1
123456789	2
666884444	3
453453453	1
453453453	2
333445555	2
333445555	3
333445555	10
333445555	20
999887777	30
999887777	10
987987987	10
987987987	30
987654321	30
987654321	20
888665555	20

SMITH_PNOS

Pno
1
2

SSNS

Ssn	
123456789	
453453453	

<u>Pno</u>	Hours
1	32.5
2	7.5
3	40.0
1	20.0
2	20.0
2	10.0
3	10.0
10	10.0
20	10.0
30	30.0
10	10.0
10	35.0
30	5.0
30	20.0
20	15.0
20	NULL
	1 2 3 1 2 2 3 10 20 30 10 30 30 20

OPERATION	PURPOSE	NOTATION
SELECT	Selects all tuples that satisfy the selection condition from a relation R .	$\sigma_{< \text{selection condition}>}(R)$
PROJECT	Produces a new relation with only some of the attributes of <i>R</i> , and removes duplicate tuples.	$\pi_{< ext{attribute list}>}(R)$
THETA JOIN	Produces all combinations of tuples from R_1 and R_2 that satisfy the join condition.	$R_1 \bowtie_{< \text{join condition}>} R_2$
EQUIJOIN	Produces all the combinations of tuples from R_1 and R_2 that satisfy a join condition with only equality comparisons.	$R_1 \bowtie_{< \text{join condition}>} R_2$, OR $R_1 \bowtie_{(< \text{join attributes 1}>)}$, (< join attributes 2>) R_2
NATURAL JOIN	Same as EQUIJOIN except that the join attributes of R_2 are not included in the resulting relation; if the join attributes have the same names, they do not have to be specified at all.	R_1 * <poin condition=""> R_2, OR R_1* (<poin 1="" attributes="">), (<poin 2="" attributes="">) R_2 OR R_1* R_2</poin></poin></poin>
UNION	Produces a relation that includes all the tuples in R_1 or R_2 or both R_1 and R_2 ; R_1 and R_2 must be union compatible.	$R_1 \cup R_2$
INTERSECTION	Produces a relation that includes all the tuples in both R_1 and R_2 ; R_1 and R_2 must be union compatible.	$R_1 \cap R_2$
DIFFERENCE	Produces a relation that includes all the tuples in R_1 that are not in R_2 ; R_1 and R_2 must be union compatible.	$R_1 - R_2$
CARTESIAN PRODUCT	Produces a relation that has the attributes of R_1 and R_2 and includes as tuples all possible combinations of tuples from R_1 and R_2 .	$R_1 \times R_2$
DIVISION	Produces a relation $R(X)$ that includes all tuples $t[X]$ in $R_1(Z)$ that appear in R_1 in combination with every tuple from $R_2(Y)$, where $Z = X \cup Y$.	$R_1(Z) \div R_2(Y)$

Complete set

- $\{\sigma, \pi, U, \rho, -, \times\}$: Any other relational operation can be expressed as a combination of these
- How can we express the *intersection* operation?
- The division operation can be expressed as a sequence of π , \times , and –

Example

R			
Α	В		
a1	b1		
a2	b1		
аЗ	b1		
a4	b1		
a1	b2		
аЗ	b2		
a2	b3		
аЗ	b3		
a4	b3		
a1	b4		

 $T_1 \leftarrow \pi_Z(R)$ $T_2 \leftarrow \pi_Z \left((S \times T_1) - R \right)$ $T \leftarrow T_1 - T_2$

 $S \times T_1$

b2

b3

b4

 b_2 a_1

 $a_1 b_1$

 a_1 b_3

 a_1 b_4

 a_2

 a_2 b_2

 $a_2 b_3$

 a_2

 a_3 b_1

 $a_3 b_2$

 $a_3 b_3$

 a_3 b_4

 $(S \times T_1) - R$

 $a_1 b_3$

 $a_2 b_2$

 T_2

В

 b_2

 b_1

 b_4

 $X = \{A,B\}, Y = \{A\}, Z = \{B\}$

b4

b4

Thank you!