

Query optimization for the quantifier ONLY

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In this note we will show how an **only** quantifier gets translated and optimized. Consider the following query:

‘Find the pairs (s, b) where s is the sid of a student and b is the bookno of a book such that student s **only** bought the books that cite book b .’

$$\{(s.sid, b.bookno) | Student(s) \wedge Book(b) \wedge \{t.bookno | \exists t (Buys(t) \wedge t.sid = s.sid)\} \subseteq \{c.bookno | \exists c (Cites(c) \wedge c.citedbook = b.bookno)\}\}$$

We expect the final RA expression to be

$$\pi_{sid}(Student) \times \pi_{bookno}(Book) - \pi_{sid, bookno}(Buys \times \pi_{bookno}(Book) - \pi_{sid, bookno, citedbookno}(Buys \bowtie Cites))$$

1 Formulation and translation in SQL

```
select s.sid, b.bookno
from   student s, book b
where  not exists (select 1
                  from   buys t
                  where   t.sid = s.sid and
                        not exists (select 1
                                  from   cites c
                                  where   t.bookno = c.bookno and
                                            c.citedbookno = b.bookno));
```

Translation in SQL of the **not exists** set predicates

```
select sid, bookno
from student, book
except
select sid, bookno
from (select s.*, b.*, t.sid as tsid, t.bookno as tbno
      from student s, book b, buys t
      where s.sid = t.sid
      except
      select s.*, b.*, t.*
      from student s, book b, buys t, cites c
      where s.sid = t.sid and t.bookno = c.bookno and c.citedbookno = b.bookno) q;
```

2 Translation to RA as expressed in SQL

```
select sid, bookno
from student cross join book
except
select sid, bookno
from (select s.*, b.*, t.sid as tsid, t.bookno as tbno
      from student s
      join buys t on (s.sid = t.sid)
      cross join book b
      except
      select s.*, b.*, t.*
      from student s
      join buys t on (s.sid = t.sid)
      join cites c on (t.bookno = c.bookno)
      join book b on (c.citedbookno = b.bookno)) q;
```

3 Optimization

3.1 Using functional constraints

Observe that we have the following functional constraints.

$$\begin{array}{ll} sid \rightarrow (sname) & \text{In Student relation} \\ bookno \rightarrow (title, price) & \text{In Book relation} \end{array}$$

These constraints permit us the "project-out" the sname, title and price attributes since they are not part of the output of the query. This gets us to the following query:¹

```
select sid, bookno
from   (select sid from student) s
       cross join (select bookno from book) b
except
select sid, bookno
from   (select s.sid, b.bookno, t.sid as tsid, t.bookno as tbno
       from   (select sid from student) s
              join buys t on (t.sid = s.sid)
              cross join (select bookno from book) b
       except
       select s.sid, b.bookno, t.sid, t.bookno
       from   (select sid from student) s
              join buys t on (s.sid = t.sid)
              natural join cites c
              join (select b.bookno from book b) b on c.citedbookno = b.bookno) q;
```

¹We have left out the `distinct` clause to aid with the readability of the query.

3.2 Utilizing foreign key constraints

Observe that we have the following foreign key constraints:

sid foreign key in *Buys* referencing primary key in *Student*
citedbookno foreign key in *Cites* referencing primary key in *Book*

Using these constraints, we can optimize the query to become the following.

```
select sid, bookno
from   (select sid from student) s
       cross join (select bookno from book) b
except
select sid, bookno
from   (select t.sid, t.bookno as tbno, b.bookno
       from   buys t
             cross join (select bookno from book) b
       except
       select t.sid, t.bookno, c.citedbookno
       from   buys t
             natural join cites c) q;
```

We get the RA expression:

$$\pi_{sid}(Student) \times \pi_{bookno}(Book) - \pi_{sid,bookno}(Buys \times \pi_{bookno}(Book) - \pi_{sid,bookno,citedbookno}(Buys \bowtie Cites)).$$