CS 747, Autumn 2020: Week 1, Lecture 1

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Autumn 2020

Multi-armed Bandits

- 1. The exploration-exploitation dilemma
- 2. Definitions: Bandit, Algorithm
- 3. ϵ -greedy algorithms
- 4. Evaluating algorithms: Regret

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Coin 1



 $\mathbb{P}\{\text{heads}\} = p_1$

Coin 2



 $\mathbb{P}\{\text{heads}\} = \textcolor{red}{p_2}$

Coin 3



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- p_1 , p_2 , and p_3 are unknown.
- You are given a total of 20 tosses.
- Maximise the total number of heads!

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Coin 3



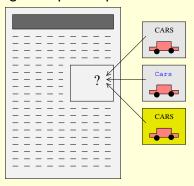
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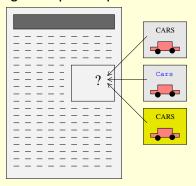
Let's play!

- Now we know: $p_1 = 0.6$, $p_2 = 0.3$, $p_3 = 0.8$.
- If you knew p_1, p_2, p_3 beforehand, how would you have played? How many heads would you have got in 20 tosses?

On-line advertising: Template optimisation

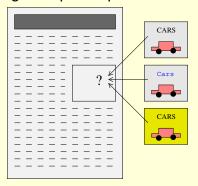


On-line advertising: Template optimisation



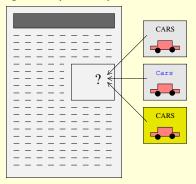
Clinical trials

On-line advertising: Template optimisation



- Clinical trials
- Packet routing in communication networks

On-line advertising: Template optimisation

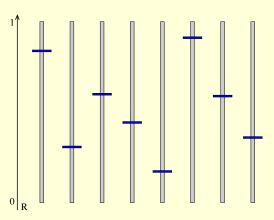


- Clinical trials
- Packet routing in communication networks
- Game playing and reinforcement learning

Multi-armed Bandits

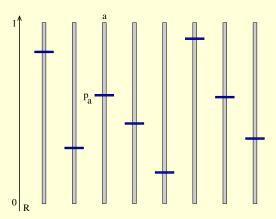
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Stochastic Multi-armed Bandits



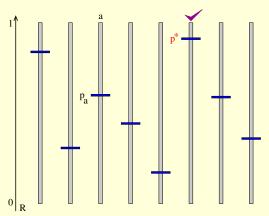
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- Highest mean is p*.

One-armed Bandits



[1]

1. https://pxhere.com/en/photo/942387.

For
$$t = 0, 1, 2, ..., T - 1$$
:

- Given the history $h^t = (a^0, r^0, a^1, r^1, a^2, r^2, \dots, a^{t-1}, r^{t-1}),$
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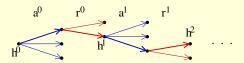
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- Formally: a deterministic algorithm is a mapping from the set of all histories to the set of all arms.

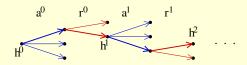
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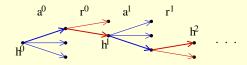
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- Note: The algorithm picks the arm to pull; the bandit instance returns the reward.

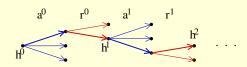




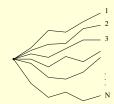
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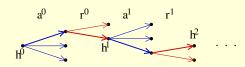


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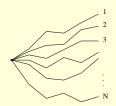


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How many histories possible if the algorithm is deterministic and rewards 0–1?

Multi-armed Bandits

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- If $t \le \epsilon T$, sample an arm uniformly at random.
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● *ϵ*G3

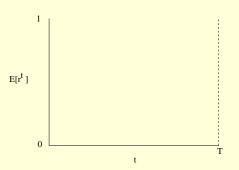
- With probability ϵ , sample an arm uniformly at random; with probability $1 - \epsilon$, sample an arm with the highest empirical mean.

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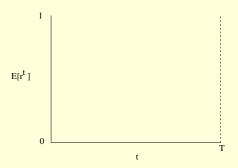
Visualising Performance

• Consider a plot of $\mathbb{E}[r^t]$ against t.



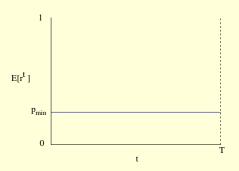
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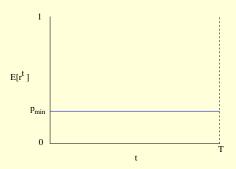


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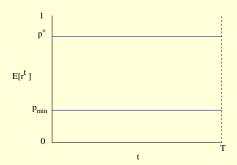
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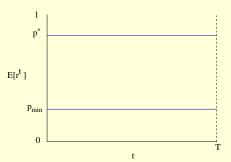
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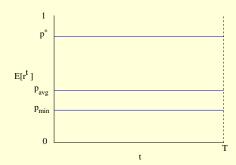


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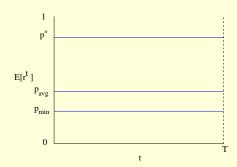


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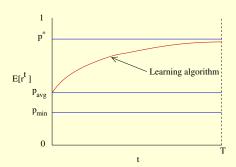
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- How will the graph look for a reasonable learning algorithm?



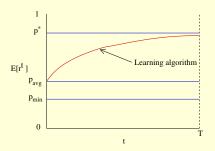
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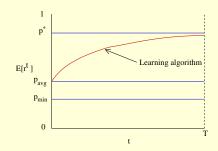
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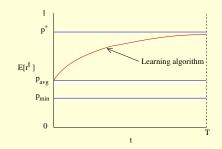
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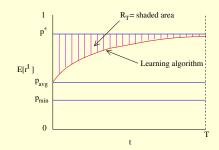
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 The (expected cumulative) regret of the algorithm for horizon T is the difference

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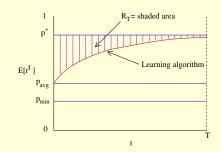
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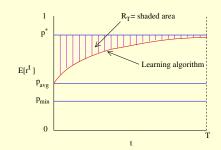


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• We would like R_T to be small, in fact for $\lim_{T\to\infty} \frac{R_T}{T} = 0$. Does this happen for ϵ G1, ϵ G2, ϵ G3?

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