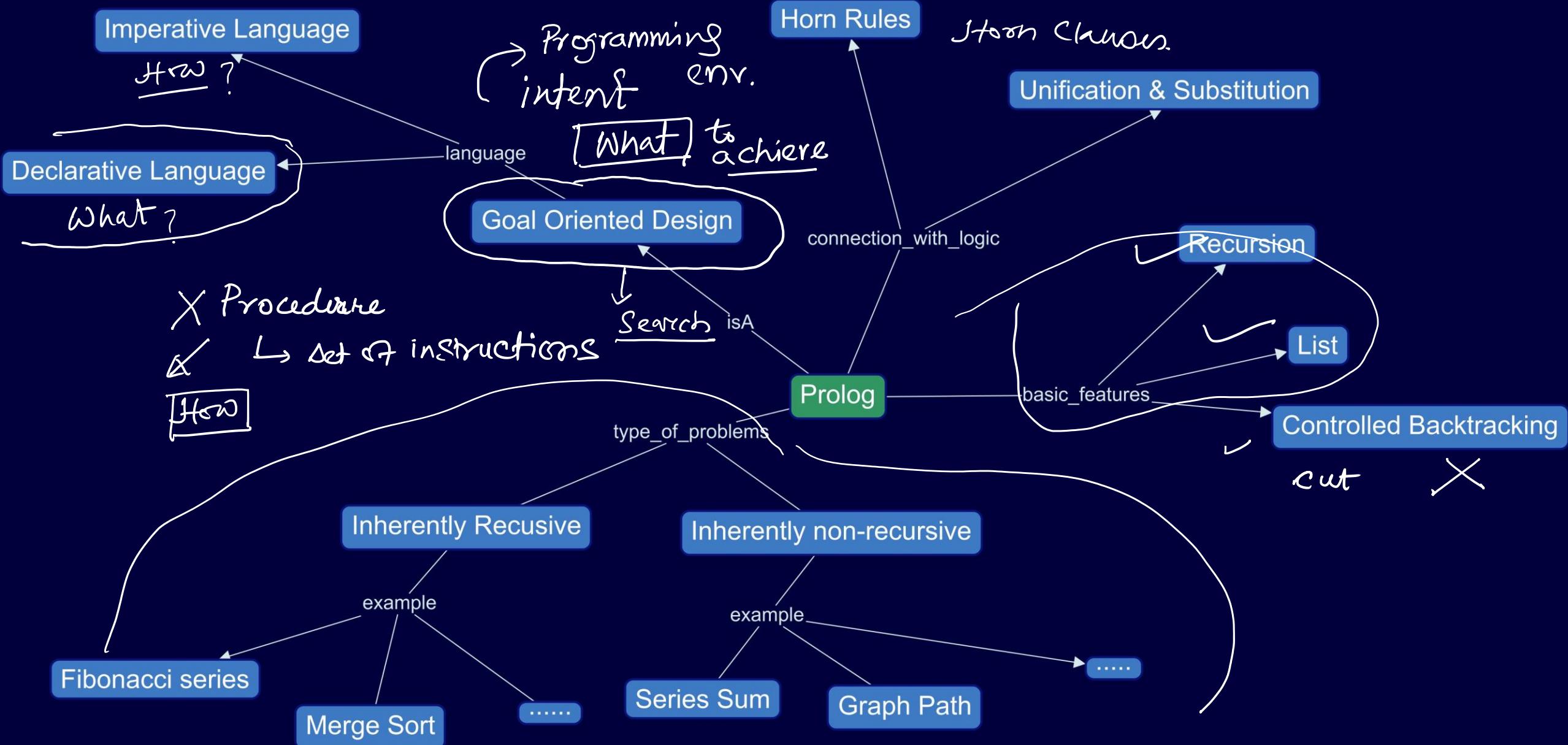


Prolog: Language for AI

Programming
Logic

AIFA, AI61005, 2021 Autumn

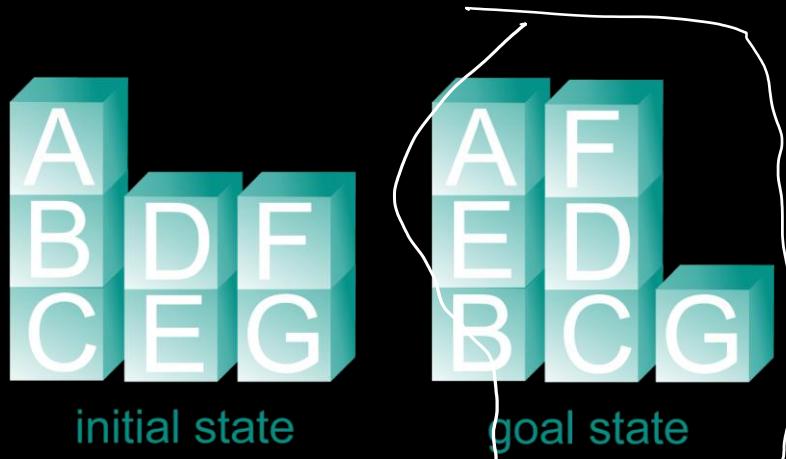
Plaban Kumar Bhowmick



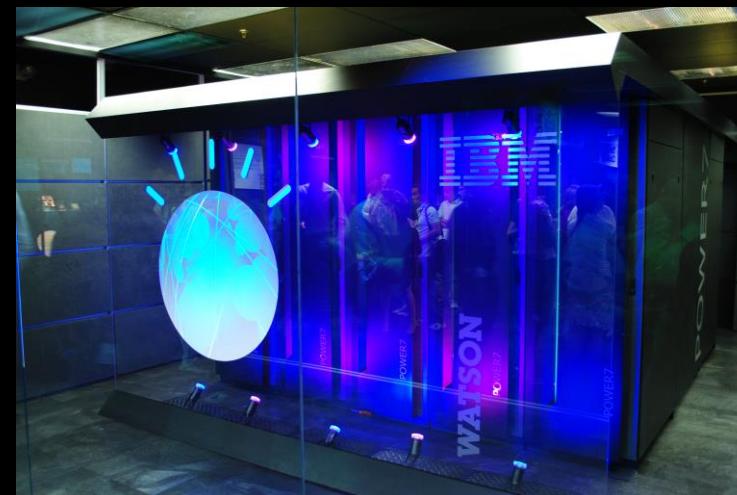
Prolog Programming for Artificial Intelligence: Ivan Bratko

<https://swish.swi-prolog.org/>

Goal Oriented Programming



Goal



Language Choice

- “Known” languages like FORTRAN, C/C++, Java, python
 - Imperative: How-type language
- Goal Oriented Languages (Declarative)
 - Declarative: What-type language
 - LISP → *functional*
 - *“Goal oriented programming is like reading Shakespeare in language other than English”* - Patrick Winston
 - ProLog: Truly what-type language

Imperative vs. Declarative

Programming languages

Imperative: Comprises a sequence of commands

— Imperative

 └ Procedural →
 └ Object-oriented

— Declarative

 └ Functional List
 └ Logic Piping
 └ Query SQL

Declarative: Declare what result we want and leave the language to come up with the procedure to produce them

Prolog and First-Order Logic

- Prolog language syntax

- Horn Clause: CNF with implicit quantifiers and with at most one positive literal

- $\text{child}(x) \wedge \text{male}(x) \Rightarrow \boxed{\text{boy}(x)}$

$$\neg \text{child}(\alpha) \vee \neg \text{male}(\alpha) \vee \underline{\text{boy}(\alpha)}$$

- Prolog proof procedure

- Resolution Principle

- Prolog goal matching

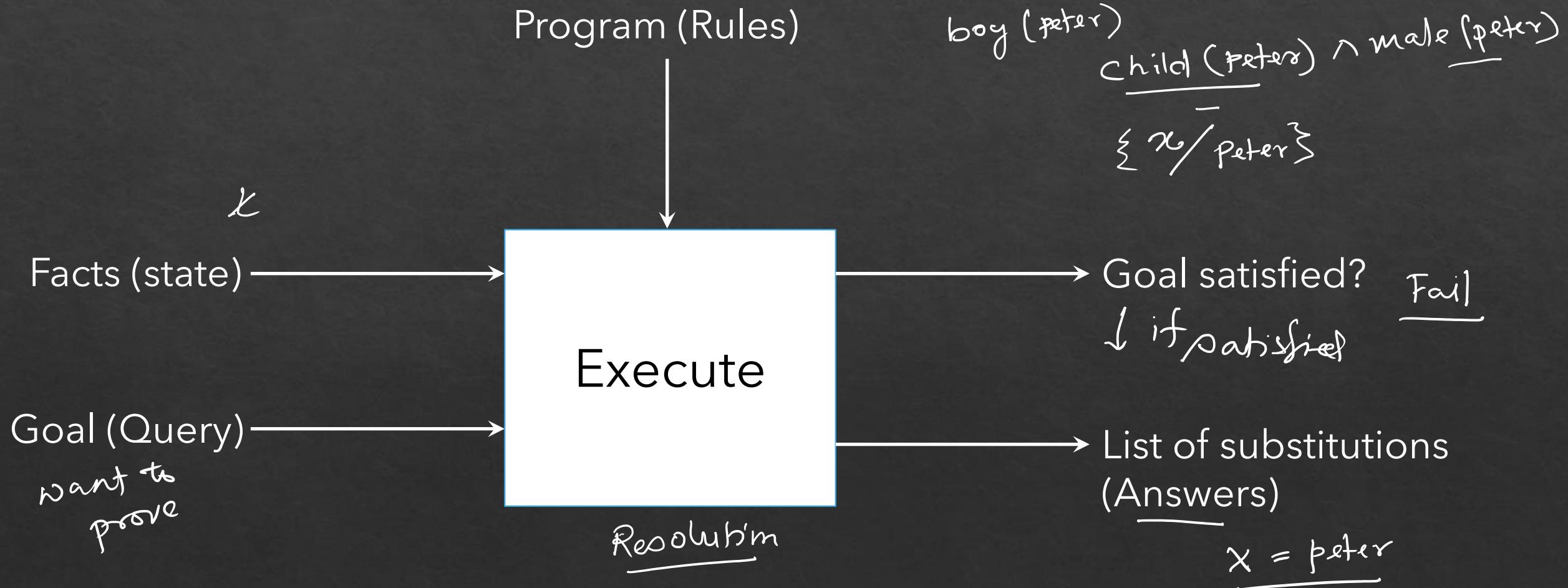
- Unification and substitution

Solving problem in Prolog
⇒ prove unit goal
 $\text{boy}(\text{peter}) \quad \left\{ \frac{x}{\text{peter}} \right\}$

A Confusion

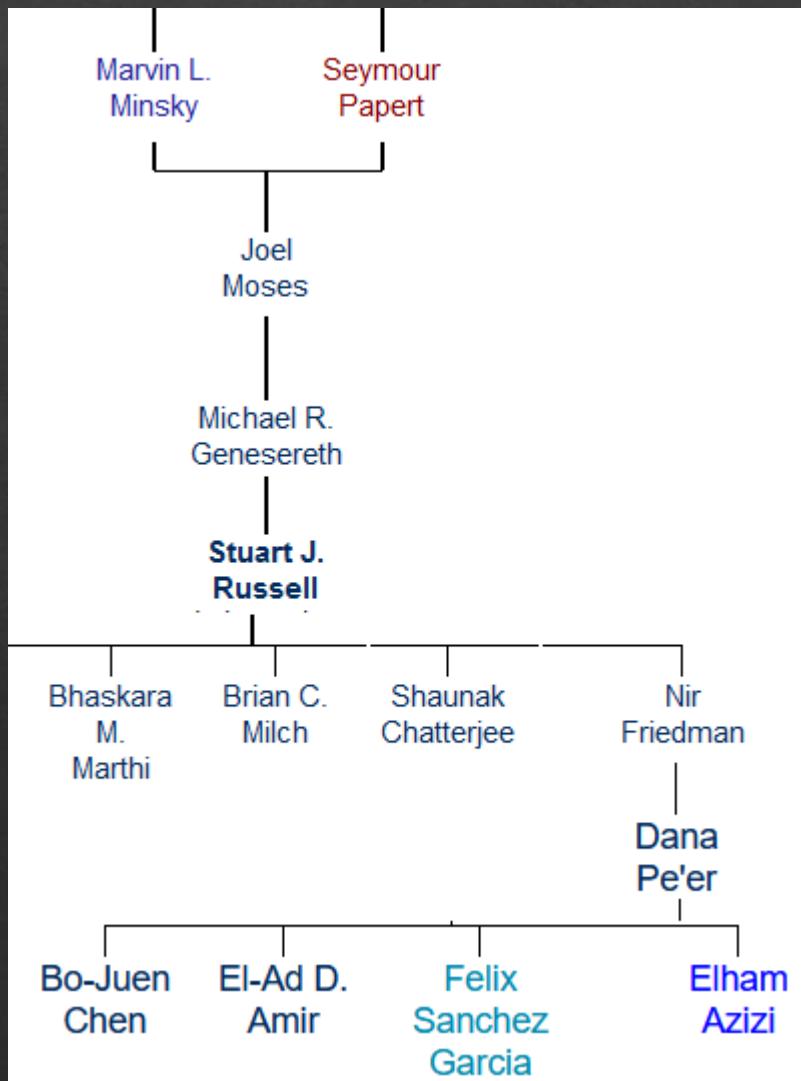
$\text{boy}(\text{peter}) \rightarrow \text{True}$
 Subgoal $\hookrightarrow \frac{\text{child}(\text{peter})}{\wedge \text{male}(\text{peter})}$
 $\rightarrow \text{True}$

Prolog Computation Model



Prolog Facts

✓



```
advisor(minsky, moses).  
advisor(papert, moses).  
advisor(moses, genesereth).  
advisor(genesereth, russell).  
advisor(russell, bhaskara).  
advisor(russell, milch).  
advisor(russell, shaunak).  
advisor(russell, friedman).  
advisor(friedman, dana).  
advisor(dana, felix).  
advisor(dana, chen).  
advisor(dana, amir).  
advisor(dana, azizi).  
  
male(felix).  
female(dana).
```

$\text{adviser}(x,y) \wedge \text{adviser}(y,z) \Rightarrow \text{grand_adviser}(x,z)$



Prolog Rules

```
{ grand_advisor(X, Z) :- advisor(X, Y), advisor(Y, Z).  
          _____|_____  
          head      body  
          (consequent) (antecedent )
```

IF there is a Y such that X is advisor of Y AND Y is advisor of Z
THEN X is a grand advisor of Z

Prolog rules are Horn Clauses:

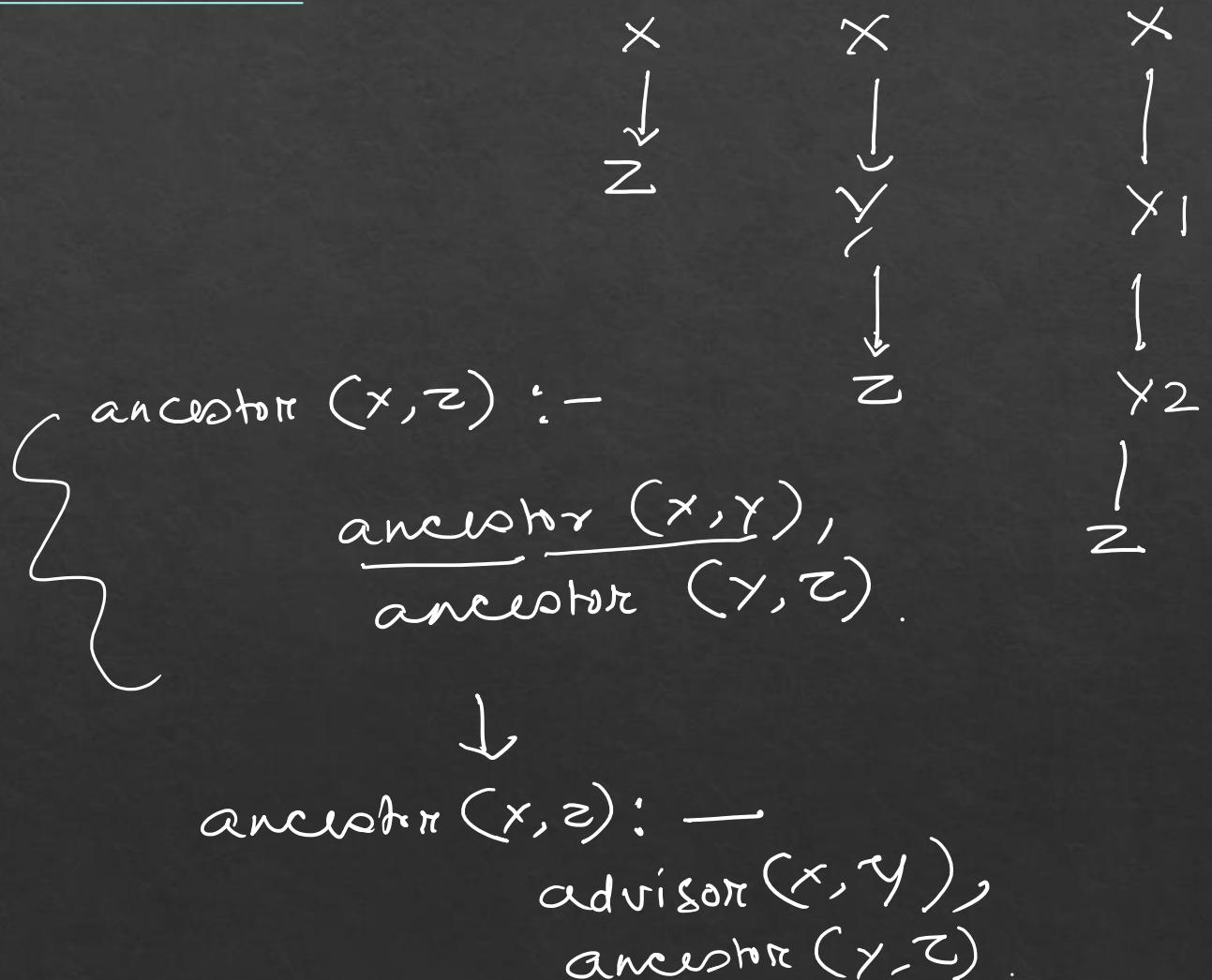
body $(P_{11} \overline{\vee} P_{12} \vee \dots P_{1m}) \overline{\wedge} \dots \wedge (P_{n1} \vee P_{n2} \vee \dots P_{nr}) \Rightarrow Q$ Generic

$Q : \neg P_{11} \vee P_{12}; \dots; P_{1m} \vee \dots \vee P_{n1}; P_{n2}; \dots; P_{nr}$

$\mathcal{Q} := \vdash \rightarrow \text{disjunct.}$

Prolog Rules: Recursion

```
ancestor(X, Z) :-  
    advisor(X, Z). ✓  
  
ancestor(X, Z) :-  
    advisor(X, Y),  
    advisor(Y, Z).  
  
ancestor(X, Z) :-  
    advisor(X, Y1),  
    advisor(Y1, Y2),  
    advisor(Y2, Z),
```



Prolog Rules: Recursion

```
ancestor(X, Z) :-  
    advisor(X, Z).
```

```
ancestor(X, Z) :-  
    advisor(X, Y),  
    ancestor(Y, Z).
```

X is an ancestor of Z if
X is an advisor of Y AND
Y is an ancestor of Z

How Prolog Answers?

\downarrow

\downarrow

\checkmark

$\text{ancestor}(X, Z) :- \text{advisor}(X, Z).$

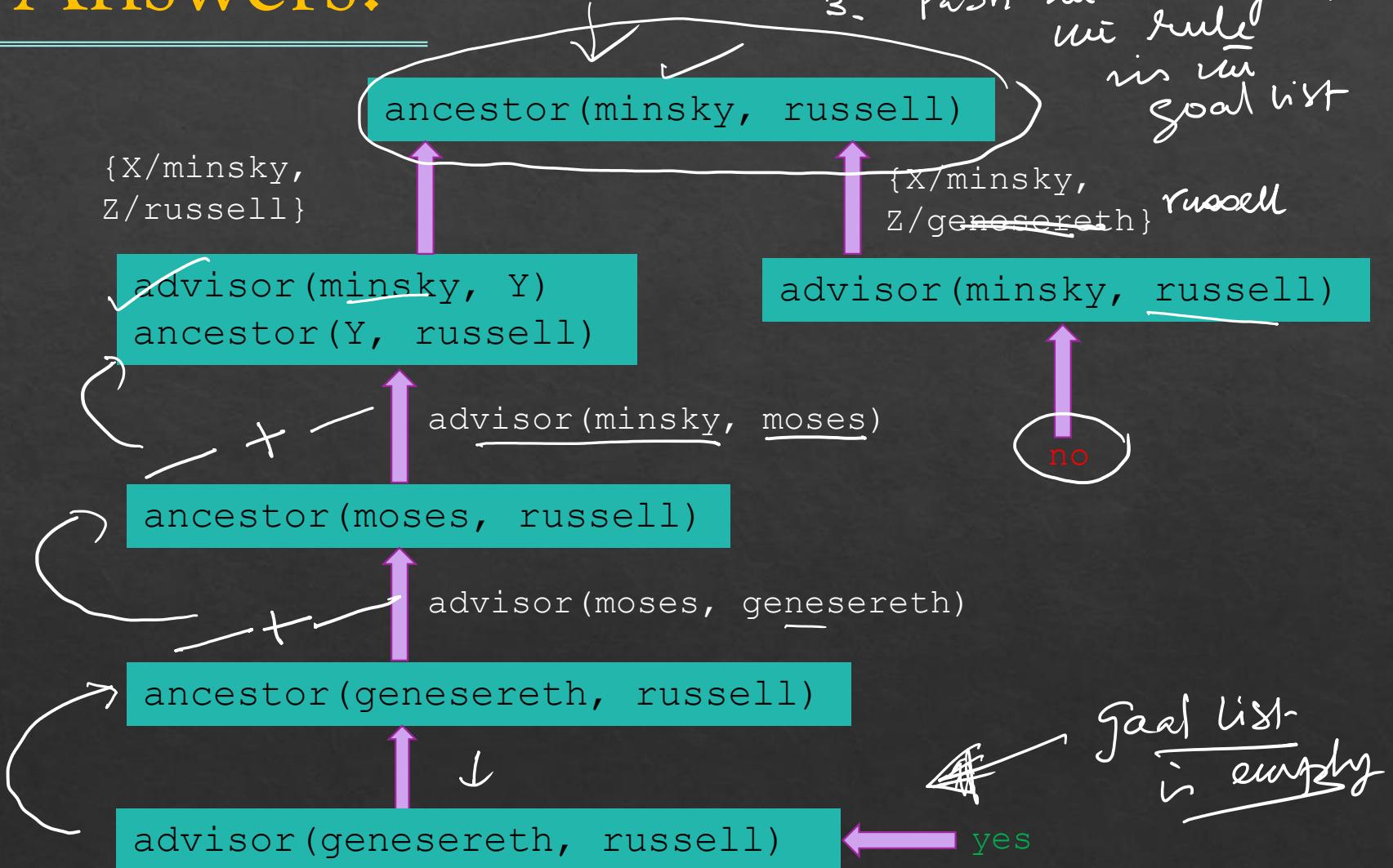
$\checkmark R \backslash$

$\text{ancestor}(X, Z) :-$

$\text{advisor}(X, Y), \checkmark$

$\text{ancestor}(Y, Z).$

?- $\text{ancestor}(\text{minsky}, \text{russell})$



1. Start with a goal list
2. match in goal with rule head.
3. push in body of rule in goal list

~~Goal list
is empty~~

yes

How Prolog Answers?

function EXECUTE(program, GoalList) **returns** [success, instance]/failure

if empty(GoalList) **then return** [true, instance]

else

goal = head(GoalList)

other_goals = tail(GoalList)

satisfied = false ✓

while not satisfied and more clauses in program **then**

✓ $H : - B_1, \dots, B_n$ //Next clause C of the program

$H' : - B'_1, \dots, B'_n$ //Variant C' of clause C →

[match_OK, instance] = match(goal, H')

if match_OK then

new_goals = append([B'_1, \dots, B'_n], other_goals)

new_goals = substitute(instance, other_goals)

[satisfied, instance] = EXECUTE(program, new_goals)

return [satisfied, instance]

$\overline{\text{anc}(m, y), \text{anc}(d, c)}$

$\overline{\text{anc}(x, z) :- \text{adv}(x, z)}$

?anc(m, y), anc(d, c)

single goal

goal → anc(m, y)

other-goals → anc(d, c)

new-goals ← $[B'_1, \dots, B'_n]$ + $\overbrace{\text{adv}(x, z)}$

new goals ← $\overbrace{\text{adv}(m, y)}$

Reordering of Clauses

Q: Whether reordering of clauses or goals have any effect in execution?

V1

```
ancestor1(X, Z) :-  
    advisor(X, Z).  
  
ancestor1(X, Z) :-  
    advisor(X, Y),  
    ancestor1(Y, Z).
```

V2

```
ancestor2(X, Z) :-  
    advisor(X, Y),  
    ancestor2(Y, Z).  
  
ancestor2(X, Z) :-  
    advisor(X, Z).
```

Clause
swap

Original



V3

```
ancestor3(X, Z) :-  
    advisor(X, Z).  
  
ancestor3(X, Z) :-  
    ancestor3(Y, Z),  
    advisor(X, Y).
```

Goal
swap

V4

```
ancestor4(X, Z) :-  
    ancestor4(Y, Z),  
    advisor(X, Y).  
  
ancestor4(X, Z) :-  
    advisor(X, Z).
```

Clause
& Goal
swap

Reordering of Clauses

```
↓  
ancestor1(X, Z) :-  
    advisor(X, Z).
```

```
ancestor1(X, Z) :-  
    advisor(X, Y),  
    ancestor1(Y, Z).
```

Original

?- ancestor(dana, azizi)

Call: *ancestor1(dana, azizi)*

Call: *advisor(dana, azizi)*

Exit: *advisor(dana, azizi)*

Exit: *ancestor1(dana, azizi)*

←

← True

← True

goal satisfied

Reordering of Clauses

R1

```
ancestor2(X, Z) :-  
    advisor(X, Y),  
    ancestor2(Y, Z).
```

R2

```
ancestor2(X, Z) :-  
    advisor(X, Z).
```

Clause swap

Success

```
Call: ancestor2(dana, azizi)  
Call: advisor(dana, _4612)  
Exit: advisor(dana, felix)  
Call: ancestor2(felix, azizi)  
Call: advisor(felix, _4614)  
Fail: advisor(felix, _4614)  
Redo: ancestor2(felix, azizi)  
Call: advisor(felix, azizi)  
Fail: advisor(felix, azizi)  
Fail: ancestor2(felix, azizi)  
Redo: advisor(dana, _4612)  
Exit: advisor(dana, chen)  
Call: ancestor2(chen, azizi)  
Call: advisor(chen, _4614)  
Fail: advisor(chen, _4614)  
Redo: ancestor2(chen, azizi)  
Call: advisor(chen, azizi)  
Fail: advisor(chen, azizi)  
Fail: ancestor2(chen, azizi)  
Redo: advisor(dana, _4612)  
Exit: advisor(dana, amir)  
Call: ancestor2(amir, azizi)  
Call: advisor(amir, _4614)  
Fail: advisor(amir, _4614)  
Redo: ancestor2(amir, azizi)  
Call: advisor(amir, azizi)  
Fail: advisor(amir, azizi)  
Fail: ancestor2(amir, azizi)  
Redo: advisor(dana, _4612)  
Exit: advisor(dana, azizi)  
Call: ancestor2(azizi, azizi)  
Call: advisor(azizi, _4614)  
Fail: advisor(azizi, _4614)
```

?- ancestor(dana, azizi) *not instantiated*

↓

advisor(dana, _4612),
advisor(dana, felix)

↙
True

ancestor(_4612, azizi)

?- ancestor(felix, azizi)
false

advisor(felix, _4652)

↙
false

ancestor(_4652, azizi)

?- ancestor(felix, azizi)

R2

↙
false

advisor(felix, azizi)

Reordering of Clauses

↳

```
ancestor3(X, Z) :-  
    advisor(X, Z).  
  
ancestor3(X, Z) :-  
    \ ancestor3(Y, Z),  
    advisor(X, Y).
```

Goal swap

?- ancestor3(bhaskara, felix)

```
ancestor4(X, Z) :-  
    ancestor4(Y, Z),  
    advisor(X, Y).  
  
ancestor4(X, Z) :-  
    advisor(X, Z).
```

Quick Solution

Infinite Loop

Infinite Loop

Clause & Goal swap

Takeaways from Ordering

- Try **simplest idea** first (practical heuristics in problem solving)
 - ancestor₁ being the simplest, ancestor₄ being the most complex

- Check your clause ordering to avoid **infinite recursion**

Efficiency or decidability

Procedural aspect is also important along with declarative ordering → sequencing of instructions.

$$N_1 = 10 \quad N = \overbrace{10+1}^{\text{Not } N=11} \quad \text{between } (10, 20)$$

Some Example Programs

variable assignment

$N = N_1 + 1$

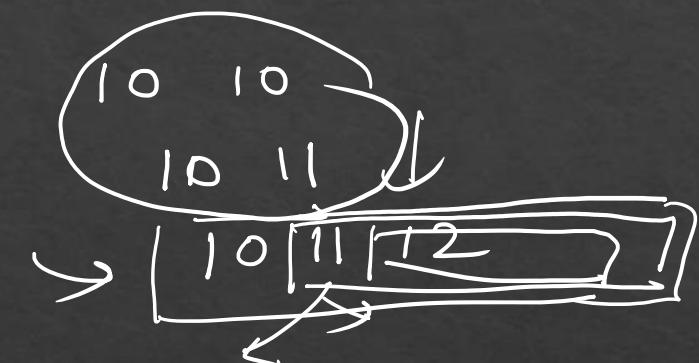
N is $N_1 + 1$

Numbers between two numbers:

```
between_number(N1, N2) :-  
    N1 < N2 - 1, N is N1 + 1,  
    print(N), nl, NN1 is N1 + 1,  
    between_number(NN1, N2).
```

10 11 12 13 15 18 20

↓ ↓ ↓ ↓ ↓ ↓ ↓



$\text{sum}(10, 10) \rightarrow 10 \rightarrow \text{base cond!}$

Sum of the numbers in a range:

$\text{sum}(10, 20) \rightarrow 10 + 11 + 12 + \dots + 20$

series_sum(N, N, N).

series_sum(N1, N2, Sum) :- N1 < N2, N is N1 + 1,

$\text{series_sum}(N, N2, \text{SumInter}),$

Sum is SumInter + N1.

$[N1]$ sumInter

Sum

Prolog List Data Structure

List Data Structure:

$$\text{empty list} \quad [1 | \underline{[2, 3, 4, 5]}] = [1, 2, 3, 4, 5]$$

$$[] \text{ OR } [\underline{\text{Head}} | \text{Tail}] \text{ OR } [\text{Item}_1, \text{Item}_2, \dots | \text{Others}]$$

$$[1, 2, 3 | \text{others}] = [1, 2, 3, 4, 5]$$

Examples:

$$\underline{\text{Color1}} = [\text{maroon}, \text{green}] .$$

$$\text{others} = [4, 5]$$

$$\underline{\text{Color2}} = [\text{red}, \text{yellow}] . \quad \text{Nested List}$$

$$\underline{\text{Clubs}} = [\underline{\text{mohanB}}, \underline{\text{Color1}}, \underline{\text{eastB}}, \underline{\text{Color2}}] .$$

$$\underline{L} = [\underline{X1}, \underline{X2} | [\underline{X3}, \underline{X4}, \underline{X5}]] .$$

Prolog List Data Structure

- Concatenation of two lists

$\text{conc}([], L, L).$

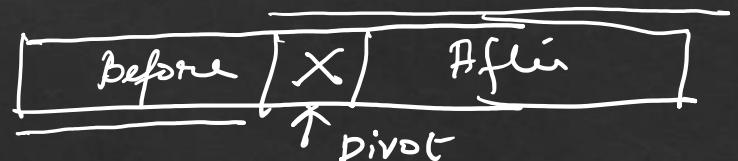
$\text{conc}([x|L_1], L_2, [x|L_3]) :- \text{conc}(L_1, L_2, L_3).$

- Membership in list

base step : $x [x| _]$ → satisfied

recursive : $x \text{ mem}(x, [-| \text{Tail}]) :- \text{Mem}(x, \text{Tail}).$

- Partition a list wrt a pivot



divide $[x|L]$

$\text{conc}(_X),$



$L_1 + L_2$

$$= \underbrace{[L_1 | L_2]}_{23}$$

$\text{conc}(L_1, L_2, L_3)$

$(x|L_1| \dots |L_3)$



$\text{conc}(\underline{\text{Before}}, [\underline{x} | \underline{\text{After}}], \underline{\text{List}})$

$\text{conc}(_, [B, X, A | _], \underline{\text{List}})$

Prolog List Data Structure

- Delete from list

base step: $x [x \ L1] \Rightarrow [L1]$

recursive step: $x [y \ L1] \Rightarrow [y \ L2] \gg \text{del } x \ L1 \ L2$

- A list is ordered or not

base step: $[] [x] \quad \text{Tail} = \text{Con}(H1, Tail1)$

recursive step: $[H \ Tail] \quad [H \ H1 \ Tail1] \quad H \leq H1 \quad \text{ordered}([H1 \ Tail1])$

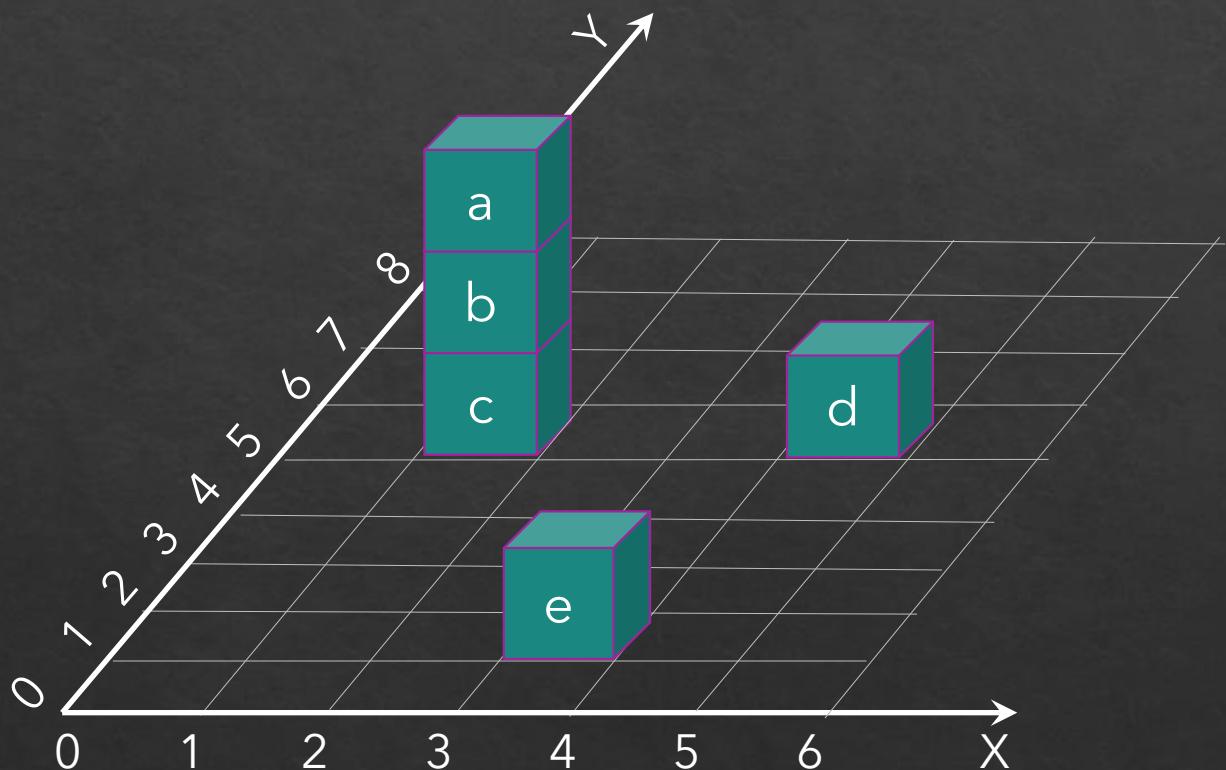
- Max in a list

base step: $[x] \Rightarrow x \text{ max}$

recursive step: $[H \ Tail]$

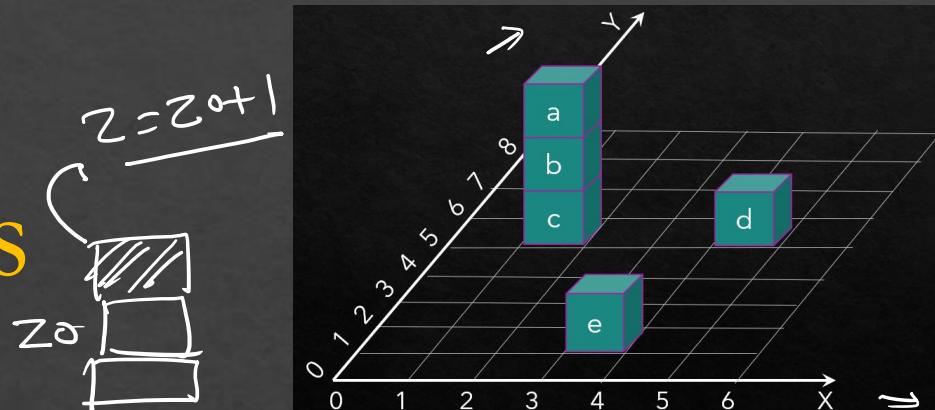
$\max(N1, Tail)$
If $H > N1$, M is H
if $H \leq N1$, M is $N1$

A Robot Playing with Blocks



- Robot sees from the top
- Robot can name the visible blocks and x-y coordinates.
 - ✓ see (a, 1, 5) .
 - see (d, 4, 5) .
 - see (e, 3, 1) .
- Positional relations
 - on (a, b) .
 - on (b, c) .
 - on (c, table) .
 - on (d, table) .
 - on (e, table) .

on (A, _), \+ see (A, _, _)



A Robot Playing with Blocks

- What are blocks in this world?

?- on (Block, _).

Blocks

- Pairs of blocks having the same y-coordinate

?- see (B1, _, Y),
see (B2, _, Y),
B1 \= B2.

- Boxes that are not visible

?- on (B, _),
\+ see (B, _, _, _).

- Leftmost visible block

✓ ?- see (B, X, _),
\+ (see (B2, X2, _), (X2 < X)).

? - Find the Z-coordinate of a block

→ z (B, 0) :- on (B, table).
→ z (B, Z) :- on (B, B0), z (B0, Z0),
Z is Z0+1.

✓ - Find blocks b/w two blocks

```
recursive_on(B1, B2) :- on(B1, B2).  
recursive_on(B1, B2) :- on(B1, BX),  
print(BX), recursive_on(BX, B2).
```

Inherently Recursive Problems

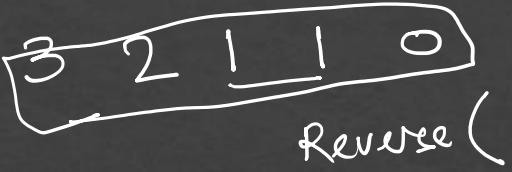
Fibonacci Series

fib series

```

✓ fib_seq(S, N) :- N > 1, reverse(SR, S).
→ fib_seq_(N, SR, 1, [1, 0]), inbuilt

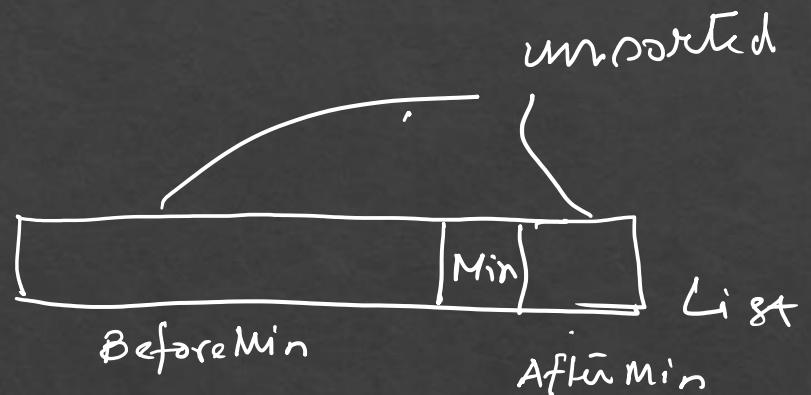
fib_seq_(N, Seq, N, Seq).
fib_seq_(N, Seq, N0, [B, A | FS]) :- N > N0,
→ N1 is N0+1, C is A+B,
↓
fib_seq_(N, Seq, N1, [C, B, A | FS]).
```



$L = [H | Tail]$
 $= [H_1, H_2 | Tail]$ Reversing
 $i = 1 \quad [0, 1] \quad [1, 0]$
 $i = 2 \quad [1, 1, 0 | []]$
 $i = 3 \quad [2, 1, 1 | [0]]$
 $[2, 1 | [1, 0]]$

Simple Sort

```
✓ min(A, A, B) :- A <= B.  
✓ min(B, A, B) :- B <= A.  
  
✓ smallest(A, [A|[]]).  
✓ smallest(Min, [A|B]) :- smallest(SB, B),  
                           min(Min, A, SB).  
  
sorted([], []).  
- -  
sorted([Min|RestSorted], List) :-  
    smallest(Min, List),  
    append(BeforeMin, [Min|AfterMin], List),  
    append(BeforeMin, AfterMin, RestUnsorted),  
    sorted(RestSorted, RestUnsorted).  
  
Bind →
```



con(BeforeMin, [Min] AfterMin,
List)

con(Beforemin, Aftermin,
unsorted)

sort(RestSorted, unsorted)

[Min | RestSorted]

Merge Sort



$[] + \text{List} \Rightarrow \text{List}$

$\text{List} + [] \Rightarrow \text{List}$

$[M_1] \text{ RestList 1 }$

$[M_2] \text{ RestList 2 }$

if $M_1 \leq M_2$ rest merged

$[M_1] \quad [\underbrace{\text{RestList 1}}_{\text{rest merged}}] \quad [M_2] \text{ RestList 2 }$

if $M_1 > M_2$

MergeSort

A diagram showing a list node consisting of a rectangle with a vertical line extending downwards from its bottom edge to a horizontal line, representing the connection to the next node in the list.

\rightarrow
A diagram showing two list nodes connected by a horizontal line, representing the merging of two sorted lists into one.

Merge Sort

```
merge(List, List, []).
```

```
merge(List, [], List).
```

```
merge([MinList1|RestMerged], [MinList1|RestList1],  
[MinList2|RestList2]) :-
```

```
MinList1 <= MinList2,
```

```
merge(RestMerged, RestList1, [MinList2|RestList2]).
```

```
merge([MinList2|RestMerged], [MinList1|RestList1],  
[MinList2|RestList2]) :-
```

```
MinList2 <= MinList1,
```

```
merge(RestMerged, [MinList1|RestList1], RestList2).
```

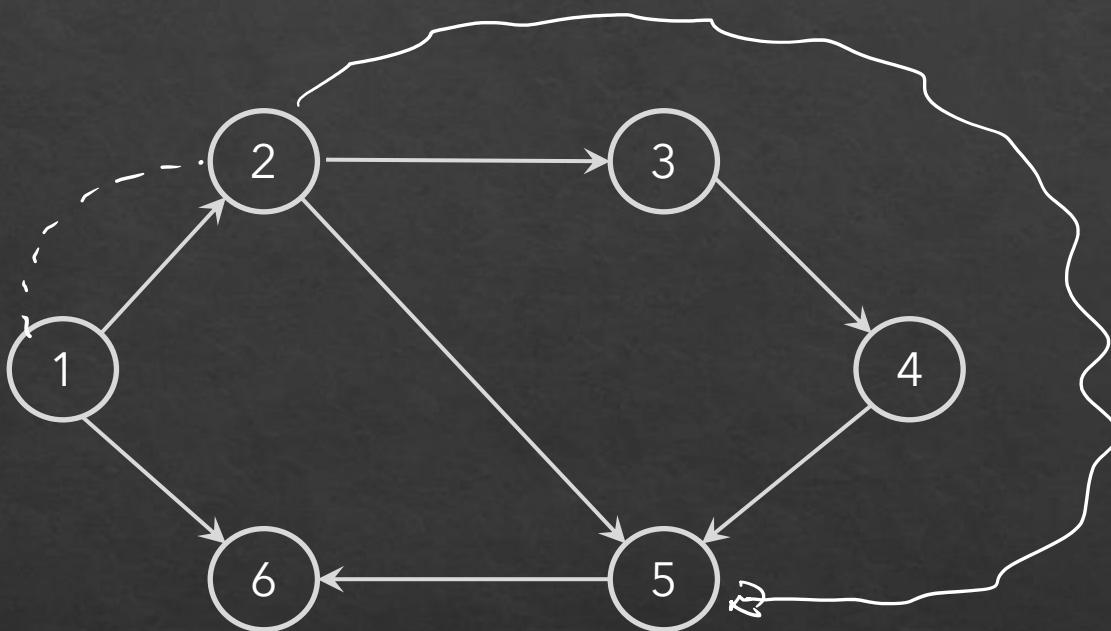
Merge Sort

```
mergeSort([], []).  
mergeSort([A], [A]).
```

```
mergeSort(Sorted, List) :-  
    length(List, N), FirstLength is // (N, 2),  
    SecondLength is N - FirstLength,  
    length(FirstUnsorted, FirstLength),  
    length(SecondUnsorted, SecondLength),  
    append(FirstUnsorted, SecondUnsorted, List),  
    mergeSort(FirstSorted, FirstUnsorted),  
    mergeSort(SecondSorted, SecondUnsorted),  
    merge(Sorted, FirstSorted, SecondSorted).
```

Inherently Iterative Problems

Path Finding in Graph



path1 (N, N) .

path1 (S, D) :- edge (S, N),
path1 (N, D) .

path2 (N, N, [N]) .

path2 (S, D, L) :- edge (S, N),
path2 (N, D, T),
conc ([S], T, L) .