

Appendix D

The ARM Processor

D.1. (a) R8, R9, and R10, contain 1, 2, and 3, respectively.

(b) The values 20 and 30 are pushed onto a stack pointed to by R1 by the two Store instructions, and they occupy memory locations 1996 and 1992, respectively. They are then popped off the stack into R8 and R9. Finally, the Subtract instruction results in 10 ($30 - 20$) being stored in R10. The stack pointer R1 is returned to its original value of 2000.

D.2. (b) A memory operand cannot be referenced in a Subtract instruction.

(d) The immediate value 257 is 100000001 in binary, and is thus too long to fit in the 8-bit immediate field. Note that it cannot be generated by the rotation of any 8-bit value.

D.3. Use register R0 as a counter register and R1 as a work register. A first possible program uses a rotation operation, RRX (Rotate right extended), that was not explained in Appendix D. The operation RRX rotates the contents of a register, extended by the C flag, to the right by one bit position.

	MOV	R0, #32	Load R0 with count value 32.
LOOP	MOVS	R2, R2, LSL #1	Shift contents of R2 left one bit position, setting the high-order bit into the C flag.
	MOV	R1, R1, RRX	Rotate R1 right one bit position, including the C flag, as shown in Figure 2.25d.
	SUBS	R0, R0, #1	Check if finished.
	BNE	LOOP	
	MOV	R2, R1	Load reversed pattern back into R2.

Next, we give a second possible program that does not use the RRX operation.

	MOV	R0, #32	Load R0 with the count value 32.
LOOP	MOV	R1, R1, LSR #1	Shift contents of R1 right by one bit position, setting the high-order bit to 0.
	CMP	R2, #0	Set condition codes to show status of R2 contents.
	ORRMI	R1, R1, #&80000000	Set high-order position of R1 to 1 if high-order position of R2 is 1.
	MOV	R2, R2, LSL #1	Shift contents of R2 left one bit position, moving next bit into the sign-bit position.
	SUBS	R0, R0, #1	Loop back if all 32 bits have not been processed.
	BNE	LOOP	
	MOV	R2, R1	Write reversed bits into R2.

D.4. Program trace:

TIME	R0	R1	R2
after 1st execution of BGT	3	4	NUM1 + 4
after 2nd execution of BGT	-14	3	NUM1 + 8
after 3rd execution of BGT	13	2	NUM1 + 12

D.5. Assume bytes are unsigned 8-bit values.

	LDR	R0, N	R0 is list counter.
	ADR	R1, X	R1 points to X list.
	ADR	R2, Y	R2 points to Y list.
	ADR	R3, LARGER	R3 points to LARGER list.
LOOP	LDRB	R4, [R1], #1	Load X list byte into R4.
	LDRB	R5, [R2], #1	Load Y list byte into R5.
	CMP	R4, R5	Compare bytes.
	STRHSB	R4, [R3], #1	Store X byte if larger or same.
	STRLOB	R5, [R3], #1	Store Y byte if larger.
	SUBS	R0, R0, #1	Check if finished.
	BGT	LOOP	

D.6. This solution assumes that the last number in the series of n numbers can be represented in a 32-bit word, and that $n > 2$.

	MOV	R0, N	Use R0 for loop count for numbers
	SUB	R0, R0, #2	generated after 1.
	ADR	R1, MEMLOC	Use R1 as memory pointer.
	MOV	R2, #0	Store first two numbers,
	STR	R2, [R1], #4	0 and 1, from R2
	MOV	R3, #1	and R3 into memory.
	STR	R3, [R1], #4	
LOOP	ADD	R3, R2, R3	Starting with number $i - 1$
	STR	R3, [R1], #4	in R2 and i in R3, compute
			and place $i + 1$ in R3
			and store in memory.
	SUB	R2, R3, R2	Recover old i and place
			in R2.
	SUBS	R0, R0, #1	Branch back if all numbers
	BGT	LOOP	have not been computed.

D.7. Inspection of Table 1.1 of the ASCII characters shows that lowercase letters can be converted to uppercase letters by changing bit b_5 from 1 to 0.

A possible program for the conversion follows. Note that the last three instructions are executed conditionally (condition = NE), as described in Sections D.1.1 and D.9.

	LDR	R1, =WORD	Load address WORD.
NEXT	LDRB	R0, [R1]	Load character byte.
	CMP	R0, #&20	Check for space.
	BICNE	R0, R0, #&20	If not space, clear
	STRNEB	R0, [R1], #1	bit b_5 , store back,
	BNE	NEXT	and process next character.

- D.8. Memory word location J contains the number of tests, j , and memory word location N contains the number of students, n . The list of student marks begins at memory word location LIST in the format shown in Figure 2.10. The parameter $\text{Stride} = 4(j + 1)$ is the distance in bytes between scores on a particular test for adjacent students in the list.

The Post-indexed addressing mode [R2], R3, LSL #2 is used to access the successive scores on a particular test in the inner loop. The value in register R2 before each entry to the inner loop is the address of the score on a particular test for the first student. Register R3 contains the value $j + 1$. Therefore, register R2 is incremented by the Stride parameter on each pass through the inner loop.

	LDR	R3, J	Load $j + 1$ into R3 to
	ADD	R3, R3, #1	be used as an address offset.
	ADR	R4, SUM	Initialize R4 to the sum
			location for test 1.
	ADR	R5, LIST	Load address of test 1 score
	ADD	R5, R5, #4	for student 1 into R5.
	LDR	R6, J	Initialize outer loop counter
			R6 to j .
OUTER	LDR	R7, N	Initialize inner loop
			counter R7 to n .
	MOV	R2, R5	Initialize base register R2
			to location of student 1 test
			score for next inner loop
			sum computation.
	MOV	R0, #0	Clear sum accumulator
			register R0.
INNER	LDR	R1, [R2], R3, LSL #2	Load test score into R1
			and increment R2 by Stride to
			point to next test score.
	ADD	R0, R0, R1	Accumulate score into R0.
	SUBS	R7, R7, #1	Check if all student scores
	BGT	INNER	for current test are added.
	STR	R0, [R4], #4	Store sum in memory.
	ADD	R5, R5, #4	Increment R5 to next test
			score for student 1.
	SUBS	R6, R6, #1	Check if sums for all test
	BGT	OUTER	scores have been accumulated.

- D.9. A possible program that computes $\text{SUM} = 580 + 68400 + 80000$.

LDR	R0, =580	Load 580.
LDR	R1, =68400	Load 68400.
ADD	R0, R0, R1	Generate $580 + 68400$ in R0.
LDR	R1, =80000	Load 80000.
ADD	R0, R0, R1	Generate the final sum in R0.
STR	R0, SUM	Store the sum.

D.10. Program that computes $\text{ANSWER} = A \times B + C \times D$.

	AREA	CODE	
	ENTRY		
	LDR	R0, A	Load operand A.
	LDR	R1, B	Load operand B.
	MUL	R0, R1, R0	Generate $A \times B$ in R0.
	LDR	R1, C	Load operand C.
	LDR	R2, D	Load operand D.
	MULA	R0, R1, R2, R0	Generate $C \times D + [R0]$ in R0.
	STR	R0, ANSWER	Store the answer.
	AREA	DATA	
A	DCD	100	Test data.
B	DCD	50	
C	DCD	20	
D	DCD	400	
ANSWER	DCD	0	Space for the answer.

D.11. The following program determines the number of negative integers in a list.

	AREA	CODE	
	ENTRY		
	LDR	R1, =NUMBERS	Load address of data list.
	LDR	R2, N	Load size of list.
	MOV	R0, #0	Clear negative number counter.
LOOP	LDR	R3, [R1], #4	Load next data item .
	CMP	R3, #0	Set condition code flags.
	ADDMI	R0, R0, #1	Increment counter if data item negative.
	SUBS	R2, R2, #1	Decrement loop counter
	BNE	LOOP	and branch back if not done.
	STR	R0, NEGNUM	Store result.
	AREA	DATA	
N	DCD	6	Size of list.
NEGNUM	DCD	0	Count of negative numbers.
NUMBERS	DCD	25, -5, -128	Test data.
	DCD	44, -23, -9	

- D.12. A possible program for byte sorting as described in Example 2.5 follows. It has the same general structure as the solution in Example 2.5. The conditional execution feature of the ARM instruction set is used to advantage in the inner loop when $LIST(k)$ must be interchanged with $LIST(j)$. The three-instruction sequence STR, STR, MOV is only executed if $LIST(k)$ is greater than $LIST(j)$, as indicated by the GT suffixes. The forward conditional branch to NEXT in the solution to Example 2.5 is not needed in the ARM program.

	ADR	R4, LIST	Load list pointer register R4,
	LDR	R10, N	and initialize outer loop base
	ADD	R2, R4, R10	register R2 to $LIST + n$.
	ADD	R5, R4, #1	Load $LIST + 1$ into R5.
OUTER	LDRB	R0, [R2, #-1]!	Load $LIST(j)$ into R0.
	MOV	R3, R2	Initialize inner loop base register
			R3 to $LIST + n - 1$.
INNER	LDRB	R1, [R3, #-1]!	Load $LIST(k)$ into R1.
	CMP	R1, R0	Compare $LIST(k)$ to $LIST(j)$.
	STRGTB	R1, [R2]	If $LIST(k) > LIST(j)$, interchange
	STRGTB	R0, [R3]	$LIST(k)$ and $LIST(j)$, and
	MOVGT	R0, R1	move (new) $LIST(j)$ into R0.
	CMP	R3, R4	If $k > 0$, repeat
	BNE	INNER	inner loop.
	CMP	R2, R5	If $j > 1$, repeat
	BNE	OUTER	outer loop.

- D.13. Assume that register R0 contains the address BOTTOM and that register R1 contains the address TOP.

SAFEPOPUSH	CMP	R5, R1
	BEQ	FULLERROR
	STR	R2, [R5, #-4]!
SAFEPOP	CMP	R5, R0
	BHI	EMPTYERROR
	LDR	R2, [R5], #4

- D.14. No program is needed for the solution to this problem; so the solution is the same as that given for Problem 2.24.

D.15. See the solution to Problem 2.24 for the procedures needed to perform the append and remove operations.

Register assignment:

R0 – Data byte to append to or remove from queue

R1 – IN pointer

R2 – OUT pointer

R3 – Address of first queue byte location

R4 – Address of last queue byte location ($= [R3] + k - 1$)

R5 – Auxiliary register for address of next appended byte.

Initially, the queue is empty with $[R1] = [R2] = [R3]$

APPEND routine:

MOV	R5, R1	
ADD	R1, R1, #1	Increment R1 Modulo k .
CMP	R1, R4	
MOVGT	R1, R3	
CMP	R1, R2	Check if queue is full.
MOVEQ	R1, R5	If queue full, restore
BEQ	QUEUEFULL	IN pointer and send
		message that queue is full.
STRB	R0, [R5]	If queue not full,
		append byte and continue.

REMOVE routine:

CMP	R1, R2	Check if queue is empty.
BEQ	QUEUEEMPTY	If empty, send message.
LDRB	R0, [R2], #1	Otherwise, remove byte
CMP	R2, R4	and increment R2.

D.16. The weights stored in memory are not actually needed because multiplication of inputs by the weights $1/8$, $1/4$, $1/2$ can be accomplished by arithmetic right shifts of the inputs by 3, 2, and 1 bit positions, respectively.

	LDR	R0, =IN	Load addresses of first three
	ADD	R1, R0, #4	inputs into R0, R1, and R2.
	ADD	R2, R1, #4	
	LDR	R5, =OUT	Load address of first output into R5.
	LDR	R6, N	Number of outputs to be computed.
LOOP	LDR	R4, [R0], #4	Load first input.
	MOV	R4, R4, ASR #3	Multiply it by $1/8$.
	LDR	R5, [R1], #4	Load second input.
	ADD	R4, R4, R5, ASR #2	Multiply it by $1/4$ and
			add to first input.
	LDR	R5, [R2], #4	Load third input.
	ADD	R4, R4, R5, ASR #1	Multiply it by $1/2$ and
			add it to sum.
	STR	R4, [R5], #4	Store output value.
	SUBS	R6, R6, #1	Decrement counter and branch
	BGT	LOOP	back if not done.

- D.17. The following program performs the copying operation as specified in the problem statement. However, if the *from* and *to* areas do not overlap, copying can actually be done in either direction. If the problem statement is relaxed to allow that possibility, then the five instructions following the STMFD instruction could be replaced by the two instructions

```
CMP    R2, R1
BLO    FORWARD
```

and the subroutine would perform the copying operation correctly. Also, we assume that the *from* and *to* addresses are not the same; but if they are, the given program does the copying in the BACKWARD loop.

Calling program:

```
LDR    R0, N          Load the length of the byte sequence.
LDR    R1, =FROM      Load the from address.
LDR    R2, =TO        Load the to address.
BL     MEMCOPY        Call the subroutine.
next instruction
```

Subroutine:

MEMCOPY	STMFD	R13!, {R3, R14}	Stack R3 and the link register.
	ADD	R3, R0, R1	Load R3 with <i>from</i> + <i>length</i> .
	CMP	R2, R1	Scan/copy forward if <i>to</i> is outside <i>from</i> area.
	BLO	FORWARD	
	CMP	R2, R3	
	BHS	FORWARD	
BACKWARD	ADD	R1, R1, R0	Otherwise (if <i>to</i> is inside <i>from</i> area),
	ADD	R2, R2, R0	scan/copy backward from end of sequence.
	LDRB	R3, [R1, #-1]!	
	STRB	R3, [R2, #-1]!	
	SUBS	R0, R0, #1	Branch back if not finished copying.
	BGT	BACKWARD	
	B	DONE	
FORWARD	LDRB	R3, [R1], #1	Scan/copy forward.
	STRB	R3, [R2], #1	
	SUBS	R0, R0, #1	Branch back if not done.
	BGT	FORWARD	
DONE	LDMFD	R13!, {R3, R15}	Return.

- D.18. Assume that the values *first* address, *second* address, and *length* of the sequence are passed to the subroutine in registers R1, R2, and R3, respectively. Use register R0 to return the count of the number of comparisons that do not match. The subroutine also uses registers R4 and R5, so they need to be saved/retored at the start/end of the subroutine.

Subroutine:

MEMCMP	STMFD	R13!, {R4, R5, R14}	Stack registers R4, R5, and the link register.
	MOV	R0, #0	Clear “no-match” count register.
LOOP	LDRB	R4, [R1], #1	Load byte from first list.
	LDRB	R5, [R2], #1	Load byte from second list.
	CMP	R4, R5	Compare bytes.
	ADDNE	R0, R0, #1	Increment counter if bytes don’t match.
	SUBS	R3, R3, #1	Branch back if not done.
	BGT	LOOP	
	LDMFD	R13!, {R4, R5, R15}	Return from subroutine.

- D.19. Assume that the address value STRNG is passed to the subroutine in register R0, and that the subroutine uses registers R1 and R2 as work registers that need to be saved/restored by the subroutine.

EXCLAIM	STMFD	R13!, {R1, R2, R14}	Save R1, R2, and the link register.
	MOV	R2, #&21	Load ASCII code for '!'.
LOOP	LDRB	R1, [R0], #1	Check for ASCII code '.' and
	CMP	R1, #&2E	replace with '!' if yes.
	STREQB	R2, [R0, #-1]	
	CMP	R1, #&00	Check for ASCII code 'NUL'
	BNE	LOOP	and read more characters if no.
	LDMFD	R13!, {R1, R2, R15}	Restore registers and return.

- D.20. This problem is similar to Problem D.7. In that problem, it is assumed that all of the input text characters are lowercase letters. Here, the text can consist of any ASCII characters. Hence, a check must be made to determine whether or not each character scanned is a lowercase letter.

A possible subroutine program ALLCAPS follows. The subroutine assumes that the address STRNG is passed to it in register R0. The ASCII code for a lowercase letter can be converted to the corresponding uppercase letter by changing bit b_5 from 1 to 0. The lowercase letter codes are in the numerical range 97 through 122.

Subroutine:

ALLCAPS	STMFD	R13!, {R1, R14}	
SCAN	LDRB	R1, [R0], #1	Load character.
	CMP	R1, #&00	Check for NUL character.
	BEQ	RETURN	Return on NUL.
	CMP	R1, #97	Check for lowercase.
	BLT	SCAN	
	CMP	R1, #122	
	BGT	SCAN	
	BIC	R1, R1, #&20	If lowercase, clear bit b_5
	STRB	R1, [R0, #-1]	and store back.
	B	SCAN	Continue scanning.
RETURN	LDMFD	R13!, {R1, R15}	Return.

- D.21. The calling program passes the address value STRNG to the subroutine in register R0, and the word count is returned in register R1. A possible subroutine is as follows.

Subroutine:

WORDS	STMFD	R13!, {R2, R14}	Stack work register R2 and the link register.
	MOV	R1, #0	Clear word counter.
LOOP	LDRB	R2, [R0], #1	Check for space character and
	CMP	R2, #&20	increment counter if yes.
	ADDEQ	R1, R1, #1	
	CMP	R2, #&00	Check for NUL character and
	BNE	LOOP	continue scanning if no.
	LDMFD	R13!, {R2, R15}	Restore R2 and return.

- D.22. Assume that the starting address of the ordered list of numbers, the length of the list (in words), and the new number to be inserted, are passed to the subroutine INSERT in registers R1, R2, and R3. Also assume that there may be sequences of two or more numbers that are the same at different places in the list, and that the list contains at least one number.

In the solution given, the subroutine moves through the list to locate the position of the first number that is larger than the new number to be inserted. This is the position where the new number is to be inserted after moving the sublist that starts at that position up by one word position.

Register R4 is first set to hold the address of the word after that of the last number in the list; register R1, which initially contains the starting address for the list, is used to step through the list searching for the correct insertion position; and the current list number to be compared to the new number is loaded into register R5. If the new number is equal to a number (or string of numbers) in the list, it will be inserted after that number (or string).

Subroutine:

INSERT	STMFD	R13!, {R4, R5, R14}	
	ADD	R4, R1, R2, LSL #2	Set R4 to first word address beyond the end of the list.
	LDR	R5, [R4, #-4]	Check if the new number should be stored after the last number in the list.
	CMP	R3, R5	
	BGE	STORE	
LOOP1	LDR	R5, [R1], #4	Load current list number into R5, increment pointer, and continue through list if new number is greater than or equal to the current number.
	CMP	R3, R5	
	BGE	LOOP1	
	SUB	R1, R1, #4	Set R1 to starting address of numbers to be moved up.
LOOP2	LDR	R5, [R4, #-4]	Shift numbers up, starting from end of list.
	STR	R5, [R4], #-4	
	CMP	R4, R1	
	BGT	LOOP2	
STORE	STR	R3, [R4]	Store new number in list.
	LDMFD	R13!, {R4, R5, R15}	

- D.23. The address OLDLIST, the length of the list in words, and the address NEWLIST, are passed to the subroutine INSERTSORT in registers R6, R7, and R8, respectively. Assume that the unordered list contains at least one number.

Subroutine:

INSERTSORT	STMFD	R13!, {R1, R2, R3, R14}	
	MOV	R1, R8	Initialize R1 to address NEWLIST.
	LDR	R3, [R6], #4	Load first number from old list into R3 and increment pointer.
	STR	R3, [R1]	Store this number in new list.
	SUBS	R7, R7, #1	Decrement old list count.
	BEQ	RETURN	Return if old list had only one number.
	MOV	R2, #1	Otherwise, load R2 and R3 before calling INSERT.
	MOV	R3, [R6], #4	(See Problem 2.32 for a description of INSERT.)
	BL	INSERT	
	MOV	R1, R8	Set up R1, R2, and R3 before calling INSERT again.
	ADD	R2, R2, #1	
	LDR	R3, [R6], #4	
	SUBS	R7, R7, #1	Check if all numbers have been inserted into the new list.
	BNE	LOOP	
	LDMFD	R13!, {R1, R2, R3, R15}	If yes, Return.

- D.24. The following program is similar to that in Figure D.17. There are two differences: A specified number of characters, $n = [N]$, are read from the keyboard, instead of reading up to a carriage return character; and the characters are pushed onto a user stack instead of being written into memory before they are displayed.

	LDR	R0, N	Use R0 as the loop counter for reading n characters.
READ	LDRB	R3, [R1, #4]	Load KBD_STATUS byte and wait for character.
	TST	R3, #2	
	BEQ	READ	
ECHO	LDRB	R3, [R1]	Read character and push onto stack.
	STRB	R3, [R6, #-1]!	
	LDRB	R4, [R2, #4]	Load DISP_STATUS byte and wait for display.
	TST	R4, #4	
	BEQ	ECHO	
	STRB	R3, [R2]	Send character to display.
	SUBS	R0, R0, #1	Repeat until n characters read.
	BGT	READ	

- D.25. Assume that almost all of the time between successive characters being struck is spent in the three-instruction wait loop that starts at location READ. The BEQ READ instruction is executed once every 15 ns (3×5 ns) while this loop is being executed. There are $10^9/10 = 10^8$ ns between successive characters. Therefore, the BEQ READ instruction is executed $10^8/15 = 6.7 \times 10^6$ times per character entered.

D.26. Main Program:

READLINE	BL	GETCHAR	Load character into R3.
	STRB	R3, [R0], #1	Store character in memory.
	BL	PUTCHAR	Display character.
	TEQ	R3, #CR	Check for carriage return character.
	BNE	READLINE	

Subroutine GETCHAR:

GETCHAR	LDRB	R3, [R1, #4]	Wait for character.
	TST	R3, #2	
	BEQ	GETCHAR	
	LDRB	R3, [R1]	Load character into R3.
	MOV	R15, R14	Return.

Subroutine PUTCHAR:

PUTCHAR	STMFD	R13!, {R4, R14}	Save R4 and Link register.
DISPLAY	LDRB	R4, [R2, #4]	Wait for display.
	TST	R4, #4	
	BEQ	DISPLAY	
	STRB	R3, [R2]	Send character to display.
	LDMFD	R13!, {R4, R15}	Restore R4 and Return.

D.27. Address KBD_DATA is passed to GETCHAR on the stack; the character read is passed back in the same stack position. The character to be displayed and the address DISP_DATA are passed to PUTCHAR by being pushed onto the stack in that order. Register R0 contains the address of the memory buffer area.

Main Program:

READLINE	LDR	R5, =KBD_DATA	Load address KBD_DATA
	STR	R5, [SP, #-4]!	into R5 and push onto stack.
	BL	GETCHAR	Call subroutine to read character.
	LDRB	R5, [SP]	Load character from top of
	STRB	R5, [R0], #1	stack, leaving it on the stack,
			and store it in memory.
	LDR	R5, =DISP_DATA	Load address DISP_DATA into
	STR	R5, [SP, #-4]!	R5, and push onto stack.
	BL	PUTCHAR	Call character display subroutine.
	ADD	SP, SP, #8	Remove parameters from stack.
	TEQ	R5, #CR	Check for carriage return character.
	BNE	READLINE	

Subroutine GETCHAR:

GETCHAR	STMFD	SP!, {R1, R3, LR}	Save registers.
	LDR	R1, [SP, #12]	Load address KBD_DATA into R1.
READ	LDRB	R3, [R1, #4]	Wait for character.
	TST	R3, #2	
	BEQ	READ	
	LDRB	R3, [R1]	Load character into R3
	STRB	R3, [SP, #12]	and overwrite address KBD_DATA
			on stack.
	LDMFD	SP!, {R1, R3, PC}	Restore registers and Return.

Subroutine PUTCHAR:

PUTCHAR	STMFD	SP!, {R2-R4, LR}	Save registers.
	LDR	R2, [SP, #16]	Load address DISP_DATA into
	LDR	R3, [SP, #20]	R2 and character into R3.
DISPLAY	LDR	R4, [R2, #4]	Wait for display.
	TST	R4, #4	
	BEQ	DISPLAY	
	STRB	R3, [R2]	Send character to display.
	LDMFD	SP!, {R2-R4, PC}	Restore registers and Return.

D.28. The first program section reads the characters, stores them in a 3-byte area beginning at CHARSTR, and echoes them to a display. The second section does the conversion to binary and stores the result in BINARY. The I/O device addresses KBD_DATA and DISP_DATA are in registers R1 and R2.

READ	LDR	R0, =CHARSTR	Initialize memory pointer
	MOV	R5, #3	R0 and counter R5.
	LDRB	R3, [R1, #4]	Read a character and
	TST	R3, #2	store it in memory.
ECHO	BEQ	READ	
	LDRB	R3, [R1]	
	STRB	R3, [R0], #1	
	LDRB	R4, [R2, #4]	Echo the character
CONVERT	TST	R4, #4	to the display.
	BEQ	ECHO	
	STRB	R3, [R2]	
	SUBS	R5, R5, #1	Check if all three
	BGT	READ	characters have been read.
	LDR	R0, =CHARSTR	Initialize memory pointers
	LDR	R1, =HUNDREDS	R0, R1, and R2.
	LDR	R2, =TENS	
	LDRB	R3, [R0], #1	Load high-order BCD digit
	AND	R3, R3, #&F	into R3.
	LDR	R4, [R1, R3, LSL #2]	Load binary value
			corresponding to decimal
			hundreds value into
			accumulator register R4.
	LDRB	R3, [R0], #1	Load middle BCD digit
	AND	R3, R3, #&F	into R3.
	LDR	R3, [R2, R3, LSL #2]	Load binary value
			corresponding to
			decimal tens value
			into register R3.
	ADD	R4, R4, R3	Accumulate into R4.
	LDRB	R3, [R0], #1	Load low-order BCD digit
	AND	R3, R3, #&F	into R3.
	ADD	R4, R4, R3	Accumulate into R4.
	STR	R4, BINARY	Store converted value
			into location BINARY.

- D.29. (a) The names FP, SP, LR, and PC, are used for registers R12, R13, R14, and R15 (frame pointer, stack pointer, link register, and program counter). The 3-byte memory area for the characters begins at address CHARSTR; and the converted binary value is stored at BINARY.

The first subroutine, labeled READCHARS, is patterned after the program in Figure D.17. It echoes the characters back to a display as well as reading them into memory. The second subroutine is labeled CONVERT.

The stack frame format used is like Figure D.12.

A possible program, including subroutines, is:

Main program:

	LDR	R10, =CHARSTR	Load parameters into
	LDR	R11, =BINARY	R10 and R11 and
	STMFD	SP!, {R10, R11}	push them onto the stack.
	BL	READCHARS	Branch to first subroutine.
RTNADDR	ADD	SP, SP, #8	Remove two parameters
	...		from stack and continue.

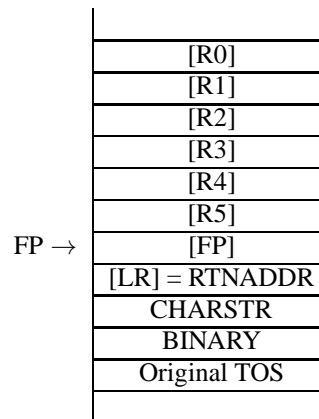
First subroutine READCHARS:

READCHARS	STMFD	SP!, {R0–R5, FP, LR}	Save registers on stack.
	ADD	FP, SP, #28	Set up frame pointer.
	LDR	R0, [FP, #4]	Load R0, R1,
	LDR	R1, =KBD_DATA	and R2 with
	LDR	R2, =DISP_DATA	parameters.
	MOV	R5, #3	Load counter R5.
READ	LDRB	R3, [R1, #4]	Read a character and
	TST	R3, #2	store it in memory.
	BEQ	READ	
	LDRB	R3, [R1]	
	STRB	R3, [R0], #1	
ECHO	LDR	R4, [R2]	Echo the character
	TST	R4, #4	to the display.
	BEQ	ECHO	
	STRB	R3, [R2]	
	SUBS	R5, R5, #1	Check if all three
	BGT	READ	characters have been read.
	LDR	R0, [FP, #8]	Load R0,R1,R2
	LDR	R5, [FP, #12]	and R5 with
	LDR	R1, =HUNDREDS	parameters.
	LDR	R2, =TENS	
	BL	CONVERT	Call second subroutine.
	LDMFD	SP!, {R0–R5, FP, PC}	Return to Main program.

Second subroutine CONVERT:

CONVERT	STMFD	SP!, {R3, R4, FP, LR}	Save registers on stack.
	ADD	FP, SP, #8	Set up frame pointer.
	LDRB	R3, [R0], #1	Load high-order BCD digit into R3.
	AND	R3, R3, #&F	
	LDR	R4, [R1, R3, LSL #2]	Load binary value corresponding to decimal hundreds value into accumulator register R4.
	LDRB	R3, [R0], #1	Load middle BCD digit into R3.
	AND	R3, R3, #&F	
	LDR	R3, [R2, R3, LSL #2]	Load binary value corresponding to decimal tens value into register R3.
	ADD	R4, R4, R3	Accumulate into R4.
	LDRB	R3, [R0], #1	Load low-order BCD digit into R3.
	AND	R3, R3, #&F	
	ADD	R4, R4, R3	Accumulate into R4.
	STR	R4, [R5]	Store binary number.
	LDMFD	SP!, {R3, R4, FP, PC}	Return to first subroutine.

(b) The contents of the top of the stack immediately after the call to the CONVERT subroutine are:



- D.30. A 16-entry table of ASCII character codes, starting at memory location TABLE, for the characters 0, 1, ... , 9, A, B, C, D, E, F, is used by the following program to generate the hexadecimal character codes, two per byte, for the ten bytes starting at memory location LOC.

The display device has the interface depicted in Figure 3.3*b*.

Assume that the character for the high-order four bits of each byte is displayed first. The two character codes for each byte are loaded into registers R1 and R2. Register R0 is initialized to the count value 10.

Program-controlled I/O is used. A display buffer memory area of 30 bytes, starting at DISPBUFF, is used to assemble the two characters, followed by a space character, for each of the ten bytes to be displayed.

	AREA	CODE	
	ENTRY		
	MOV	R0, #10	Load byte count.
	LDR	R3, =TABLE	Load addresses.
	LDR	R4, =LOC	
	LDR	R5, =DISPBUFF	
	MOV	R7, #&20	Load space character code.
LOOP	LDRB	R6, [R4], #1	Load byte to be displayed.
	MOV	R1, R6, LSR #4	Shift high-order digit into low-order digit position.
	LDRB	R1, [R3, R1]	Load hex digit code
	STRB	R1, [R5], #1	and store it in buffer.
	AND	R2, R6, #&F	Clear high-order digit.
	LDRB	R2, [R3, R2]	Load hex digit code
	STRB	R2, [R5], #1	and store it in buffer.
	STRB	R7, [R5], #1	Store space character in next byte.
	SUBS	R0, R0, #1	Check if all 10 bytes processed.
	BNE	LOOP	
DISPLAY	MOV	R0, #30	Load count of characters to be displayed.
	LDR	R5, =DISPBUFF	Load buffer address.
	LDR	R2, =DISP_DATA	Load I/O device address.
WRITEWAIT	LDRB	R4, [R2, #4]	Display 30 characters.
	TST	R4, #4	
	BEQ	WRITEWAIT	
	LDRB	R3, [R5], #1	
	STRB	R3, [R2]	
	SUBS	R0, R0, #1	
	BNE	WRITEWAIT	
		next instruction	
	AREA	DATA	
TABLE	DCB	&30, &31, &32, &33 &34, &35, &36, &37 &38, &39, &41, &42 &43, &44, &45, &46	
DISPBUFF	SPACE	30	

D.31. Assume (big-endian assumption) that the 16 bits to be displayed are in the high-order half of the word at memory location BINARY, and that the high-order bits are to be displayed first.

	LDR	R2, =DISP_DATA	Load I/O device address.
	LDR	R1, =&80000000	R1 has a single 1 in the high-order bit position.
	LDR	R7, BINARY	16-bit pattern is in high-order half of register R7.
	MOV	R0, #16	R0 is the counter for displayed bits.
	MOV	R5, #&30	Load ASCII code for 0.
	MOV	R6, #&31	Load ASCII code for 1.
LOOP	TST	R7, R1	Check if high-order bit of R7 is 1.
	MOVNE	R3, R6	If it is, load ASCII code for 1 into R3.
	MOVEQ	R3, R5	If bit is 0, load ASCII code for 0 into R3.
WRITEWAIT	LDRB	R4, [R2, #4]	Display bit character.
	TST	R4, #4	
	BEQ	WRITEWAIT	
	STRB	R3, [R2]	
	MOV	R7, R7, LSL #1	Move next bit into high-order position of register R7.
	SUBS	R0, R0, #1	Check if done.
	BGT	LOOP	

D.32. Assume that the 7-segment display device is driven by an 8-bit interface register at address location SEVEN, and that the first register in the timer interface is at address location TIM_STATUS, as shown in Figure 3.14. The number of 100-MHz clock cycles in a 1-second time period is given by the hexadecimal number 5F5E100. The following possible program uses polling to detect when the timer reaches the end of each 1-second period.

	AREA	CODE	
	ENTRY		
	LDR	R2, =SEVEN	Load device interface addresses.
	LDR	R3, =TIM_STATUS	
	LDR	R8, =TABLE	Load address of
			display patterns.
	LDR	R4, =&5F5E100	Load 1-second timer count.
	STR	R4, [R3, #8]	Store count in device register.
	MOV	R4, #6	Start timer in continuous
	STRB	R4, [R3, #4]	down counting mode.
	MOV	R5, #0	Clear digit counter.
LOOP	LDRB	R6, [R3]	Wait for timer to reach
	TST	R6, #2	end of 1-second period.
	BEQ	LOOP	
	LDRB	R7, [R8, R5]	Look up 7-segment pattern
	STRB	R7, [R2]	and display it.
	ADD	R5, R5, #1	Increment digit counter and
	CMP	R5, #10	loop back if less than 10.
	BLT	LOOP	
	MOV	R5, #0	Otherwise, reset counter to 0
	B	LOOP	and loop back.
	AREA	DATA	
TABLE	DCB	&7E, &30, &6D, &79	
	DCB	&33, &5B, &5F, &70	
	DCB	&7F, &7B	

- D.33. Assume that the two 7-segment display devices are driven by a 16-bit interface register at address location SEVENS2, and that the first register in the timer interface is at address location TIM_STATUS, as shown in Figure 3.14. The number of 100-MHz clock cycles in a 1-second time period is given by the hexadecimal number 5F5E100. The following possible program uses polling to detect when the timer reaches the end of each 1-second period.

	AREA	CODE	
	ENTRY		
	LDR	R2, =SEVENS2	Load device interface addresses.
	LDR	R3, =TIM_STATUS	
	LDR	R10, =TABLE	Load address of display patterns.
	LDR	R4, =&5F5E100	Load 1-second timer count.
	STR	R4, [R3, #8]	Store count in timer register.
	MOV	R4, #6	Start timer in continuous
	STRB	R4, [R3, #4]	down counting mode.
	MOV	R5, #0	Clear low-order digit counter.
	MOV	R6, #0	Clear high-order digit counter.
LOOP	LDRB	R7, [R3]	Wait for timer to reach
	TST	R7, #2	end of 1-second period.
	BEQ	LOOP	
	LDRB	R8, [R10, R5]	Load 7-segment pattern
			for low-order digit.
	LDRB	R9, [R10, R6]	Load 7-segment pattern
			for high-order digit.
	ORR	R8, R8, R9, LSL #8	Place high-order digit
			pattern in high-order byte
			of half word in register R8.
	STRH	R8, [R2]	Display two digit patterns.
	ADD	R5, R5, #1	Increment low-order digit counter.
	CMP	R5, #10	Loop back if count is
	BLT	LOOP	less than 10.
	MOV	R5, #0	Otherwise, reset counter to 0.
	ADD	R6, R6, #1	Increment high-order digit counter.
	CMP	R6, #10	Loop back if count is
	BLT	LOOP	less than 10.
	MOV	R6, #0	Otherwise, reset counter to 0
	B	LOOP	and loop back.
	AREA	DATA	
TABLE	DCB	&7E, &30, &6D, &79	
	DCB	&33, &5B, &5F, &70	
	DCB	&7F, &7B	

- D.34. Assume that the four 7-segment display devices are driven by the bytes of a 32-bit device interface register at address location SEVENS4, and that the first register in the timer interface is at address location TIM_STATUS, as shown in Figure 3.14. The number of 100-MHz clock cycles in a 1-second time period is given by the hexadecimal number 5F5E100. Register pair R1,R0 is used to hold the minute count (00 to 59), and register pair R3,R2 is used to hold the hour count (00 to 23). The following possible program uses polling to detect when the timer reaches the end of each 1-second period.

	AREA	CODE	
	ENTRY		
	LDR	R4, =TABLE	Load address of table of 7-segment patterns for the digits 0 through 9.
	LDR	R5, =SEVENS4	Load display device address.
	LDR	R6, =TIM_STATUS	Load address of timer device.
	LDR	R7, =&5F5E100	Load 1-second timer count.
	MOV	R0, #0	Clear minute count.
	MOV	R1, #0	
	MOV	R2, #0	Clear hour count.
	MOV	R3, #0	
	STR	R7, [R6, #8]	Load timer count and mode for
	MOV	R7, #6	continuous down counting operation.
	STRB	R7, [R6, #4]	
DELAY	MOV	R7, #0	Delay for 1 minute.
LOOP	LDRB	R8, [R6]	Wait for 1 second.
	TST	R8, #2	
	BEQ	LOOP	
	ADD	R7, R7, #1	Increment second count.
	CMP	R7, #60	Check if 60 seconds.
	BLT	LOOP	Loop back if not at 60.
	BL	DISPLAY	Otherwise, display time.
INCRCLK	ADD	R0, R0, #1	Increment real time clock.
	CMP	R0, #10	First, increment minutes from 00 to 59,
	BLT	DELAY	delay for 1 minute, then display time.
	MOV	R0, #0	
	ADD	R1, R1, #1	
	CMP	R1, #6	
	BLT	DELAY	
	MOV	R1, #0	
	CMP	R3, #2	Then, increment hours from 00 to 23.
	BEQ	TWOON3	Low-order digit of hours cycles
NOT2ON3	ADD	R2, R2, #1	from 0 through 9 while
	CMP	R2, #10	high-order digit is less than 2.
	BLT	DELAY	
	MOV	R2, #0	
	ADD	R3, R3, #1	
	B	DELAY	
TWOON3	ADD	R2, R2, #1	When high-order digit is 2,
	CMP	R2, #4	increment R2 only to 3.
	BLT	DELAY	
	MOV	R2, #0	Then, hours are reset to 00 on
	MOV	R3, #0	the minute increment after 23 59
	B	DELAY	has been reached, and time is displayed.

DISPLAY	STMFD	R13!, {R10, R11, R14}	
	LDRB	R10, [R4, R0]	Pack the four 7-segment
	LDRB	R11, [R4, R1]	patterns, corresponding to the
	ORR	R10, R10, R11, LSL #8	hours and minutes digits, into R10.
	LDRB	R11, [R4, R2]	
	ORR	R10, R10, R11, LSL #16	
	LDRB	R11, [R4, R3]	
	ORR	R10, R10, R11, LSL #24	
	STR	R10, [R5]	Send packed patterns to display.
	LDMFD	R13!, {R10, R11, R15}	
	AREA	DATA	
TABLE	DCB	&7E, &30, &6D, &79	
	DCB	&33, &5B, &5F, &70	
	DCB	&7F, &7B	

D.35. Assume that the main program starts in the Supervisor mode. This allows the execution of the privileged instruction MSR, which manipulates bits in the status register, CPSR.

Also assume that the timer device is the only device enabled to raise interrupts, and that its interrupt request signal is connected to the IRQ request line. The first register in the timer interface is at address location TIM_STATUS, as shown in Figure 3.14. The number of 100-MHz clock cycles in a 10-second time period is given by the hexadecimal number 3B9ACA00.

Main program:

	LDR	R1, =TIM_STATUS	Load timer device address.
	LDR	R2, =&3B9ACA00	LOAD 10-second count
	STR	R2, [R1, #8]	and store in timer device.
	LDR	R2, #7	Load timer mode for continuous
	STRB	R2, [R1, #4]	down counting operation
			with interrupt enabled.
	MOV	R2, #&50	Switch processor to User mode
	MSR	CPSR, R2	and clear IRQ disable bit.
COMPUTE		next instruction	

Interrupt-service routine:

IRQLOC	STMFD	R13!, {R0}	Save register R0.
	LDRB	R0, [R1]	Clear TIRQ and ZERO bits
			in timer device status register.
	BL	DISPLAY	Display updated results
			from COMPUTE routine.
RETURN	LDMFD	R13!, {R0}	Restore register R0
	SUBS	R15, R14, #4	and return.

D.36. An ARM program corresponding to Figure 3.18 is given as follows:

	AREA	CODE	
	ENTRY		
	LDR	R1, =&800	Load address DIGIT.
	LDR	R7, =&4030	Load address of 7-segment display.
	LDR	R6, =TABLE	Load address of table of
			7-segment patterns.
	LDRB	R2, [R1]	Load ASCII-encoded digit.
	AND	R3, R2, #&F0	Extract high-order bits of code.
	AND	R2, R2, #&F	Extract BCD number.
	CMP	R3, #&30	Check if high-order bits
	BEQ	HIGH3	are 0011.
	MOV	R2, #&F	If no, load index to blank.
HIGH3	LDRB	R5, [R6, R2]	Load 7-segment pattern.
	STRB	R5, [R7]	Display the digit.
	AREA	DATA	
TABLE	DCB	&7E, &30, &6D, &79	
	DCB	&33, &5B, &5F, &70	
	DCB	&7F, &7B, &00, &00	
	DCB	&00, &00, &00, &00	

D.37. The prompt is “Name: ”, after which the user types a name, followed by carriage return, which is echoed back on the same line and stored in memory, starting at location NAME. The next line leads off with the message “Name Reversed: ”, after which the user name is displayed in reverse.

The two subroutines for reading and displaying characters used by the following ARM program use registers R1, R2, R3, and R4. These registers are not used in the main program, so they are not saved/restored on subroutine entry/exit. The subroutines are patterned after the routines in Section D.8.1.

The ASCII characters for the prompt and the message are stored in byte strings starting at memory locations PROMPT and MESSAGE, respectively. Each of these strings is terminated by a space character.

Main program:

	AREA	CODE	
	ENTRY		
	LDR	R0, =NAME	Load address for user name.
	LDR	R1, =KBD_DATA	See Figure 3.3 for I/O
	LDR	R2, =DISP_DATA	device addresses.
	MOV	R10, #&20	Load ASCII code for space.
	MOV	R11, #&0D	Load ASCII code for carriage return.
	LDR	R6, =PROMPT	R6 points to prompt characters.
LOOP1	LDRB	R3, [R6], #1	Display prompt, which is
	BL	WRITECHAR	terminated by a space character.
	CMP	R3, R10	Check for space.
	BNE	LOOP1	If not, loop back.
LOOP2	BL	READCHAR	Read user name and store
	STRB	R3, [R0], #1	at NAME.
	BL	WRITECHAR	Echoback name.
	CMP	R3, R11	Check for carriage return.
	BNE	LOOP2	If not, loop back.
	LDR	R6, =MESSAGE	R6 points to message characters.
LOOP3	LDRB	R3, [R6], #1	Display message, which is
	BL	WRITECHAR	terminated by a space character.
	CMP	R3, R10	Check for space.
	BNE	LOOP3	If not, loop back.
	LDR	R6, =NAME	Load address for user name.
	SUB	R0, R0, #1	Decrement R0 to point to
			carriage return following user name.
LOOP4	LDRB	R3, [R0, #-1]!	Display user name in reverse.
	BL	WRITECHAR	
	CMP	R6, R0	Has first letter of name
			been displayed?
	BNE	LOOP4	If not, loop back.
			next instruction

Subroutines:

READCHAR	LDRB	[R1, #4]	Subroutine to read character.
	TST	R3, #2	
	BEQ	READCHAR	
	LDRB	R3, [R1]	
	MOV	R15, R14	
			Return.
WRITECHAR	LDRB	R4, [R2, #4]	Subroutine to display character.
	TST	R4, #4	
	BEQ	WRITECHAR	
	STRB	R3, [R2]	
	MOV	R15, R14	
			Return.
	AREA	DATA	
PROMPT	DCB	&4E, &61, &6D, &65	
		&3A, &20	
MESSAGE	DCB	&4E, &61, &6D, &65	
		&20, &72, &65, &76	
		&65, &72, &73, &65	
		&64, &3A &20	
NAME	SPACE	40	

D.38. The following possible program assumes that there are at least two characters in the word entered by the user. The characters are stored starting at the memory location WORD. Subroutines are used to read and display characters using the routines shown in Section D.8.1. They use registers R1 through R4, which are not saved/restored in the subroutines. The prompt character is ‘*’; and the result of the palindrome check is ‘Y’ (yes) or ‘N’ (no).

Main program:

	LDR	R1, =KBD_DATA	Load addresses for I/O devices.
	LDR	R2, =DISP_DATA	
	LDR	R5, =WORD	
	MOV	R6, R5	R6 is a pointer that is used to step through the character locations.
	MOV	R3, #&2A	Display the prompt character,
	BL	WRITECHAR	‘*’ (&2A), followed by
	MOV	R3, #&0D	a carriage return character (&0D).
	BL	WRITECHAR	
GETCHARS	BL	READCHAR	Read a character typed by the user.
	STRB	R3, [R6], #1	Store it in memory and increment address pointer.
	BL	WRITECHAR	Echo character back to display.
	CMP	R3, #&0D	Check for carriage return and
	BNE	GETCHARS	proceed to palindrome analysis if yes; otherwise, continue reading.
	SUBS	R6, R6, #2	Decrement R6 to point to last character of word entered and start palindrome analysis.
	MOV	R3, #&59	Anticipate “yes” answer.
COMPCHARS	LDRB	R7, [R5], #1	Working in from both ends of
	LDRB	R8, [R6], #-1	word, check if characters match.
	CMP	R7, R8	
	MOVNE	R3, #&4E	If they do not, change “yes” to “no”.
	CMP	R6, R5	See Note below.
	BGT	COMPCHARS	If not finished checking, branch back to continue.
	BL	WRITECHAR	Otherwise, display answer.
		next instruction	

Note: Character checking has been completed when the pointers become equal (for an odd length word) or they cross each other (for an even length word). Otherwise, checking continues while R6 is greater than R5.

Subroutines:

READCHAR	LDRB	R3, [R1, #4]	
	TST	R3, #2	
	BEQ	READCHAR	
	LDRB	R3, [R1]	
	MOV	R15, R14	Return.
WRITECHAR	LDRB	R4, [R2, #4]	
	TST	R4, #4	
	BEQ	WRITECHAR	
	STRB	R3, [R2]	
	MOV	R15, R14	Return.

D.39. The first part of the main program below determines the length of the sample text to be displayed and writes the length in register R0. The expression

$$[80 - ([R0] + 2)]/2 = 39 - [R0]/2$$

gives the number of leading space characters needed to center the display.

Three subroutines are used. LENGTH determines the length of the given character string located at memory location STRING. DISPA displays the upper or lower line of the display box. WRITECHAR displays a single character, using the routine given in Section D.8.1.

Main program:

	LDR	R1, =STRING	Load address of string.
	MOV	R0, #0	Clear register R0.
	BL	LENGTH	Compute and write length of string into R0.
	CMP	R0, #78	If length greater than 78,
	MOVGT	R0, #78	truncate to 78.
			Length of sample text is now in R0.
	MOV	R1, R0, ASR #1	Write [R0]/2 into R1.
	RSB	R1, R1, #39	Number of leading spaces
			is now in R1.
	MOV	R10, R0	Load length of sample text into R10.
	MOV	R11, R1	Load number of leading spaces into R11.
	BL	DISPA	Display upper line of bounding box,
			followed by carriage return.
	MOV	R10, R0	Display sample text between
	MOV	R11, R1	vertical lines.
	MOV	R3, #&20	Load space character
LOOP1	SUBS	R11, R11, #1	and display 0 or more leading spaces.
	BLT	DISP2	
	BL	WRITECHAR	
	B	LOOP1	
DISP2	MOV	R3, #&7C	Display vertical character.
	BL	WRITECHAR	
	LDR	R5, =STRING	Display sample text.
LOOP2	LDRB	R3, [R5], #1	
	BL	WRITECHAR	
	SUBS	R10, R10, #1	
	BNE	LOOP2	
	MOV	R3, #&7C	Display vertical character.
	BL	WRITECHAR	
	MOV	R3, #&0D	Send carriage return.
	BL	WRITECHAR	
	MOV	R10, R0	Display bottom line of
	MOV	R11, R1	bounding box.
	BL	DISPA	
			next instruction

Subroutines:

LENGTH	LDRB	R5, [R1], #1	Determine length of string
	CMP	R5, #&00	and return in R0.
	ADDNE	R0, R0, #1	
	BNE	LENGTH	
	MOV	R15, R14	Return.
DISPA	MOV	R3, #&20	ASCII code for space.
LOOP3	SUBS	R11, R11, #1	Display zero or more
	BLT	DISP1	leading space characters.
	BL	WRITECHAR	
	B	LOOP3	
DISP1	MOV	R3, #&2B	Display '+' character.
	BL	WRITECHAR	
	MOV	R3, #&2D	Display one or more '-' characters.
LOOP4	BL	WRITECHAR	
	SUBS	R10, R10, #1	
	BNE	LOOP4	
	MOV	R3, #&2B	Display '+' character.
	BL	WRITECHAR	
	MOV	R3, #&0D	Send carriage return.
	BL	WRITECHAR	
	MOV	R15, R14	Return.
WRITECHAR	LDRB	R4, [R2, #4]	Assume R2 contains DISP_DATA
	TST	R4, #4	and the next word contains DISP_STATUS.
	BEQ	WRITECHAR	
	STRB	R3, [R2]	R3 contains character to be displayed.
	MOV	R15, R14	Return.

D.40. Assume that the text to be displayed is stored in the memory starting at location TEXT. Also, assume that words consist of from 1 to 80 characters, and that the text contains at least one word.

The following program is a possible solution. The program does a complete scan of the text, replacing some of the space characters with carriage returns to meet the stated requirements. During this preprocessing of the text, no lines are sent to the display. After the preprocessing has been completed, the text is displayed. The NUL character is replaced by a carriage return when the text is displayed.

Registers R2, R3, and R4 are used by the single character output subroutine, patterned after the routine shown in Section D.8.1, assuming the interface arrangement shown in Figure 3.3b.

	LDR	R2, =DISP_DATA	Load I/O device address.
	LDR	R0, =TEXT	Register R0 points to start of text.
	MOV	R1, #&0D	R1 contains carriage return character.
RESET	MOV	R5, #0	Clear character counter.
SCAN	LDRB	R6, [R0], #1	Start/continue scanning text.
	ADD	R5, R5, #1	Increment character count.
	CMP	R6, #00	If character is NUL,
	BEQ	DISPLAY	display text.
	CMP	R5, #81	If not at 81st character,
	BLT	SCAN	continue scanning.
	CMP	R6, #&20	If 81st character is a space,
	STREQB	R1, [R0, #-1]	replace with carriage return,
	BEQ	RESET	reset counter, and continue scanning.
BACKUP	LDRB	R6, [R0, #-1]!	Otherwise, backup until
	CMP	R6, #&20	space is encountered.
	BNE	BACKUP	
	STRB	R1, [R0], #1	When space is encountered,
	B	RESET	replace it with a carriage return,
			reset counter, and continue scanning.
DISPLAY	LDR	R0, =TEXT	Starting at TEXT, display
DISPLAY1	LDRB	R3, [R0], #1	characters until NUL
	CMP	R3, #& 00	is reached.
	BEQ	FINISH	
	BL	WRITECHAR	
	B	DISPLAY1	
FINISH	MOV	R3, R1	When NUL is reached,
	BL	WRITECHAR	display carriage return.
		next instruction	
WRITECHAR	LDRB	R4, [R2, #4]	Subroutine to display character.
	TST	R4, #4	
	BEQ	WRITECHAR	
	STRB	R3, [R2]	
	MOV	R15, R14	Return.