

1.

Suppose  $C \cong C'$ . Then there is an isomorphism  $f : C \rightarrow C'$  with inverse  $f^{-1}$ . Thus  $Ff : FC \rightarrow FC'$  is an isomorphism with inverse  $Ff^{-1}$ . Conversely, suppose  $FC \cong FC'$ . Then there is an isomorphism  $h : FC \rightarrow FC'$ . Since  $F$  is full, there are  $f : C \rightarrow C', g : C' \rightarrow C$  such that  $Ff = h, Fg = h^{-1}$ . Then  $1_{FC} = h \circ h^{-1} = Ff \circ Fg = F(f \circ g)$ . Since  $F$  is faithful,  $f \circ g = 1_C$  and thus  $g = f^{-1}$ , meaning  $C \cong C'$ .

2.

Let  $P, Q \in \mathbf{Sets}^{\mathbf{C}^{\text{op}}}$  and  $\varphi, \psi : P \rightarrow Q$  such that  $\varphi \circ \vartheta = \psi \circ \vartheta$  for  $\vartheta : yC \rightarrow P$  for each  $C \in \mathbf{C}$ . We can define a map

$$\varphi_{C*} = \psi_{C*} : \text{Hom}(yC, P) \rightarrow \text{Hom}(yC, Q)$$

by  $\varphi_{C*}(\vartheta) = \psi_{C*}(\vartheta) = \varphi \circ \vartheta = \psi \circ \vartheta$ . By Yoneda lemma, we can deduces  $\text{Hom}(yC, P) \cong PC$  and  $\text{Hom}(yC, Q) \cong QC$ , making a commutative diagram

$$\begin{array}{ccc} \text{Hom}(yC, P) & \xrightarrow{\cong} & PC \\ \psi_{C*} \downarrow \varphi_{C*} & & \psi_C \downarrow \varphi_C \\ \text{Hom}(yC, Q) & \xrightarrow{\cong} & QC \end{array} .$$

But since  $\varphi_{C*} = \psi_{C*}$ , we conclude that  $\varphi_C = \psi_C$  for all  $C$ , meaning  $\varphi = \psi$ .

3.

$$\begin{aligned} \text{Hom}(X, A^B \times A^C) &\cong \text{Hom}(X, A^B) \times \text{Hom}(X, A^C) \\ &\cong \text{Hom}(X \times B, A) \times \text{Hom}(X \times C, A) \\ &\cong \text{Hom}(X \times B + X \times C, A) \\ &\cong \text{Hom}(X \times (B + C), A) \\ &\cong \text{Hom}(X, A^{B+C}). \end{aligned}$$

By Yoneda lemma,  $A^B \times A^C \cong A^{B+C}$ . Similarly,  $(A \times B)^C \cong A^C \times B^C$ .

4.

Let  $f : P \rightarrow Q$ . Then  $S(f) = \text{Hom}(-, f) : \text{Hom}(-, P) \rightarrow \text{Hom}(-, Q)$  and hence  $S$  is a functor. This functor is faithful. From text, simplicial nerve is defined by for any  $n$  and a poset  $P$ ,

$$S(P)_n = \{(p_1, \dots, p_n) \in P^n \mid p_1 \leq \dots \leq p_n\},$$

which is just a set of monotone map

$$S(P)(n) = \text{Hom}([n], P).$$

5.

Define  $S : \mathbf{Cat} \rightarrow \mathbf{Cat}^{\Delta^{\text{op}}}$  by  $S(\mathbf{C}) = \text{Fun}(-, \mathbf{C})$ . Then it generalizes the concept of simplicial nerve.

6.

Consider a diagram  $\mathbf{J} \rightarrow \mathbf{D}^{\mathbf{C}}$ . Then for some  $C \in \mathbf{C}_0$ , we have an induced diagram  $\mathbf{J} \rightarrow \mathbf{D}$  defined by  $i \mapsto F_i(C)$  and  $(i \rightarrow j) \mapsto (F_i(C) \rightarrow F_j(C))$ . Since  $\mathbf{D}$  is complete, we have a limit  $\lim_{\leftarrow \mathbf{J}} F_i(C)$ . Thus from the diagram, we have the limit  $\lim_{\leftarrow \mathbf{J}} F_i$  defined by  $(\lim_{\leftarrow \mathbf{J}} F_i)(C) = \lim_{\leftarrow \mathbf{J}} (F_i(C))$ . Colimit of  $\mathbf{D}^{\mathbf{C}}$  is the limit of  $(\mathbf{D}^{\text{op}})^{\mathbf{C}^{\text{op}}} = (\mathbf{D}^{\mathbf{C}})^{\text{op}}$ .

7.

For product,

$$\begin{aligned} y(A \times B)(C) &\cong \text{Hom}(C, A \times B) \\ &\cong \text{Hom}(C, A) \times \text{Hom}(C, B) \\ &\cong yA(C) \times yB(C) \\ &\cong (yA \times yB)(C). \end{aligned}$$

Thus,  $yA \times yB \cong y(A \times B)$ . For exponential,

$$\begin{aligned} (yB)^{(yA)}(C) &\cong \text{Hom}(yC, (yB)^{(yA)}) \\ &\cong \text{Hom}(yC \times yA, yB) \\ &\cong \text{Hom}(y(C \times A), yB) \\ &\cong \text{Hom}(C \times A, B) \\ &\cong \text{Hom}(C, B^A) \\ &\cong y(B^A)(C). \end{aligned}$$

Thus  $(yB)^{(yA)} \cong y(B^A)$ .

8.

Let  $C, C' \in P$  and  $AC = x, AC' = x'$ . In  $\int_P A$ , define  $(x, C) \leq (x', C')$  if  $C \leq C'$  and there is a map  $f : x \mapsto x'$ . Then it is a poset and the projection  $\pi$  is monotone. Finally, consider a natural transformation  $\alpha : A \rightarrow B$ . Then we can define  $\int_P \alpha : \int_P A \rightarrow \int_P B$  by  $(AC, C) \mapsto (BC, C)$  for all  $C \in P$ . It clearly defines the functor

$$\int_P : \mathbf{Sets}^{P^{\text{op}}} \rightarrow \mathbf{Pos}/P.$$

9,10.

skip

11.

Product of  $a : A \rightarrow X$  and  $b : B \rightarrow X$  is a pullback of them. Since  $X$  is a set,  $\mathbf{Sets}^X$  is equivalent to  $\mathbf{Sets}/X$ . Therefore, it is enough to find the exponential of  $\mathbf{Sets}^X$  to show that  $\mathbf{Sets}/X$  is cartesian closed. As a category,  $X$  is a discrete category. Yoneda embedding  $y : X \rightarrow \mathbf{Sets}^X$  is defined by  $y(x) = \text{Hom}_X(-, x) : X \rightarrow \mathbf{Sets}$  for  $x \in X$ . Since  $X$  is discrete,  $y(x)(z) = \emptyset$  for all  $z \neq x$  and  $y(x)(x) = \{1_x\}$ . Now for  $P, Q \in \mathbf{Sets}^X$ , exponential  $Q^P$  can be defined by  $Q^P(x) = \text{Hom}(y(x) \times P, Q)$ , which exists. For a category  $\mathbf{C}$ , the subobject classifier is the pair  $(\Omega, t : 1 \rightarrow \Omega)$  where  $\Omega(C)$  is a set of sieves on  $C$  and  $t_C$  is a total sieve. Since  $X$  is discrete,  $\Omega(x) = \{1_x\}$  and  $t_x = \{1_x\}$  for all  $x \in X$ .

**12.**

- (a) We know that  $\Omega(C) = \{S \mid S \text{ sieve on } C\}$ . For  $p \in P^{\text{op}}$  where  $P$  is a poset, then  $S$  is the set of arrows  $\{p \leq x\}$ , which is identical to  $\{x \mid p \leq x\}$ . Therefore, for  $\mathbf{P} = \mathbf{2} = \{0 \leq 1\}$ ,  $\Omega : \mathbf{2}^{\text{op}} \rightarrow \mathbf{Sets}$  where  $\Omega(0) = \{\{0, 1\}, \{1\}\}$ ,  $\Omega(1) = \{\{1\}\}$  and  $t : 1 \rightarrow \Omega$  where  $t_0(*) = \{0, 1\}$ ,  $t_1(*) = \{1\}$ . For  $\mathbf{P} = \omega$ , the subobject classifier is defined by

$$\begin{aligned}\Omega(0) &= \{\{0, 1, 2, \dots\}, \{1, 2, 3, \dots\}, \{2, 3, 4, \dots\}, \dots\} \\ \Omega(1) &= \{\{1, 2, 3, \dots\}, \{2, 3, 4, \dots\}, \{3, 4, 5, \dots\}, \dots\} \\ &\vdots \\ \Omega(n) &= \{\{n, n+1, \dots\}, \{n+1, n+2, \dots\}, \dots\}\end{aligned}$$

and  $t_n = \{n, n+1, \dots\}$ .

- (b) Since all topos operations (pullbacks, exponentials, subobject classifiers) are finite, it is also a topos.

**13.**

Recall that the category of graphs is equivalent to the functor category  $\mathbf{Sets}^{\Gamma^{\text{op}}}$  where  $\Gamma$  is a category

$$C \rightrightarrows D.$$

Thus the category of graphs is a topos. Then the subobject classifier  $\Omega$  is defined by  $\Omega(C) = \{1_C\}$  and  $\Omega(D) = \{1_D, s : C \rightarrow D, t : C \rightarrow D\}$ . point  $1 \rightarrow \Omega$  is the same.