The Identity Problem in the special affine group of \mathbb{Z}^2

Ruiwen Dong

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Markov (1940s): is (semigroup) Membership Problem decidable?

Input: Set of square matrices $\mathcal{G} = \{A_1, \dots, A_K\}$, target matrix T. **Output:** Is there a sequence $B_1, B_2, \dots, B_m \in \mathcal{G}$, s.t. $B_1B_2 \cdots B_m = T$?

Image: A matrix and a matrix

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Special case: is the Identity Problem decidable?

Input: Set of square matrices $\mathcal{G} = \{A_1, \dots, A_K\}$. **Output:** Is there a sequence $B_1, B_2, \dots, B_m \in \mathcal{G}$, s.t. $B_1B_2 \cdots B_m = I$? i.e. whether $I \in \langle \mathcal{G} \rangle$?

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Known results.

 $SL(n,\mathbb{Z})$: the group of $n \times n$ integer matrices of determinant one.

group type	Membership: $T \in \langle \mathcal{G} \rangle$?	Identity Prob: $I \in \langle \mathcal{G} \rangle$?
$SL(2,\mathbb{Z})$	NP-complete [BHP23]	NP-complete [BHP17]
SL(3,ℤ)	?	?
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* There exist groups where Membership Problem is undecidable but Identity Problem is decidable.

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 $\mathsf{SA}(2,\mathbb{Z})$: the Special Affine group

$$\left\{M = \begin{pmatrix} * & * & * \\ * & * & * \\ 0 & 0 & 1 \end{pmatrix} \middle| \det(M) = 1\right\}$$

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Let $\mathcal{G} = \{(A_1, a_1), \dots, (A_K, a_K)\}$. Goal: decide whether $(I, \mathbf{0}) \in \langle \mathcal{G} \rangle$.

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Elements of SA(2, \mathbb{Z}) are (A, a). Group law (A, a)(B, b) = (AB, Ab + a). Let $\mathcal{G} = \{(A_1, a_1), \dots, (A_K, a_K)\}$. Goal: decide whether $(I, \mathbf{0}) \in \langle \mathcal{G} \rangle$.

Step 1: for s = 1, ..., K, check if $A_s^{-1} \in \langle A_1, ..., A_K \rangle$. If $A_s^{-1} \notin \langle A_1, ..., A_K \rangle$, then

$$(A_i, \boldsymbol{a}_i) \cdots (A_s, \boldsymbol{a}_s) \cdots (A_{i'}, \boldsymbol{a}_{i'}) \neq (I, \boldsymbol{0}).$$

So we can delete (A_s, a_s) from \mathcal{G} .

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Theorem (Bell, Hirvensalo, Potapov)

It is decidable in NP whether $A_s^{-1} \in \langle A_1, \ldots, A_K \rangle$.

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Theorem (Bell, Hirvensalo, Potapov)

It is decidable in NP whether $A_s^{-1} \in \langle A_1, \ldots, A_K \rangle$.

We can perform Step 1 iteratively, until $A_s^{-1} \in \langle A_1, \ldots, A_K \rangle$ for all s. So the semigroup $\langle A_1, \ldots, A_K \rangle$ becomes a group.

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$$\mathcal{G} = \{ (A_1, \boldsymbol{a}_1), \dots, (A_K, \boldsymbol{a}_K) \}.$$

Additionally, $H = \langle A_1, \ldots, A_K \rangle \leq SL(2, \mathbb{Z})$ is a group.

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Theorem (Tits alternative)

Let *H* be a subgroup of $SL(n, \mathbb{Z})$. Then

- either H contains a non-abelian free subgroup,
- I or H contains a solvable subgroup of finite index.

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In SL(2, \mathbb{Z}) this means:

- either H contains two matrices A, B that are not simultaneously triangularizable,
- **(2)** or *H* contains a finite-index subgroup that is isomorphic to \mathbb{Z} or $\{I\}$.

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In SL(2, \mathbb{Z}) this means:

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• or *H* contains a finite-index subgroup that is isomorphic to \mathbb{Z} or $\{I\}$. Furthermore, the two cases can be distinguished in PTIME (Beals 1999).

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Proposition

Suppose $\langle A_1, \ldots, A_K \rangle \leq SL(2, \mathbb{Z})$ is a group containing two matrices A, B that are not simultaneously triangularizable, then $(I, \mathbf{0}) \in \langle \mathcal{G} \rangle$.

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Suppose $\langle A_1, \ldots, A_K \rangle \leq SL(2, \mathbb{Z})$ is a group containing two matrices A, B that are not simultaneously triangularizable, then $(I, \mathbf{0}) \in \langle \mathcal{G} \rangle$.

Proof idea:

Since $\langle A_1, \ldots, A_K \rangle$ is a group containing A and B, it also contains some Y such that AYB = I. In particular $\langle \mathcal{G} \rangle$ contains some elements $(A, \mathbf{a}), (Y, \mathbf{y}), (B, \mathbf{b})$ such that $(A, \mathbf{a})(Y, \mathbf{y})(B, \mathbf{b}) = (I, \mathbf{x})$ for some $\mathbf{x} \in \mathbb{Z}^2$.

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Suppose A has eigenspaces V_A , W_A , and B has eigenspaces V_B , W_B , since A and B are not simultaneously triangularizable, we can suppose V_A , W_A , V_B , W_B pairwise distinct.



We have $(A, \boldsymbol{a}), (Y, \boldsymbol{y}), (B, \boldsymbol{b}) \in \langle \mathcal{G} \rangle$ s.t. $(A, \boldsymbol{a})(Y, \boldsymbol{y})(B, \boldsymbol{b}) = (I, \boldsymbol{x}).$



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Consider

$$\underbrace{(A, \boldsymbol{a})(Y, \boldsymbol{y})(A, \boldsymbol{a})(Y, \boldsymbol{y})\cdots(A, \boldsymbol{a})(Y, \boldsymbol{y})}_{m \text{ times}}(B, \boldsymbol{b})^m = (I, \boldsymbol{x}_1) \in \langle \mathcal{G} \rangle$$



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When $m \to \infty$, the vector \mathbf{x}_1 tends towards V_B .



We have $(A, \boldsymbol{a}), (Y_1, \boldsymbol{y}_1), (B, \boldsymbol{b}) \in \langle \mathcal{G} \rangle$ s.t. $(A, \boldsymbol{a})(Y_1, \boldsymbol{y})(B, \boldsymbol{b}) = (I, \boldsymbol{x}_1).$



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We have $(A, a), (Y_1, y_1), (B, b) \in \langle \mathcal{G} \rangle$ s.t. $(A, a)(Y_1, y_1)(B, b) = (I, x_1)$. Consider

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When $m \to \infty$, the vector \boldsymbol{x}_2 tends towards W_A .



We have $(A, a), (Y_2, y_2), (B, b) \in \langle \mathcal{G} \rangle$ s.t. $(A, a)(Y_2, y_2)(B, b) = (I, x_2)$.



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$$(A, \boldsymbol{a})(Y_2, \boldsymbol{y}_2)(B, \boldsymbol{b})^m \underbrace{(Y_2, \boldsymbol{y}_2)(A, \boldsymbol{a}) \cdots (Y_2, \boldsymbol{y}_2)(A, \boldsymbol{a})}_{m \text{ times}} (B, \boldsymbol{b}) = (I, \boldsymbol{x}_3) \in \langle \mathcal{G} \rangle$$



We have $(A, a), (Y_2, y_2), (B, b) \in \langle \mathcal{G} \rangle$ s.t. $(A, a)(Y_2, y_2)(B, b) = (I, x_2)$. Consider

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When $m \to \infty$, the vector \boldsymbol{x}_3 tends towards W_B .



Continue like this, we obtain $(I, \mathbf{x}_1), (I, \mathbf{x}_2), \dots, (I, \mathbf{x}_6) \in \langle \mathcal{G} \rangle$.



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There exist **positive** integers n_1, \ldots, n_6 such that $n_1 \mathbf{x}_1 + \cdots + n_6 \mathbf{x}_6 = \mathbf{0}$.



Continue like this, we obtain $(I, \mathbf{x}_1), (I, \mathbf{x}_2), \dots, (I, \mathbf{x}_6) \in \langle \mathcal{G} \rangle$.

There exist **positive** integers n_1, \ldots, n_6 such that $n_1 x_1 + \cdots + n_6 x_6 = \mathbf{0}$. Therefore

$$(I,\mathbf{0})=(I,\mathbf{x}_1)^{n_1}(I,\mathbf{x}_2)^{n_2}\cdots(I,\mathbf{x}_6)^{n_6}\in\langle\mathcal{G}
angle.$$



$$\mathcal{G} = \{(A_1, \boldsymbol{a}_1), \dots, (A_K, \boldsymbol{a}_K)\}.$$

We have proved the first case of the dichotomy:

Proposition

Suppose $\langle A_1, \ldots, A_K \rangle \leq SL(2, \mathbb{Z})$ is a group containing two matrices A, B that are not simultaneously triangularizable, then $(I, \mathbf{0}) \in \langle \mathcal{G} \rangle$.

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We can also prove the second case of the dichotomy:

Proposition

Suppose $\langle A_1, \ldots, A_K \rangle$ is a group containing a finite-index subgroup that is isomorphic to \mathbb{Z} or $\{I\}$. Then it is decidable in PTIME whether or not $(I, \mathbf{0}) \in \langle \mathcal{G} \rangle$.

Let $\mathcal{G} = \{(A_1, a_1), \dots, (A_K, a_K)\}$, we want to decide if $(I, \mathbf{0}) \in \langle \mathcal{G} \rangle$. We defined $H = \langle A_1, \dots, A_K \rangle$.

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Let $\mathcal{G} = \{(A_1, a_1), \dots, (A_K, a_K)\}$, we want to decide if $(I, \mathbf{0}) \in \langle \mathcal{G} \rangle$. We defined $H = \langle A_1, \dots, A_K \rangle$.

Step 1: narrowing down to the case where H is a group is done in NP.

- Step 2: distinguishing dichotomy is in PTIME.
- Step 3: first dichotomy case, always true.
- Step 4: second dichotomy case, complexity is PTIME.

In total, complexity is in NP.

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NP-hardness comes from the NP-hardness in $SL(2,\mathbb{Z}) \leq SA(2,\mathbb{Z})$.

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NP-hardness comes from the NP-hardness in $SL(2,\mathbb{Z}) \leq SA(2,\mathbb{Z})$.

Theorem

The Identity Problem in $SA(2, \mathbb{Z})$ is NP-complete.

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Let $\mathcal{G} = \{(A_1, a_1), \dots, (A_K, a_K)\}$, we want to decide if $(I, \mathbf{0}) \in \langle \mathcal{G} \rangle$. We defined $H = \langle A_1, \dots, A_K \rangle$.

Step 1: narrowing down to the case where H is a group is done in NP.

- Step 2: distinguishing dichotomy is in PTIME.
- Step 3: first dichotomy case, always true.
- Step 4: second dichotomy case, complexity is PTIME.

In total, complexity is in NP.

NP-hardness comes from the NP-hardness in $SL(2, \mathbb{Z}) \leq SA(2, \mathbb{Z})$.

Theorem

The Identity Problem in $SA(2, \mathbb{Z})$ is NP-complete.

Open Problem

Is Membership Problem in $SA(2,\mathbb{Z})$ decidable?

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