

# The Identity Problem for nilpotent groups of bounded class

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Markov (1940s): is *Semigroup Membership* decidable?

## Definition (Semigroup Membership)

**Input:** Set of square matrices  $S = \{A_1, \dots, A_K\}$ , target matrix  $T$ .

**Output:** Is there a sequence  $A_{i_1}, A_{i_2}, \dots, A_{i_m} \in S$ , s.t.  $A_{i_1} A_{i_2} \cdots A_{i_m} = T$ ?

# Membership Problems

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## Theorem (Bell, Potapov 2010)

*Identity Problem is undecidable, even when  $S \subseteq \text{SL}(4, \mathbb{Z})$ .*

Membership Problems has been proven to be decidable in groups with additional structures:

**Theorem (Babai et al. 1996)**

*Semigroup Membership is decidable (and NP) for commutative matrices.  
Identity Problem is decidable (and PTIME) for commutative matrices.*

“proof”: suppose we work with  $(\mathbb{Z}^n, +)$  instead of multiplication of commutative matrices. Suppose  $S = \{a_1, \dots, a_K\} \subset \mathbb{Z}^n$  and  $t \in \mathbb{Z}^n$ .

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$t$  is in semigroup generated by  $S \iff$  there exist  $n_1, \dots, n_K \in \mathbb{N}$ , such that  $n_1 \cdot a_1 + \dots + n_K \cdot a_K = t$ .

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$e$  is in semigroup generated by  $S \iff$  there exist  $n_1, \dots, n_K \in \mathbb{N}$ , not all zero, such that  $n_1 \cdot a_1 + \dots + n_K \cdot a_K = 0^n$ .

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Solving linear equations over  $\mathbb{Q}_{\geq 0}$ : PTIME (linear programming).

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## Definition

The **lower central series** of a group  $G$  is the sequence of subgroups

$$G = G_1 \geq G_2 \geq G_3 \geq \cdots,$$

in which  $G_k = [G, G_{k-1}]$ . ( $[G, H]$  is the group generated by  $ghg^{-1}h^{-1}, g \in G, h \in H$ .)

$G$  is **nilpotent** if  $G_{d+1} = \{I\}$  for some  $d$ . The smallest such  $d$  is the **nilpotency class** of  $G$ .

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## Example

$G = \text{UT}(3, \mathbb{Z})$  has nilpotency class two:

$$G_1 = \left\{ \begin{pmatrix} 1 & * & * \\ 0 & 1 & * \\ 0 & 0 & 1 \end{pmatrix} \right\} \geq G_2 = \left\{ \begin{pmatrix} 1 & 0 & * \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \right\} \geq G_3 = \left\{ \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \right\}$$

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$\text{UT}(n, \mathbb{Z})$  has nilpotency class  $n - 1$ , so does  $\text{UT}(n, \mathbb{Z})^k$ .

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# Semigroup Membership and Identity Problem in nilpotent groups

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*For each  $d > 10$ , subject to a conjecture  $P_d$ , the Identity Problem is decidable (and PTIME) in all finitely generated nilpotent groups of class  $d$ .  
For each  $d$ , the conjecture  $P_d$  can be verified by computer algebra software in case it is true.*

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Denote by  $\langle S \rangle_{\text{semigrp}}$  the semigroup generated by  $S$ .

## Lemma (Very easy lemma)

*We have  $1 \in \langle S \rangle_{\text{semigrp}}$  if and only if there exists a non-empty subset  $H \subseteq S$ , such that the semigroup  $\langle H \rangle_{\text{semigrp}}$  is a group.*

# Identity Problem in nilpotent groups: main idea

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Recall that  $[G, G]$  denotes the (normal) subgroup of  $G$  generated by  $ghg^{-1}h^{-1}$ ,  $g, h \in G$ . In particular, the quotient group  $G/[G, G]$  is abelian.

## Proposition (Very difficult proposition)

*For  $d \leq 10$ , let  $G$  be a class- $d$  nilpotent group. Let  $S$  be the generators of  $G$  as a group, then  $\langle S \rangle_{\text{semigrp}} = G$  if and only if  $\langle S[G, G] \rangle_{\text{semigrp}} = G/[G, G]$ .*

## Example with $UT(4, \mathbb{Z})$

Let's illustrate with  $G := UT(4, \mathbb{Z})$ .

$$G = \left\{ \begin{pmatrix} 1 & * & * & * \\ 0 & 1 & * & * \\ 0 & 0 & 1 & * \\ 0 & 0 & 0 & 1 \end{pmatrix} \mid * \in \mathbb{Z} \right\}, \quad [G, G] = \left\{ \begin{pmatrix} 1 & 0 & * & * \\ 0 & 1 & 0 & * \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \mid * \in \mathbb{Z} \right\}$$

$G/[G, G] \cong \mathbb{Z}^3$ : multiplication acts additively on the superdiagonal.

$$\begin{pmatrix} 1 & a_1 & * & * \\ 0 & 1 & b_1 & * \\ 0 & 0 & 1 & c_1 \\ 0 & 0 & 0 & 1 \end{pmatrix} \times \begin{pmatrix} 1 & a_2 & * & * \\ 0 & 1 & b_2 & * \\ 0 & 0 & 1 & c_2 \\ 0 & 0 & 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & a_1 + a_2 & * & * \\ 0 & 1 & b_1 + b_2 & * \\ 0 & 0 & 1 & c_1 + c_2 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

$[G, G]$  itself is also abelian:

$$\begin{pmatrix} 1 & 0 & d_1 & f_1 \\ 0 & 1 & 0 & e_1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \times \begin{pmatrix} 1 & 0 & d_2 & f_2 \\ 0 & 1 & 0 & e_2 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 & d_1 + d_2 & f_1 + f_2 \\ 0 & 1 & 0 & e_1 + e_2 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$



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Suppose  $S = \{A_1, A_2, A_3, A_4\}$ ,

$$A_1 = \begin{pmatrix} 1 & 1 & 2 & 2 \\ 0 & 1 & 1 & 3 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \end{pmatrix}, A_2 = \begin{pmatrix} 1 & -1 & 4 & -2 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix},$$

$$A_3 = \begin{pmatrix} 1 & 0 & -2 & 0 \\ 0 & 1 & -1 & 3 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}, A_4 = \begin{pmatrix} 1 & 0 & 7 & 5 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & -1 \\ 0 & 0 & 0 & 1 \end{pmatrix},$$

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$S$  generates  $G$  as a group. What is  $\langle S[G, G] \rangle_{\text{semigrp}}$ ?

$$S[G, G] = \{A_1[G, G], \dots, A_4[G, G]\} = \left\{ \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}, \begin{pmatrix} -1 \\ 0 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ -1 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 0 \\ -1 \end{pmatrix} \right\}.$$

So indeed  $\langle S[G, G] \rangle_{\text{semigrp}} = \mathbb{Z}^3 = G/[G, G]$ .

## Identity Problem in $UT(4, \mathbb{Z})$ : Example

Since  $\langle S[G, G] \rangle_{\text{semigrp}} = G/[G, G]$ , the proposition claims that  $\langle S \rangle_{\text{semigrp}} = G$ .  
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We will prove  $\langle S \rangle_{\text{semigrp}} = G$  for this example.

Since  $\langle S[G, G] \rangle_{\text{semigrp}} = G/[G, G]$ , we have  $\langle S \rangle_{\text{semigrp}} \cap [G, G] \neq \emptyset$ .

$$A_1 A_2 A_3 A_4 = \begin{pmatrix} 1 & 0 & 11 & 2 \\ 0 & 1 & 0 & 8 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \in \langle S \rangle \cap [G, G],$$

$$A_2^{100} A_3^{100} A_1^{100} A_4^{100} = \begin{pmatrix} 1 & 0 & 6050 & 77350 \\ 0 & 1 & 0 & -4250 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \in \langle S \rangle \cap [G, G],$$

$$A_2^{100} A_1^{100} A_3^{100} A_4^{100} = \begin{pmatrix} 1 & 0 & -3950 & 127350 \\ 0 & 1 & 0 & 5750 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \in \langle S \rangle \cap [G, G],$$

$$A_4^{100} A_3^{100} A_2^{100} A_1^{100} = \begin{pmatrix} 1 & 0 & -3950 & -287650 \\ 0 & 1 & 0 & -4250 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \in \langle S \rangle \cap [G, G].$$

$$\begin{pmatrix} 1 & 0 & 11 & 2 \\ 0 & 1 & 0 & 8 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}^{1880000} \begin{pmatrix} 1 & 0 & 6050 & 77350 \\ 0 & 1 & 0 & -4250 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}^{14443} \times \\
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 = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \quad (1)$$

Therefore

$$(A_1 A_2 A_3 A_4)^{1880000} (A_2^{100} A_3^{100} A_1^{100} A_4^{100})^{14443} \dots (A_4^{100} A_3^{100} A_2^{100} A_1^{100})^{11096} = I.$$

Consequently,  $A_1^{-1}, \dots, A_4^{-1} \in \langle A_1, A_2, A_3, A_4 \rangle_{\text{semigroup}}$ . i.e.  $\langle S \rangle_{\text{semigrp}} = G$ .

## The general case

In general, we might not be able to find **positive** powers (1880000, 14443 etc.) to cancel elements in  $\langle S \rangle \cap [G, G]$ . Therefore, we need to prove that such cancellations always exist.

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$$\log : A \mapsto \sum_{k=1}^n \frac{(-1)^{k-1}}{k} (A - I)^k.$$

In particular,  $\log I = 0$ .



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### Theorem (Baker-Campbell-Hausdorff formula)

Let  $B_1, \dots, B_m \in \text{UT}(n, \mathbb{Z})$ , then

$$\log(B_1 \cdots B_m) = \sum_{i=1}^m \log B_i + \sum_{k=2}^d H_k(\log B_1, \dots, \log B_m),$$

where  $H_k, k = 2, 3, \dots$ , are expressions with explicitly computable forms.

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To find cancellation for elements of  $\langle S \rangle \cap [G, G]$ : cancel  $H_2, H_3, \dots, H_d$  one by one.

We use computer algebra software to find the explicit cancellation pattern for  $H_k, k = 2, 3, \dots$ . (These are **fixed** expressions!)

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*For each  $d > 10$ , subject to a conjecture  $P_d$ , the Identity Problem is decidable (and PTIME) in all finitely generated nilpotent groups of class- $d$ .*

*For each  $d$ , the conjecture  $P_d$  can be verified by computer algebra software in case it is true.*

In particular, the conjecture  $P_d$  concerns the existence of cancellation for the term  $H_d$ , which we were able to verify up to  $d \leq 10$ .