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;
; params.asm
;
; Fall, 2019
; An example to illustrate passing parameters
via
; the stack.
; Adapted from Mr. Jason Corless's code
; Modified by Victoria Li

; The stack pointer is in I/O space, so if we
; want to use LDS and STS instructions we have
to
; use the alternate addresses.
;
; SPH is 0x3E if using IN/OUT and 0x5E if
using LDS/STS
; SPL is 0x3D if using IN/OUT and 0x5D if
using LDS/STS
;
;.equ SPH=0x5E
;.equ SPL=0x5D

;
; The Z register is the combination of R31:R30
; since ZH, ZL are defined in "m2560def.inc",
we don't need redefine them here.
;.def ZH=r31
;.def ZL=r30
.def temp=r16
.def n1=r0
.def n2=r1
.def sumH=r19
.def sumL=r18

;to read "m2560def.inc", in "Solution
Explorer" -> under project name, in this case,
"project6"
;-> "Dependencies" ->"m2560def.inc"
;in file "m2560def.inc", SPL/H are defined as
following
; ****      .equ   SPL      = 0x3d
; ****      .equ   SPH      = 0x3e
; therefore, must use IN/OUT for read and
store

                ; initialize the stack pointer
(SP), so that SP points to 0x21FF
                ldi temp, low(RAMEND) ;.equ
RAMEND = 0x21ff <- defined in line 1747 of
file "m2560def.inc"
                out SPL, temp
                ldi temp, high(RAMEND)
                out SPH, temp

                ; call the subroutine
                ; Note that it is the caller's
responsibility

                ; to push the parameters on
the stack before

                ; the call and pop the
parameters from the

                ; stack after the call
                ;
                ; push the first parameter
ldi temp, 0xEE
push temp
                ; push the second parameter
ldi temp, 0xCC
push temp
call add_num
                ; now that the subroutine has
returned

                ; pop the parameters we
previously pushed. Why? To restore the stack
to the state before the call
pop temp
pop temp
                ; At this point, the stack is
empty, which

                ; is what we want.

done:           jmp done

; note: add two 8-bit numbers the result may
be 9 bits if there is a carry,
; therefore, the sum is store in register pair
sumH:sumL - r19:r18
; add_num sumH:sumL=n1+n2  0x1BA=0xEE + 0xCC
;
; This subroutine demonstrates using the stack
to pass
; parameters and using the stack to 'protect'
registers
; that are used in the subroutine.
;
; By protecting registers this subroutine
uses, the callers
; are free to use any registers.
;
;
; After the Z register is set to be the stack
pointer, the
; stack frame looks like:
;
;
; |         | <- Z and SP
; | n2      | saved register <- register n2 (r1)
is going to be used in the subroutine,
preserve its value on stack
; | n1      | saved register <- register n1 (r0)
is going to be used in the subroutine,
preserve its value on stack
; | ZH      | saved register
; | ZL      | saved register
; | ret     | return address
; | ret     | return address
; | ret     | return address
; | 0xCC    | parameter n2 (Z + 8)
; | 0xEE    | parameter n1 (Z + 9)
;
add_num:
                ; first protect the Z
register, since we will use it
push ZL
push ZH
                ; now protect r0 and r1 since
they will be used

                ; to store the parameters
push n1 ;r0
push n2 ;r1

                ; load the value in stack
pointer into the Z register
in ZH, SPH
in ZL, SPL

                ; get the 1st parameter pushed
on the stack:
ldd n1, Z+9
                ; get the 2nd parameter pushed
on the stack:
ldd n2, Z+8

                ; sumH:sumL=n1+n2
clr sumH
mov sumL, n1
add sumL, n2
rol sumH ;carry bit is brought
into position 0 of sumH

add_num_end:   ; This is where we return from
the subroutine

                ; restore the registers
protected on entry
                ; into the subroutine
pop n2
pop n1
pop ZH
pop ZL
ret

```