

CSC 230

# Numeration

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# Outline

- Number systems
- Conversion between
- Binary arithmetic
- Computer logic
- Negative numbers
- Horner's Algorithm
- Endianness

- Numeration - The action or process of calculating or assigning a number to something (Oxford dictionary).
- Arabic numbers (digits)
- Abacus (oldest calculator)
- Binary numbers (bits)

# Integer Number Systems

## Binary

- Base: 2
- Digits: 0,1

## Octal

- Base: 8
- Digits: 0,1,2,3,4,5,6,7

## Decimal

- Base: 10
- Digits: 0,1,2,3,4,5,6,7,8,9

## Hexadecimal

- Base: 16
- Digits: 0,1,2,3,4,5,6,7,8,9,A,B,C,D,E,F

# Positional Representation

Base 10

$$\begin{array}{r} + \quad 7 \\ + \quad 80 \\ + \quad 500 \\ + \quad 3000 \\ \hline = \quad 3587 \end{array}$$

# Weighted Positional Representation

$$\text{Integer Value} = \sum_{i=0}^{n-1} d_i * b_i$$

$$\text{Decimal Value} = \sum_{i=-m}^{n-1} d_i * b_i$$

e.g.  $312.98 = 3 * 10^2 + 1 * 10^1 + 2 * 10^0 + 9 * 10^{-1} + 8 * 10^{-2}$

# Polynomial representation

## Base 10

$$\begin{aligned} 3 * 1000 + 5 * 100 + 8 * 10 + 7 &= \\ 3 * 10^3 + 5 * 10^2 + 8 * 10^1 + 7 * 10^0 \end{aligned}$$

## General Form

$$d_n * b^n + d_{n-1} * b^{n-1} + \dots + d_2 * b^2 + d_1 * b^1 + d_0 * b^0$$

where  $d_n$  is the digit in  $n^{th}$  position starting from the right  
and  $b$  is the base.

# Notation

Binary: **0b** or **%** prefix

e.g. 0b11011010 or %11011010

Octal: **0o** (zero-oh) or **0** prefix

e.g. 0o617 or 0617

Hexadecimal (hex): **0x** (zero-oh) or **\$** prefix

e.g. 0x1F or \$1F

General: use a subscript with base number

e.g.  $101_2 \neq 101_8 \neq 101_{10} \neq 101_{16}$



# Memorize this table! (not)

Decimal	Binary	Hex	Octal
0	0000	0	0
1	0001	1	1
2	0010	2	2
3	0011	3	3
4	0100	4	4
5	0101	5	5
6	0110	6	6
7	0111	7	7
8	1000	8	
9	1001	9	
10	1010	A	
11	1011	B	
12	1100	C	
13	1101	D	
14	1110	E	
15	1111	F	

# Conversion shortcut: binary $\rightarrow$ hex or octal

Group the bits:

Binary:

$$(11010110\ 10110101)_2 = 54965_{10}$$

Octal:

$$(1\ 101\ 011\ 010\ 110\ 101)_2 = 153265_8$$

Hex:

$$(1101\ 0110\ 1011\ 0101)_2 = D6B5_{16}$$

# Conversion: decimal $\rightarrow$ any base

Example: convert  $119_{10}$  to octal.

Repeated division by base:

$$119/8 = 14 * 8 + 7 \quad r = 7 \quad (\text{right most digit})$$

$$14/8 = 1 * 8 + 6 \quad r = 6 \quad (\text{second last})$$

$$1/8 = 0 * 8 + 1 \quad r = 1 \quad ()$$

stop at 0

$$119_{10} = 167_8$$

# Conversion review

Decimal  $\rightarrow$  binary, octal, hex

Binary  $\rightarrow$  decimal, octal, hex

Octal  $\rightarrow$  decimal, binary, hex

Hex  $\rightarrow$  decimal, binary, octal

# Binary Arithmetic

# Binary Addition

	1	1	1	1	1	1	1		← carry
		1	0	1	1	1	1	1	← input 1
+		0	1	0	1	0	1	0	← input 2
<hr/>									
	1	0	0	0	1	0	0	1	← result

Traditional vs long carry methods

# Binary Subtraction

	*							← borrow
	1	0	1	1	1	1	1	← input 1
-	0	1	0	1	0	1	0	← input 2
<hr/>								
	0	1	1	0	1	0	1	← result

# Binary Multiplication

					1	0	1	0	← input 1
*					1	1	0	1	← input 2
<hr/>									
					1	0	1	0	
+				0	0	0	0		
+			1	0	1	0			
+		1	0	1	0				
<hr/>									
	1	0	0	0	0	0	1	0	← result



# Binary Division

$$\begin{array}{r}
 \phantom{110}111 \leftarrow \text{quotient} \\
 \hline
 110 \, ) \, 101101 \leftarrow \text{divisor } ) \text{ dividend} \\
 \phantom{110}110 \\
 \phantom{110}- \\
 \phantom{110}\phantom{1}010 \\
 \phantom{110}\phantom{1}110 \\
 \phantom{110}\phantom{1}- \\
 \phantom{110}\phantom{1}\phantom{1}001 \\
 \phantom{110}\phantom{1}\phantom{1}110 \\
 \phantom{110}\phantom{1}\phantom{1}- \\
 \phantom{110}\phantom{1}\phantom{1}\phantom{1}110 \\
 \phantom{110}\phantom{1}\phantom{1}\phantom{1}110 \\
 \phantom{110}\phantom{1}\phantom{1}\phantom{1}- \\
 \hline
 \phantom{110}\phantom{1}\phantom{1}\phantom{1}11 \leftarrow \text{remainder}
 \end{array}$$

# Computer (Boolean) Logic

# Logic Gates Notation

NOT

$(\bar{A})$  or  $(\neg A)$  or  $(\sim A)$  or  $(!A)$  or  $(A')$

AND

$(A \cdot B)$  or  $(A \wedge B)$  or  $(A \& B)$

OR

$(A + B)$  or  $(A \vee B)$  or  $(A \parallel B)$

XOR

$(A \oplus B)$  or  $(A \underline{\vee} B)$

# Truth Table

A	B	NOT		AND	OR	XOR	NAND	NOR
0	0	1	1	0	0	0	1	1
0	1	1	0	0	1	1	1	0
1	0	0	1	0	1	1	1	0
1	1	0	0	1	1	0	0	0

AND  $\rightarrow$  1 when both are one, 0 otherwise

OR  $\rightarrow$  0 when both are 0, 1 otherwise

XOR  $\rightarrow$  0 when both are the same, 1 otherwise

# Bitwise Operations

## Rotating

## Shifting

<< left shift

- multiply by 2
- overflow can occur

>> right shift

- divide by 2
- underflow may occur

## Masking

**input:** string of bits

**mask:** string of bits (same size as input)

binary Boolean **operator:**

- AND - 1 preserves, 0 clears
- OR - 0 preserves, 1 sets
- XOR - 0 preserves, 1 toggle

# Negative Numbers

# Signed Magnitude

Binary	Decimal
0000	0
0001	1
0010	2
0011	3
0100	4
0101	5
0110	6
0111	7
1000	-0
1001	-1
1010	-2
1011	-3
1100	-4
1101	-5
1110	-6
1111	-7

# 1's Complement

Binary	Decimal
0000	0
0001	1
0010	2
0011	3
0100	4
0101	5
0110	6
0111	7
1000	-7
1001	-6
1010	-5
1011	-4
1100	-3
1101	-2
1110	-1
1111	-0



# 2's Complement

Binary	Decimal
0000	0
0001	1
0010	2
0011	3
0100	4
0101	5
0110	6
0111	7
1000	-8
1001	-7
1010	-6
1011	-5
1100	-4
1101	-3
1110	-2
1111	-1

# 2's Complement Algorithm

Convert negative value to binary:

1. Convert the absolute value to binary
2. Complement (flip all the bits)
3. Add 1 (in binary)

Convert a **negative** binary number to decimal:

1. Subtract 1 (in binary)
2. Complement (flip all the bits)
3. Convert as positive binary integer
4. Prepend a minus

## 2's Complement Advantages

- efficient
- uniform
- avoids two zeros
- only need one adder (no subtraction unit)

In Addition...

# Endianness

Byte-addressable memory:

Byte order	Big-endian	23-bit integer	Little-endian
		0A1B2C3D	
1	0A		3D
2	1B		2C
3	2C		1B
4	3D		0A

Word-addressable memory:

Byte order	Big-endian	23-bit integer	Little-endian
		0A1B2C3D	
1	0A1B		2C3D
2	2C3D		0A1B

# Horner's Method

In 1819, William George Horner developed a method for evaluation of a polynomial of degree  $n$  using at most  $n$  multiplications and  $n$  additions.

$$\begin{aligned} & d_n * b^n + d_{n-1} * b^{n-1} + \dots + d_2 * b^2 + d_1 * b^1 + d_0 * b^0 \\ &= ((\dots(d_n * b + d_{n-1}) * b + \dots + d_2) * b + d_1) * b + d_0 \end{aligned}$$

Note: evaluation starts from inner-most parentheses.

Allows for fast multiplication and division w/o hardware multiplier.