Part IV

Phases of Database Design



- Phases of Database Design
- Further Steps During Design



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- Further Steps During Design
- 3 Capacity-preserving Transformations



- Phases of Database Design
- Further Steps During Design
- Capacity-preserving Transformations
- ER-to-RM Transformation



- Phases of Database Design
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- Capacity-preserving Transformations
- ER-to-RM Transformation
- Transformation of Relationships

- Phases of Database Design
- Further Steps During Design
- Capacity-preserving Transformations
- **ER-to-RM Transformation**
- Transformation of Relationships
- **Special Cases for Transformation**



- Phases of Database Design
- Purther Steps During Design
- Capacity-preserving Transformations
- ER-to-RM Transformation
- Transformation of Relationships
- 6 Special Cases for Transformation
- Overview and Summary



Educational Objective for Today . . .

Goals and steps of the database design process



Educational Objective for Today ...

- Goals and steps of the database design process
- Rules to transform ER schemata into relational schemata



Phases of Database Design

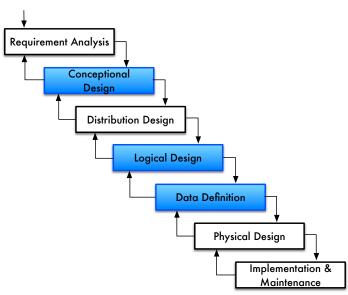


Goal of Design

- Data management for multiple application systems, for multiple years
- Therefore: special importance
- Design requirements
 - For each application, it should be possible to derive application data from stored data in the database — efficiently!
 - Only store "sensible" (actually needed) data
 - Avoid redundancies



Phase Model



Requirements Analysis

- Approach: Collecting information needs from all specialist divisions
- Result:
 - Informal description (text, tabular lists, forms, etc.) of the problem domain
 - Separation of the information about data (data analysis) from the information about functions (functional analysis)
- "Classical" DB design:
 - Only data analysis and following steps
- Functional design:
 - See methods of software engineering

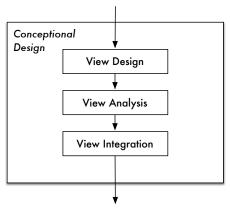


Conceptual Design

- First formal description of the problem domain
- Language means: semantical data model
- Process:
 - Modeling of views, e.g., for different specialist divisions
 - Analysis of existing views with respect to conflicts
 - Integration of views into a full schema
- Result: full conceptual schema, e.g., ER diagram



Phases of Conceptual Design



Further Steps During Design

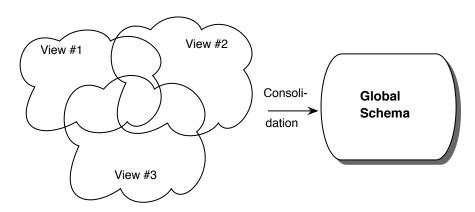
Further Steps During Design

- ER modeling of different views of the complete information, e.g., for different specialist divisions of a company ~> conceptual design
 - Analysis and integration of views
 - Result: full conceptual schema
- Distribution design when using distributed storage
- Transformation to concrete implementation model (e.g., relational model) → logical design
- Data definition, implementation and maintenance → physical design



View Integration

- Analysis of existing view with respect to conflicts
- Integration of views into a full schema



Integration Conflicts

- Naming conflicts: Homonyms / synonyms
 - Homonyms: bank (money / river); order (command / request for goods)
 - Synonyms: car, vehicle, automobile
- Typing conflicts: different structures for the same element
- Domain mismatch: different domains for an element
- Identifier conflicts: e.g., different keys for the same element
- Structural conflicts: same fact expressed in different ways



Distribution Design

- If data should be distributed to several machines, a way of distributed storage must be determined
- E.g., for a relation
 CUSTOMER (CNo, Name, Address, Zipcode, Account)
 - Horizontal distribution:
 CUSTOMER_1 (CNo, Name, Address, Zipcode, Account)
 where Zipcode < 50000
 CUSTOMER_2 (CNo, Name, Address, Zipcode, Account)
 where Zipcode >= 50000
 - Vertical distribution (connection via attribute CNo): CUSTOMER_Adr (CNo, Name, Address, Zipcode) CUSTOMER_Account (CNo, Account)

Logical Design

- Language means: Data model of the chosen "implementation"
 DBMS, e.g., relational model
- Process:
 - $\begin{tabular}{ll} \hline \textbf{0} & (Automatical) transformation of the conceptual schema, e.g., \\ ER & \rightarrow relational model \\ \hline \end{tabular}$
 - Improvement of the relational schema based on quality criteria (normalization, see Chapter 5): Design goals: avoid redundancies, . . .
- Result: logical schema, e.g., collection of relation schemata

Data Definition

- Translation of logical schema into a concrete schema
- Language means: DDL and DML of DBMS (e.g., Oracle, DB2, SQL Server)
 - Database declaration in the DDL of the DBMS
 - Realization of integrity constraints
 - Definition of views

Physical Design

- Supplement physical design with support for efficient access, e.g., by defining indexes
- Index
 - Access path: data structure for additional, key-based access to tuples (\(\langle key attribute value, tuple address \rangle\)
 - Usually implemented as a B*-tree
- Language means: storage structure (definition) language SSL

Indexes in SQL

```
create [ unique ] index indexname
  on relname (
     attrname [ asc | desc ],
     attrname [ asc | desc ],
     ...
)
```

Example

```
create index WineIdx on WINES (Name)
```

Necessity of Access Paths

- Example: Table with 100 GB of data, hard disk transfer rate of ca.
 50 MB/s
- Operation: Search for a tuple (selection)
- Implementation: Sequential search
- Cost: 102.400/50 = 2.048 sec. ≈ 34 min.

Implementation and Maintenance

- Phases of ...
 - Maintenance.
 - Further optimization of the physical layer,
 - Adaptation to new requirements or operating system platforms,
 - Porting to new database management systems
 - etc.

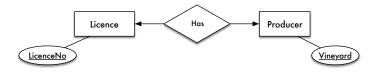
Capacity-preserving Transformations

Transformation of the Conceptual Schema

- Translation to logical schema
 - Example: ER → RM
 - Correct?
 - Quality of transformation?
- Preservation of information capacity
 - Is it possible, after the transformation, to store exactly the same data as before?
 - ... or more?
 - ... or less?



Capacity-increasing Transformation



Transformation into

$$R = \{ \texttt{LicenseNo}, \texttt{Vineyard} \}$$

with exactly one key

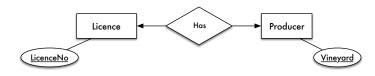
$$K = \{\{\text{LicenseNo}\}\}$$

Possible invalid relation:

Has	LicenseNo	Vineyard
	007	Helena
	42	Helena



Capacity-preserving Transformation



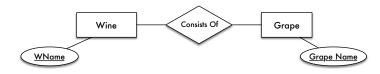
Correct instantiation

Has	LicenseNo	Vineyard
	007	Helena
	42	Mueller

Correct set of keys

$$K = \{\{\text{LicenseNo}\}, \{\text{Vineyard}\}\}$$

Capacity-decreasing Transformation



- Relation schema with one key {WName}
- Instantiation that is no longer valid:

ConsistsOf	WName	GrapeName
	Zinfandel Red Blossom	Zinfandel
	Bordeaux Blanc	Cabernet Sauvignon
	Bordeaux Blanc	Muscadelle

 Capacity-preserving when using the keys of both entity types as the new key in the relation schema

$$K = \{\{\text{WName}, \text{GrapeName}\}\}$$

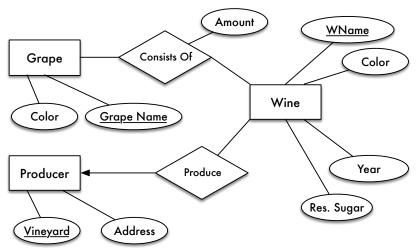


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ER-to-RM Transformation



Example Transformation ER-RM: Input



Example Transformation ER-RM: Result

- $\textbf{0} \ \, \mathsf{GRAPE} = \{\mathsf{Color}, \mathsf{GrapeName}\} \ \, \mathsf{with} \ \, K_{\mathsf{GRAPE}} = \{\{\mathsf{GrapeName}\}\}$
- $\textbf{@} \label{eq:consists0f} \textbf{ Consists0f} = \{ \textbf{GrapeName}, \textbf{WName}, \textbf{Amount} \} \ \ \, \\ K_{\textbf{Consists0f}} = \{ \{ \textbf{GrapeName}, \textbf{WName} \} \}$
- **②** PRODUCE = $\{WName, Vineyard\}$ with $K_{PRODUCE} = \{\{WName\}\}$
- **1** PRODUCER = $\{Vineyard, Address\}$ with $K_{PRODUCER} = \{\{Vineyard\}\}$

ER Transformation into Relations

- Entity types and relationship types: both transformed into relation schemata
- Attributes: attributes of the relation schema, keys are adopted
- Cardinalities of the relationships: expressed in respective relation schemata by choice of keys
- In some cases: merge of the relation schemata of entity and relationship types
- Introduce foreign key constraints between the remaining relation schemata



Transformation of Relationships



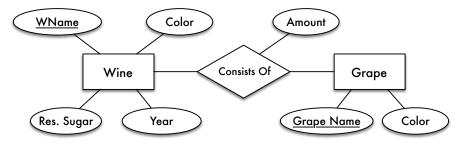
Transformation of Relationship Types

New relation schema with all attributes of the relationship type;
 additionally, adopt all primary keys of the participating entity types

Determining keys:

- m:n relationship: both primary keys together form the key in the new relation schema
- 1:n relationship: primary keys of the n-side (in the functional notation, this is the side without the arrowhead) form key in the new relation schema
- ▶ 1:1 relationship: both primary keys become a key in the new relation schema; the primary key is then chosen from these keys

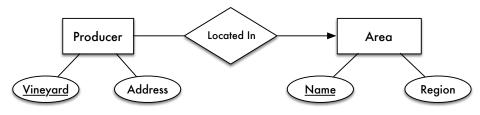
n:m Relationships



- Transformation
 - $\textbf{0} \; \; \mathsf{GRAPE} = \{\mathsf{Color}, \mathsf{GrapeName}\} \; \mathsf{with} \; K_{\mathsf{GRAPE}} = \{\{\mathsf{GrapeName}\}\}$
 - ConsistsOf = {GrapeName, WName, Amount} with
 KConsistsOf = {{GrapeName, WName}}
- Attributes GrapeName and WName together are key



1:n Relationships

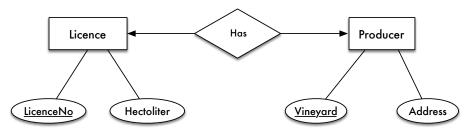


- (Preliminary) transformation
 - PRODUCER with the attributes Vineyard and Address,
 - AREA with the attributes Name and Region, and
 - LocatedIn with the attributes Vineyard and Name and the primary key of the n-side Vineyard as primary key of this schema.

Possible Merges

- Optional relationships ([0,1] or [0,n]) are not merged
- With cardinalities [1,1] or [1,n] (mandatory relationships), merge is possible:
 - ▶ 1:n relationship: the entity-relation schema of the n-side can be integrated into the relation schema of the relationship
 - ▶ 1:1 relationship: both entity-relation schemata can be integrated into the relation schema of the relationship

1:1 Relationships



- (Preliminary) transformation
 - PRODUCER with the attributes Vineyard and Address
 - LICENSE with the two attributes LicenseNo and Hectoliters
 - Has with the primary keys of both participating entity types each as key of this schema, that is LicenseNo and Vineyard

1:1 Relationships: Merge

- Transformation with merge
 - Merged relation:

PRODUCER	Vineyard	Address	LicenseNo	Hectoliters
	Rotkäppchen	Freiberg	42-007	10 000
	Vineyard Müller	Dagstuhl	42-009	250

Producers without license require null values:

PRODUCER	Vineyard	Address	LicenseNo	Hectoliters
	Rotkäppchen	Freiberg	42-007	10 000
	Vineyard Müller	Dagstuhl	上	上

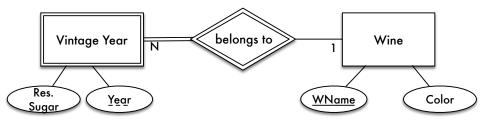
Free Licenses lead to additional null values:

PRODUCER	Vineyard	Address	LicenseNo	Hectoliters
	Rotkäppchen	Freiberg	42-007	10 000
	Vineyard Müller	Dagstuhl	上	_
			42-003	100 000

Special Cases for Transformation



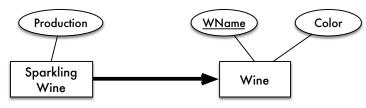
Dependent Entity Types



- Transformation

 - ② WINE = $\{\text{Color}, \text{WName}\}\$ with $K_{\text{WINE}}=\{\{\text{WName}\}\}$
 - Attribute WName in WINEVINTAGE is foreign key to relation WINE

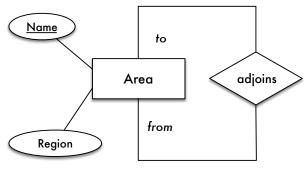
IS-A Relationship



- Transformation
 - WINE = {Color, WName} with $K_{WINE} = \{\{WName\}\}\$
 - SPARKLING_WINE = {WName, Production} with

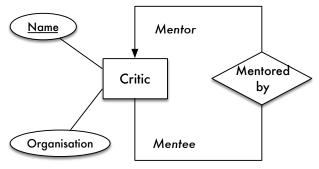
 KSPARKLING_WINE = {{WName}}
 - ▶ WName in SPARKLING_WINE is foreign key with respect to relation WINE

Recursive Relationships



- Transformation
 - AREA = {Name, Region} with $K_{ARFA} = \{\{Name\}\}\$
 - ② ADJOINS = $\{to, from\}$ with $K_{ADJOINS} = \{\{to, from\}\}$

Recursive Functional Relationships

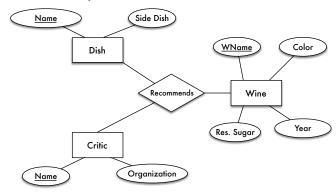


- Transformation

 - ▶ Mentorname is foreign key to attribute Name of relation CRITIC.



N-ary Relationships



- Every participating entity type is treated according to the rules stated above
- For relationship Recommends, the primary keys of the three participating entity types are included in the resulting relation schema
- Relationship has a generic type (k:m:n relationship): all primary keys together form the key



N-ary Relationships: Result

- RECOMMENDS = {WName, DName, Name} with $K_{RECOMMENDS} = \{\{WName, DName, Name\}\}$
- ② DISH = $\{DName, Side_Dish\}$ with $K_{DISH} = \{\{DName\}\}$
- CRITIC = {Name, Organization} with $K_{CRITIC} = \{\{\text{Name}\}\}$
 - The three key attributes of RECOMMENDS are foreign keys to the respective source relations (CRITIC, WINE, DISH).

Overview and Summary



Overview of Transformations

ER Concept	Is Translated into Relational Concept	
Entity type E_i	Relation schema R_i	
Attributes of E_i	Attributes of R_i	
Primary key P_i	Primary key P_i	
Relationship type	Relation schema	
	Attributes: P_1 , P_2	
Its attributes	Further attributes	
1:n	P ₂ becomes primary key of the relationship	
1:1	P_1 and P_2 become key of the relationship	
m:n	$P_1 \cup P_2$ becomes primary key of the relationship	
IS-A relationship	R_1 gets an additional key P_2	

 E_1 , E_2 : Entity types participating in a relationship,

 P_1, P_2 : Their primary keys,

1: n relationship: E_2 is n-side,

IS-A relationship: E_1 is a special entity type

Summary

Phases of database design



Summary

- Phases of database design
- Capacity-preserving transformations

Summary

- Phases of database design
- Capacity-preserving transformations
- Transformation ER → relational

 Which steps does the database design process comprise?



- Which steps does the database design process comprise?
- Which requirements do the transformations between each design step have to fulfill? Why?



- Which steps does the database design process comprise?
- Which requirements do the transformations between each design step have to fulfill? Why?
- How are concepts of the ER model translated into concepts of the relational model?



- Which steps does the database design process comprise?
- Which requirements do the transformations between each design step have to fulfill? Why?
- How are concepts of the ER model translated into concepts of the relational model?
- How are the different cardinalities of relationship types accounted for during transformation?

